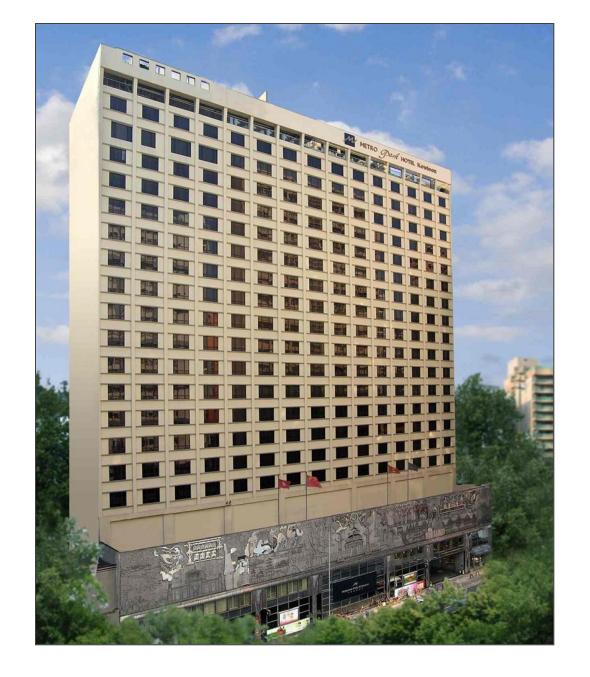
Which Animal Gave Us SARS?

Evolutionary Trees Part 1: The Neighbor-Joining Algorithm

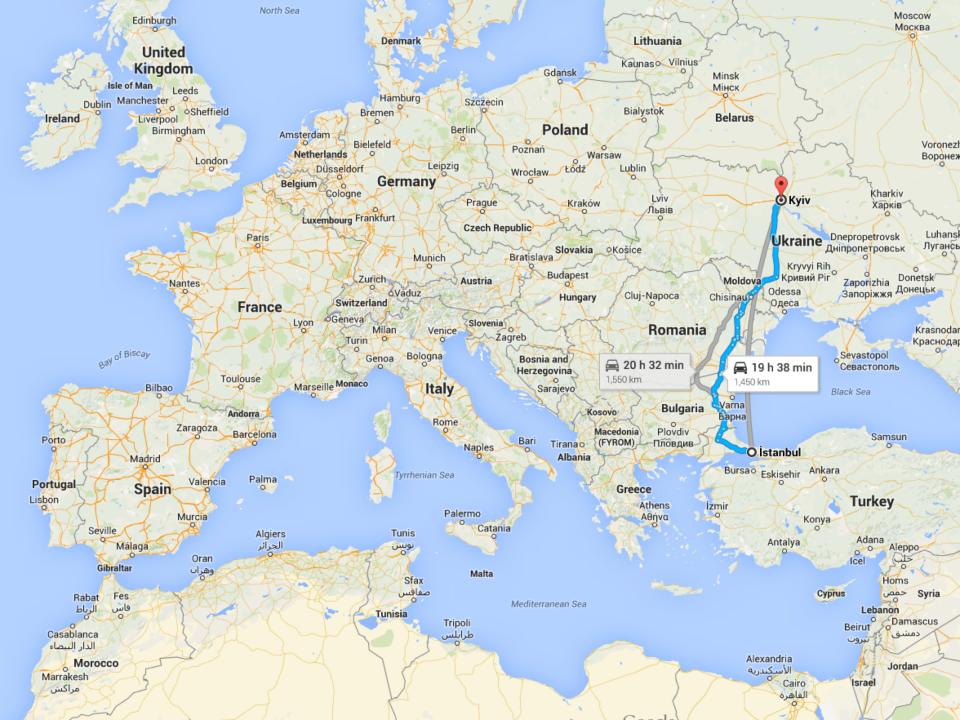
Phillip Compeau and Pavel Pevzner Bioinformatics Algorithms: An Active Learning Approach

©2018 by Compeau and Pevzner. All rights reserved.



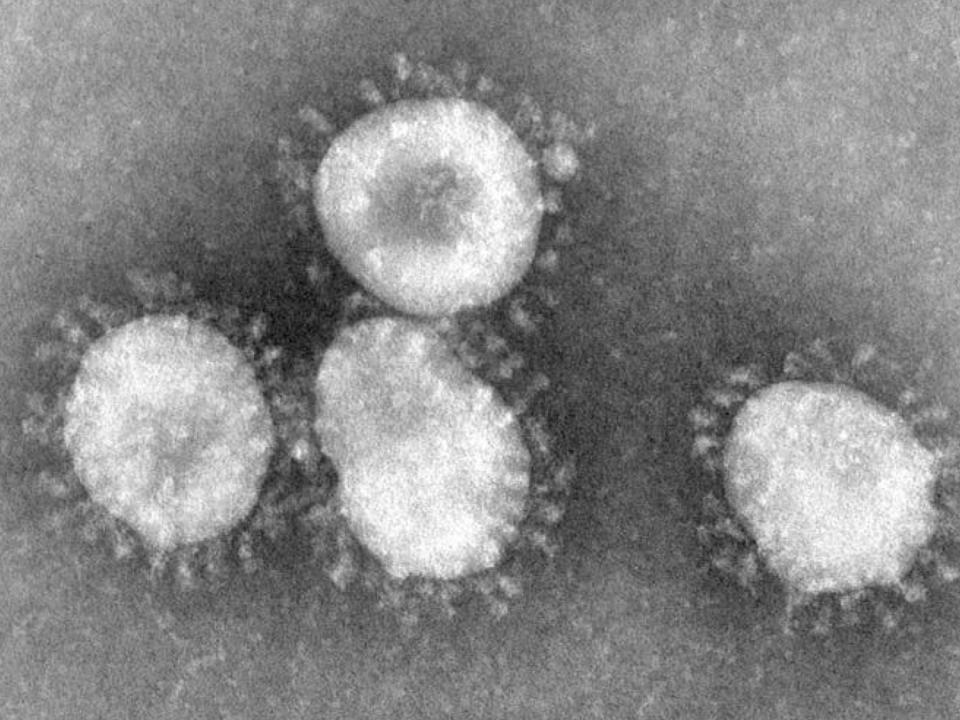
Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.













Questions about SARS

Which animal gave us SARS? How does SARS compare to other viruses and how did it mutate over time?

Questions about SARS

Which animal gave us SARS? How does SARS compare to other viruses and how did it mutate over time?

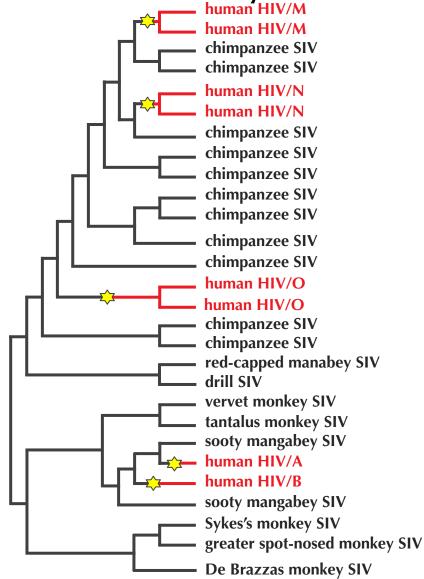
To answer these questions, we need to learn how to construct evolutionary trees (a.k.a. phylogenies).

Example: HIV Evolutionary Tree

— SIVs (monkeys)

— HIV (human)

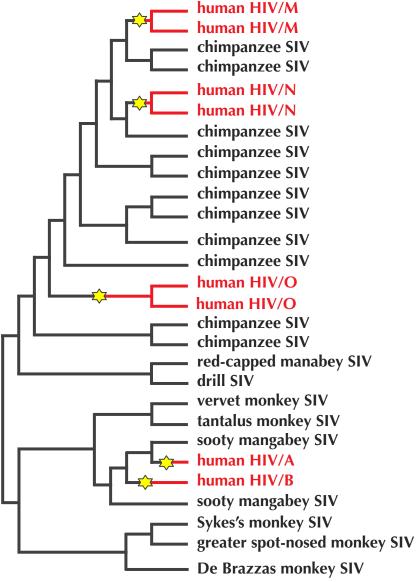
human infection



Two Computational Questions

How do we construct the tree's *structure*?

Can we infer anything about the ancestors?

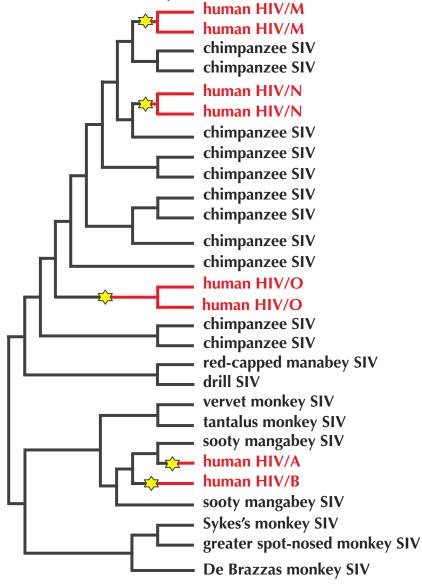


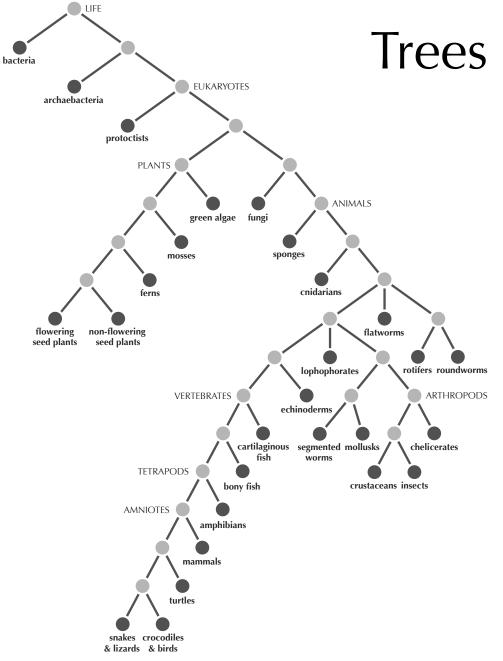
Two Computational Questions

How do we construct the tree's *structure*?

Can we infer anything about the ancestors?

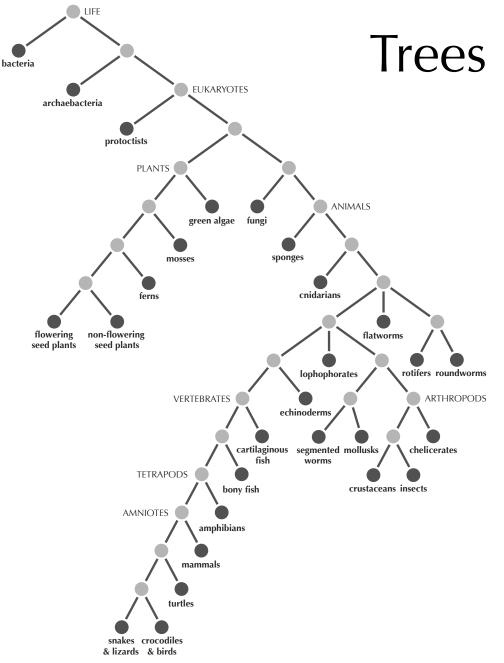
Checkpoint: Any thoughts on how we could answer either question?





Tree: Connected graph containing no cycles.

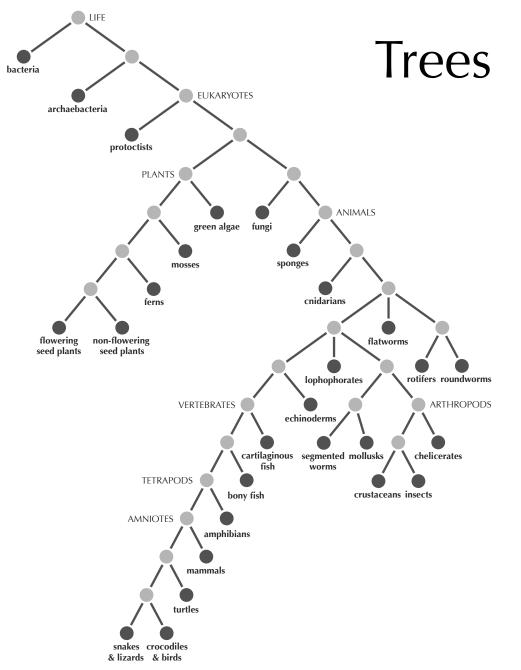
Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.



Tree: Connected graph containing no cycles.

Leaves (degree = 1): present-day species

Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.



Tree: Connected graph containing no cycles.

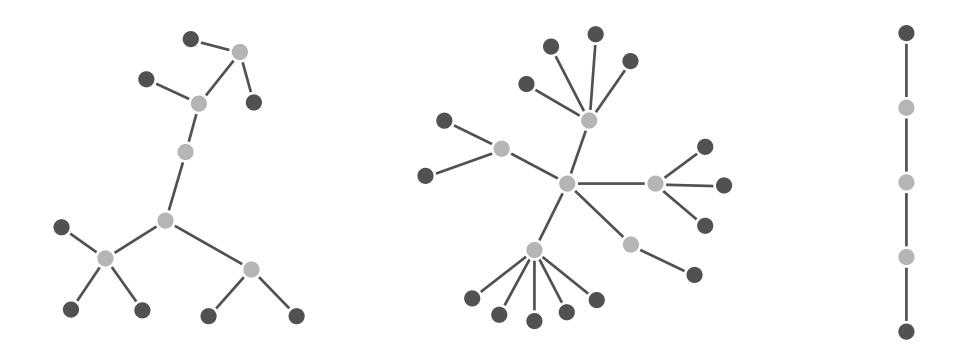
Leaves (degree = 1): present-day species

Internal nodes

(degree \geq 2): ancestral species

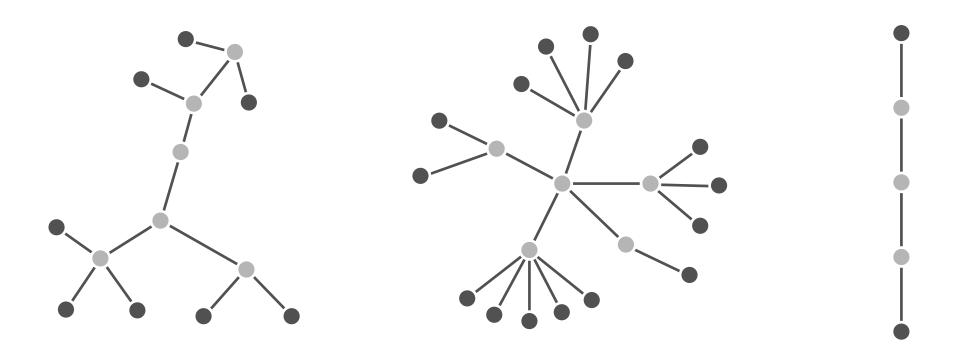
Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.

