

# 15-110 Quiz 4 Review Session

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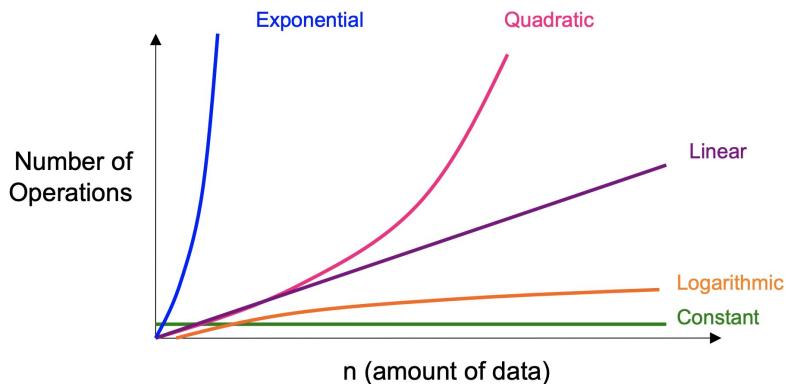
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# Gameplan

- Review the major topics
  - Big O
  - Trees
  - Graphs
  - Tractability
- Go over some practice problems
- Q + A

# Big O

- A way for programmers to group certain functions into different families based on their speed/efficiency
- Measure based on worst case
- Ignore lower order terms!



$f(n)$	Big-O
$n$	$O(n)$
$32n + 23$	$O(n)$
$5n^2 + 6n - 8$	$O(n^2)$
$18 \log(n)$	$O(\log n)$

# Common Themes in Big O

- For loops
  - Check how many times the for loop is running
    - If it's related to the input (i.e. `len(lst)`), then it's  $O(n)$
    - If it's looping a constant amount (`for i in range(20)`), it's  $O(1)$
- While loops
  - Check how the while loop is increasing/decreasing the iterator variable
    - If it's using addition/subtraction, it's  $O(n)$
    - If it's using multiplication/division, it's  $O(\log n)$
- Most built in functions/methods (`len`, `.append()`, indexing into an array, etc.) are  $O(1)$  **UNLESS OTHERWISE SPECIFIED**
  - `.find()`, `in`, and other methods/functions that search through are  $O(n)$

# Strategy to Big O

1. Go through a function line by line
2. Find the Big O value of each individual line
  - a. Most lines are  $O(1)$ !
3. Once you're done, go through to find the total Big O
  - a. If it's one the same indentation levels, add up their big O's
    - i.  $O(1) + O(1) + O(1) = O(3)$
  - b. If a line is indented, or nested, in another line, multiply their big O's
    - i.  $O(n) * O(n) = O(n^2)$
4. At the end, make sure to simplify whatever value you get to fit into a function family
  - a.  $O(3n^2 + 2n + 6) \rightarrow O(n^2)$

# Big O Example

```
# s is a string larger than 10 characters
def mur(s):
    result = ""
    for i in range(10):
        result += s[i]

    return result
```

# Big O Example

```
# L is a list of integers
def allPossiblePairs(L):
    result = []
    for x in L:
        for y in L:
            result += [[x, y]]
    return result
```

# Big O Example

```
# n is a positive integer
def foo(n):
    result = 0
    while (n > 0):
        n = n // 2
        result += 1
    return result
```



# Tractability

- A problem is said to be tractable if it has a reasonably efficient runtime:
  - $O(1)$ ,  $O(\log n)$ ,  $O(n \log n)$ ,  $O(n^2)$ ,  $O(n^{10000})$
  - ^ Polynomial time
- Intractable:
  - $O(2^n)$ ,  $O(n!)$ ,  $O(k^n)$
  - ^ Bigger than polynomial time

