

Fast Fourier Transform, convolutions, Applications.

$$A(x) = a_0 + a_1 x + \dots + a_d x^d$$

$$B(x) = b_0 + b_1 x + \dots + b_d x^d$$

Goal: compute the product $C(x)$

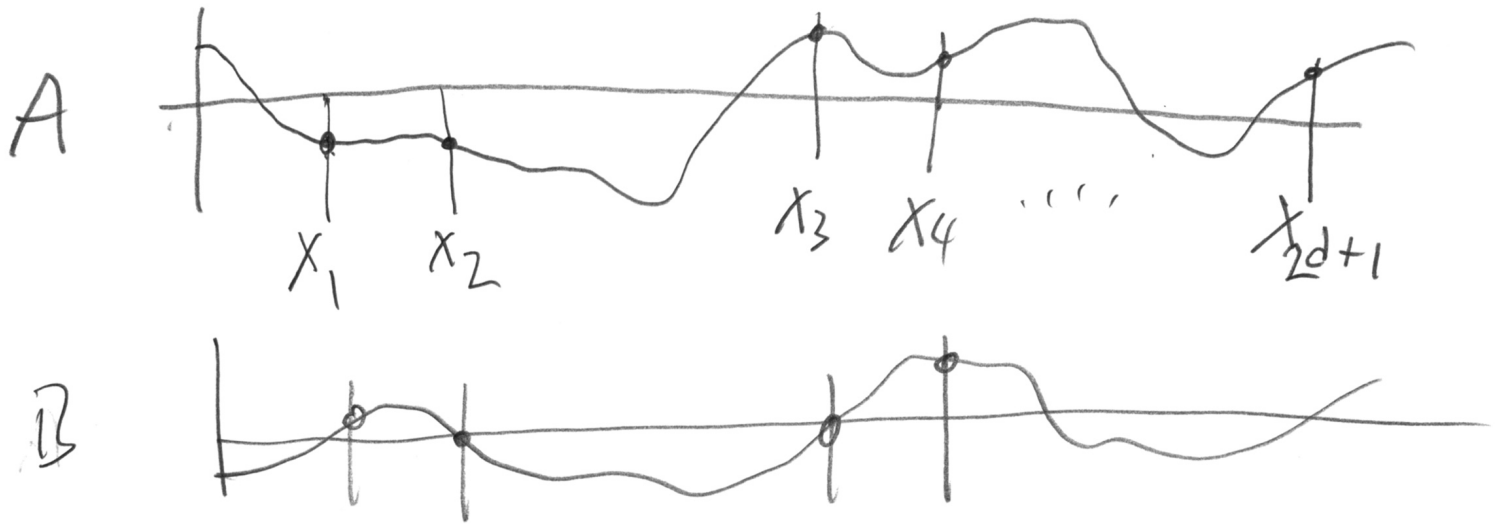
$$C(x) = c_0 + c_1 x + \dots + c_{2d} x^{2d}$$

$$c_k = \sum_{\substack{0 \leq i, j \leq k \\ i+j=k}} a_i b_j \leftarrow O(d^2)$$

convolution of $[a_0 \dots a_d] * [b_0 \dots b_d]$

Alternative Rep of Polynomials

Point-Value Rep.



$$C = A \cdot B$$

$$C(x_i) = A(x_i) \cdot B(x_i)$$

$O(d)$
time

Convolution

$O(d^2)$ Naive

$O(d^{1.58})$ Karatsuba's

today $O(d \log d)$ using FFT

Shoehaga Strassen

integer mult.

