

## UNIT 10A

### Multiprocessing & Deadlock

15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

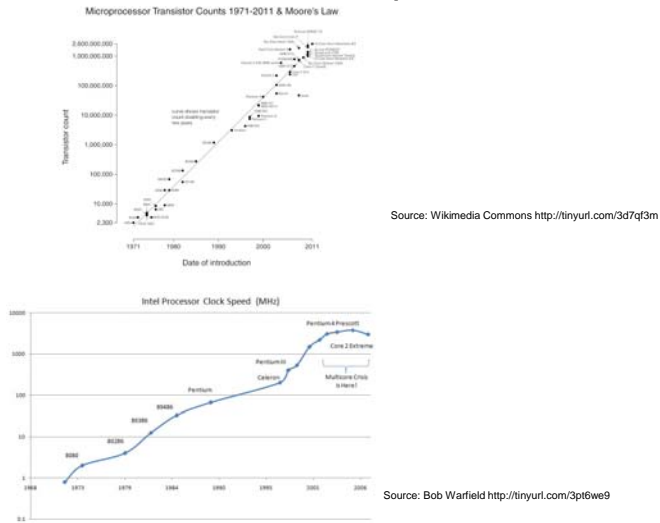
1

## Why Multiprocessing ?

- Everything happens at once in the world. Inevitably, computers must deal with that world.
  - Traffic control, process control, banking, fly by wire, etc.
- It is essential to future speed-up of any computing process.
  - Google, Yahoo, etc. use thousands of small computers, even when a job could be done with one big computer.
  - Chips can't run any faster because they would generate too much heat.
  - Moore's law will allow many processors per chip.

2

## Moore's Law vs. Clock Speed

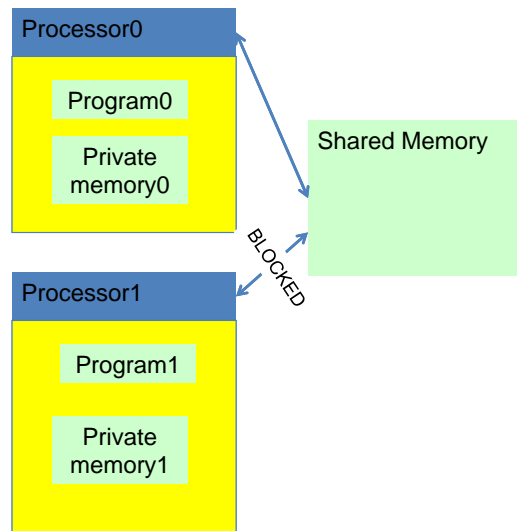


15110 Principles of Computing, Carnegie Mellon University - MORRIS

3

## A Ruby Multiprocessor Model

- The processors run independently.
- The shared memory is used for communication.
- Only one processor at a time may execute a line of Ruby that touches the shared memory. The memory hardware makes the others wait.



15110 Principles of Computing, Carnegie Mellon University - MORRIS

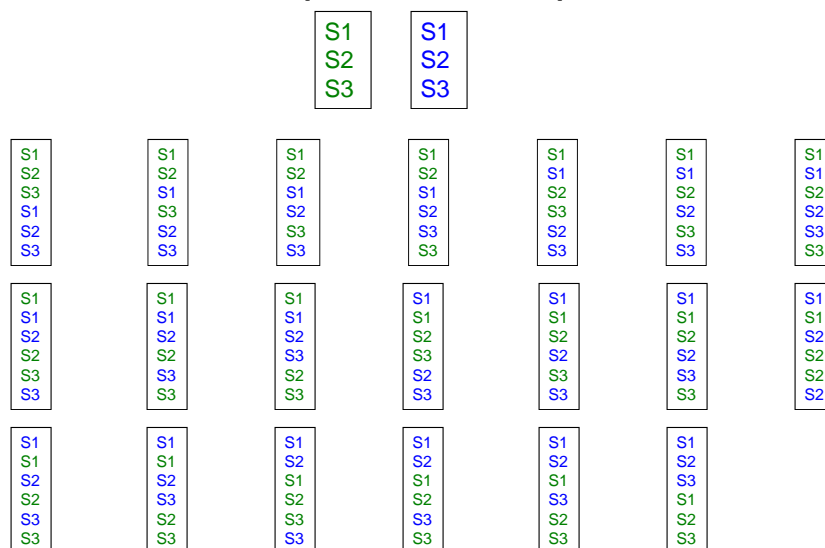
4

## Multiprocessing is *very* hard.

- Only a tiny percentage of practicing programmers can do it.
- It requires art and mathematics.
  - It's like digital hardware design.
  - It needs proofs.
- Conventional debugging doesn't work.
  - If you stop the program to observe, you change the behavior.
  - Testing is futile because the number of possible execution sequences for the same input explodes.