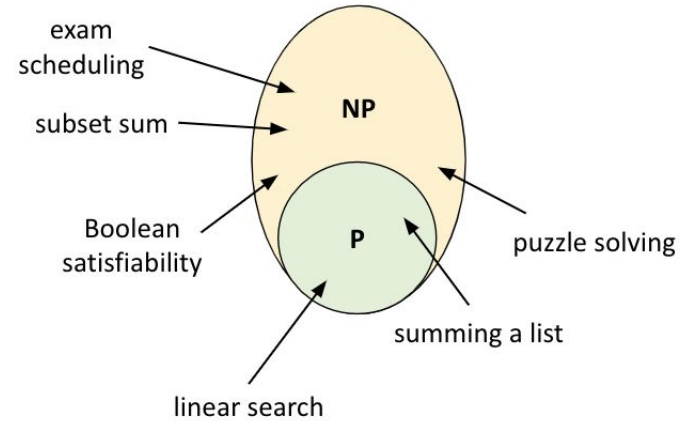
# Brute Force Algorithms

- Brute force algorithms check every possible solution.
  - Ex. Testing every subset of [1,2,3] for subset sum ==>
  - Other ones: travelling salesperson, puzzle-solving, and exam scheduling
- Generally are intractable solutions.

- []
- [1]
- [2]
- [1, 2]
- [3]
- [1, 3]
- [2, 3]
- [1, 2, 3]

# P vs. NP

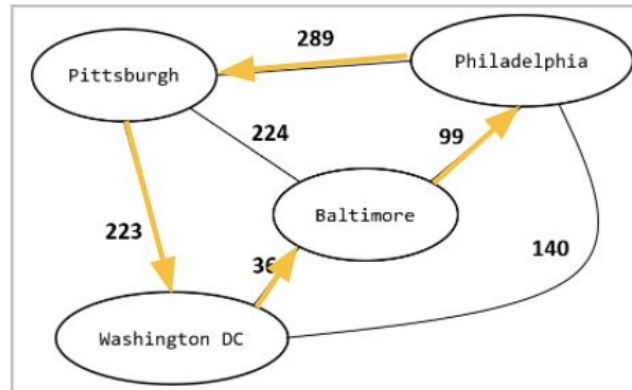
	P	NP
Verifying	Tractable	Tractable
Solving	Tractable	?



Why isn't traveling salesperson in this diagram?

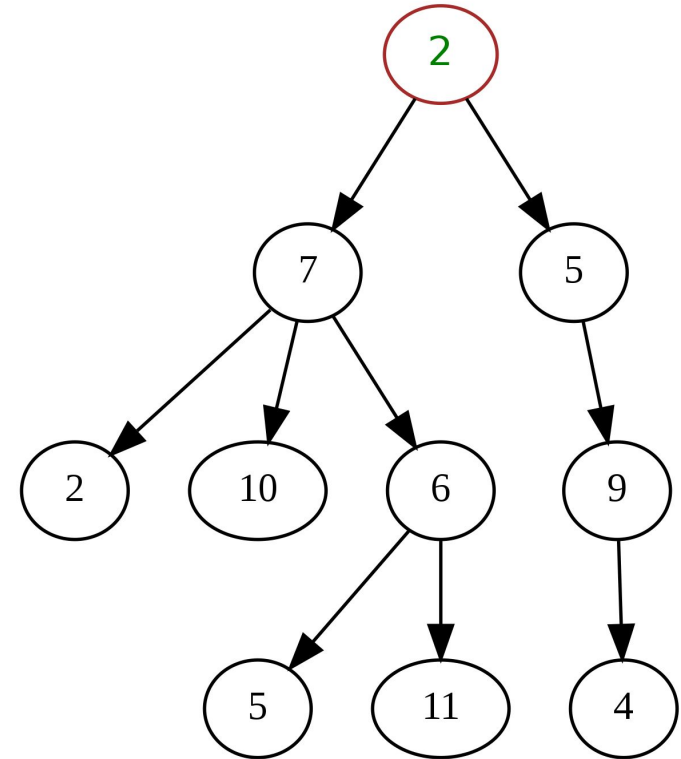
# Heuristics

- Shortcuts to find a solution that is not the best, but is close.
  - Finding the best solution often takes a really long time
  - Finding a good solution is often much easier