$O(n \log n \log \log n)$

Plan

Pick $N \geq 2d+1$ (N power of 2)

- ① Pick N points $x_0 \dots x_{N-1}$
in a special way
- ② Evaluate $A()$ at them
- ③ Evaluate $B()$ at them
- ④ compute $c(x_0) \dots c(x_{N-1})$ $O(N)$
- ⑤ "Interpolate" C back to
the coefficient rep

Picking the points (x values)

$$A(x) = a_0 + a_1x + \dots + a_7x^7$$

$$x_0 = 1, \quad x_1 = -1$$

$$\left. \begin{aligned} A(1) &= a_0 + a_1 + a_2 + a_3 + a_4 + a_5 + a_6 + a_7 \\ A(-1) &= a_0 - a_1 + a_2 - a_3 + a_4 - a_5 + a_6 - a_7 \end{aligned} \right\} 14$$

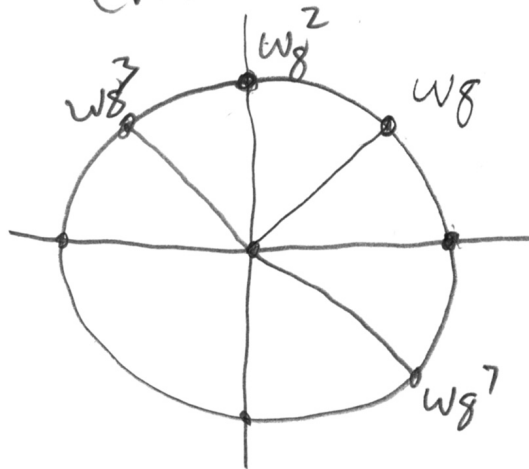
$$\left. \begin{aligned} W &= a_0 + a_2 + a_4 + a_6 \\ Z &= a_1 + a_3 + a_5 + a_7 \end{aligned} \right\} 6$$

$$\left. \begin{aligned} A(1) &= W + Z \\ A(-1) &= W - Z \end{aligned} \right\} 2$$

8

Complex Numbers

$$w_8^8 = 1$$

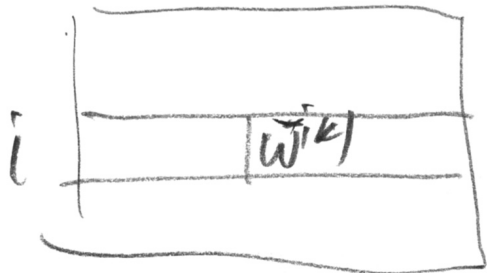


w_N = the primitive
Nth root of unity

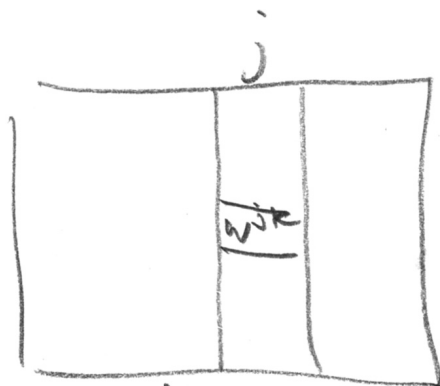
$$w_N = e^{2\pi i/N} = i \sin\left(\frac{2\pi}{N}\right) + \cos\left(\frac{2\pi}{N}\right)$$

$$w_N^N = 1$$

$$w_N^i \neq 1 \quad \text{for } 0 < i < N$$



DFT (ω, N)



DFT (ω, N)

$$\begin{pmatrix} \omega^{0 \cdot 0} & \omega^{0 \cdot 1} & \dots & \omega^{0 \cdot (N-1)} \\ \omega^{1 \cdot 0} & \omega^{1 \cdot 1} & \dots & \omega^{1 \cdot (N-1)} \\ \vdots & \vdots & \ddots & \vdots \\ \omega^{(N-1) \cdot 0} & \omega^{(N-1) \cdot 1} & \dots & \omega^{(N-1) \cdot (N-1)} \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_{N-1} \end{pmatrix} = \begin{pmatrix} A(\omega^0) \\ A(\omega^1) \\ \vdots \\ A(\omega^{N-1}) \end{pmatrix}$$

||
DFT (ω_N, N)

