

I RANDOM RELATIONSHIP

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(*Hint:* express a random variable as a polynomial)

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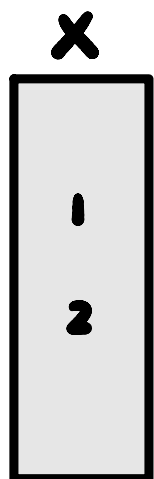
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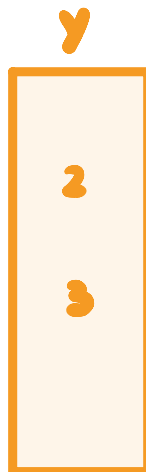
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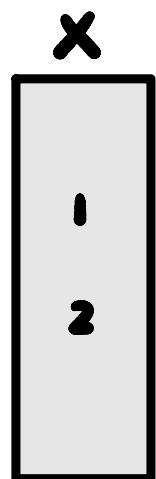
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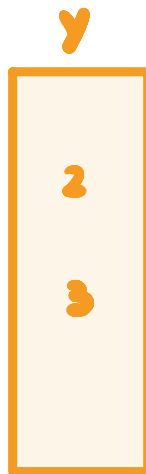
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KEY IDEA:

you can think of these r.v.'s as n -sided dice, with each side weighted by a probability.

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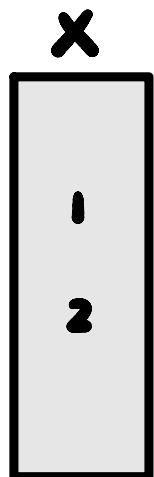
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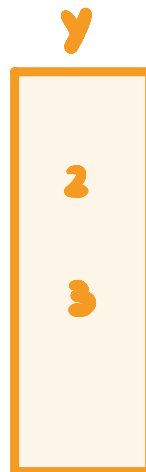
*you roll X , you roll Y
and then you add the
outcomes

$$Z = X + Y$$



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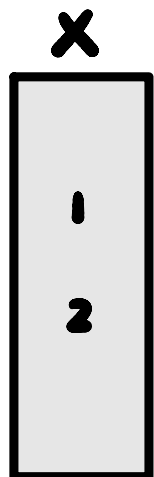
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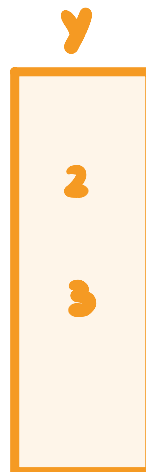
* what is the probability of each outcome?

$$Z = X + Y$$



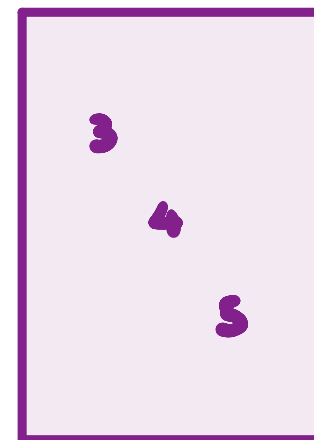
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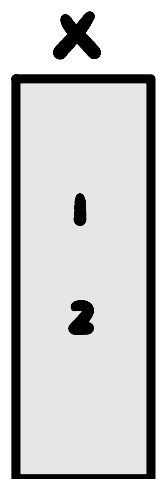
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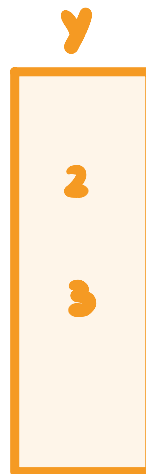
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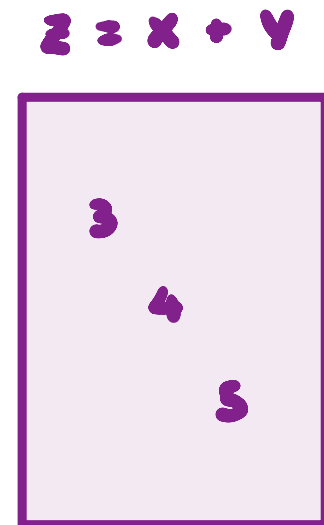
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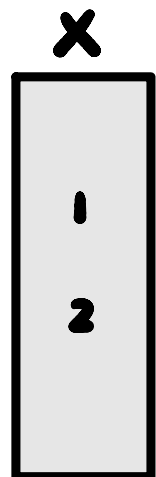
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KEY OBSERVATION:

what process did you use to derive $\mathbb{P}(Z = z)$?

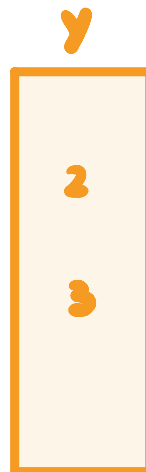
does it remind you of smth?

$$Z = X + Y$$



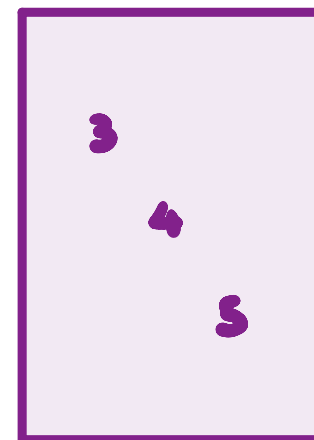
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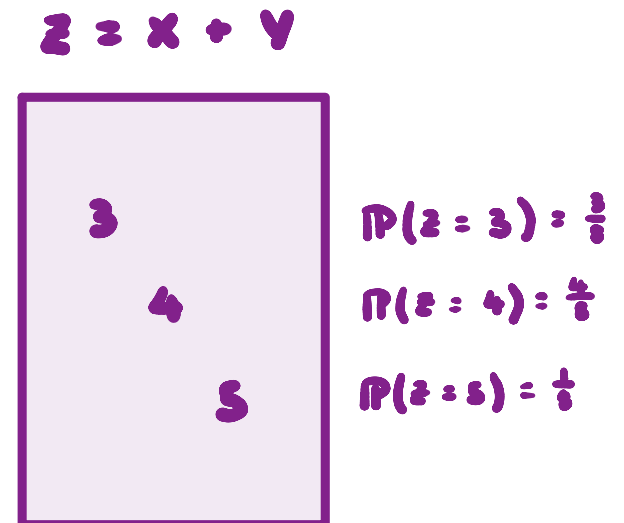
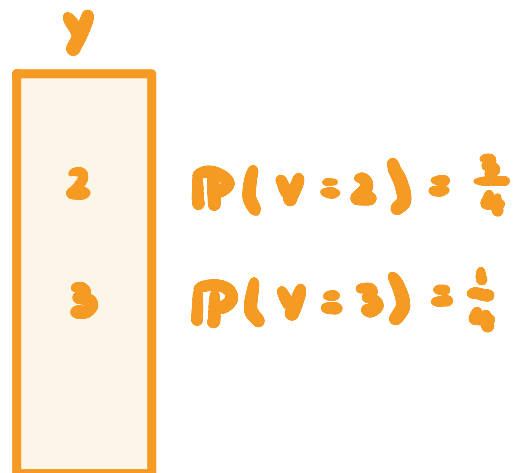
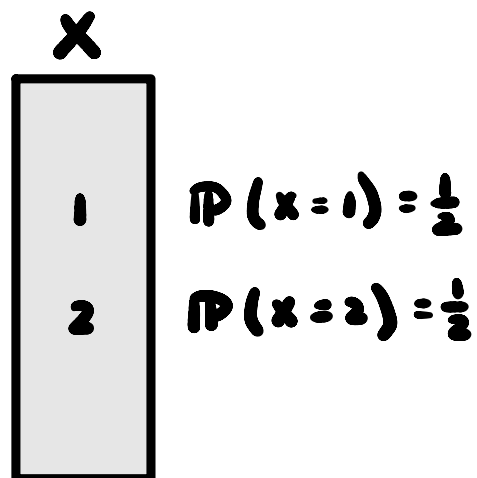
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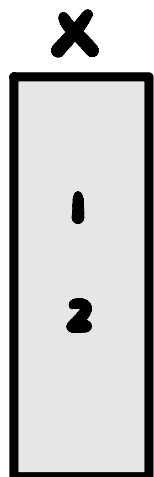
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* let's convert X and Y into polynomials and represent Z as their multiplication.