5

## There are many ways to execute two sequences in parallel.



6



# Streams: One process sends, another receives.

```
# Shared
@full = false
@box = nil
```

```
# Producer 0
while true do
  mail0 = whatever()
  while @full do #nothing
  end
  @box = mail0
  @full = true
end

# Consumer 1
while true do
  while !@full do #nothing
  end
  mail1 = @box
  @full = false
  process(mail1)
end
```

15110 Principles of Computing, Carnegie Mellon University - MORRIS

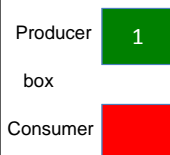
7

## A Typical Execution Pattern

```
# Shared
@full = false
@box = nil
```

```
# Producer 0
while true do
  mail0 = whatever()
  while @full do #nothing
  end
  @box = mail0
  @full = true
end

# Consumer 1
while true do
  while !@full do #nothing
  end
  mail1 = @box
  @full = false
  process(mail1)
end
```

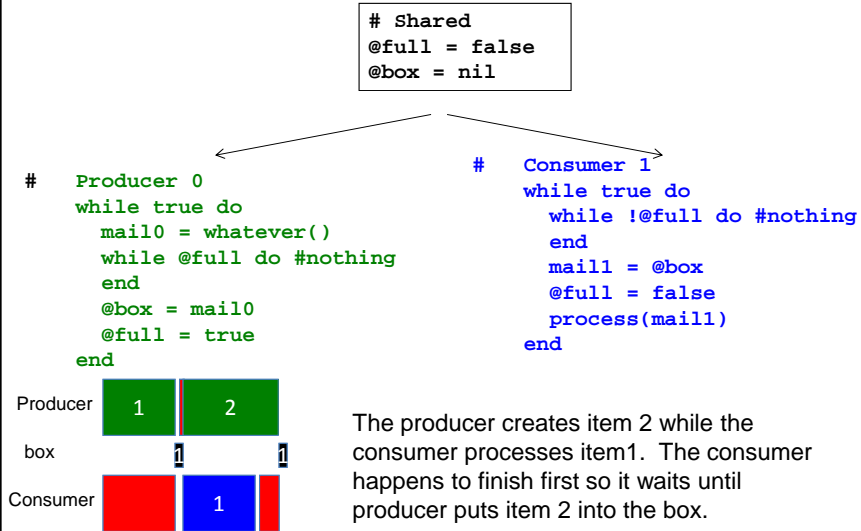


The producer creates item 1 and puts it in the box while the consumer waits.

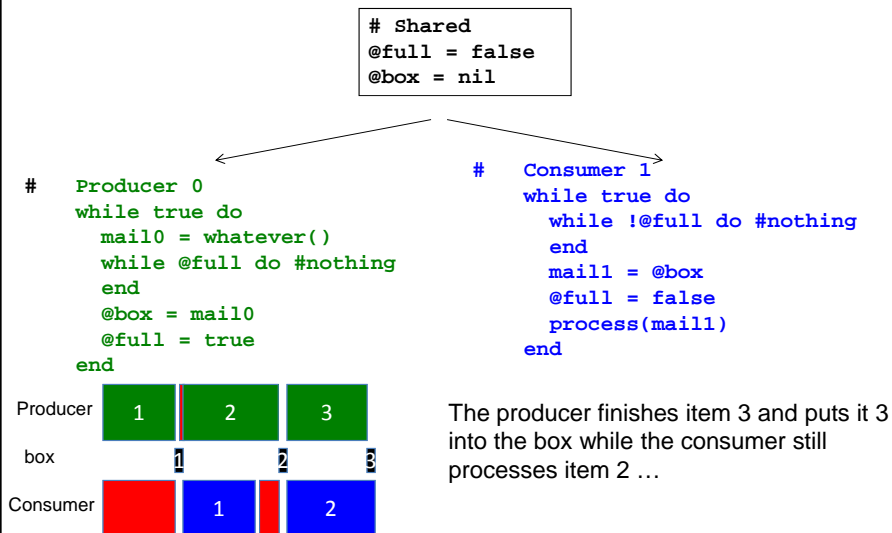
15110 Principles of Computing, Carnegie Mellon University - MORRIS