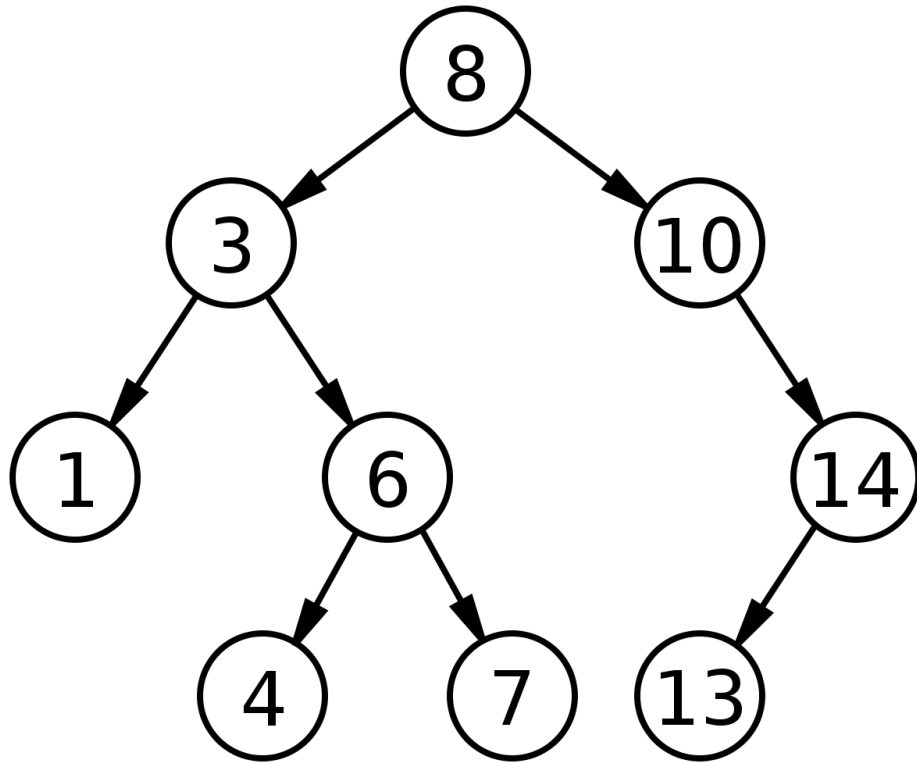


# Trees

- Another different data structure
- Hierarchical structure (up, down)
- Uses nodes, and each node has a value
- Nodes connected below a node are the children, and the node above is the parent
- Top node is root, nodes with no children are leaves



# Trees Example



Root: 8

Node: all of them! (1, 3, 4, 6...13)