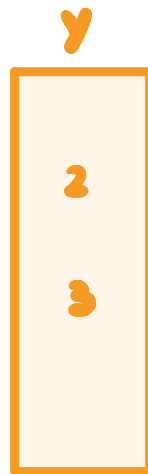
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$$Z = X + Y$$



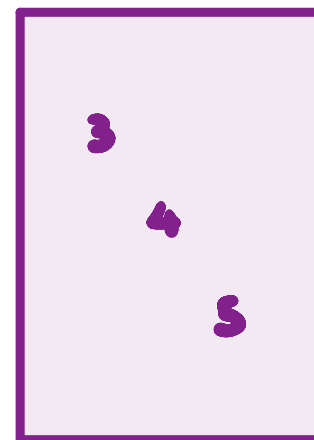
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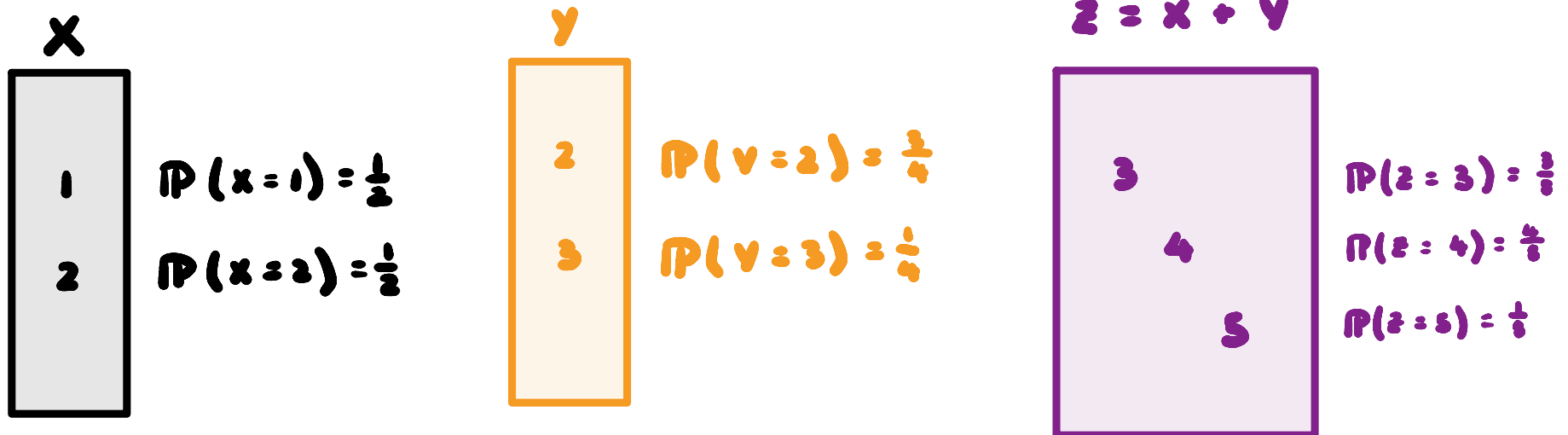
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we need degrees for our polynomial