$$F_N(A)_j = A(\omega^j)$$

$$F_N(A) = [A(\omega^0), A(\omega^1), \dots, A(\omega^{N-1})]$$

2

$$= \sum_{i=0}^{N-1} a_i \omega^{ij}$$

$$= \sum_{\substack{i=0 \\ i \text{ even}}}^{N-1} a_i \omega^{ij} + \sum_{\substack{i=1 \\ i \text{ odd}}}^{N-1} a_i \omega^{ij}$$

$$= \sum_{i=0}^{\frac{N}{2}-1} a_{2i} \omega^{2ij} + \sum_{i=0}^{\frac{N}{2}-1} a_{2i+1} \omega^{(2i+1)j}$$

$$= \sum_{i=0}^{\frac{N}{2}-1} a_{2i} (\omega^2)^{ij} + \omega^j \sum_{i=0}^{\frac{N}{2}-1} a_{2i+1} (\omega^2)^{ij}$$

$$= \sum_{i=0}^{\frac{N}{2}-1} a_{2i} (\omega_{N/2})^{i(j \bmod N/2)} + \omega_N^j \sum_{i=0}^{\frac{N}{2}-1} a_{2i+1} (\omega_{N/2})^{i(j \bmod N/2)}$$

$$F_N(A)_j = F_{N/2}(A_{\text{even}})_{j \bmod N/2} + \omega_N^j F_{N/2}(A_{\text{odd}})_{j \bmod N/2}$$

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FFT([ $a_0, \dots, a_{N-1}$ ],  $\omega, N$ )
  if  $N = 1$  then return [ $a_0$ ]
   $F_{\text{even}} \leftarrow$  FFT([ $a_0, a_2, \dots, a_{N-2}$ ],  $\omega^2, N/2$ )
   $F_{\text{odd}} \leftarrow$  FFT([ $a_1, a_3, \dots, a_{N-1}$ ],  $\omega^2, N/2$ )
   $F \leftarrow$  a new vector of length  $N$ 
   $x \leftarrow 1$ 
  for  $j = 0$  to  $N - 1$  do
     $F[j] \leftarrow F_{\text{even}}[j \bmod (N/2)] + x * F_{\text{odd}}[j \bmod (N/2)]$ 
     $x \leftarrow x * w$ 
  return  $F$ 

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$$T(N) = 2T(N/2) + O(N)$$

Let's see what happens if we multiply $\mathbf{DFT}(\omega_N^{-1}, N)$ by $\mathbf{DFT}(\omega, N)$.

$$\text{The } (i, j) \text{ entry of } \mathbf{DFT}(\omega^{-1}, N) \mathbf{DFT}(\omega, N) = \sum_{k=0}^{N-1} \omega^{-ik} \omega^{kj}$$

So let's try to evaluate the summation on the right in the two cases of $i = j$ and $i \neq j$.

If $i = j$ then we get:

$$\sum_{k=0}^{N-1} \omega^{-ik} \omega^{kj} = \sum_{k=0}^{N-1} \omega^0 = N$$

If $i \neq j$ then we get:

$$\sum_{k=0}^{N-1} \omega^{-ik} \omega^{kj} = \sum_{k=0}^{N-1} \omega^{(j-i) \cdot k} = \boxed{\sum_{k=0}^{N-1} (\omega^{j-i})^k}$$

$$1 + r + r^2 + \dots + r^{N-1} = \frac{1 - r^N}{1 - r}$$

So

$$\sum_{k=0}^{N-1} (\omega^{j-i})^k = \frac{1 - (\omega^{j-i})^N}{1 - \omega^{j-i}} = \frac{1 - 1}{1 - \omega^{j-i}} = 0$$

So, summarizing what we just learned:

$$\text{The } (i, j) \text{ entry of } \mathbf{DFT}(\omega^{-1}, N) \mathbf{DFT}(\omega, N) = \begin{cases} N & \text{if } i = j \\ 0 & \text{if } i \neq j \end{cases}$$