Trees

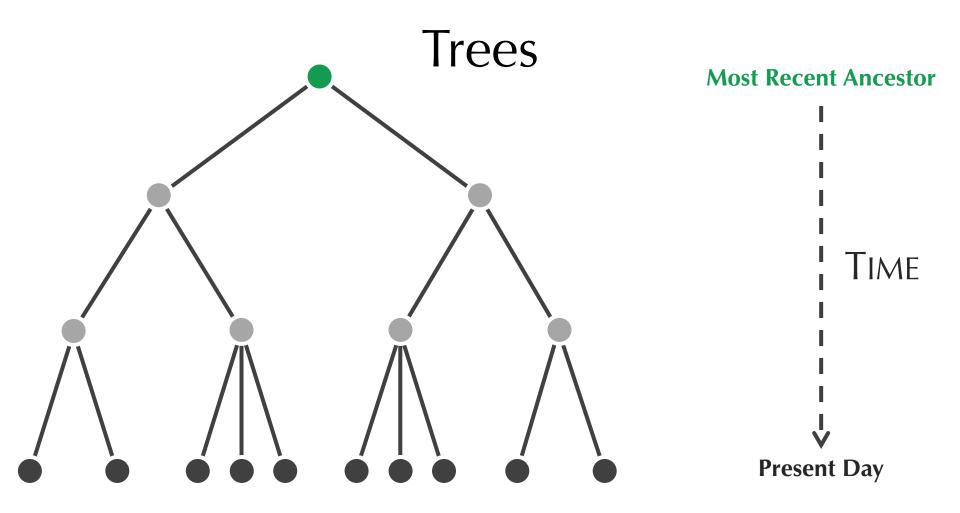


Note: We proved in a previous lecture that every tree with n nodes has exactly n-1 edges.

Trees



Exercise: Prove that there is a unique path connecting any two nodes in a tree.

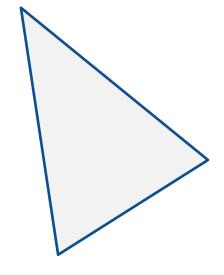


Rooted tree: one node is designated as the **root** (most recent common ancestor).

Definition of a Distance Matrix

Distance matrix: A matrix *D* representing distances between pairs of *n* organisms that satisfies three properties:

- **1. Symmetry:** $D_{i,j} = D_{i,j}$ for all pairs i, j
- **2. Non-negativity:** $D_{i,j} >= 0$ for all pairs i, j
- **3. Triangle inequality:** For all i, j, and k, $D_{i,j} + D_{j,k} >= D_{i,k}$.



SPECIES ALIGNMENT

Chimp ACGTAGGCCT

Human ATGTAAGACT

Seal TCGAGAGCAC

Whale TCGAAAGCAT

 $D_{i,j}$ = number of differing symbols between *i*-th and *j*-th rows of a multiple alignment.

SPECIES	ALIGNMENT	Distance Matrix			
		Chimp	Human	Seal	Whale
Chimp	ACGTAGGCCT	0	3	6	4
Human	ATGTAAGACT	3	0	7	5
Seal	TCGAGAGCAC	6	7	0	2
Whale	TCGAAAGCAT	4	5	2	0

 $D_{i,j}$ = number of differing symbols between i-th and j-th rows of a multiple alignment.

SPECIES	ALIGNMENT	Distance Matrix			
		Chimp	Human	Seal	Whale
Chimp	ACGTAGGCCT	0	3	6	4
Human	A T GTA A G A CT	3	0	7	5
Seal	TCGAGAGCAC	6	7	0	2
Whale	TCGAAAGCAT	4	5	2	0

Exercise: Prove that for any multiple sequence alignment, this way of defining *D* produces a distance matrix.

SPECIES	ALIGNMENT	Distance Matrix			
		Chimp	Human	Seal	Whale
Chimp	ACGTAGGCCT	0	3	6	4
Human	A T GTA A G A CT	3	0	7	5
Seal	TCGAGAGCAC	6	7	O	2
Whale	TCGAAAGCAT	4	5	2	0

Distance-Based Phylogeny

Distance-Based Phylogeny Problem.

- Input: A distance matrix.
- Output: The unrooted tree "fitting" this distance matrix.

Distance-Based Phylogeny

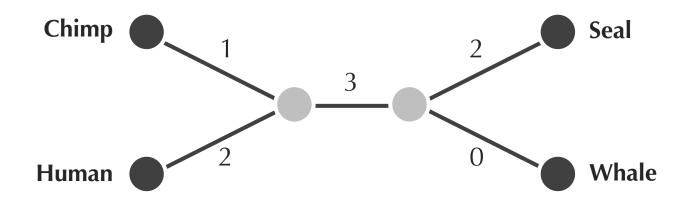
Distance-Based Phylogeny Problem.

- Input: A distance matrix.
- Output: The unrooted tree "fitting" this distance matrix.

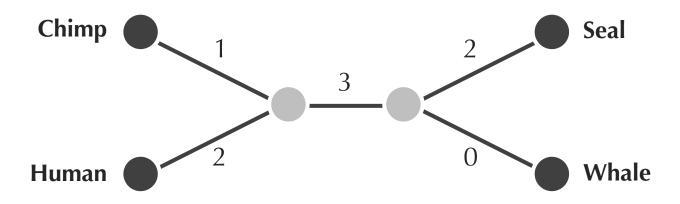
Of course, we are getting a bit ahead of ourselves – we should define what we mean by "fitting"!

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

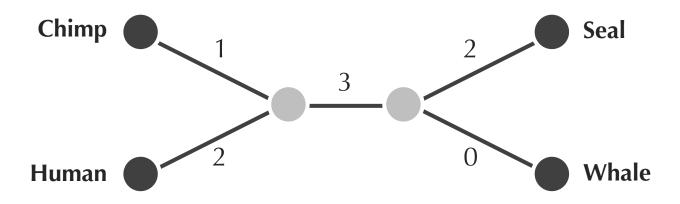


	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



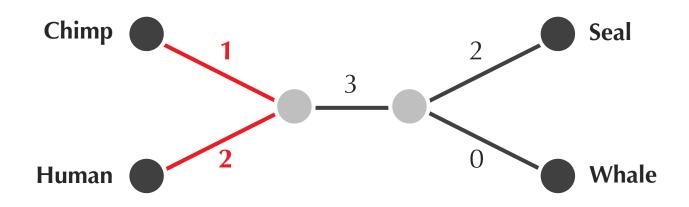
 $d_{i,j}(T)$ = distance between nodes i and j in tree T, computed by summing edge weights from i to j.

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

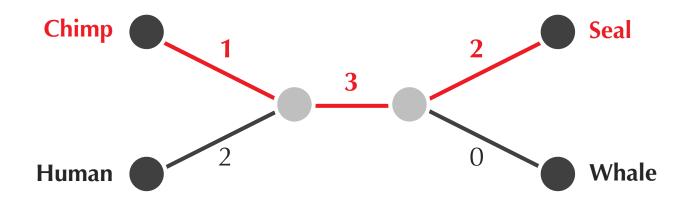


We say that T fits matrix D if for every pair i and j, $d_{i,j}(T) = D_{i,j}$.

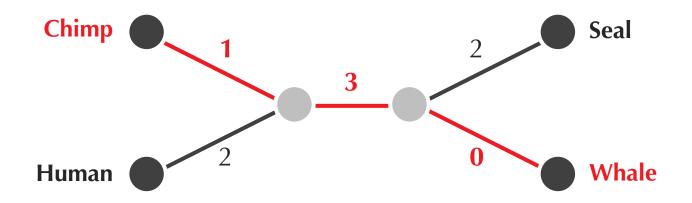
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



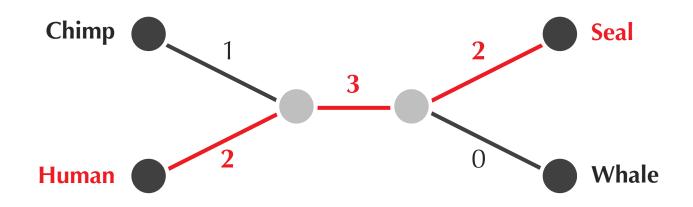
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



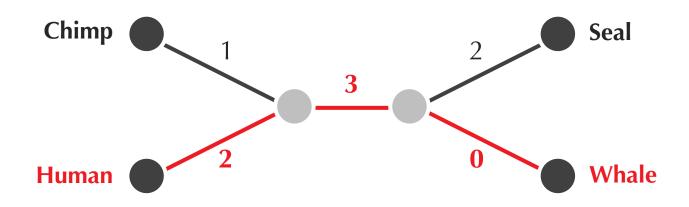
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



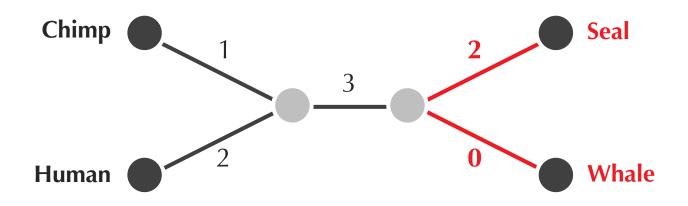
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	O



	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



Return to Distance-Based Phylogeny

Exercise: Find a tree fitting the following matrix.

```
v_1 v_2 v_3 v_4

v_1 0 3 4 3

v_2 3 0 4 5

v_3 4 4 0 2

v_4 3 5 2 0
```

Sometimes, No Tree Fits a Matrix

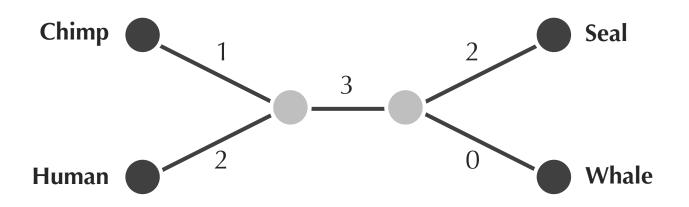
Exercise: Find a tree fitting the following matrix.

$$v_1$$
 v_2 v_3 v_4
 v_1 0 3 4 3
 v_2 3 0 4 5
 v_3 4 4 0 2
 v_4 3 5 2 0

Additive matrix: distance matrix such that there exists an unrooted tree fitting it.

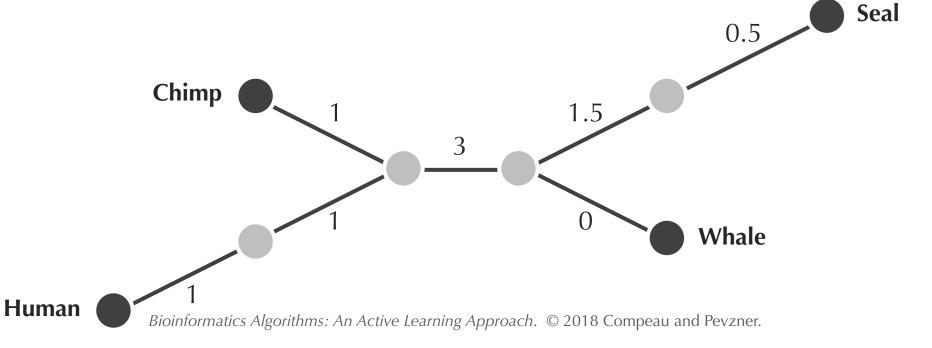
Sometimes, More Than One Tree Fits a Matrix

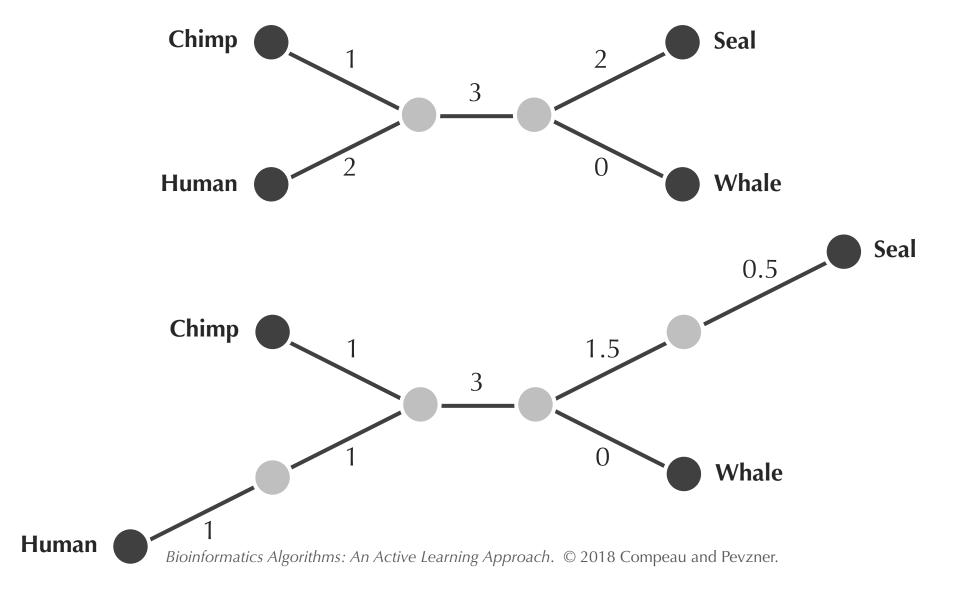
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

