Synchronization: Basics

15-213: Introduction to Computer Systems
24th Lecture, April 17, 2018

Instructor:
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Today

- Threads review
- Sharing
- Mutual exclusion
- Semaphores
Traditional View of a Process

- Process = process context + code, data, and stack

**Process context**

- **Program context:**
  - Data registers
  - Condition codes
  - Stack pointer (SP)
  - Program counter (PC)

- **Kernel context:**
  - VM structures
  - Descriptor table
  - brk pointer

**Code, data, and stack**

- Stack
  - Shared libraries
  - Run-time heap
  - Read/write data
  - Read-only code/data

PC → 0
Alternate View of a Process

- Process = thread + code, data, and kernel context

**Thread (main thread)**
- Stack
- Thread context:
  - Data registers
  - Condition codes
  - Stack pointer (SP)
  - Program counter (PC)

**Code, data, and kernel context**
- Shared libraries
- Run-time heap
- Read/write data
- Read-only code/data
- Kernel context:
  - VM structures
  - Descriptor table
  - brk pointer
A Process With Multiple Threads

Multiple threads can be associated with a process

- Each thread has its own logical control flow
- Each thread shares the same code, data, and kernel context
- Each thread has its own stack for local variables
  - but not protected from other threads
- Each thread has its own thread id (TID)

Thread 1 (main thread)  Thread 2 (peer thread)

<table>
<thead>
<tr>
<th>Thread 1 context:</th>
<th>Thread 2 context:</th>
</tr>
</thead>
<tbody>
<tr>
<td>Data registers</td>
<td>Data registers</td>
</tr>
<tr>
<td>Condition codes</td>
<td>Condition codes</td>
</tr>
<tr>
<td>$SP_1$</td>
<td>$SP_2$</td>
</tr>
<tr>
<td>$PC_1$</td>
<td>$PC_2$</td>
</tr>
</tbody>
</table>

Shared code and data

- shared libraries
- run-time heap
- read/write data
- read-only code/data
- Kernel context:
  - VM structures
  - Descriptor table
  - brk pointer
Shared Variables in Threaded C Programs

Question: Which variables in a threaded C program are shared?
- The answer is not as simple as “global variables are shared” and “stack variables are private”

Def: A variable $x$ is shared if and only if multiple threads reference some instance of $x$.

Requires answers to the following questions:
- What is the memory model for threads?
- How are instances of variables mapped to memory?
- How many threads might reference each of these instances?
Threads Memory Model: Conceptual

- Multiple threads run within the context of a single process
- Each thread has its own separate thread context
  - Thread ID, stack, stack pointer, PC, condition codes, and GP registers
- All threads share the remaining process context
  - Code, data, heap, and shared library segments of the process virtual address space
  - Open files and installed handlers
Threads Memory Model: Actual

- Separation of data is not strictly enforced:
  - Register values are truly separate and protected, but...
  - Any thread can read and write the stack of any other thread

The mismatch between the conceptual and operation model is a source of confusion and errors
Example Program to Illustrate Sharing

```c
char **ptr; /* global var */

int main(int argc, char *argv[]) {
    long i;
    pthread_t tid;
    char *msgs[2] = {
        "Hello from foo",
        "Hello from bar"
    };

    ptr = msgs;
    for (i = 0; i < 2; i++)
        Pthread_create(&tid,
                        NULL,
                        thread,
                        (void *)i);

    Pthread_exit(NULL);
}

void *thread(void *vargp) {
    long myid = (long)vargp;
    static int cnt = 0;

    printf("[%ld]:  %s (cnt=%d)\n", myid, ptr[myid], ++cnt);
    return NULL;
}

Peer threads reference main thread’s stack indirectly through global ptr variable

A common, but inelegant way to pass a single argument to a thread routine
```
Mapping Variable Instances to Memory

- **Global variables**
  - *Def:* Variable declared outside of a function
  - Virtual memory contains exactly one instance of any global variable