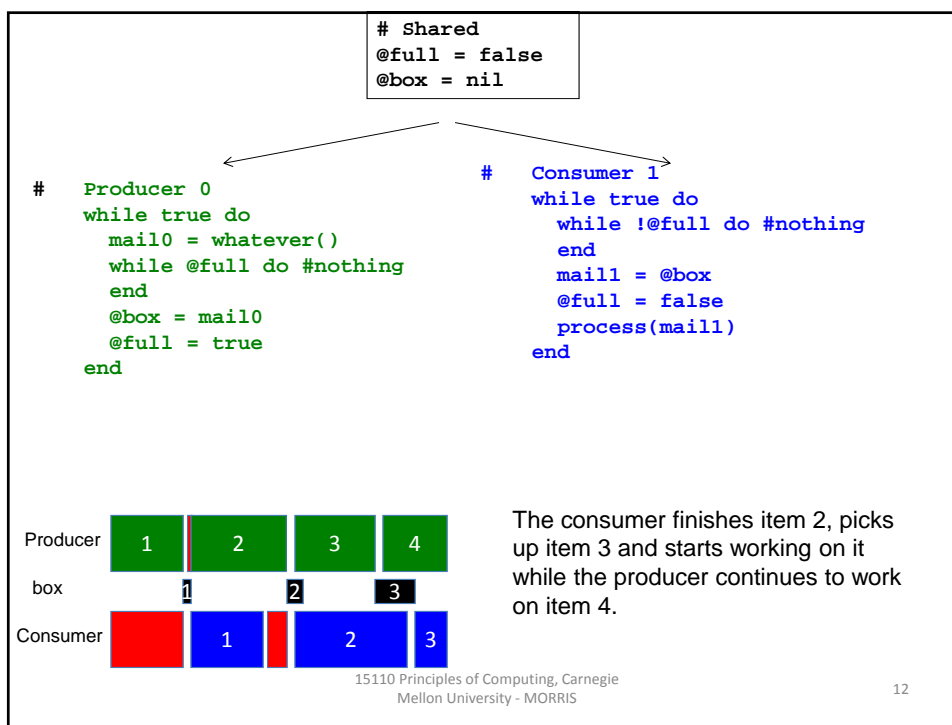
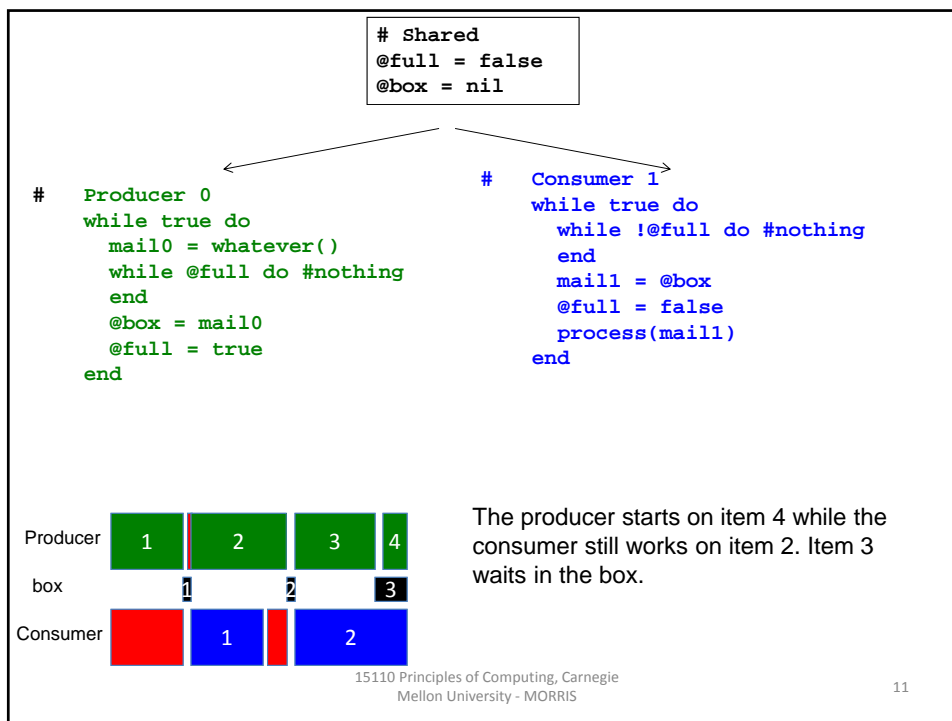
8