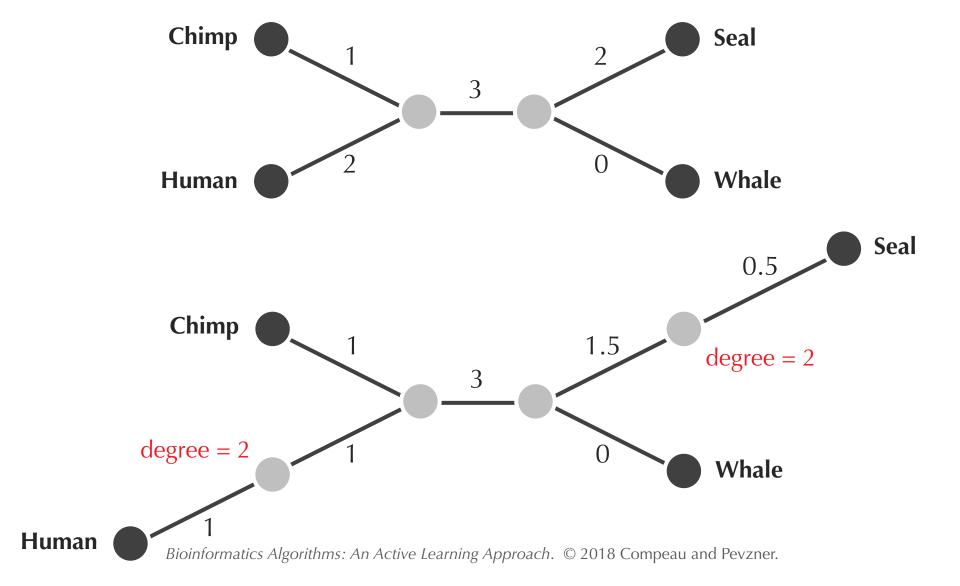


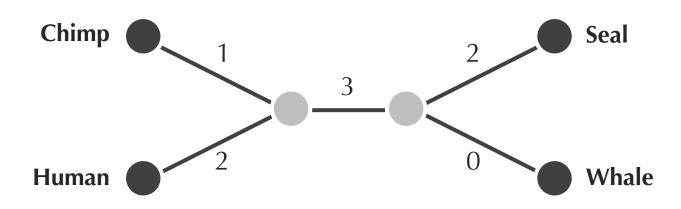
Sometimes, More Than One Tree Fits a Matrix

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

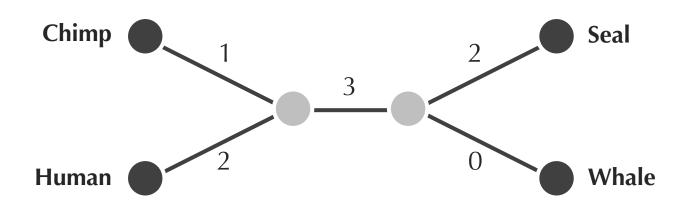








Simple tree: tree with no nodes of degree 2.



Simple tree: tree with no nodes of degree 2.

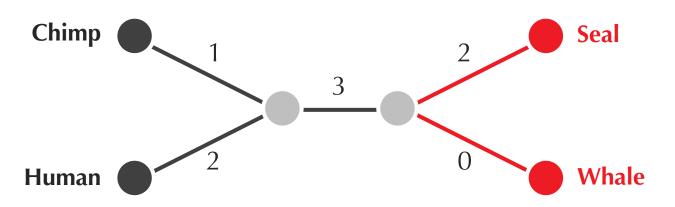
Theorem: There is a unique *simple* tree fitting an additive matrix.

Reformulating Distance-Based Phylogeny

Distance-Based Phylogeny Problem: Construct an evolutionary tree from a distance matrix.

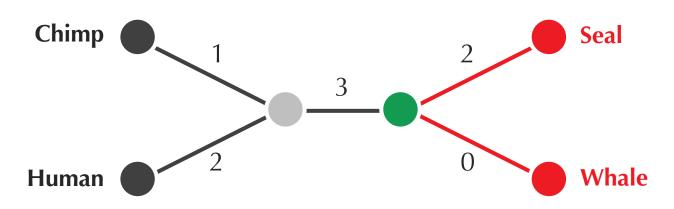
- Input: A distance matrix.
- Output: The simple tree fitting this distance matrix (if this matrix is additive).

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.

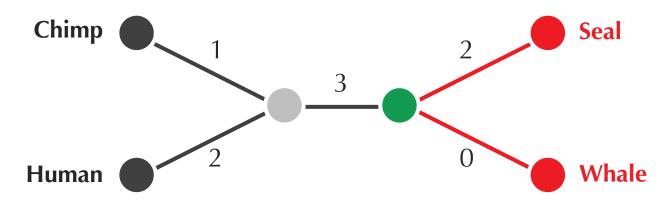
Seal and whale are **neighbors** (meaning they share the same **parent**).



Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.

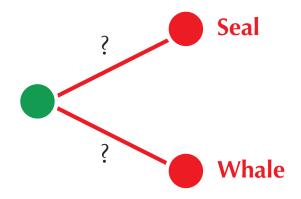
Seal and whale are **neighbors** (meaning they share the same **parent**).

Theorem: Every simple tree with at least three leaves has at least one pair of neighboring leaves.



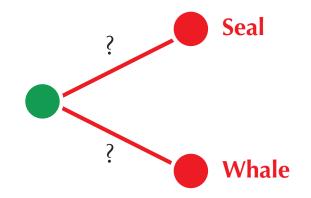
Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.

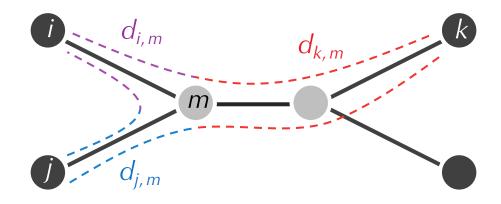
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

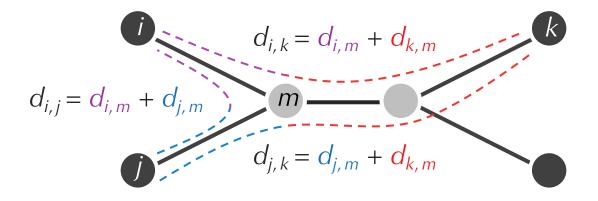


	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

Key Point: How do we compute the unknown distances?





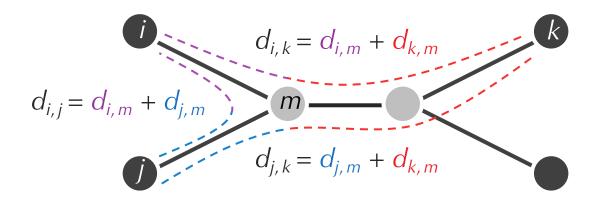


$$d_{i,k} = d_{i,m} + d_{k,m}$$

$$d_{i,k} = d_{i,m} + d_{k,m}$$

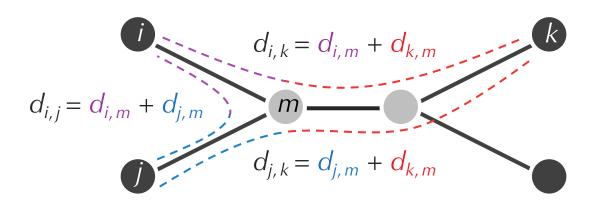
$$d_{j,k} = d_{j,m} + d_{k,m}$$

$$d_{k,m} = \left[(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m}) \right] / 2$$



$$d_{k,m} = [(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m})] / 2$$

$$d_{k,m} = (d_{i,k} + d_{j,k} - d_{i,j}) / 2$$



$$d_{k,m} = [(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m})] / 2$$

$$d_{k,m} = (d_{i,k} + d_{j,k} - d_{i,j}) / 2$$

$$d_{k,m} = (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

$$d_{i,k} = d_{i,m} + d_{k,m}$$

$$d_{i,j} = d_{i,m} + d_{j,m}$$

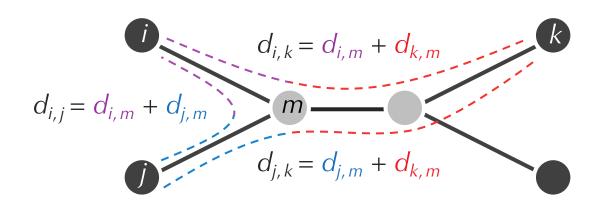
$$d_{j,k} = d_{j,m} + d_{k,m}$$

$$d_{k,m} = [(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m})] / 2$$

$$d_{k,m} = (d_{i,k} + d_{j,k} - d_{i,j}) / 2$$

$$d_{k,m} = (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

$$d_{i,m} = D_{i,k} - (D_{i,k} + D_{i,k} - D_{i,j}) / 2$$



$$d_{k,m} = [(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m})] / 2$$

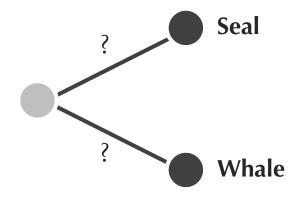
$$d_{k,m} = (d_{i,k} + d_{j,k} - d_{i,j}) / 2$$

$$d_{k,m} = (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

$$d_{i,m} = D_{i,k} - (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

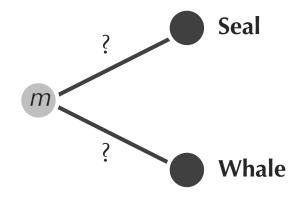
$$d_{i,m} = (D_{i,k} + D_{i,i} - D_{i,k}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



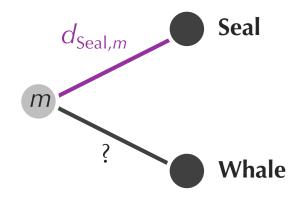
$$d_{i,m} = (D_{i,k} + D_{i,j} - D_{j,k}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



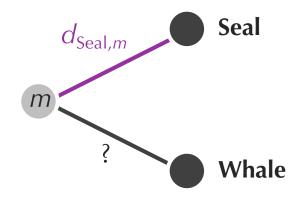
$$d_{i,m} = (D_{i,k} + D_{i,j} - D_{j,k}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



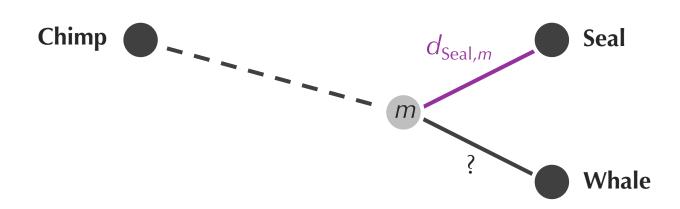
$$d_{\text{Seal},m} = (D_{\text{Seal},k} + D_{\text{Seal},j} - D_{j,k}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



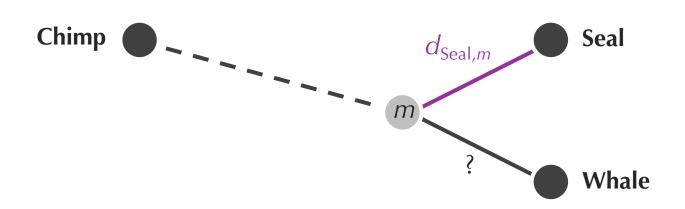
$$d_{\text{Seal,m}} = (D_{\text{Seal,k}} + D_{\text{Seal,Whale}} - D_{\text{Whale,k}}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



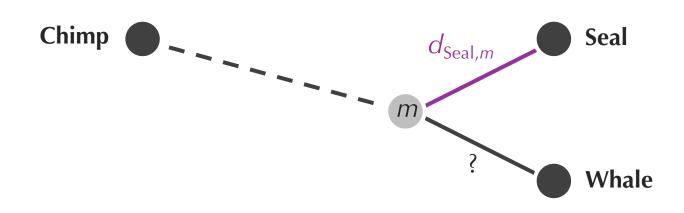
$$d_{\text{Seal,m}} = (D_{\text{Seal,Chimp}} + D_{\text{Seal,Whale}} - D_{\text{Whale,Chimp}}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



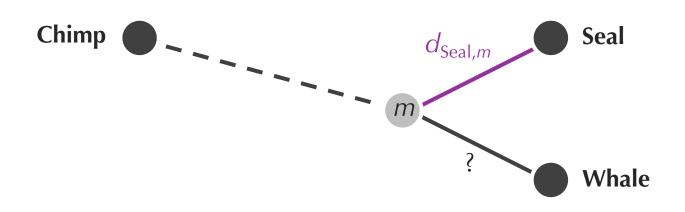
$$d_{\text{Seal,m}} = (6 + D_{\text{Seal,Whale}} - D_{\text{Whale,Chimp}}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



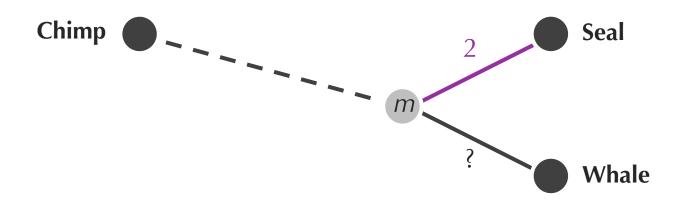
$$d_{\text{Seal},m} = (6 + 2 - D_{\text{Whale,Chimp}}) / 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



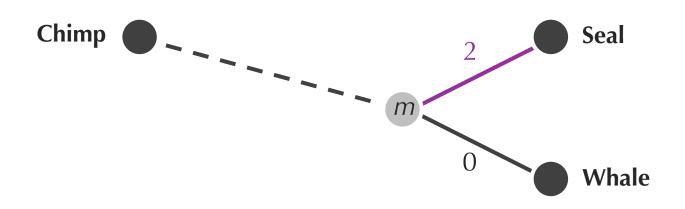
$$d_{\text{Seal},m} = (6 + 2 - 4)/2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



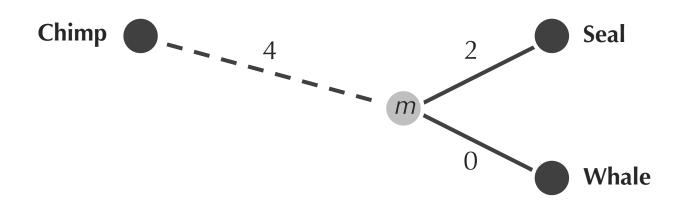
$$d_{\text{Seal},m} = 2$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	O