- **Local variables**
  - *Def:* Variable declared inside function without `static` attribute
  - Each thread stack contains one instance of each local variable

- **Local static variables**
  - *Def:* Variable declared inside function with the `static` attribute
  - Virtual memory contains exactly one instance of any local static variable.
Mapping Variable Instances to Memory

**Global var:** 1 instance (ptr [data])

```c
char **ptr; /* global var */
int main(int main, char *argv[])
{
    long i;
    pthread_t tid;
    char *msgs[2] = {
        "Hello from foo",
        "Hello from bar"
    };
    ptr = msgs;
    for (i = 0; i < 2; i++)
        Pthread_create(&tid, NULL, thread, (void *)i);
    Pthread_exit(NULL);
}
void *thread(void *vargp)
{
    long myid = (long)vargp;
    static int cnt = 0;
    printf("[%ld]: %s (cnt=%d)\n", myid, ptr[myid], ++cnt);
    return NULL;
}
```

**Local vars:** 1 instance (i.m, msgs.m, tid.m)

**Local var:** 2 instances (myid.p0 [peer thread 0’s stack], myid.p1 [peer thread 1’s stack])

**Local static var:** 1 instance (cnt [data])
Shared Variable Analysis

- Which variables are shared?

<table>
<thead>
<tr>
<th>Variable instance</th>
<th>Referenced by main thread?</th>
<th>Referenced by peer thread 0?</th>
<th>Referenced by peer thread 1?</th>
</tr>
</thead>
<tbody>
<tr>
<td>ptr</td>
<td>yes</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>cnt</td>
<td>no</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>i.m</td>
<td>yes</td>
<td>no</td>
<td>no</td>
</tr>
<tr>
<td>msgs.m</td>
<td>yes</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>myid.p0</td>
<td>no</td>
<td>yes</td>
<td>no</td>
</tr>
<tr>
<td>myid.p1</td>
<td>no</td>
<td>no</td>
<td>yes</td>
</tr>
</tbody>
</table>

char **ptr;  /* global var */
int main(int main, char *argv[]) {
    long i; pthread_t tid;
    char *msgs[2] = {"Hello from foo", "Hello from bar"};

    ptr = msgs;
    for (i = 0; i < 2; i++)
        Pthread_create(&tid,
                        NULL, thread,(void *)i);
    Pthread_exit(NULL);}

void *thread(void *vargp)
{
    long myid = (long)vargp;
    static int cnt = 0;

    printf("[%ld]:  %s (cnt=%d)\n", myid, ptr[myid], ++cnt);
    return NULL;
}
## Shared Variable Analysis

### Which variables are shared?

<table>
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<th>Referenced by main thread?</th>
<th>Referenced by peer thread 0?</th>
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<tbody>
<tr>
<td>ptr</td>
<td>yes</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>cnt</td>
<td>no</td>
<td>yes</td>
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<td>no</td>
<td>yes</td>
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</tr>
<tr>
<td>myid.p1</td>
<td>no</td>
<td>no</td>
<td>yes</td>
</tr>
</tbody>
</table>

### Answer: A variable \( x \) is shared iff multiple threads reference at least one instance of \( x \). Thus:

- \( \text{ptr, cnt, and msgs are shared} \)
- \( \text{i and myid are not shared} \)
Synchronizing Threads

- Shared variables are handy...

- ...but introduce the possibility of nasty synchronization errors.
badcnt.c: Improper Synchronization

/* Global shared variable */
volatile long cnt = 0; /* Counter */

int main(int argc, char **argv)
{
    long niters;
    pthread_t tid1, tid2;
    niters = atoi(argv[1]);
    Pthread_create(&tid1, NULL,
                   thread, &niters);
    Pthread_create(&tid2, NULL,
                   thread, &niters);
    Pthread_join(tid1, NULL);
    Pthread_join(tid2, NULL);

    /* Check result */
    if (cnt != (2 * niters))
        printf("BOOM! cnt=%ld\n", cnt);
    else
        printf("OK cnt=%ld\n", cnt);
    exit(0);
}

/* Thread routine */
void *thread(void *vargp)
{
    long i, niters =
        *((long *)vargp);

    for (i = 0; i < niters; i++)
        cnt++;
    return NULL;
}

linux> ./badcnt 10000
OK cnt=20000
linux> ./badcnt 10000
BOOM! cnt=13051
linux>

cnt should equal 20,000.