and we need coefficients

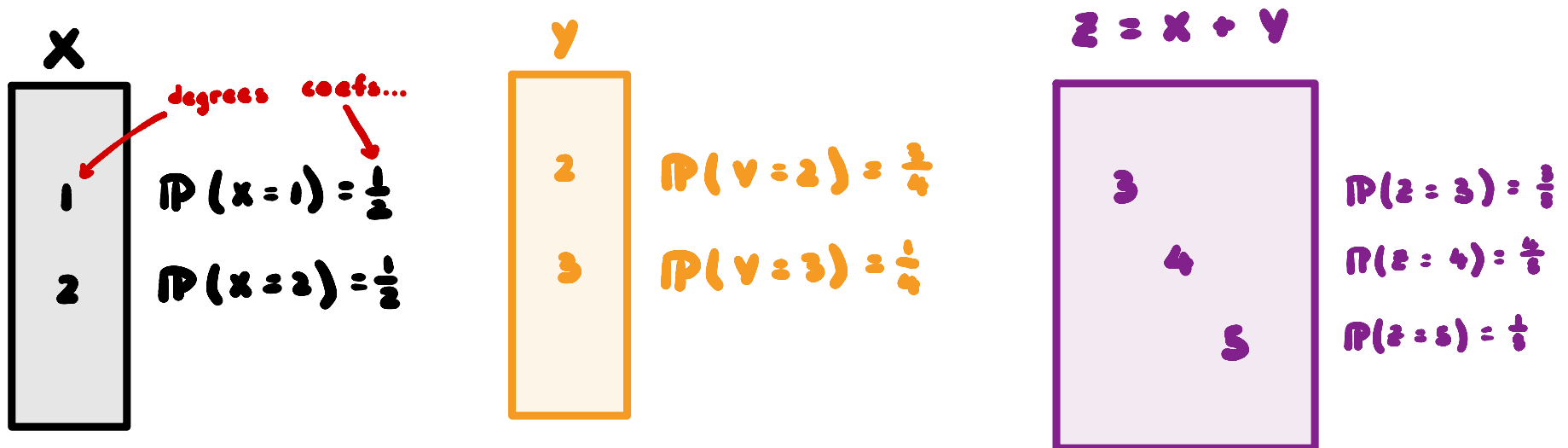
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