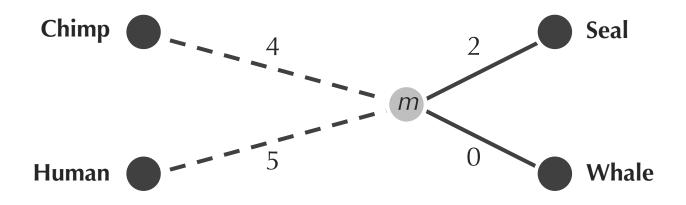


$$d_{\text{Seal},m} = 2$$

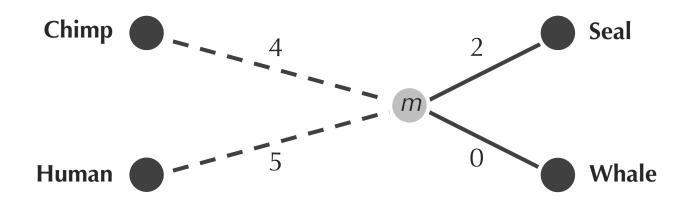
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



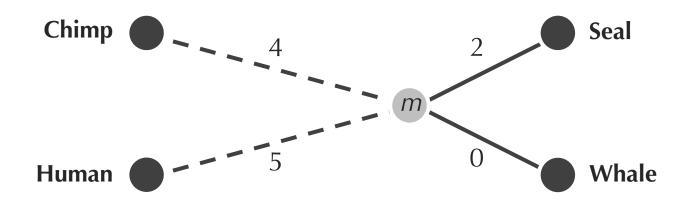
	Chimp	Human	Seal	Whale	m
Chimp	0	3	6	4	4
Human	3	0	7	5	5
Seal	6	7	0	2	2
Whale	4	5	2	0	0
m	4	5	2	0	0



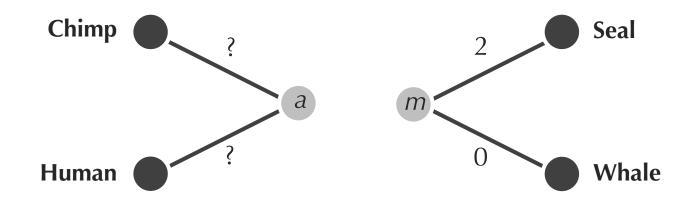
	Chimp	Human	Seal	Whale	m
Chimp	0	3	6	4	4
Human	3	0	7	5	5
Seal	6	7	0	2	2
Whale	4	5	2	0	0
m	4	5	2	0	0



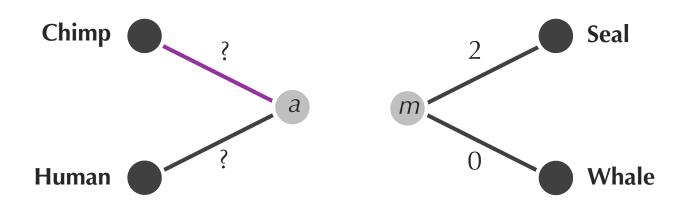
	Chimp	Human	m
Chimp	0	3	4
Human	3	0	5
m	4	5	0



	Chimp	Human	m
Chimp	0	3	4
Human	3	0	5
m	4	5	0

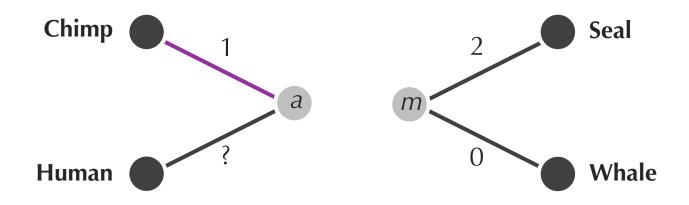


	Chimp	Human	m
Chimp	0	3	4
Human	3	0	5
m	4	5	0



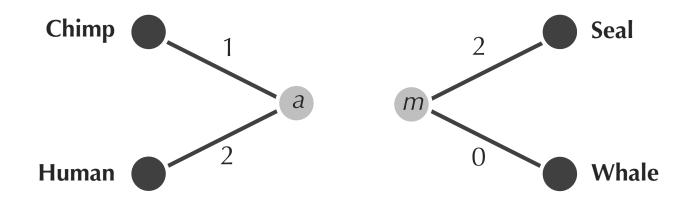
$$d_{\text{Chimp},a} = (D_{\text{Chimp},m} + D_{\text{Chimp},\text{Human}} - D_{\text{Human},m}) / 2$$

	Chimp	Human	m
Chimp	0	3	4
Human	3	0	5
m	4	5	0

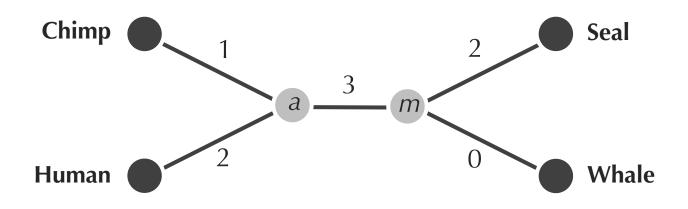


$$d_{\text{Chimp},a} = 1$$

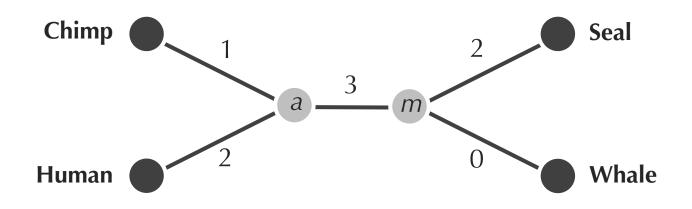
	Chimp	Human	m
Chimp	0	3	4
Human	3	0	5
m	4	5	0



	Chimp	Human	m
Chimp	0	3	4
Human	3	0	5
m	4	5	0



	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0



```
0
1
2
3
0
0
13
21
22
1
13
0
12
13
2
21
12
0
13
3
22
13
13
0
```

Exercise: Apply this recursive approach to this distance matrix.

```
v_1 v_2 v_3 v_4

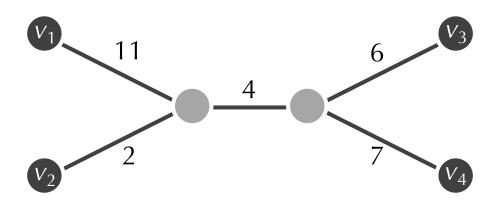
v_1 0 13 21 22

v_2 13 0 12 13

v_3 21 12 0 13

v_4 22 13 13 0
```

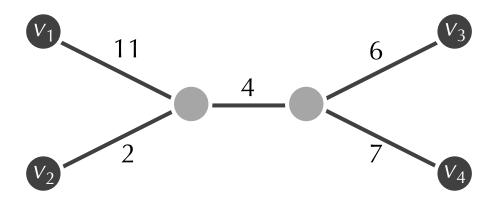
$$v_1$$
 v_2 v_3 v_4 v_1 0132122 v_2 1301213 v_3 2112013 v_4 2213130



Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.

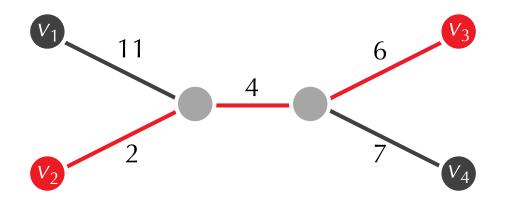
```
v_1v_2v_3v_4v_10132122v_21301213v_32112013v_42213130
```

minimum element is **D**_{2,3}



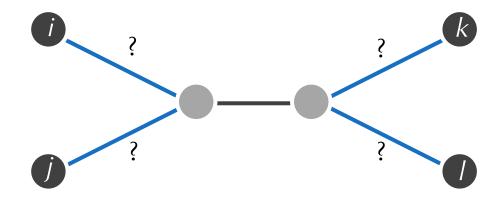
```
v_1v_2v_3v_4v_10132122v_21301213v_32112013v_42213130
```

minimum element is **D**_{2,3}



v₂ and v₃ are
not neighbors!

From Neighbors to Limbs



Rather than trying to infer **neighbors**, let's instead try to compute the length of **limbs**, the edges attached to leaves.

From Neighbors to Limbs

$$d_{i,k} = d_{i,m} + d_{k,m}$$

$$d_{i,k} = d_{i,m} + d_{k,m}$$

$$d_{j,k} = d_{j,m} + d_{k,m}$$

$$d_{k,m} = [(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m})] / 2$$

$$d_{k,m} = (d_{i,k} + d_{j,k} - d_{i,j}) / 2$$

$$d_{k,m} = (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

$$d_{i,m} = D_{i,k} - (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

$$d_{i,m} = (D_{i,k} + D_{i,j} - D_{j,k}) / 2$$

Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.

From Neighbors to Limbs

$$d_{i,k} = d_{i,m} + d_{k,m}$$

$$d_{i,j} = d_{i,m} + d_{j,m}$$

$$d_{j,k} = d_{j,m} + d_{k,m}$$

$$d_{k,m} = \left[(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m}) \right] / 2$$

$$d_{k,m} = (d_{i,k} + d_{j,k} - d_{i,j}) / 2$$

$$d_{k,m} = (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

$$\therefore d_{i,m} = D_{i,k} - (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$

$$d_{i,m} = (D_{i,k} + D_{i,j} - D_{j,k}) / 2$$

Assumes that *i* and *j* are *neighbors*...

Bioinformatics Algorithms: An Active Learning Approach. © 2018 Compeau and Pevzner.

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

$$(D_{\text{chimp, human}} + D_{\text{chimp, seal}} - D_{\text{human, seal}}) / 2 = (3 + 6 - 7) / 2 = 1$$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

$$(D_{\text{chimp, human}} + D_{\text{chimp, seal}} - D_{\text{human, seal}}) / 2 = (3 + 6 - 7) / 2 = 1$$

 $(D_{\text{chimp, human}} + D_{\text{chimp, whale}} - D_{\text{human, whale}}) / 2 = (3 + 4 - 5) / 2 = 1$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

$$(D_{\rm chimp,\ human} + D_{\rm chimp,\ seal} - D_{\rm human,\ seal}) \ / \ 2 = (3 + 6 - 7) \ / \ 2 = 1$$

$$(D_{\rm chimp,\ human} + D_{\rm chimp,\ whale} - D_{\rm human,\ whale}) \ / \ 2 = (3 + 4 - 5) \ / \ 2 = 1$$

$$(D_{\rm chimp,\ whale} + D_{\rm chimp,\ seal} - D_{\rm whale,\ seal}) \ / \ 2 = (6 + 4 - 2) \ / \ 2 = 4$$

$$(D_{\rm chimp,\ whale} + D_{\rm chimp,\ seal} - D_{\rm whale,\ seal}) \ / \ 2 = (6 + 4 - 2) \ / \ 2 = 4$$

$$(D_{\rm chimp,\ whale} + D_{\rm chimp,\ seal} - D_{\rm whale,\ seal}) \ / \ 2 = (6 + 4 - 2) \ / \ 2 = 4$$

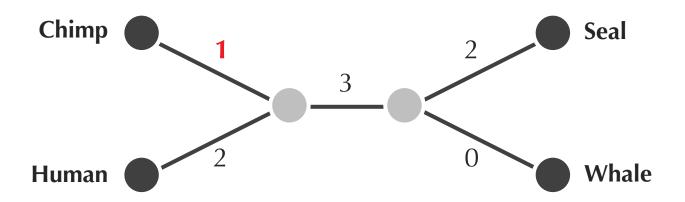
$$(D_{\rm chimp,\ whale} + D_{\rm chimp,\ seal} - D_{\rm whale,\ seal}) \ / \ 2 = (6 + 4 - 2) \ / \ 2 = 4$$

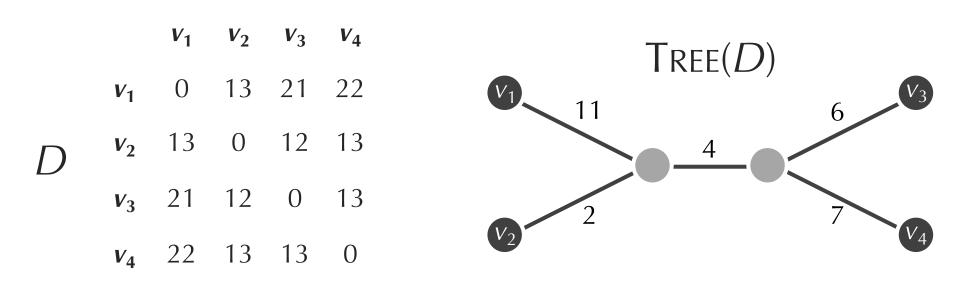
	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0

$$(D_{\text{chimp, human}} + D_{\text{chimp, seal}} - D_{\text{human, seal}}) / 2 = (3 + 6 - 7) / 2 = 1$$

 $(D_{\text{chimp, human}} + D_{\text{chimp, whale}} - D_{\text{human, whale}}) / 2 = (3 + 4 - 5) / 2 = 1$
 $(D_{\text{chimp, whale}} + D_{\text{chimp, seal}} - D_{\text{whale, seal}}) / 2 = (6 + 4 - 2) / 2 = 4$

	Chimp	Human	Seal	Whale
Chimp	0	3	6	4
Human	3	0	7	5
Seal	6	7	0	2
Whale	4	5	2	0





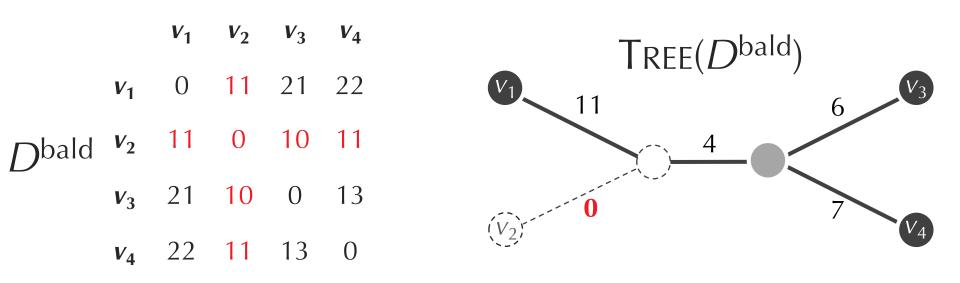
$$v_1$$
 v_2 v_3 v_4
 v_1 0 13 21 22
 v_2 13 0 12 13
 v_3 21 12 0 13
 v_4 22 13 13 0