What went wrong?
Assembly Code for Counter Loop

C code for counter loop in thread i

```c
for (i = 0; i < niters; i++)
    cnt++;
```

Asm code for thread i

```
movq (%rdi), %rcx
testq %rcx,%rcx
jle .L2
movl $0, %eax
.L3:
    movq  cnt(%rip),%rdx
    addq  $1, %rdx
    movq  %rdx, cnt(%rip)
    addq  $1, %rax
    cmpq  %rcx, %rax
    jne   .L3
.L2:
```

\( H_i \): Head
\( L_i \): Load cnt
\( U_i \): Update cnt
\( S_i \): Store cnt
\( T_i \): Tail

Hi: Head
Ti: Tail
Li: Load cnt
Ui: Update cnt
Si: Store cnt
Concurrent Execution

**Key idea:** In general, any sequentially consistent interleaving is possible, but some give an unexpected result!

- $I_i$ denotes that thread $i$ executes instruction $I$
- $\%rdx_i$ is the content of $\%rdx$ in thread $i$’s context

<table>
<thead>
<tr>
<th>$i$ (thread)</th>
<th>$instr_i$</th>
<th>$%rdx_1$</th>
<th>$%rdx_2$</th>
<th>cnt</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>$H_1$</td>
<td>-</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>$L_1$</td>
<td>0</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>$U_1$</td>
<td>1</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>$S_1$</td>
<td>1</td>
<td>-</td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>$H_2$</td>
<td>-</td>
<td>-</td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>$L_2$</td>
<td>-</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>$U_2$</td>
<td>-</td>
<td>2</td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>$S_2$</td>
<td>-</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>2</td>
<td>$T_2$</td>
<td>-</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>1</td>
<td>$T_1$</td>
<td>1</td>
<td>-</td>
<td>2</td>
</tr>
</tbody>
</table>

**OK**
Concurrent Execution

**Key idea:** In general, any sequentially consistent interleaving is possible, but some give an unexpected result!

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<th>$\text{cnt}$</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>$H_1$</td>
<td>-</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>$L_1$</td>
<td>0</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>$U_1$</td>
<td>1</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>$S_1$</td>
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<td>-</td>
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</tr>
<tr>
<td>2</td>
<td>$H_2$</td>
<td>-</td>
<td>-</td>
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</tr>
<tr>
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<td>2</td>
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<td>-</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>2</td>
<td>$T_2$</td>
<td>-</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>1</td>
<td>$T_1$</td>
<td>1</td>
<td>-</td>
<td>2</td>
</tr>
</tbody>
</table>

Thread 1

**critical section**

Thread 2

**critical section**

*OK*
Concurrent Execution (cont)

- Incorrect ordering: two threads increment the counter, but the result is 1 instead of 2