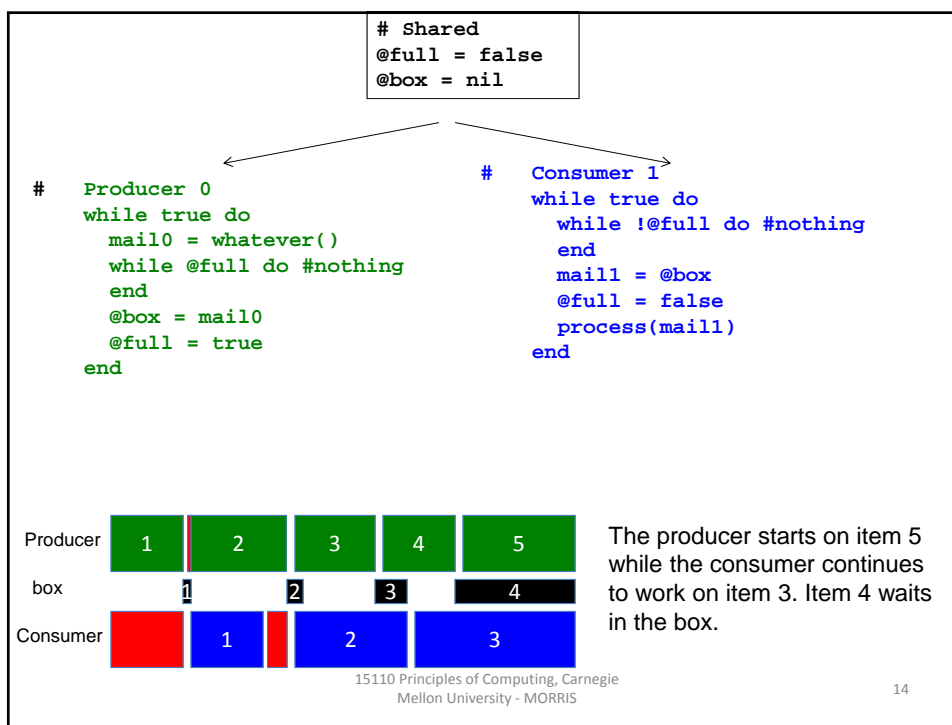
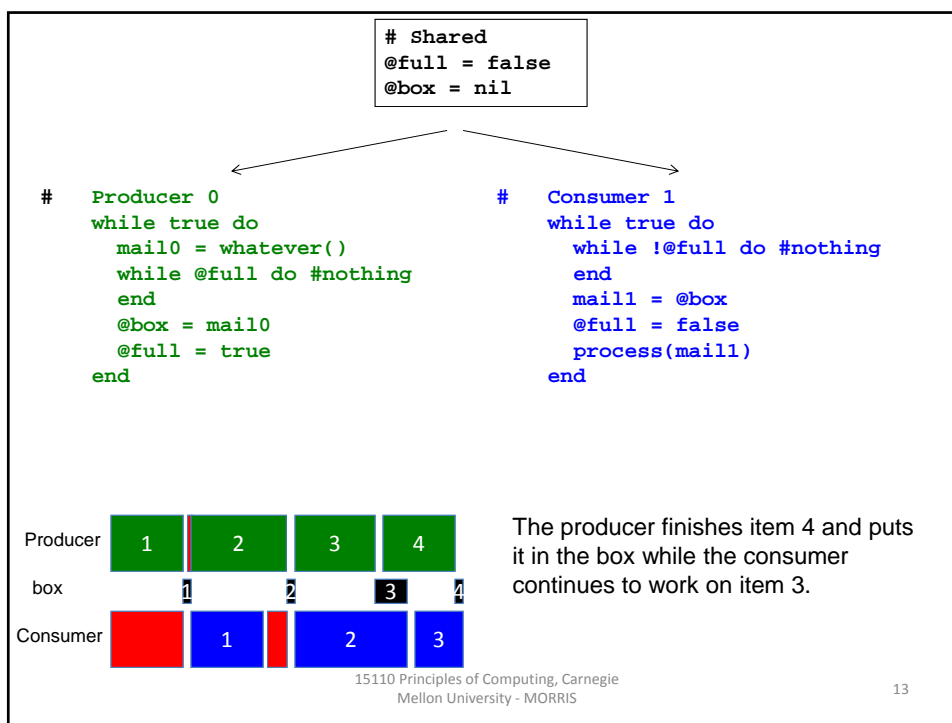
## A Typical Execution Pattern