1. Pick an arbitrary leaf j (say, $j = v_2$).

$$v_1$$
 v_2 v_3 v_4
 v_1 0 13 21 22
 v_2 13 0 12 13
 v_3 21 12 0 13
 v_4 22 13 13 0

LimbLength(v_2) = 2

2. Compute its limb length, *LimbLength(j)*.



3. Subtract LimbLength(j) from each row and column to produce D^{bald} in which j is a **bald limb** (length 0).

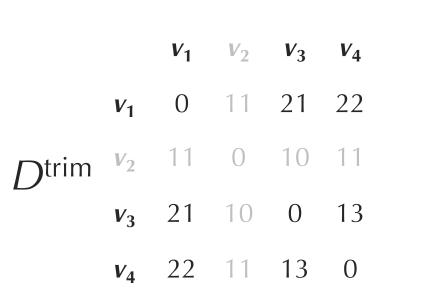
$$v_1$$
 v_2 v_3 v_4
 v_1 0 11 21 22

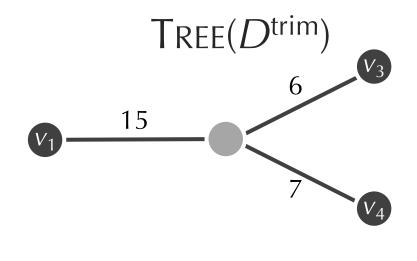
Dtrim v_2 11 0 10 11

 v_3 21 10 0 13

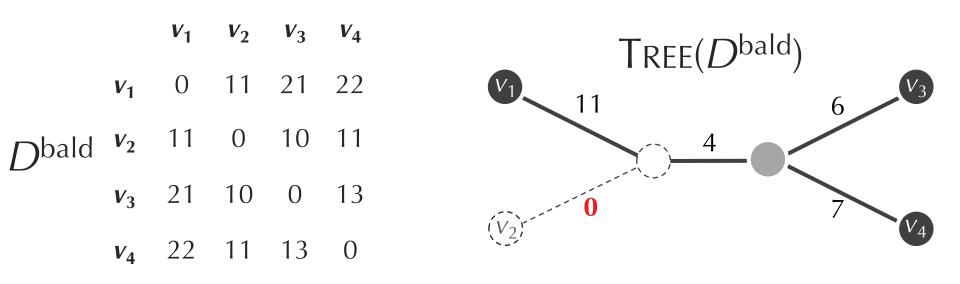
 v_4 22 11 13 0

4. Remove the *j*-th row and column of the matrix to form the $(n - 1) \times (n - 1)$ matrix D^{trim} .

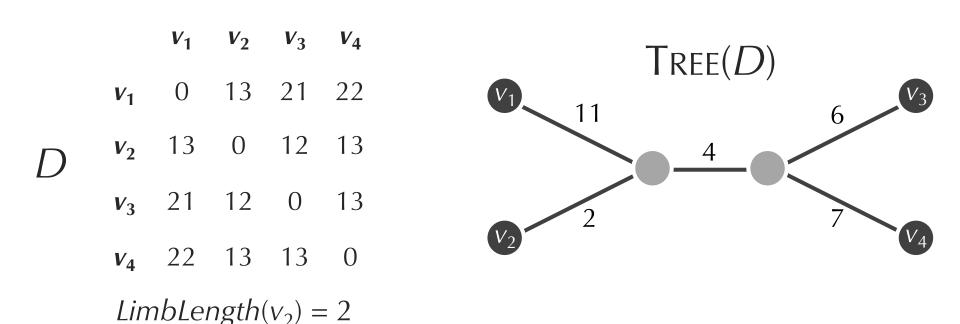




5. Call AdditivePhylogeny recursively to construct *Tree*(*D*^{trim}).



6. Identify the point in $Tree(D^{trim})$ where leaf j should be attached.



7. Attach *j* by an edge of length *LimbLength*(*j*) in order to form *Tree*(*D*).

AdditivePhylogeny

AdditivePhylogeny(*D*):

- 1. Pick an arbitrary leaf *j*.
- 2. Compute its limb length, *LimbLength(j)*.
- 3. Subtract LimbLength(j) from each row and column to produce D^{bald} in which j is a bald limb (length 0).
- 4. Remove the *j*-th row and column of the matrix to form the $(n 1) \times (n 1)$ matrix D^{trim} .
- 5. Recursively call **AdditivePhylogeny**(D^{trim}) to obtain $Tree(D^{\text{trim}})$.
- 6. Identify the point in $Tree(D^{trim})$ where leaf j should be attached.
- 7. Attach j by an edge of length LimbLength(j) in order to form Tree(D).

Attaching a leaf to the tree after the recursive step is the difficult part ...

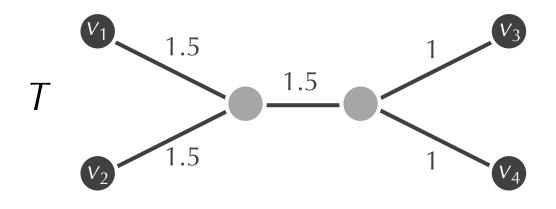
AdditivePhylogeny

AdditivePhylogeny(*D*):

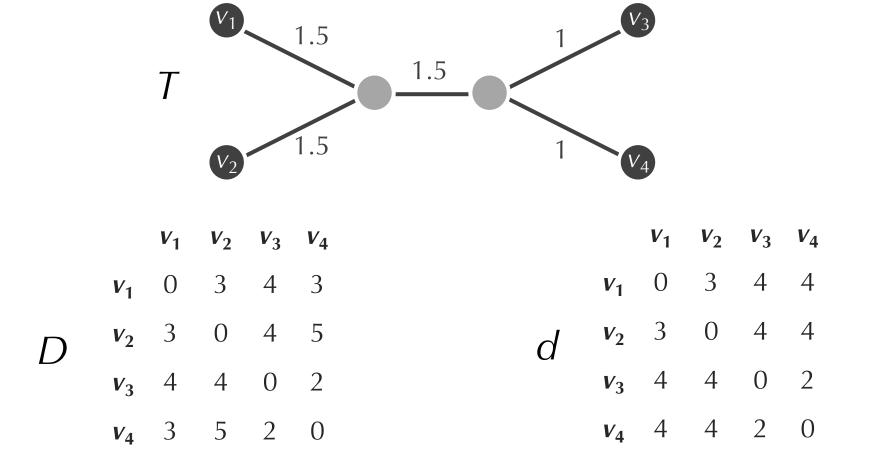
- 1. Pick an arbitrary leaf *j*.
- 2. Compute its limb length, *LimbLength(j)*.
- 3. Subtract LimbLength(j) from each row and column to produce D^{bald} in which j is a bald limb (length 0).
- 4. Remove the *j*-th row and column of the matrix to form the $(n 1) \times (n 1)$ matrix D^{trim} .
- 5. Recursively call **AdditivePhylogeny**(D^{trim}) to obtain $Tree(D^{trim})$.
- 6. Identify the point in $Tree(D^{trim})$ where leaf j should be attached.
- 7. Attach j by an edge of length LimbLength(j) in order to form Tree(D).

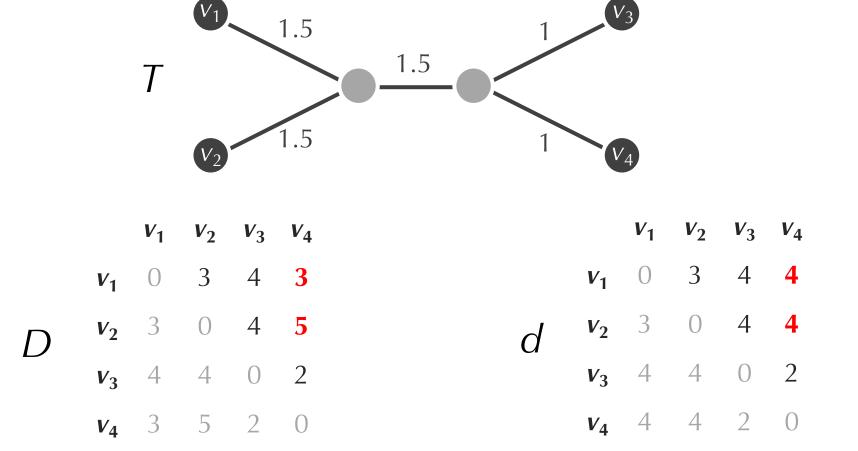
Attaching a leaf to the tree after the recursive step is the difficult part ...

... and what do we do about non-additive matrices?

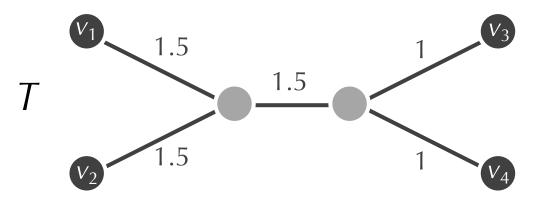


$$v_1$$
 v_2 v_3 v_4
 v_1 0 3 4 3
 v_2 3 0 4 5
 v_3 4 4 0 2
 v_4 3 5 2 0





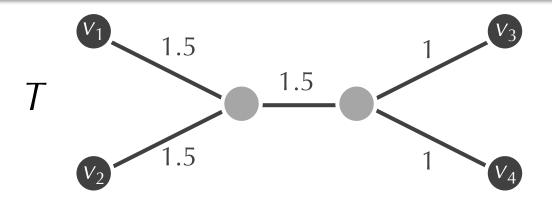
$$Discrepancy(T, D) = \sum_{1 \leq i < j \leq n} (d_{i,j}(T) - D_{i,j})^2$$



		V_1	V_2	V_3	V_4		ı	[′] 1
	V_1	0	3	4	3	<i>V</i> ₁	0	
)	V_2	3	0	4	5	V_2	3	
			4				4	
	V_4	3	5	2	0	V_4	4	

Discrepancy(T, D) =
$$\sum_{1 \le i < j \le n} (d_{i,j}(T) - D_{i,j})^2$$

= $1^2 + 1^2 = 2$





Least-Squares Phylogeny

Least-Squares Distance-Based Phylogeny Problem:

Given a distance matrix, find the tree that minimizes the sum of squared errors.

- Input: An n x n distance matrix D.
- Output: A weighted tree T with n leaves
 minimizing Discrepancy(T, D) over all weighted
 trees with n leaves.

Least-Squares Phylogeny

Least-Squares Distance-Based Phylogeny Problem:

Given a distance matrix, find the tree that minimizes the sum of squared errors.

- Input: An n x n distance matrix D.
- Output: A weighted tree T with n leaves
 minimizing Discrepancy(T, D) over all weighted
 trees with n leaves.

Unfortunately, this problem is *NP*-Complete...

The Critical Insight

Our first idea of using neighbors to construct the tree was a good one! The problem was not with our idea, but with the distance matrix itself!

Given an $n \times n$ distance matrix D, its **neighbor-joining** matrix is the matrix D^* defined as

 $D^*_{i,j} = (n-2) \cdot D_{i,j} - TotalDistance_D(i) - TotalDistance_D(j)$

Given an $n \times n$ distance matrix D, its **neighbor-joining** matrix is the matrix D^* defined as

$$D^*_{i,j} = (n-2) \cdot D_{i,j} - TotalDistance_D(i) - TotalDistance_D(j)$$

Given an $n \times n$ distance matrix D, its **neighbor-joining** matrix is the matrix D^* defined as

$$D^*_{i,j} = (n-2) \cdot D_{i,j} - TotalDistance_D(i) - TotalDistance_D(j)$$

		<i>v</i> ₁	v_2	v_3	V_4	TotalDistance _D
D	v_1	0	13	21	22	56
	v_2	13	0	12	13	38
	v_3	21	12	12 0	13	46
	V_4	22	13	13	0	48

Given an $n \times n$ distance matrix D, its **neighbor-joining** matrix is the matrix D^* defined as

$$D^*_{i,j} = (n-2) \cdot D_{i,j} - TotalDistance_D(i) - TotalDistance_D(j)$$

Given an $n \times n$ distance matrix D, its **neighbor-joining** matrix is the matrix D^* defined as

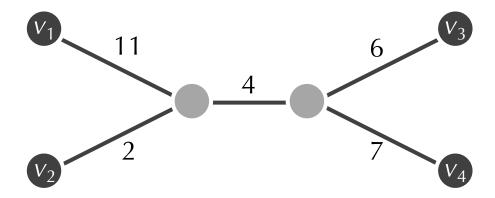
$$D^*_{i,j} = (n-2) \cdot D_{i,j} - TotalDistance_D(i) - TotalDistance_D(j)$$

Given an $n \times n$ distance matrix D, its **neighbor-joining** matrix is the matrix D^* defined as

$$D^*_{i,j} = (n-2) \cdot D_{i,j} - TotalDistance_D(i) - TotalDistance_D(j)$$

Note: What does *D** do to outliers?

$$D = \begin{bmatrix} v_1 & v_2 & v_3 & v_4 \\ v_1 & 0 & 13 & 21 & 22 \\ v_2 & 13 & 0 & 12 & 13 \\ v_3 & 21 & 12 & 0 & 13 \\ v_4 & 22 & 13 & 13 & 0 \end{bmatrix}$$



Neighbor-Joining Theorem: If D is additive, then the smallest element of D^* corresponds to neighboring leaves in Tree(D)!