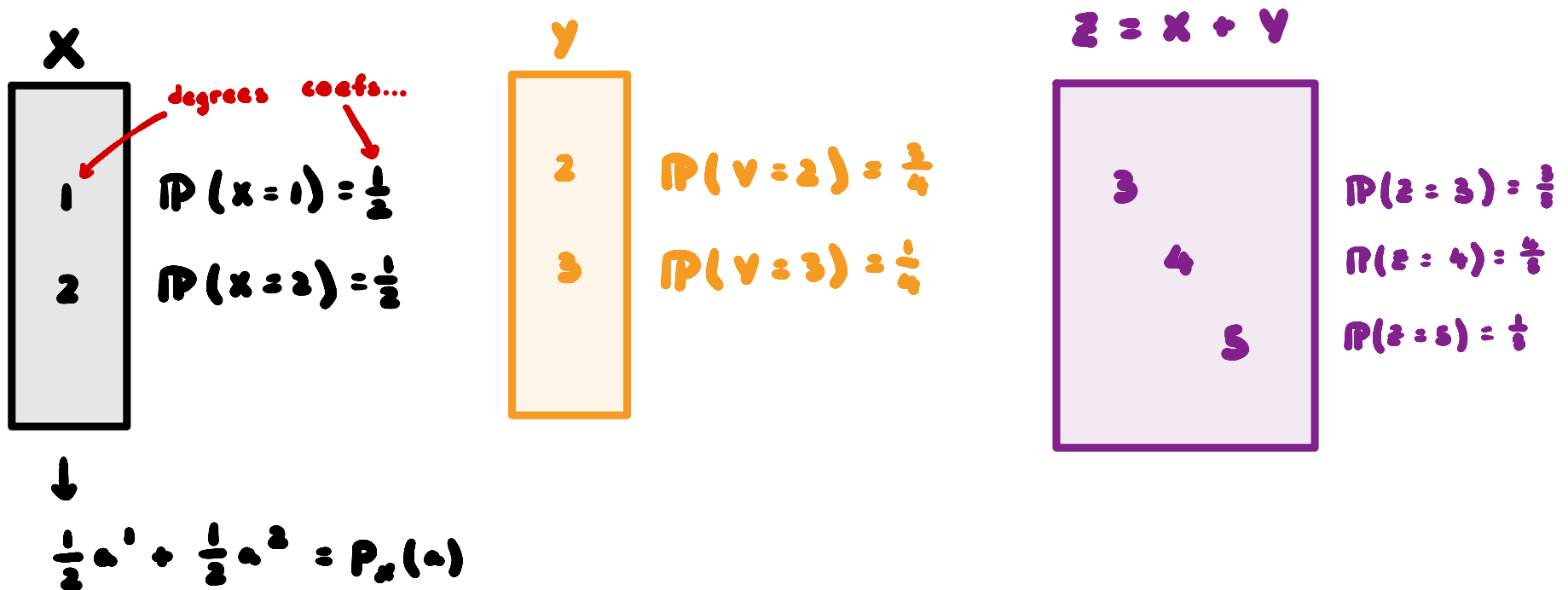
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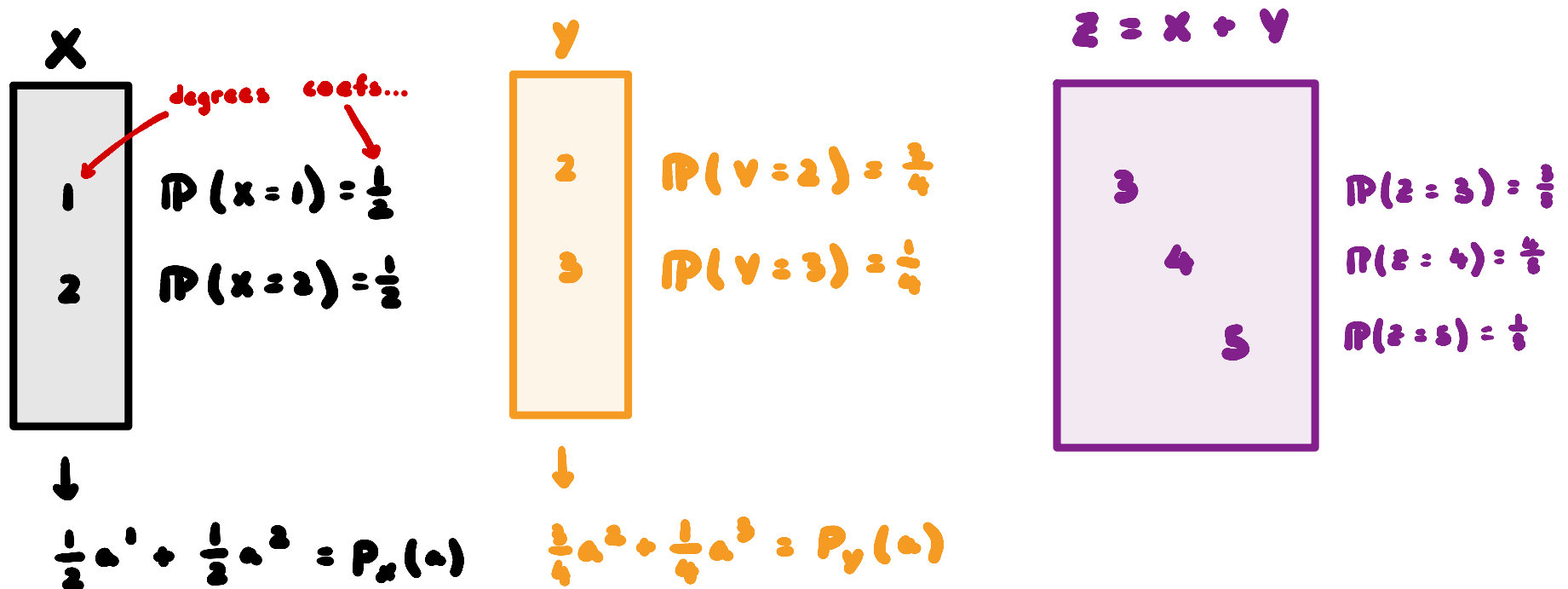