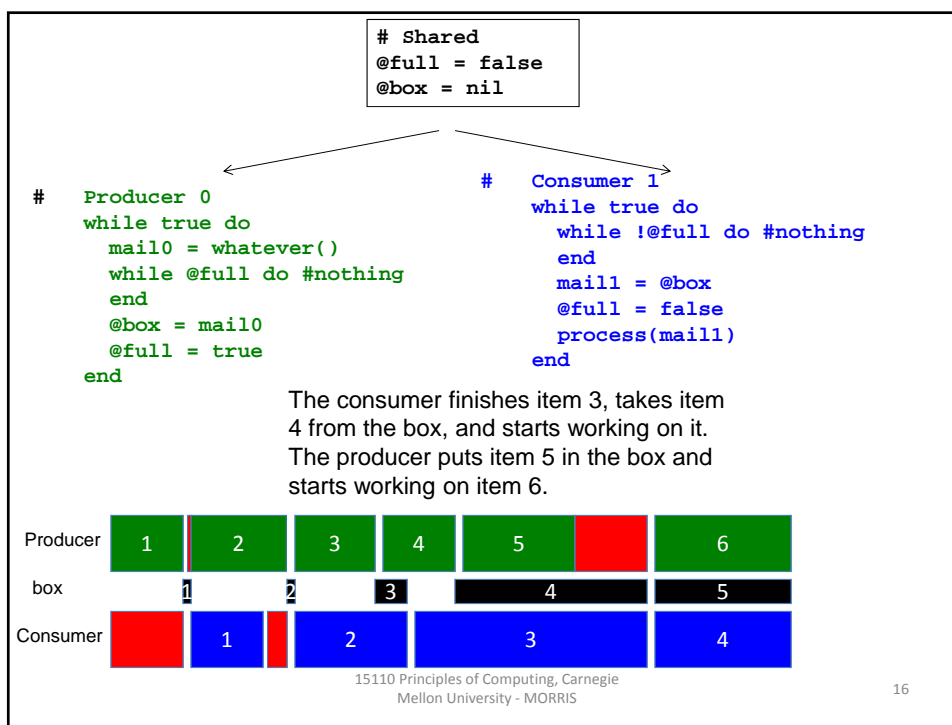
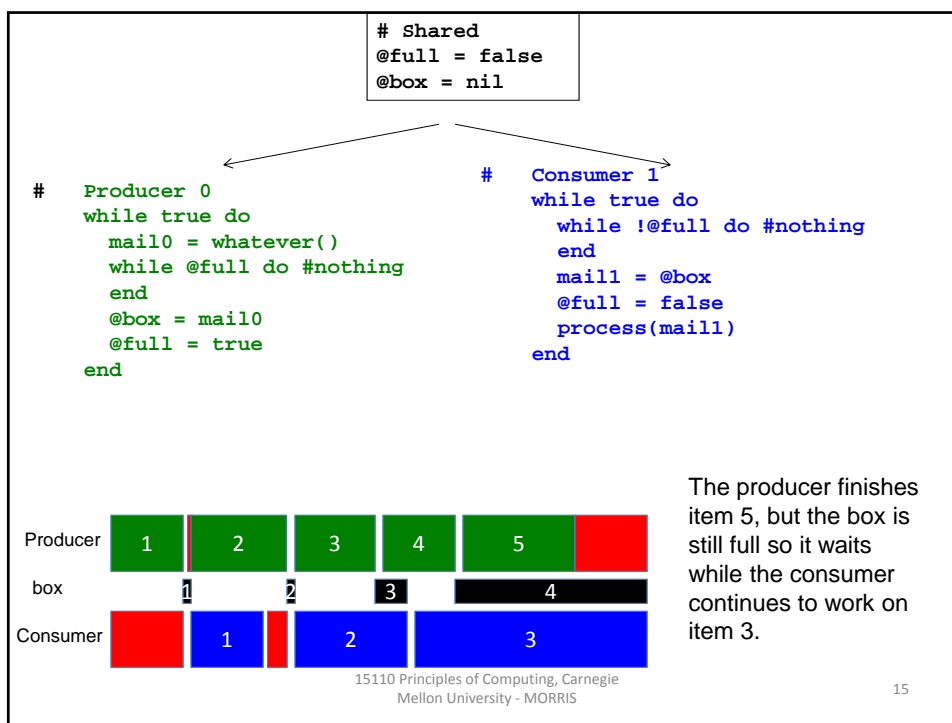