Neighbor-Joining Theorem: If D is additive, then the smallest element of D^* corresponds to neighboring leaves in Tree(D)!

		v_1	v_2	V_3	V_4	TotalDistance _D
	v_1	0	-68	-60	-60	56
D^*	v_2	-68	0	-60	-60	38
	V_3	-60	-60	0	-68	46
	V_4	-60	-60	-68	0	48

1. Construct neighbor-joining matrix D^* from D.

$$v_1$$
 v_2 v_3 v_4 TotalDistance_D
 v_1 0 -68 -60 -60 56

 D^* v_2 -68 0 -60 -60 38
 v_3 -60 -60 0 -68 46
 v_4 -60 -60 -68 0 48

2. Find a minimum element $D^*_{i,j}$ of D^* .

$$v_1$$
 v_2 v_3 v_4 TotalDistance_D
 v_1 0 -68 -60 -60 56

 D^* v_2 -68 0 -60 -60 38
 v_3 -60 -60 0 -68 46
 v_4 -60 -60 -68 0 48

2. Find a minimum element $D^*_{i,j}$ of D^* .

$$v_1$$
 v_2 v_3 v_4 TotalDistance_D
 v_1 0 -68 -60 -60 56

 D^* v_2 -68 0 -60 -60 38 $\Delta_{i,j} = (56 - 38) / (4 - 2)$
 v_3 -60 -60 0 -68 46 = 9

 v_4 -60 -60 -68 0 48

3. Compute
$$\Delta_{i,j} = (TotalDistance_D(i) - TotalDistance_D(j)) / (n - 2).$$

$$v_1$$
 v_2 v_3 v_4 TotalDistance_D
 v_1 0 13 21 22 56

 v_2 13 0 12 13 38 $\Delta_{i,j} = (56 - 38) / (4 - 2)$
 v_3 21 12 0 13 46 = 9

 v_4 22 13 13 0 48

LimbLength(i) =
$$\frac{1}{2}(13 + 9) = 11$$

LimbLength(i) = $\frac{1}{2}(13 - 9) = 2$

4. Set LimbLength(i) equal to $\frac{1}{2}(D_{i,j} + \Delta_{i,j})$ and LimbLength(j) equal to $\frac{1}{2}(D_{i,j} - \Delta_{i,j})$.

4. Set LimbLength(i) equal to $\frac{1}{2}(D_{i,j} + \Delta_{i,j})$ and LimbLength(j) equal to $\frac{1}{2}(D_{i,j} - \Delta_{i,j})$.

Wait ... where do these formulas come from?

$$m$$
 v_3 v_4 $TotalDistance_D$
 m 0 10 11 21
 D' v_3 10 0 13 23
 v_4 11 13 0 24

5. Form a matrix D' by removing i-th and j-th row/column from D and adding an m-th row/column such that for any k, $D_{k,m} = (D_{i,k} + D_{j,k} - D_{i,j}) / 2$.

Flashback: Computation of $d_{k,m}$

$$d_{i,k} = d_{i,m} + d_{k,m}$$

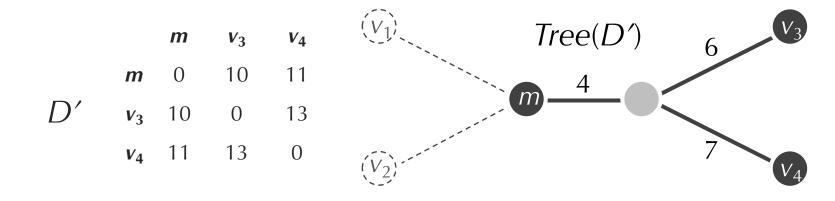
$$d_{i,k} = d_{i,m} + d_{k,m}$$

$$d_{j,k} = d_{j,m} + d_{k,m}$$

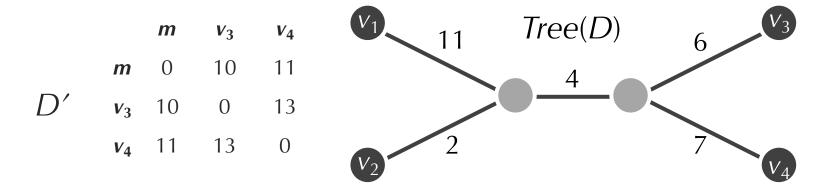
$$d_{k,m} = [(d_{i,m} + d_{k,m}) + (d_{j,m} + d_{k,m}) - (d_{i,m} + d_{j,m})] / 2$$

$$d_{k,m} = (d_{i,k} + d_{j,k} - d_{i,j}) / 2$$

$$d_{k,m} = (D_{i,k} + D_{j,k} - D_{i,j}) / 2$$



6. Apply **NeighborJoining** recursively to D' to obtain Tree(D').



$$LimbLength(v_1) = \frac{1}{2}(13 + 9) = 11$$

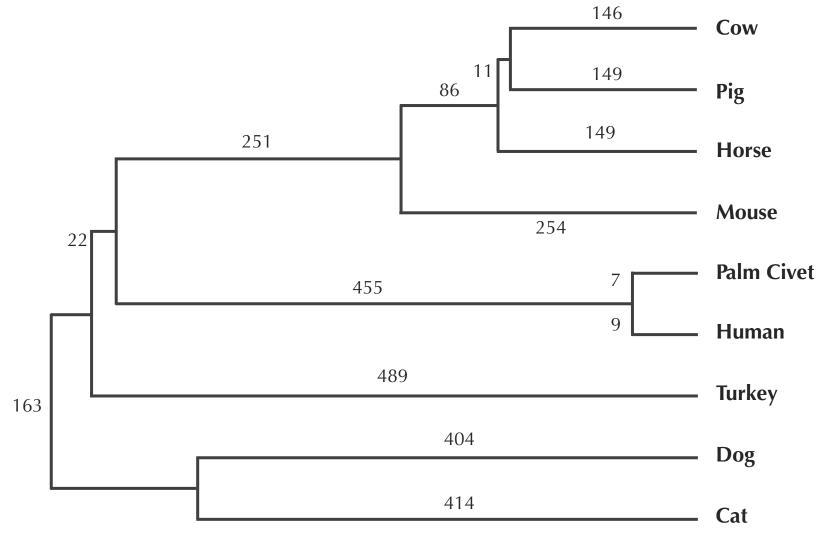
 $LimbLength(v_2) = \frac{1}{2}(13 - 9) = 2$

7. Reattach limbs of i and j to obtain Tree(D).

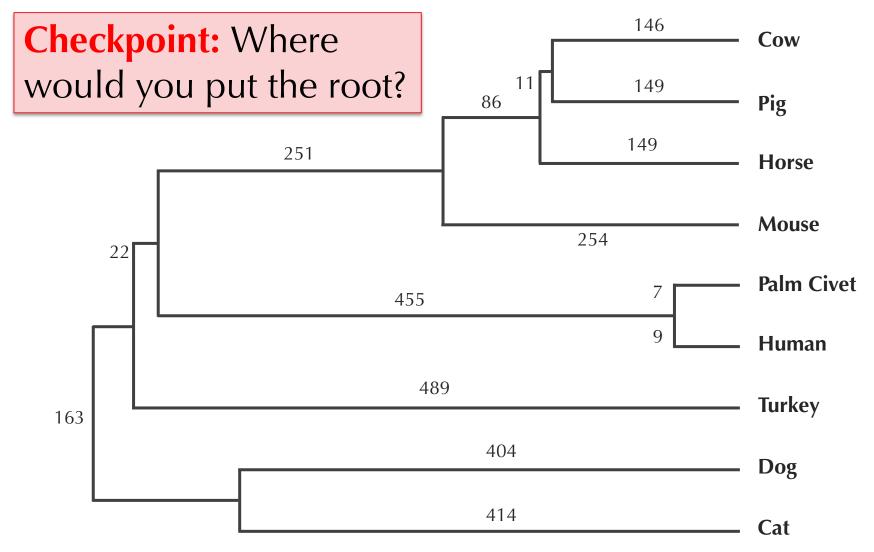
NeighborJoining(*D*):

- 1. Construct neighbor-joining matrix D^* from D.
- 2. Find a minimum element $D^*_{i,j}$ of D^* .
- 3. Compute $\Delta_{i,j} = (TotalDistance_D(i) TotalDistance_D(j)) / (n 2).$
- 4. Set LimbLength(i) equal to $\frac{1}{2}(D_{i,j} + \Delta_{i,j})$ and LimbLength(j) equal to $\frac{1}{2}(D_{i,j} \Delta_{i,j})$.
- 5. Form a matrix D' by removing i-th and j-th row/column from D and adding an m-th row/column such that for any k, $D_{k,m} = (D_{k,i} + D_{k,j} D_{i,j}) / 2$.
- 6. Apply **NeighborJoining** recursively to D' to obtain Tree(D').
- 7. Reattach limbs of i and j to obtain Tree(D).

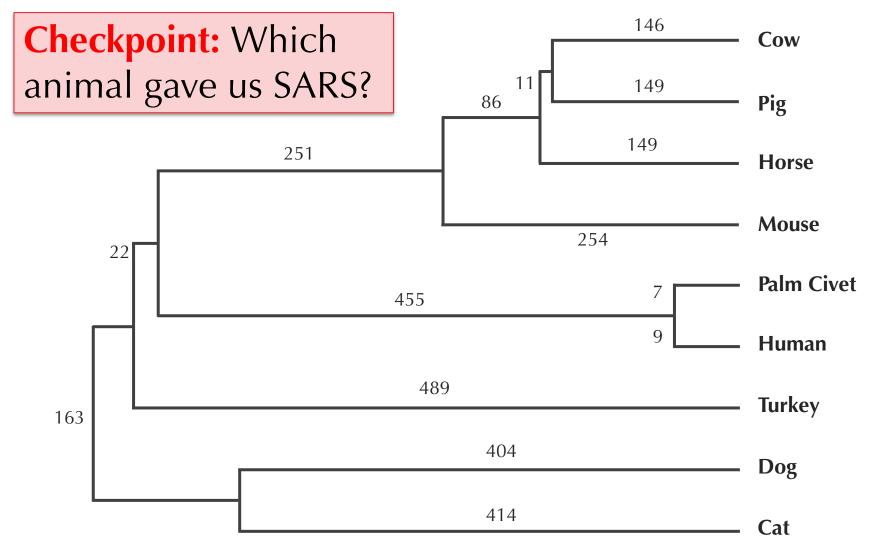
Neighbor-Joining on Coronavirus "Spike Proteins"



Neighbor-Joining on Coronavirus "Spike Proteins"



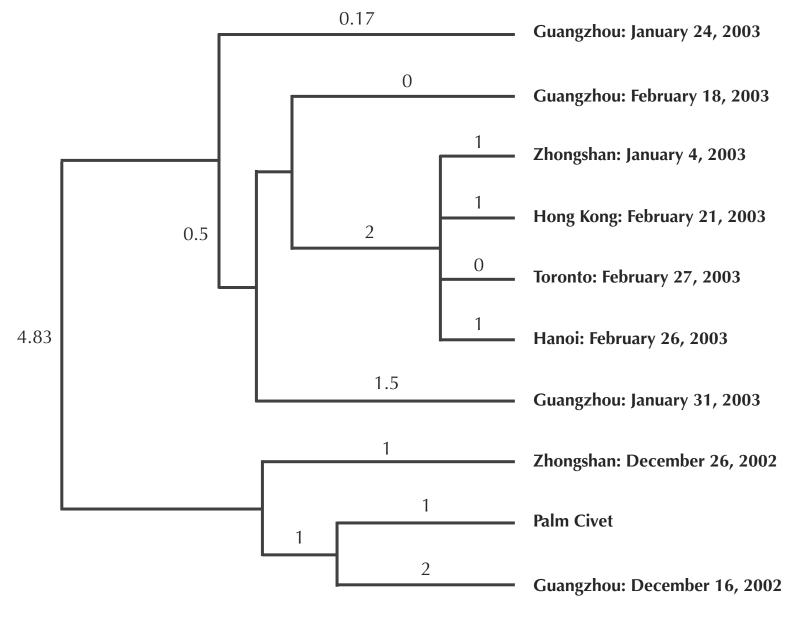
Neighbor-Joining on Coronavirus "Spike Proteins"







Neighbor-Joining on Coronaviruses



Weakness of Distance-Based Methods

Distance-based algorithms for evolutionary tree reconstruction say nothing about ancestral states at internal nodes.

Weakness of Distance-Based Methods

Distance-based algorithms for evolutionary tree reconstruction say nothing about ancestral states at internal nodes.

We *lost* information when we converted a multiple alignment to a distance matrix...

SPECIES	ALIGNMENT	Distance Matrix			
		Chimp	Human	Seal	Whale
Chimp	ACGTAGGCCT	0	3	6	4
Human	ATGTAAGACT	3	0	7	5
Seal	TCGAGAGCAC	6	7	0	2
Whale	TCGAAAGCAT	4	5	2	0

Which Animal Gave Us SARS?

Evolutionary Trees Part 2: The Small Parsimony Algorithm for Inferring Ancestral States

Phillip Compeau and Pavel Pevzner Bioinformatics Algorithms: An Active Learning Approach

©2018 by Compeau and Pevzner. All rights reserved.



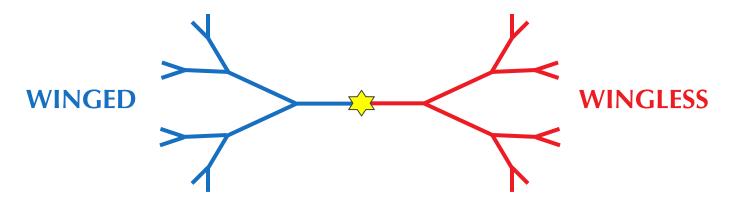
Biologists used to form trees based on anatomical or physiological properties called **characters**.

Biologists used to form trees based on anatomical or physiological properties called **characters**.

Checkpoint: Say you wanted to construct an evolutionary tree of all insects. What might be the first thing you would do?

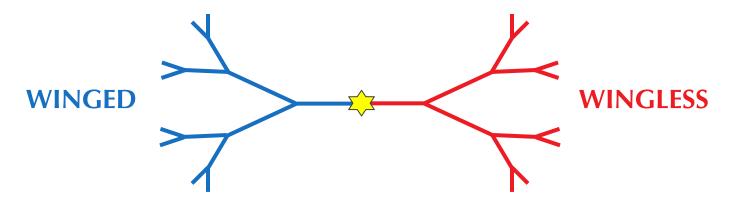
Biologists used to form trees based on anatomical or physiological properties called **characters**.