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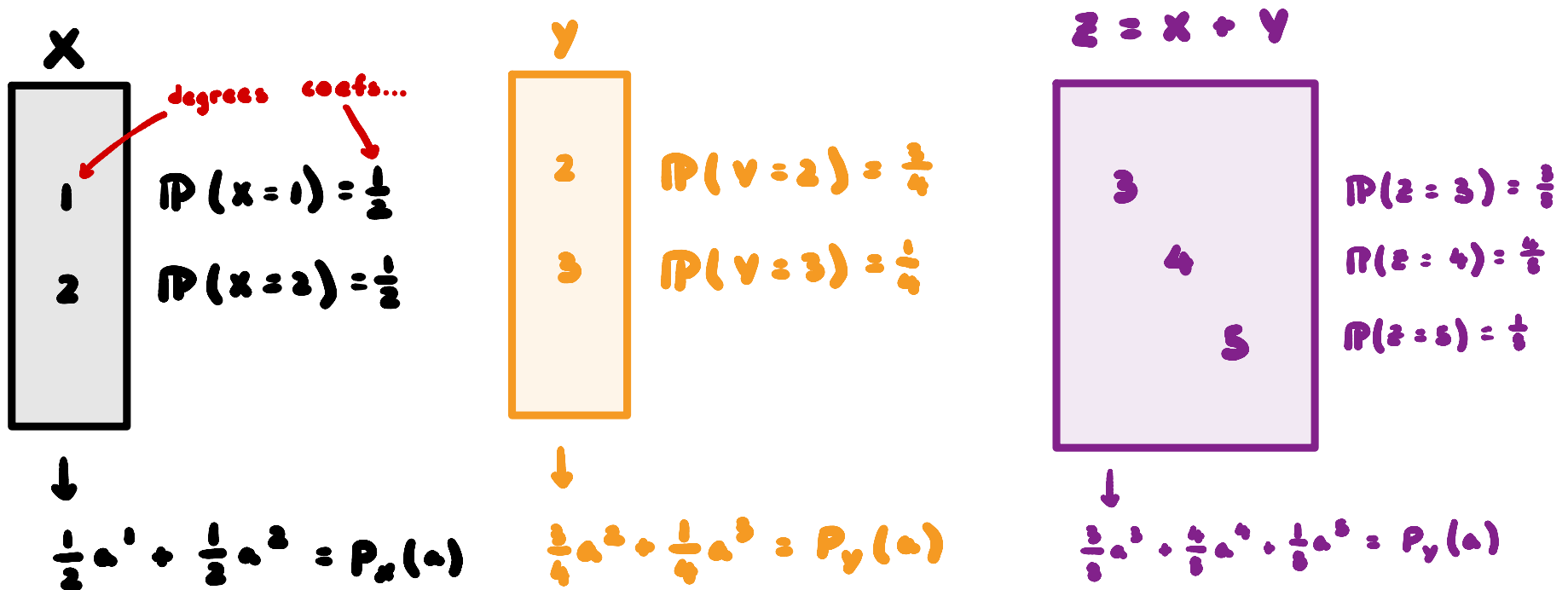
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X