<table>
<thead>
<tr>
<th>i (thread)</th>
<th>instr&lt;sub&gt;i&lt;/sub&gt;</th>
<th>%rdx&lt;sub&gt;1&lt;/sub&gt;</th>
<th>%rdx&lt;sub&gt;2&lt;/sub&gt;</th>
<th>cnt</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>H&lt;sub&gt;1&lt;/sub&gt;</td>
<td>-</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>L&lt;sub&gt;1&lt;/sub&gt;</td>
<td>0</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>U&lt;sub&gt;1&lt;/sub&gt;</td>
<td>1</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>2</td>
<td>H&lt;sub&gt;2&lt;/sub&gt;</td>
<td>-</td>
<td>-</td>
<td>0</td>
</tr>
<tr>
<td>2</td>
<td>L&lt;sub&gt;2&lt;/sub&gt;</td>
<td>-</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>S&lt;sub&gt;1&lt;/sub&gt;</td>
<td>1</td>
<td>-</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>T&lt;sub&gt;1&lt;/sub&gt;</td>
<td>1</td>
<td>-</td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>U&lt;sub&gt;2&lt;/sub&gt;</td>
<td>-</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>S&lt;sub&gt;2&lt;/sub&gt;</td>
<td>-</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>T&lt;sub&gt;2&lt;/sub&gt;</td>
<td>-</td>
<td>1</td>
<td>1</td>
</tr>
</tbody>
</table>

*Oops!*
Concurrent Execution (cont)

- How about this ordering?

<table>
<thead>
<tr>
<th>i (thread)</th>
<th>instr_i</th>
<th>%rdx_1</th>
<th>%rdx_2</th>
<th>cnt</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>H_1</td>
<td>0</td>
<td></td>
<td>0</td>
</tr>
<tr>
<td>1</td>
<td>L_1</td>
<td>0</td>
<td></td>
<td>0</td>
</tr>
<tr>
<td>2</td>
<td>H_2</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>L_2</td>
<td>0</td>
<td></td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>U_2</td>
<td>1</td>
<td></td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>S_2</td>
<td>1</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>U_1</td>
<td>1</td>
<td></td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>S_1</td>
<td>1</td>
<td></td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>T_1</td>
<td></td>
<td></td>
<td>1</td>
</tr>
<tr>
<td>2</td>
<td>T_2</td>
<td></td>
<td></td>
<td>1</td>
</tr>
</tbody>
</table>

Oops!

- We can analyze the behavior using a progress graph
A progress graph depicts the discrete execution state space of concurrent threads.

Each axis corresponds to the sequential order of instructions in a thread.

Each point corresponds to a possible execution state (Inst₁, Inst₂).

E.g., (L₁, S₂) denotes state where thread 1 has completed L₁ and thread 2 has completed S₂.
A trajectory is a sequence of legal state transitions that describes one possible concurrent execution of the threads.

Example:

H1, L1, U1, H2, L2, S1, T1, U2, S2, T2
A trajectory is a sequence of legal state transitions that describes one possible concurrent execution of the threads.

Example:

H1, L1, U1, H2, L2, S1, T1, U2, S2, T2
Critical Sections and Unsafe Regions

L, U, and S form a critical section with respect to the shared variable \( c \).

Instructions in critical sections (wrt some shared variable) should not be interleaved.

Sets of states where such interleaving occurs form unsafe regions.
Critical Sections and Unsafe Regions

Def: A trajectory is safe iff it does not enter any unsafe region

Claim: A trajectory is correct (wrt cnt) iff it is safe
badcnt.c: Improper Synchronization

```c
/* Global shared variable */
volatile long cnt = 0; /* Counter */

int main(int argc, char **argv)
{
    long niters;
    pthread_t tid1, tid2;

    niters = atoi(argv[1]);
    Pthread_create(&tid1, NULL,
                   thread, &niters);
    Pthread_create(&tid2, NULL,
                   thread, &niters);
    Pthread_join(tid1, NULL);
    Pthread_join(tid2, NULL);

    /* Check result */
    if (cnt != (2 * niters))
        printf("BOOM! cnt=%ld\n", cnt);
    else
        printf("OK cnt=%ld\n", cnt);
    exit(0);
}

/* Thread routine */
void *thread(void *vargp)
{
    long i, niters =
        *((long *)vargp);

    for (i = 0; i < niters; i++)
        cnt++;
    return NULL;
}
```