Checkpoint: Say you wanted to construct an evolutionary tree of all insects. What might be the first thing you would do?

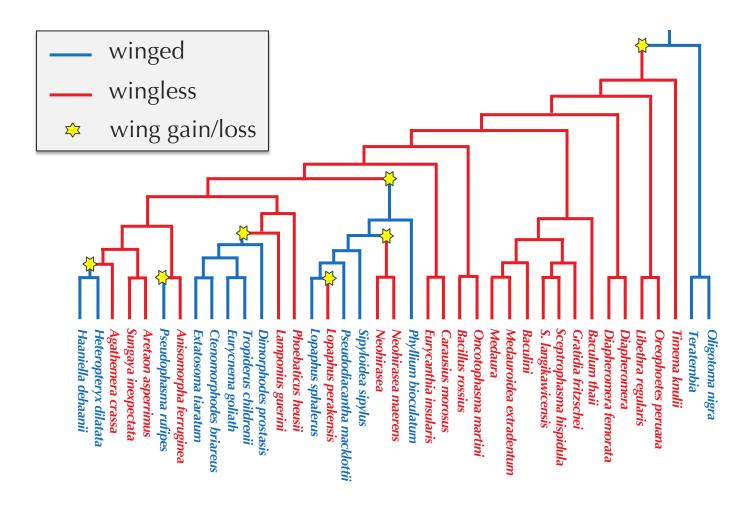


Dollo's principle of irreversibility (1893): evolution doesn't reinvent the same organ (e.g. insect wings).

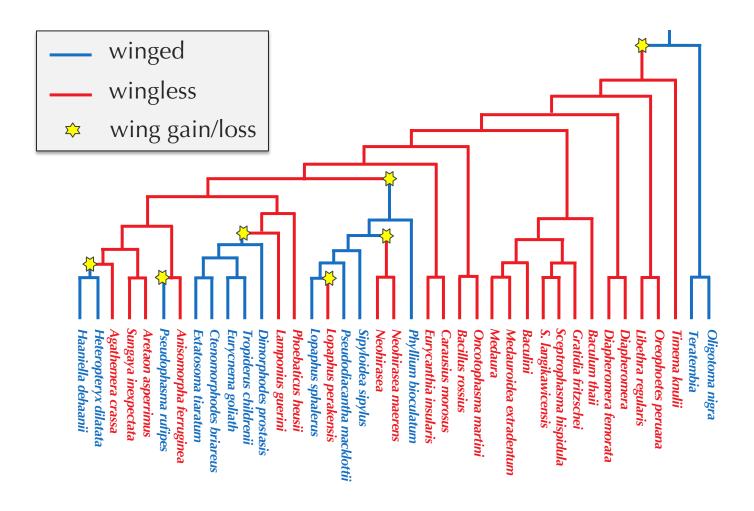
Checkpoint: Say you wanted to construct an evolutionary tree of all insects. What might be the first thing you would do?



Dollo's Principle Violated in Stick Insect Phylogeny

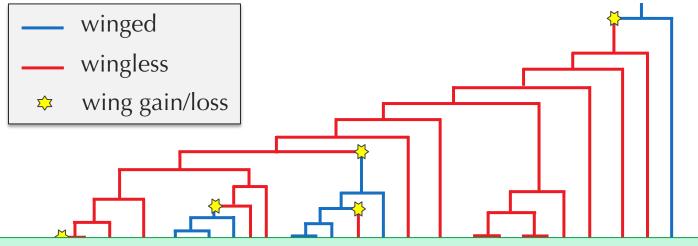


Dollo's Principle Violated in Stick Insect Phylogeny



Checkpoint: What do you think happened?

Dollo's Principle Violated in Stick Insect Phylogeny



Key Point: This approach won't help reconstruct the tree *structure*, but perhaps we can infer *ancestral* characters after constructing a tree using another method (e.g., neighbor-joining).

Checkpoint: What do you think happened?

An Alignment Is a Character Table

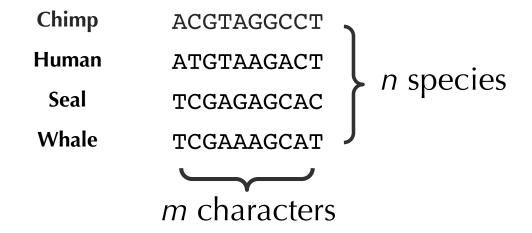
Chimp ACGTAGGCCT

Human ATGTAAGACT

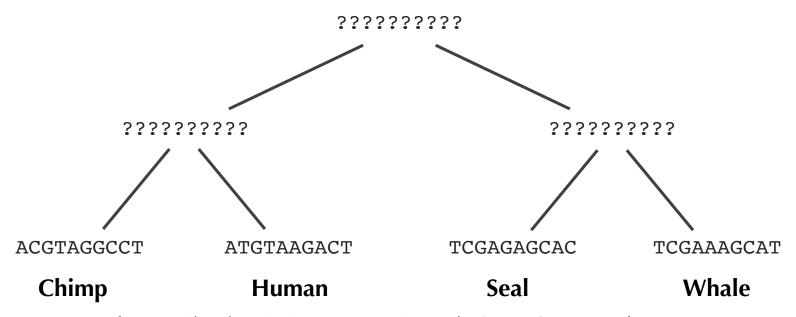
Seal TCGAGAGCAC

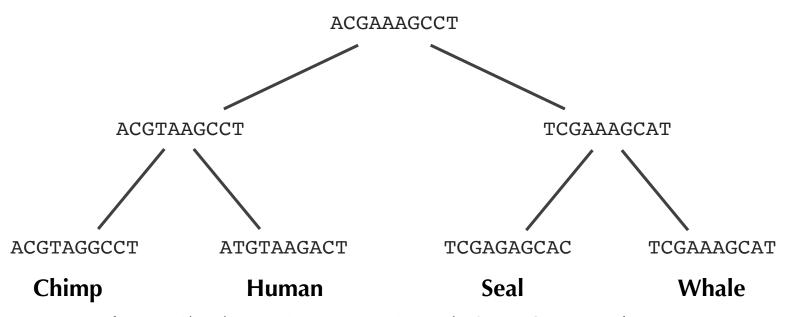
Whale TCGAAAGCAT

An Alignment Is a Character Table

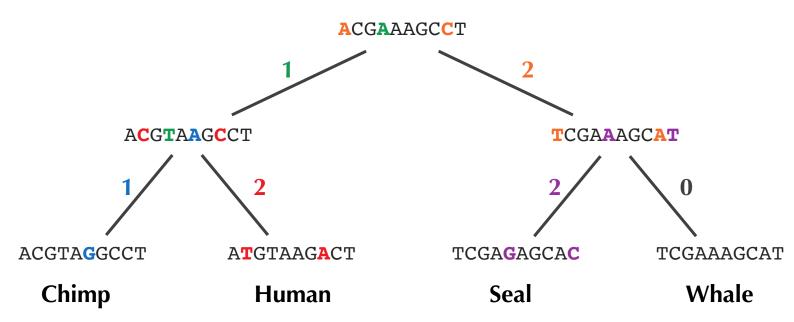


Chimp ACGTAGGCCT
Human ATGTAAGACT
Seal TCGAGAGCAC
Whale TCGAAAGCAT



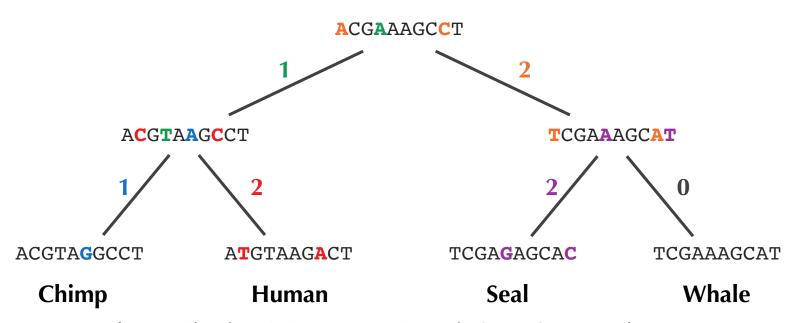


Parsimony score: sum of Hamming distances (total mismatches) along each edge.



Parsimony score: sum of Hamming distances (total mismatches) along each edge.

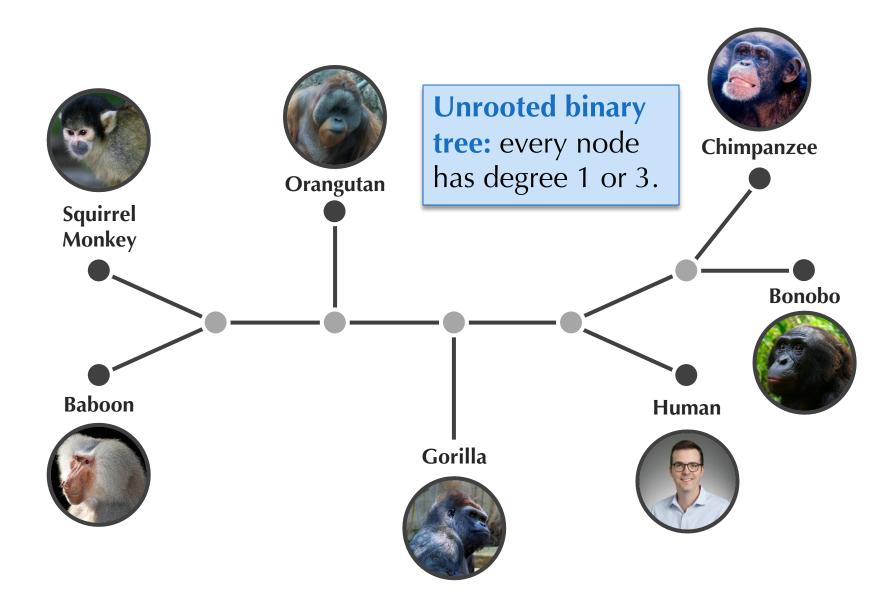
Parsimony Score: 8



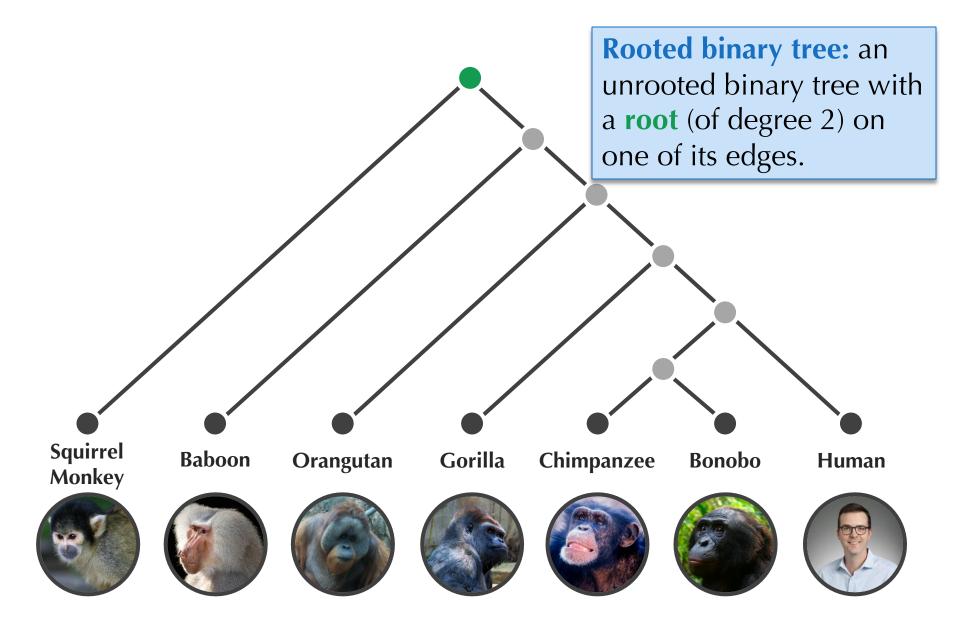
Modeling Speciations

Researchers often assume that all internal nodes correspond to **speciations**, where one species splits into two.

Modeling Speciations



Modeling Speciations



Small Parsimony Problem: Find the most parsimonious labeling of the internal nodes of a rooted tree.

- Input: A rooted binary tree with each leaf labeled by a string of length m.
- Output: A labeling of all other nodes of the tree by strings of length m that minimizes the tree's parsimony score.

Small Parsimony Problem: Find the most parsimonious labeling of the internal nodes of a rooted tree.

- Input: A rooted binary tree with each leaf labeled by a string of length m.
- **Output:** A labeling of all other nodes of the tree by strings of length *m* that minimizes the tree's parsimony score.

Checkpoint: Is there any way we can simplify this problem statement?

Small Parsimony Problem: Find the most parsimonious labeling of the internal nodes of a rooted tree.

- Input: A rooted binary tree with each leaf labeled by a single symbol.
- Output: A labeling of all other nodes of the tree by single symbols that minimizes the tree's parsimony score.

Small Parsimony Problem: Find the most parsimonious labeling of the internal nodes of a rooted tree.

- Input: A rooted binary tree with each leaf labeled by a single symbol.
- Output: A labeling of all other nodes of the tree by single symbols that minimizes the tree's parsimony score.

Checkpoint: Why is this an acceptable simplification?

Small Parsimony Problem: Find the most parsimonious labeling of the internal nodes of a rooted tree.

- Input: A rooted binary tree with each leaf labeled by a single symbol.
- Output: A labeling of all other nodes of the tree by single symbols that minimizes the tree's parsimony score.

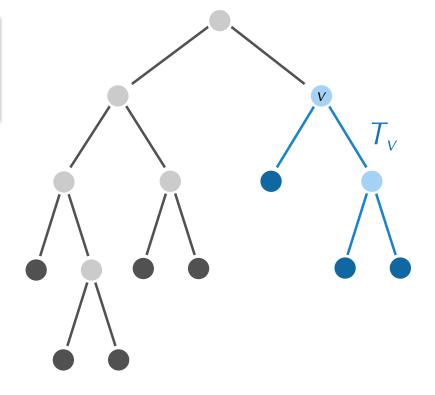
Answer: We may choose to assume that the characters are *independent*.