degrees coefs...

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from the defn of poly mult. in notes

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substitute terms

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• try all possible combos of these... $O(n^2)$

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- what should we convert to a polynomial?
- what are the coeffs and degrees?

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Given a binary string S of length n . We wish to determine whether there exists three evenly spaced ones within S . For example, 11100000, 110110010 both have three evenly spaced 1s, while 1011 does not.

(b) Derive an algorithm with $O(n \log n)$ complexity that uses polynomial multiplication and convolutions.

* uh... so in the last problem, we converted r.v.'s to polynomials. let's not reinvent the wheel...

- what should we convert to a polynomial?
- what are the coeffs and degrees?

• the binary string!

• coeffs = 1 or 0

• degree = index

ex. 1011

$$\downarrow \\ 1x^0 + 0x^1 + 1x^2 + 1x^3$$

$$= x^0 + x^2 + x^3$$

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• how does this help us? ... not so clear, let's keep exploring...

• what should we multiply this poly by?
(hint: 2-SUM)

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 - what should we multiply this poly by?
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 - itself!

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10101 $\rightarrow x^0 + 2x^2 + 3x^3 + 2x^4 + x^5$

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$$\begin{aligned} & \downarrow \\ & 1x^0 + 0x^1 + 1x^2 + 1x^3 \\ & = x^0 + x^2 + x^3 \end{aligned}$$

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0101 $\rightarrow x^1 + 2x^2 + x^3$

10101 $\rightarrow x^0 + 2x^2 + 3x^3 + 2x^4 + x^6$

10111 $\rightarrow x^0 + 2x^2 + 2x^3 + 3x^4 + 2x^5 + 3x^6 + 2x^7 + x^8$

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* notice anything interesting about the magnitude of the coeffs?

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* notice anything interesting about the magnitude of the coeffs?

• NO \rightarrow coeffs ≤ 2

• YES $\rightarrow \exists$ coeffs ≥ 3

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* notice anything interesting about the magnitude of the coeffs?

• NO \rightarrow coeffs ≤ 2

• YES \rightarrow coeffs $\geq 3 \rightarrow$ located at degree DOUBLE the center of ones!

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* notice anything interesting about the magnitude of the coeffs?

• NO \rightarrow coeffs ≤ 2 ^{why is this happening?}

• YES \rightarrow coeffs $\geq 3 \rightarrow$ located at degree DOUBLE the center of ones!

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* notice anything interesting about the magnitude of the coeffs?

• NO \rightarrow coeffs ≤ 2 ^{suppose indices (i, j, k) are evenly spaced}
 $\Rightarrow j - i = k - j = k - i$

• YES $\rightarrow \exists$ coeffs $\geq 3 \rightarrow$ located at degree DOUBLE the center of ones!

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ex. 10111

$$(x^0 + x^2 + x^3 + x^4)$$

$$(x^0 + x^2 + x^3 + x^4)$$

$$x^0 + 2x^2 + 2x^3 + 3x^4 + 2x^5 + 3x^6 + 2x^7 + x^8$$

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- (b) Derive an algorithm with $O(n \log n)$ complexity that uses polynomial multiplication and convolutions.

ex. 10111

* these multiplications in the FOIL give rise to 3 (bc evenly-spaced)

$$\begin{array}{r} (x^0 + x^2 + x^3 + x^4) \\ \hline (x^0 + x^2 + x^3 + x^4) \\ \hline = \\ x^0 + 2x^2 + 2x^3 + \underline{3x^4} + 2x^5 + 3x^6 + 2x^7 + x^8 \end{array}$$

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***so, what's the algorithm?**

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- (b) Derive an algorithm with $O(n \log n)$ complexity that uses polynomial multiplication and convolutions.

***so, what's the algorithm?**

1. make polynomial for string $O(n)$

2. square it $O(n \log n)$

3. for degree $d \in P(x)$ ^{*degrees with 1s}

1. check coefficient of $2d$ in $P(x)^2$. if $\neq 0$, return TRUE

4. return FALSE

THANK YOU. 😊

*** this is my last recitation ever...
y'all mean the world to me**