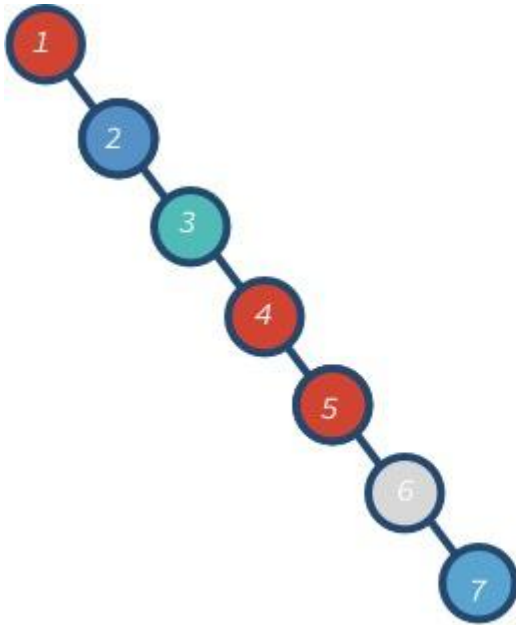
Children (of 8): 3, 10

Leaves: 1, 4, 7, 13

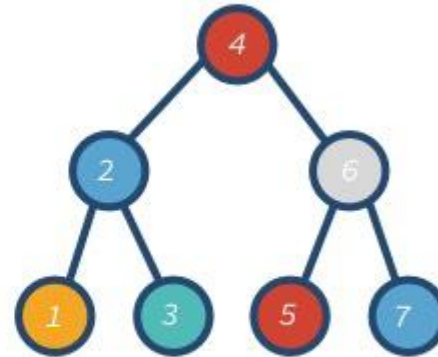
# Binary Search Trees

- A specific type of tree which allows for easy search of values (binary search! It's in the name!)
- The main rule for BSTs
  - Everything to the **left** of a node must be **less** than the value at that node, and everything to the **right** of the node must be **greater** than the value at that node

# Balanced vs. Unbalanced Search Trees



AdrianMejia.c





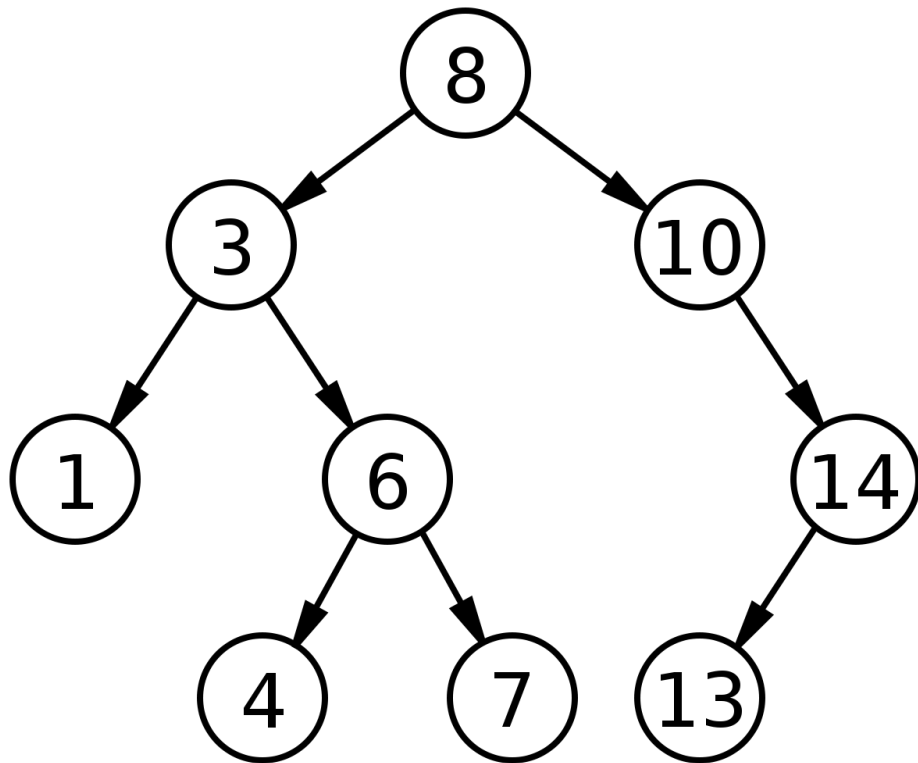
# Trees in Code

- Store in a dictionary!
  - Three keys: contents, left, right
    - Contents is the value at that node
    - Left and right are either another dictionary containing the same three keys, or None, meaning there is no child.

```
{ "contents" : nodeValue,  
  "left"      : leftChildSubtree,  
  "right"     : rightChildSubtree }
```