<table>
<thead>
<tr>
<th>Variable</th>
<th>main</th>
<th>thread1</th>
<th>thread2</th>
</tr>
</thead>
<tbody>
<tr>
<td>cnt</td>
<td>yes*</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>niters.m</td>
<td>yes</td>
<td>no</td>
<td>no</td>
</tr>
<tr>
<td>tid1.m</td>
<td>yes</td>
<td>no</td>
<td>no</td>
</tr>
<tr>
<td>i.1</td>
<td>no</td>
<td>yes</td>
<td>no</td>
</tr>
<tr>
<td>i.2</td>
<td>no</td>
<td>no</td>
<td>yes</td>
</tr>
<tr>
<td>niters.1</td>
<td>no</td>
<td>yes</td>
<td>no</td>
</tr>
<tr>
<td>niters.2</td>
<td>no</td>
<td>no</td>
<td>yes</td>
</tr>
</tbody>
</table>
Break Time!

**flimflam**: “nonsensical or insincere talk"

Check out:

Quiz: day 24: Synchronization Basic

https://canvas.cmu.edu/courses/3822
Enforcing Mutual Exclusion

- **Question:** How can we guarantee a safe trajectory?

- **Answer:** We must *synchronize* the execution of the threads so that they can never have an unsafe trajectory.
  - i.e., need to guarantee *mutually exclusive access* for each critical section.

- **Classic solution:**
  - Semaphores (Edsger Dijkstra)

- **Other approaches (out of our scope)**
  - Mutex and condition variables (Pthreads)
  - Monitors (Java)
Semaphores

- **Semaphore**: non-negative global integer synchronization variable. Manipulated by $P$ and $V$ operations.

- **$P(s)$**
  - If $s$ is nonzero, then decrement $s$ by 1 and return immediately.
    - Test and decrement operations occur atomically (indivisibly)
  - If $s$ is zero, then suspend thread until $s$ becomes nonzero and the thread is restarted by a $V$ operation.
  - After restarting, the $P$ operation decrements $s$ and returns control to the caller.

- **$V(s)$**:
  - Increment $s$ by 1.
    - Increment operation occurs atomically
  - If there are any threads blocked in a $P$ operation waiting for $s$ to become non-zero, then restart exactly one of those threads, which then completes its $P$ operation by decrementing $s$.

- **Semaphore invariant**: $(s \geq 0)$
Semaphores

- **Semaphore**: non-negative global integer synchronization variable

- **Manipulated by $P$ and $V$ operations:**
  - $P(s)$: `[ while (s == 0) wait(); s--; ]`
    - Dutch for "Proberen" (test)
  - $V(s)$: `[ s++; ]`
    - Dutch for "Verhogen" (increment)

- **OS kernel guarantees that operations between brackets [ ] are executed indivisibly**
  - Only one $P$ or $V$ operation at a time can modify $s$.
  - When `while` loop in $P$ terminates, only that $P$ can decrement $s$

- **Semaphore invariant**: $(s >= 0)$
C Semaphore Operations

Pthreads functions:

```c
#include <semaphore.h>

int sem_init(sem_t *s, 0, unsigned int val); /* s = val */
int sem_wait(sem_t *s); /* P(s) */
int sem_post(sem_t *s); /* V(s) */
```

CS:APP wrapper functions:

```c
#include "csapp.h"

void P(sem_t *s); /* Wrapper function for sem_wait */
void V(sem_t *s); /* Wrapper function for sem_post */
```
badcnt.c: Improper Synchronization

/* Global shared variable */
volatile long cnt = 0; /* Counter */

int main(int argc, char **argv)
{
    long niters;
    pthread_t tid1, tid2;
    
    niters = atoi(argv[1]);
    Pthread_create(&tid1, NULL,
            thread, &niters);
    Pthread_create(&tid2, NULL,
            thread, &niters);
    Pthread_join(tid1, NULL);
    Pthread_join(tid2, NULL);
    
    /* Check result */
    if (cnt != (2 * niters))
        printf("BOOM! cnt=%ld\n", cnt);
    else
        printf("OK cnt=%ld\n", cnt);
    exit(0);
}

How can we fix this using semaphores?