## A Typical Execution Pattern

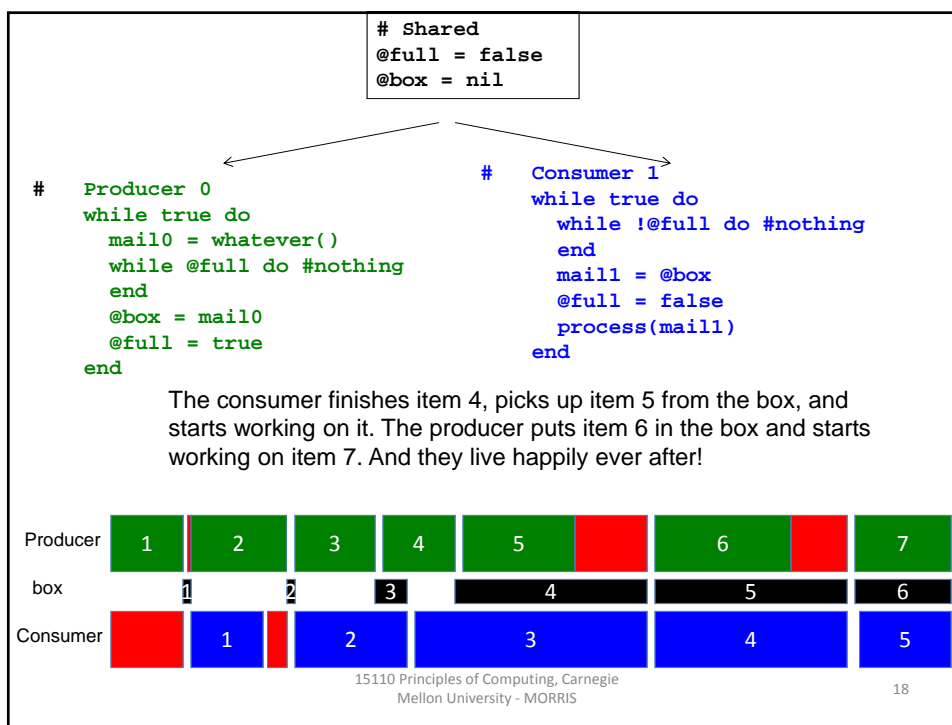
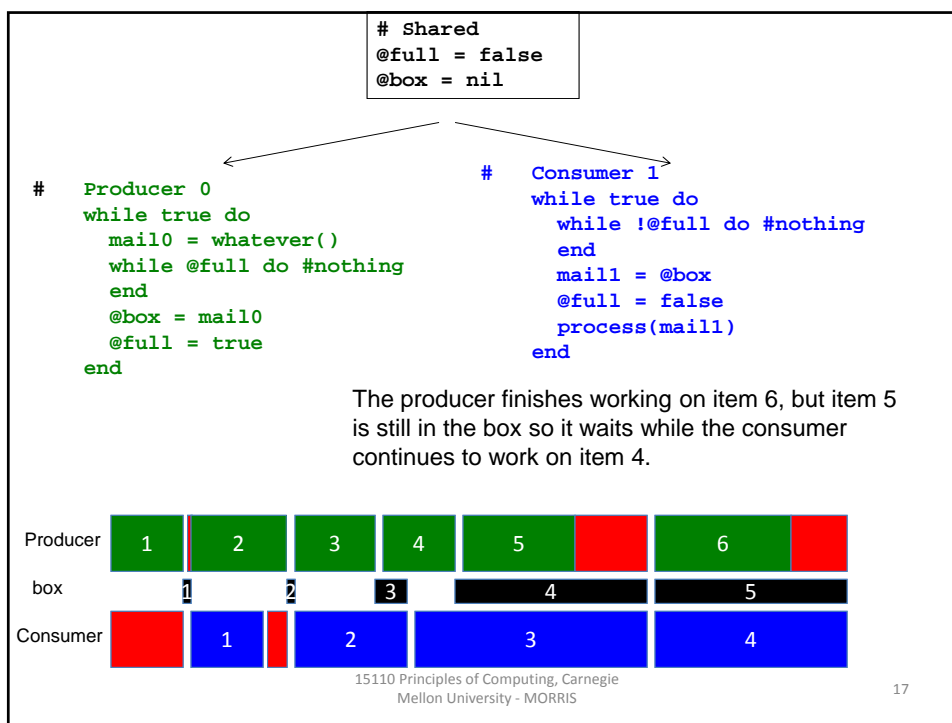












## Streams with a *Race Condition*

The order of accesses to @box and @full is very important!

```
# Producer
while true do
  mail0 = whatever()
  while @full do #nothing
  end
  @full = true
  @box = mail0
end

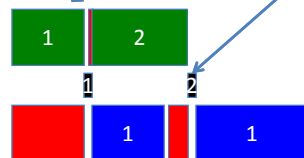
# Consumer
while true do
  while !@full do #nothing
  end
  mail1 = @box
  @full = false
  process(mail1)
end
```

BUG!

```
@full = true
@box = mail0 = 1
while !@full
  mail1 = @box = 1
  @full = false
end

@full = true
while !@full
  mail1 = @box = 1
  @full = false
  @box = mail0 = 2
end
```

An unfortunate interleaving of process steps leads to a mistake.



Item 1 is processed twice!