# Trees Example

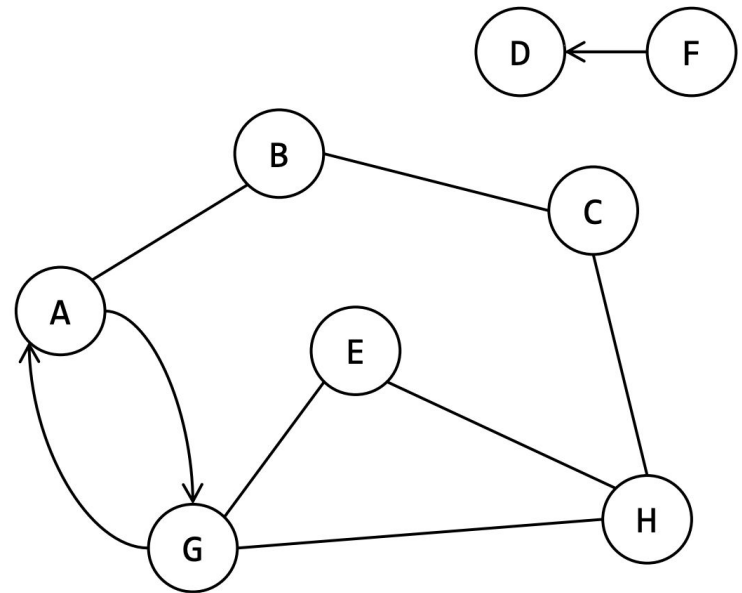
Say we are given a tree, where each node has two children (not necessarily a BST, but you can think of it like that). Write a recursive function `addOdds(tree)` that adds all of the odd leaves and returns the sum of all the odd leaves.



```
def addOdds(tree):  
    if tree["left"] == None and tree["right"] == None:  
        # Base case: We are at a leaf node  
        if tree["contents"] % 2 == 1:  
            return tree["contents"]  
        else:  
            return 0  
    else:  
        # Recursive case: We are not at a leaf node  
        result = 0  
        if tree["left"] != None:  
            result += addOdds(tree["left"])  
  
        if tree["right"] != None:  
            result += addOdds(tree["right"])  
  
        return result
```

# Graphs

- Similar to trees – nodes connected to other nodes
- Less restrictions – any node can be connected to any other node, no longer follows a hierarchical structure
- Graphs have **edges** – edges are the connections between the nodes
  - Sometimes, edges can have **weights**, which is just a number associated with the edge
  - Edges can also be **directed** or **undirected**



# Graphs in Code

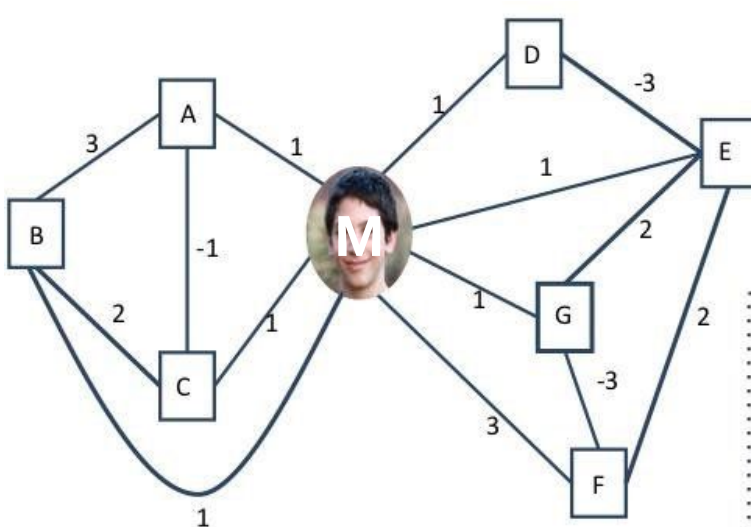
```
{ node : [ [neighborValue, weight] ],  
  ... }
```

- Treated as a dictionary!
  - Stored slightly differently if the graph is weighted or not
  - Loop through all nodes with a for each loop (same as keys in dictionary)

```
{ nodeValue : [ neighborValue ],  
  ... }
```

# Graph Problem

Say we are given a certain person in our graph, in this case, Michael. Write a function `findFriendsList(person, g)` that takes a person and a graph, and returns all the friends of that person in the form of a list.