/* Thread routine */
void *thread(void *vargp)
{
    long i, niters =
            *((long *)vargp);
    
    for (i = 0; i < niters; i++)
        cnt++;
    
    return NULL;
}
Using Semaphores for Mutual Exclusion

- **Basic idea:**
  - Associate a unique semaphore **mutex**, initially 1, with each shared variable (or related set of shared variables).
  - Surround corresponding critical sections with \( P(mutex) \) and \( V(mutex) \) operations.

- **Terminology:**
  - **Binary semaphore**: semaphore whose value is always 0 or 1
  - **Mutex**: binary semaphore used for mutual exclusion
    - P operation: “locking” the mutex
    - V operation: “unlocking” or “releasing” the mutex
    - “Holding” a mutex: locked and not yet unlocked.
  - **Counting semaphore**: used as a counter for set of available resources.
goodcnt.c: Proper Synchronization

- Define and initialize a mutex for the shared variable `cnt`:

```
volatile long cnt = 0;  /* Counter */
sem_t mutex;            /* Semaphore that protects cnt */
sem_init(&mutex, 0, 1); /* mutex = 1 */
```

- Surround critical section with `P` and `V`:

```
for (i = 0; i < niters; i++) {
    P(&mutex);
    cnt++;
    V(&mutex);
}
```

Warning: It’s orders of magnitude slower than `badcnt.c`. 
goodcnt.c: Proper Synchronization

- Define and initialize a mutex for the shared variable cnt:

```c
volatile long cnt = 0;  /* Counter */
sem_t mutex;            /* Semaphore that protects cnt */
sem_init(&mutex, 0, 1); /* mutex = 1 */
```

- Surround critical section with P and V:

```c
for (i = 0; i < niters; i++) {
P(&mutex);
cnt++;
V(&mutex);
}
```

```
linux> ./goodcnt 10000
OK cnt=20000
linux> ./goodcnt 10000
OK cnt=20000
linux>
```

Warning: It's orders of magnitude slower than badcnt.c.

```
goodcnt.c
OK cnt=2000000 BOOM! cnt=1036525 Slowdown
real 0m0.138s 0m0.007s 20X
user 0m0.120s 0m0.008s 15X
sys 0m0.108s 0m0.000s NaN
```

And slower means much slower!
Why Mutexes Work

Provide mutually exclusive access to shared variable by surrounding critical section with $P$ and $V$ operations on semaphore $s$ (initially set to 1)

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Provide mutually exclusive access to shared variable by surrounding critical section with *P* and *V* operations on semaphore *s* (initially set to 1).

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Binary Semaphores

- Mutex is special case of semaphore
  - Value either 0 or 1
- Pthreads provides pthread_mutex_t
  - Operations: lock, unlock
- Recommended over general semaphores when appropriate
goodmcnt.c: Mutex Synchronization

- Define and initialize a mutex for the shared variable `cnt`:

```c
volatile long cnt = 0; /* Counter */
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL); // No special attributes
```

- Surround critical section with `lock` and `unlock`:

```c
for (i = 0; i < niters; i++) {
    pthread_mutex_lock(&mutex);
    cnt++;
    pthread_mutex_unlock(&mutex);
}
```

<table>
<thead>
<tr>
<th>Function</th>
<th>badcnt</th>
<th>goodcnt</th>
<th>goodmcnt</th>
</tr>
</thead>
<tbody>
<tr>
<td>Time (ms) niters = 10^6</td>
<td>12</td>
<td>450</td>
<td>214</td>
</tr>
<tr>
<td>Slowdown</td>
<td>1.0</td>
<td>37.5</td>
<td>17.8</td>
</tr>
</tbody>
</table>
Summary

- Programmers need a clear model of how variables are shared by threads.

- Variables shared by multiple threads must be protected to ensure mutually exclusive access.

- Semaphores are a fundamental mechanism for enforcing mutual exclusion.