19

## Critical Sections

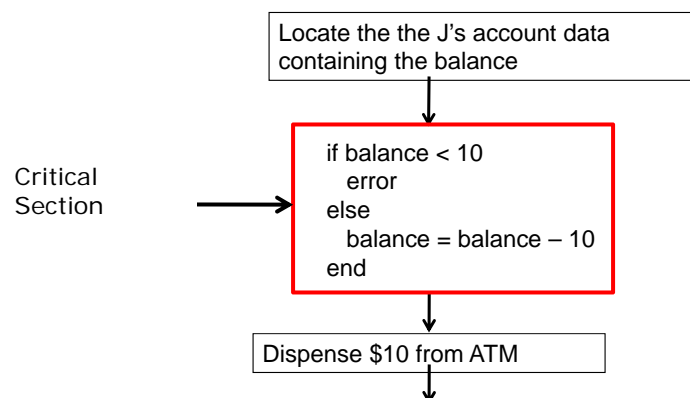
- Often, a process really needs exclusive access to some data for more than one line.
- A **critical section** is a sequence of two or more lines that need exclusive access to the shared memory.
- Real Life Examples
  - Crossing a traffic intersection
  - A bank with many ATMs
  - Making a ticket reservation

## Critical Section Example

- Consider a bank with multiple ATM's.
- At one, Mr. J requests a withdrawal of \$10.
- At another, Ms. J requests a withdrawal of \$10 from the same account.
- The bank's computer executes:
  1. For Mr. J, verify that the balance is big enough.
  2. For Ms. J, verify that the balance is big enough.
  3. Subtract 10 from the balance for Mr. J.
  4. Subtract 10 from the balance for Ms. J.
- The balance went negative if it was less than \$20!

21

## Critical Sections in Ruby



What can we do to prevent one processor from entering the critical section while another is in it?

15110 Principles of Computing, Carnegie Mellon University - MORRIS

22



## Careful Driver Method

Don't enter the intersection unless it's empty.

In shared memory:    @free = true    #initially unlocked

```
#Process 1
while true do
  Non-Critical_Section
  while !@free do #nothing
    end
  @free = false
  Critical_Section
  @free = true
end

#Process 2
while true do
  Non-Critical_Section
  while !@free do #nothing
    end
  @free = false
  Critical_Section
  free = true
end
```

*Interference is possible!*

15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

23

## Computers vs. Real Life

- The careful driver method works in real life because
  - The number of times in your life you cross the intersection is low. Twice a day for forty years is about 29,000.
  - The chance of two drivers arriving at the intersection simultaneously is low.
  - Cars move slowly enough that if you don't see anyone coming, you'll get across before anyone comes.

15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

24

## The Probability of a Collision

```

while true do
  Non-Critical_Section
  while !@free do #nothing
  end
  @free = false
  Critical_Section
  @free = true
end

```

Average time to perform Non-Critical Section: 1,000 nanoseconds  
 Average time to perform Critical Section: 10 nanoseconds  
 Average time to test and change @free: 3 nanoseconds

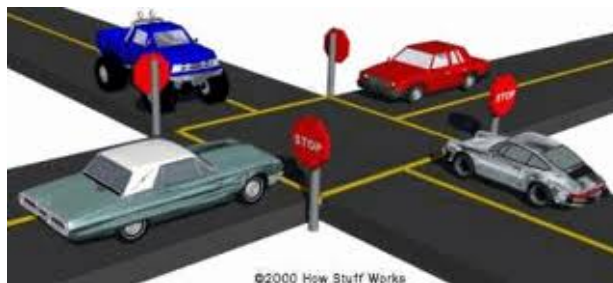


Probability of one collision:  $1/1,000 = 0.001$   
 Iterations of outer loop in one second:  $1,000,000,000/1,013 = 987,166$   
 Probability of *no* collisions in 1 second:  $(1-0.001)^{987,166} = (0.999)^{987,166} = 0$

15110 Principles of Computing, Carnegie Mellon University - MORRIS

25

## The Stop Sign Method



1. Signal your intention (by stopping).
2. Wait until cross road has no one waiting or crossing.
3. Cross intersection.
4. Renounce intention (by leaving intersection).

15110 Principles of Computing, Carnegie Mellon University - MORRIS

26

## The Stop and Look Method

```

# Shared Memory
@free[0] = true    #P0 is not stopped at
sign
@free[1] = true    #P1 is not stopped at
sign
# Process 0
while true do
  Non-Critical_Section
  @free[0] = false
  while !@free[1] do
    end
  Critical_Section
  @free[0] = true
end
# Process 1
while true do
  Non-Critical_Section
  @free[1] = false
  while !@free[0] do
    end
  Critical_Section
  @free[1] = true
end

```

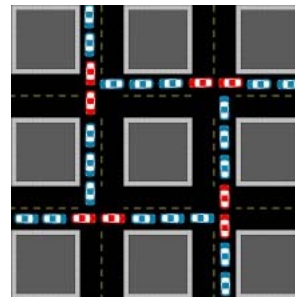
Deadlock is possible!