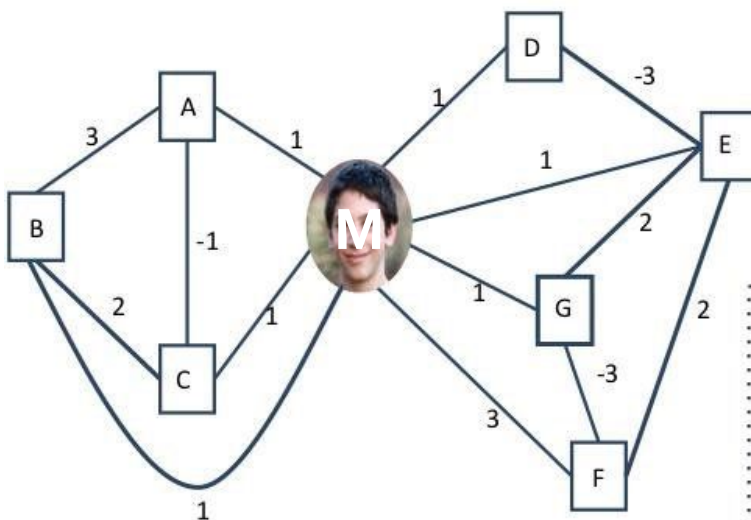
Michael

- Friends = 1
- Dating = 3
- Hate = -3
- Dislike = -1
- One has a crush on the other = 2
- No relationship stated = 0

```
def getFriends(person, g):  
    friendList = []  
    for relation in g[person]:  
        if relation[1] == 1:  
            # This is a friend :)  
            friendList += [relation[0]]  
    return friendList
```

# Graph Problem

Write a function `getAllCouples(g)` that takes in a graph, and returns a list of all the couples in the graph, stored together as a 2d list.



Michael

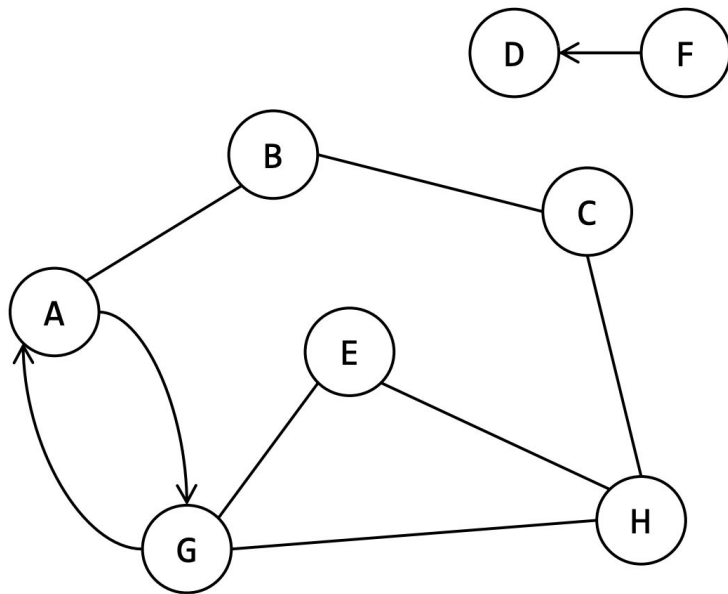
- Friends = 1
- Dating = 3
- Hate = -3
- Dislike = -1
- One has a crush on the other = 2
- No relationship stated = 0



```
def getAllCouples(g):  
    couples = []  
    for person in g:  
        for relation in g[person]:  
            if relation[1] == 3:  
                # They are dating  
                couple = [person, relation[0]]  
                couples += [couple]  
  
    # NOTE: This will contain duplicates of each couple  
    return couples
```

# Breadth vs. Depth First Search

- Different ways of searching for a given node on a graph
  - BFS: start at a node, go through every one of its neighbors, and then go to the neighbors after, etc. until found/looked through everything
  - DFS: start at a node, keep searching down one path until you can/can't find something, start from start again and repeat



# That's it!

Any questions?