Small Parsimony Problem: Find the most parsimonious labeling of the internal nodes of a rooted tree.

- Input: A rooted binary tree with each leaf labeled by a single symbol.
- Output: A labeling of all other nodes of the tree by single symbols that minimizes the tree's parsimony score.

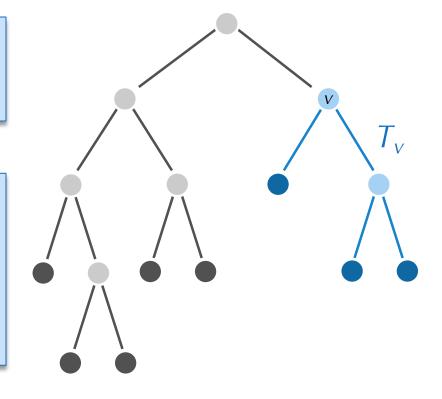
Checkpoint: Any thoughts on what approach we might use to solve this problem?

Let T_v denote the subtree of T whose root is v.



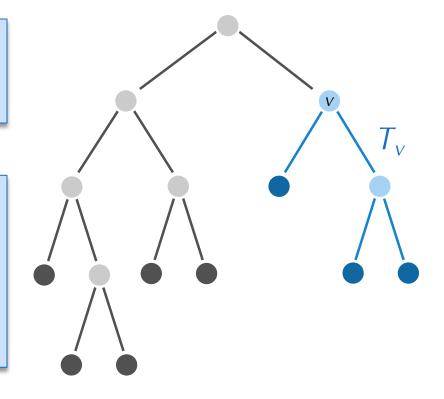
Let T_v denote the subtree of T whose root is v.

Define $s_k(v)$ as the minimum parsimony score of T_v over all labelings of T_v , assuming that v is labeled by k.



Let T_v denote the subtree of T whose root is v.

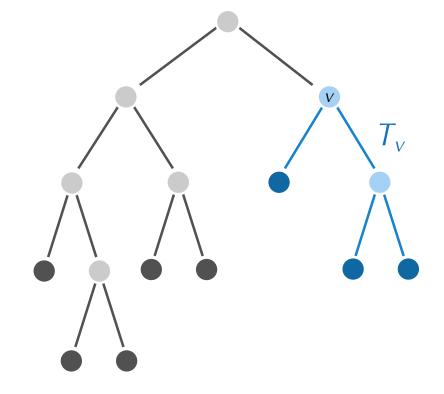
Define $s_k(v)$ as the minimum parsimony score of T_v over all labelings of T_v , assuming that v is labeled by k.



The minimum parsimony score for the tree is equal to the minimum value of $s_k(root)$ over all symbols k.

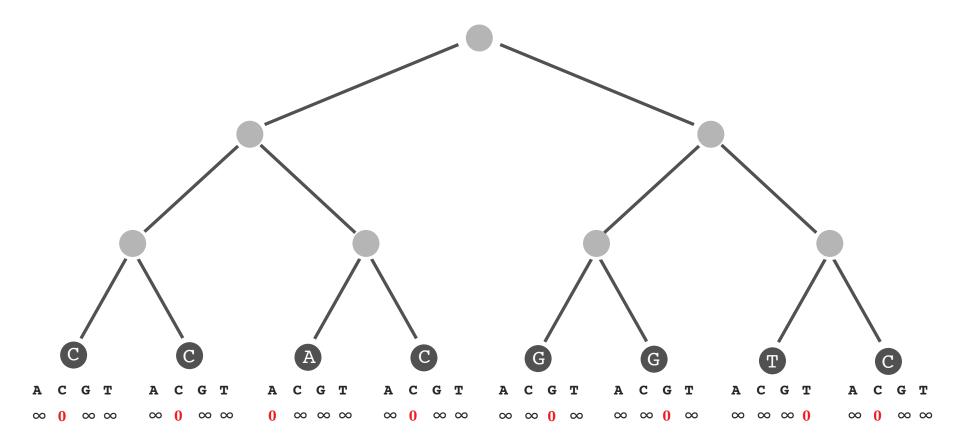
For symbols *i* and *j*, define

- $\alpha_{i,j} = 0$ if i = j
- $\alpha_{i,j} = 1$ otherwise.

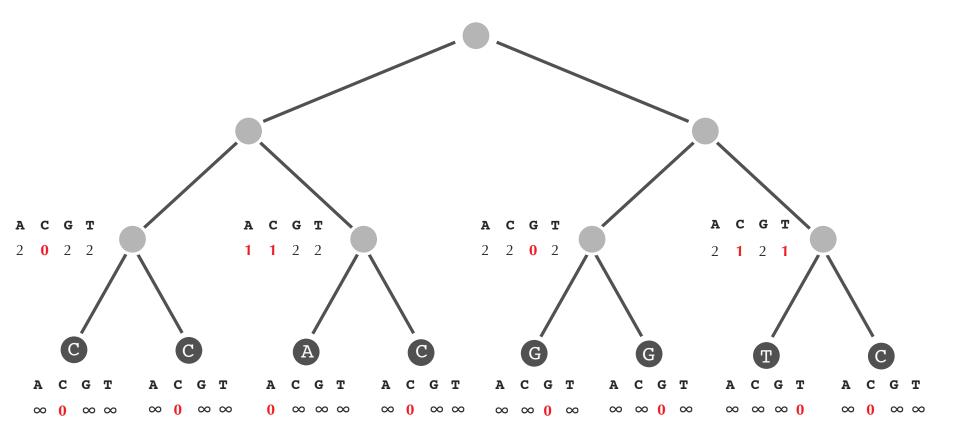


Theorem: The following recurrence relation holds:

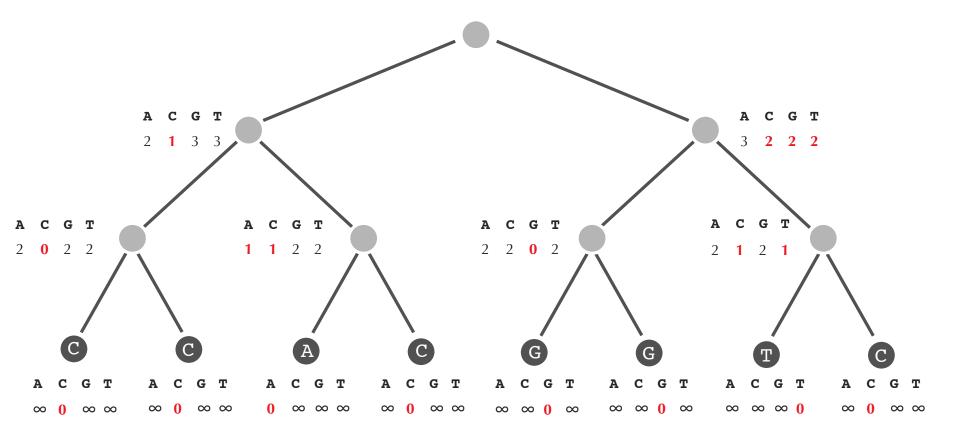
 $s_k(v) = \min_{\text{all symbols } i} \{s_i(Daughter(v)) + \delta_{i,k}\} + \min_{\text{all symbols } i} \{s_i(Son(v)) + \delta_{j,k}\}$



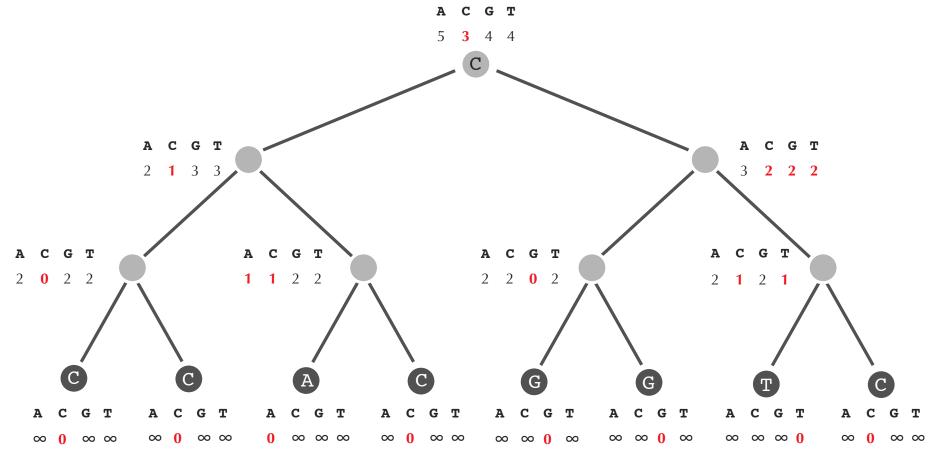
$$s_k(v) = \min_{\text{all symbols } i} \{ s_i(Daughter(v)) + \delta_{i,k} \} + \min_{\text{all symbols } i} \{ s_i(Son(v)) + \delta_{i,k} \}$$



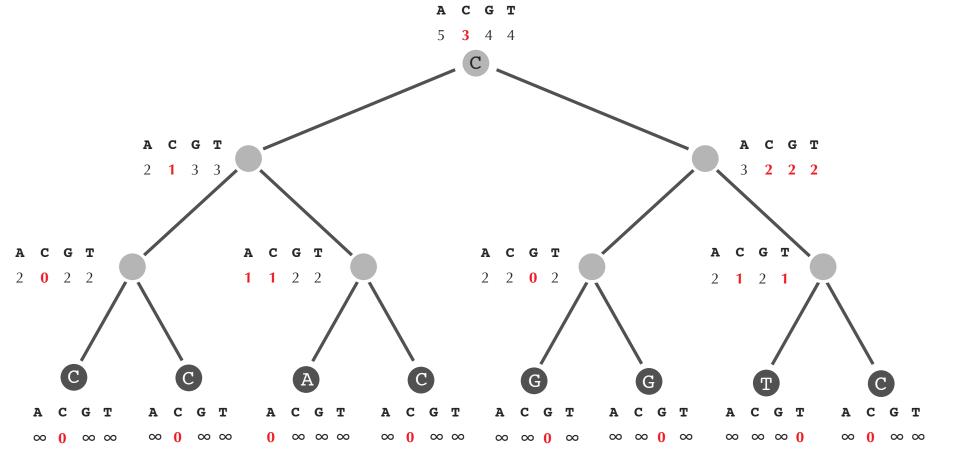
$$s_k(v) = \min_{\text{all symbols } i} \{ s_i(Daughter(v)) + \delta_{i,k} \} + \min_{\text{all symbols } i} \{ s_i(Son(v)) + \delta_{i,k} \}$$



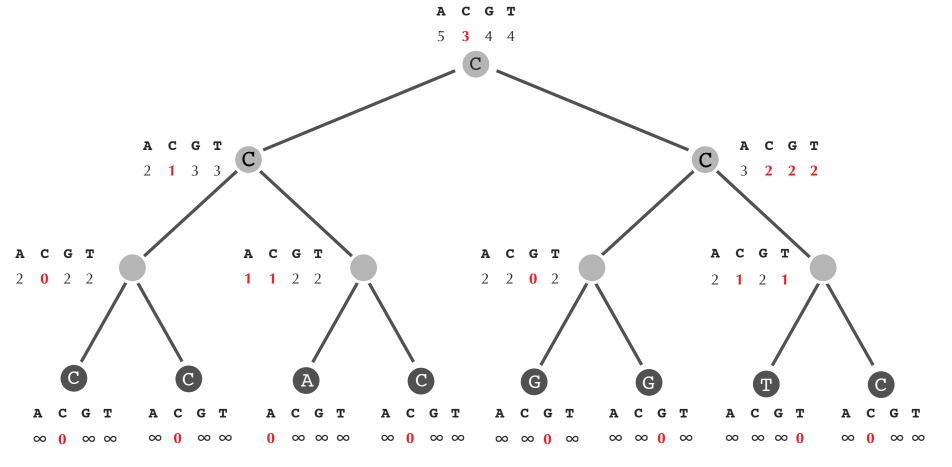
$$s_k(v) = \min_{\text{all symbols } i} \{ s_i(Daughter(v)) + \delta_{i,k} \} + \min_{\text{all symbols } i} \{ s_i(Son(v)) + \delta_{i,k} \}$$



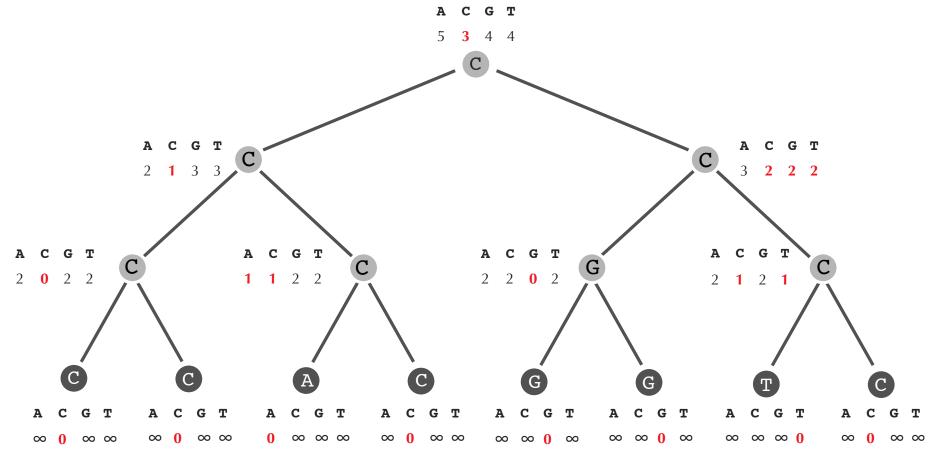
$$s_k(v) = \min_{\text{all symbols } i} \{s_i(Daughter(v)) + \delta_{i,k}\} + \min_{\text{all symbols } i} \{s_i(Son(v)) + \delta_{j,k}\}$$



Checkpoint: How should we "backtrack" to fill in the remaining nodes of the tree?



Checkpoint: How should we "backtrack" to fill in the remaining nodes of the tree?



Checkpoint: How should we "backtrack" to fill in the remaining nodes of the tree?