15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

27

## Deadlock

- Deadlock is the condition when two or more processes are all waiting for some shared resource, but no process actually has it to release, so all processes to wait forever without proceeding.
- It's like gridlock in real traffic.



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

28

## The Stop Sign Method with Tie Breaking



1. Signal your intention (by stopping).
2. Wait until cross road has no one else waiting or crossing.
3. If two of you are both waiting, yield to the car to your right.
4. Cross intersection.
5. Renounce intention (by leaving intersection).

29

## Stop Sign with Tie Breaking

```

@free[1] = true
@free[2] = true

# Process 1
while true
  Non-Critical_Section1
  @free[1] = false
  while !@free[2] do
    end
  Critical_Section1
  @free[1] = true
end

# Polite-Process 2
while true do
  Non-Critical_Section2
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    while !@free[1] do
      end
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end

```

Process 2 backs off when it detects a conflict.

30

## Types of Race Condition Bugs

In decreasing order of seriousness:

1. Interference: multiple process in critical section.
2. Deadlock: two processes idle forever, neither entering their critical or non-critical sections.
3. Starvation: one process needlessly idles forever while the other stays in its non-critical section.
4. Unfairness: a process has lower priority for no reason.

31

## Peterson's algorithm avoids all bugs!

```

@free[0] = true
@free[1] = false
priority = 0

# Process 0
while true do
  Non-Critical_Section0
  @free[0] = false
  priority = 1
  while !@free[1] and
    priority==1 do
  end
  Critical_Section0
  @free[0] = true
end

# Process 1
while true do
  Non-Critical_Section1
  @free[1] = false
  priority = 0
  while !@free[0] and
    priority==0 do
  end
  Critical_Section1
  @free[1] = true
end

```



## A Probabilistic Approach

```
# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end
```



## A Probabilistic Approach

```
# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end
```



## A Probabilistic Approach

```
# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end
```



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

35

## A Probabilistic Approach

```
# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end
```



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

36

## A Probabilistic Approach

```

# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end

```



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

37

## A Probabilistic Approach

```

# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end

```



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

38

## A Probabilistic Approach

```
# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end
```



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

39

## A Probabilistic Approach

```
# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end
```



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

40

## A Probabilistic Approach

```
# Process 1
while true
  Non_Critical_Section1
  n1 = 0.000001 #microsecond
  @free[1] = false
  while !@free[2] do
    @free[1] = true
    sleep(rand(n1))
    n1 = 2 * n1
    @free[1] = false
  end
  Critical_Section1
  @free[1] = true
end

# Process 2
while true
  Non_Critical_Section2
  n2 = 0.000001
  @free[2] = false
  while !@free[1] do
    @free[2] = true
    sleep(rand(n2))
    n2 = 2 * n2
    @free[2] = false
  end
  Critical_Section2
  @free[2] = true
end
```



The probability of dithering is vanishingly small, proportional to  $1/2^N$  for  $N$  collisions.

15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

41

## New Vocabulary

- Stream: A programming pattern in which one process sends data to another process sequentially
- Race Condition: A multiprocessing bug in which proper functioning depends upon luck
- Deadlock: A condition in which all processes are stalled waiting for each other
- Starvation: A condition in which a process is needlessly stalled

15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

42

## Takeaways

- Multiprocessing is very hard because controlling events in the real world is very hard.
- Sequential programming was a great invention because it made controlling simple things very easy.
- Leave it to the Engineers and hope they get it right.

## Afterthoughts

Some counter-intuitive ideas about bugs and risks.

This man removed all the traffic lights and signs!



15110 Principles of Computing, Carnegie  
Mellon University - MORRIS

45

Why did Jared Diamond sleep under a tree when his aborigine companion wouldn't?



## Why is a 1% chance of a bug biting better than a 0.1% chance?

- If there is a 1% chance of error, the bug will show up during 100 days of testing.
- If there is a 0.1% chance, the bug will show up after three years when the system is deployed.

15110 Principles of Computing, Carnegie Mellon University - MORRIS

47

## Economics as Multiprocessing

A national economy could be looked at a system with 1B independent processes representing buyers and sellers of goods. Consider the following economic maladies:

- A. Depression
- B. Bubbles
- C. Income Inequality
- D. Wasted productive resources

How do these problems correspond to the four multiprocessing problems?

- 1. Interference
- 2. Deadlock
- 3. Starvation
- 4. Unfairness

Hint: Think of entering a critical section as buying a good.

48