This is just the identity matrix times N . So what we've just proven is that:

$$\mathbf{DFT}(\omega, N)^{-1} = \frac{1}{N} \mathbf{DFT}(\omega^{-1}, N)$$

A = 1 2 3 4 5 6 7 8 0 0 0 0 0 0 0 0
 B = 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
 conv A B = 1 2 3 4 5 6 7 8 0 0 0 0 0 0 0 0

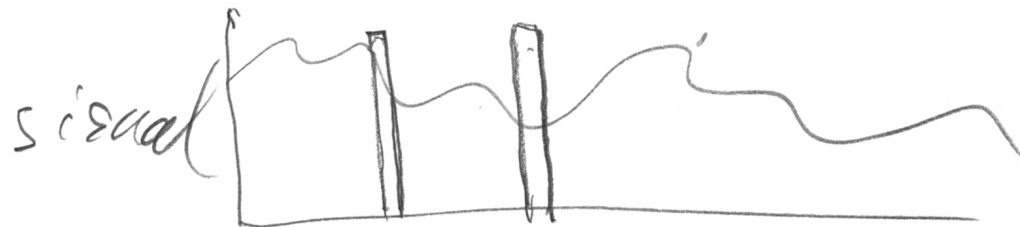
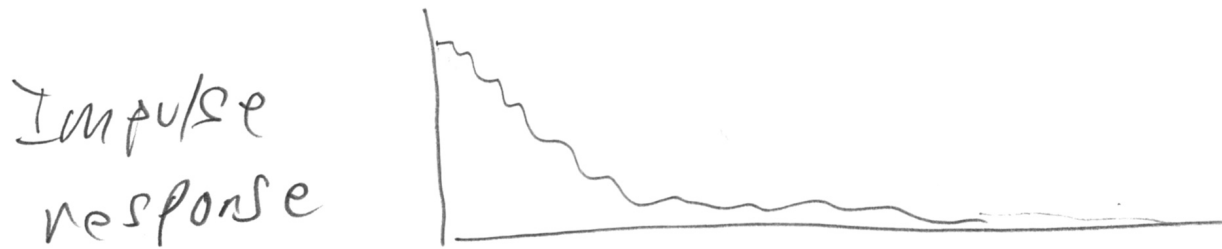
A = 1 2 3 4 5 6 7 8 0 0 0 0 0 0 0 0
 B = 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0
 conv A B = 0 0 0 0 0 0 0 1 2 3 4 5 6 7 8 0

A = 1 2 3 4 5 6 7 8 0 0 0 0 0 0 0 0
 B = 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 0
 conv A B = 4 5 6 7 8 0 0 0 0 0 0 0 0 1 2 3

A = 1 2 3 4 5 6 7 8 0 0 0 0 0 0 0 0
 B = 0 0 0 0 0 1 0 0 0 0 1 0 0 0 0 0
 conv A B = 7 8 0 0 0 1 2 3 4 5 7 9 11 4 5 6

A = 1 2 3 4 5 6 7 8 9 10 11 12 0 0 0 0
 B = 0 0 0 0 0 1 0 0 0 0 1 0 0 0 0 0
 conv A B = 19 8 9 10 11 13 2 3 4 5 7 9 11 13 15 17

Digital Signal Processing



Evenly Spaced Ones

S is a binary string

$$S = \overset{0}{1} \overset{1}{1} \overset{2}{0} \overset{3}{1} \overset{4}{1} 0 0 1 0$$

Brute force: $O(N^2)$

FFT method

$$P(x) = x^0 + x^1 + x^3 + x^4 + x^7$$

$$P^2(x) = x^0 + 2x^1 + x^2 + 2x^3 + 4x^4 + \dots + 3x^8 + \dots$$

the coefficient of x^t is the number of pairs (a, c) with ones and $a+c=t$
position $b=4$ x^{2b} term

