15-853: Algorithms in the Real World

Computational Biology II
- Sequence Alignment
- Database searches

Extending LCS for Biology

The LCS/Edit distance problem is not a "practical" model for comparing DNA or proteins.

Why?

Good example of the simple model failing.

What we will talk about today

Extensions
- **Sequence Alignment**: a generalization of LCS to account for the closeness of different elements
- **Gap Models**: More sophisticated models for accounting for the cost of adjacent insertions or deletions
- **Local Alignment**: Finding parts of one sequence in parts of another sequence.

Applications
- **FASTA** and **BLAST**: The most common sequence matching tools used in Molecular Biology.
Sequence Alignment

A generalization of LCS / Edit Distance

**Extension:** A' is an extension of A if it is A with spaces _ added.

**Alignment:** An alignment of A and B is a pair of extensions A' and B' such that |A'| = |B'|

**Example:**
- A = a b a c d a
- B = a a d c d d c
- A' = _ a b a c d a _
- B' = a a d _ c d d c

The Score (Weight)

Σ⁺ = alphabet including a "space" character

**Scoring Function:** σ(x,y), x,y ∈ Σ⁺

**Alignment score:** \( W(A', B') = \sum_{i=1}^{[A]} \sigma(A_i', B_i') \)

**Optimal alignment:** An alignment \((A', B')\) of \((A, B)\) such that \(W(A', B')\) is maximized. We will denote this optimized score as \(W(A, B)\).

Same as |LCS| when: \( \sigma(x, y) = \begin{cases} 1 & \text{if } x = y \neq _ \\ 0 & \text{otherwise} \end{cases} \)

Scores vs. Distances

**Maximizing vs. Minimizing.**

**Scores:**
- Can be positive, zero, or negative. We try to maximize scores.

**Distances:**
- Must be non-negative, and typically we assume they obey the triangle inequality (i.e., they are a metric). We try to minimize distances.

Scores are more flexible, but distances have better mathematical properties. The local alignment method we will use requires scores.
σ(x,y) for Protein Matching

How is the function/matrix derived?

- **Identity**: entries are 0 or 1, either same or not
- **Genetic code**: number of DNA changes, Remember that each amino acid is coded with a 3 bp codon. Changes can be between 0 and 3
- **Chemical Similarity**: size, shape, or charge
- **Experimental**: see how often mutations occur from one amino acid to another in real data.

This is what is used in practice.
- PAM (or dayhoff) Matrix
- BLOSUM (BLOcks SUbstitution Matrix)

Optimal Alignment Recursive Solution

**Edit Distance, from last lecture:**

- D(A, ε) = |A|
- D(ε, B) = |B|
- D(A:x, B:x) = D(A, B)
- D(A:a, B:b) = \min(I + D(A:a, B), I + D(A, B:b))

**The optimal alignment problem:**

- W(ε, ε) = 0
- W(ε, B:b) = σ(ε, b) + W(ε, B)
- W(A:a, ε) = σ(a, ε) + W(A, ε)
- W(A:a, B:b) = max(σ(a, b) + W(A, B), σ(ε, b) + W(A:a, B), σ(a, ε) + W(A, B:b))

BLOSUM62 Matrix

| A | B | C | D | E | F | G | H | I | K | L | M | N | P | Q | R | S | T | V | W | Y |
| A | 4 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | 21 |
| B | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |
| C | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 |
| D | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
| E | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| F | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 |
| G | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| H | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |
| I | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 |
| K | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| L | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| M | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| N | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| P | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Q | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| R | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
| S | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 | 4 |
| T | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 | 3 |
| V | 18 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 | 2 |
| W | 19 | 18 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | 1 |
| Y | 20 | 19 | 18 | 17 | 16 | 15 | 14 | 13 | 12 | 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 |

W_{i,j} = \max \{ W_{i-1,j} + \sigma(a_{i-1}, b_{j-1}) \}

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Dynamic programming

for i = 1 to n
M[i,1] = σ(A[i], _ )
for j = 1 to m
M[1,j] = σ(_ , B[i]);
for i = 1 to n
for j = 1 to m
M[i,j] = max3(σ(A[i],B[j]) + M[i-1,j-1],
              σ(A[i], _) + M[i-1,j],
              σ(_ , B[j]) + M[i , j-1]);

Example

<table>
<thead>
<tr>
<th>B</th>
<th>a</th>
<th>t</th>
<th>c</th>
<th>a</th>
<th>c</th>
<th>a</th>
<th>c</th>
</tr>
</thead>
<tbody>
<tr>
<td>t</td>
<td>0</td>
<td>-1</td>
<td>-2</td>
<td>-3</td>
<td>-4</td>
<td>-5</td>
<td>-6</td>
</tr>
<tr>
<td>c</td>
<td>-1</td>
<td>1</td>
<td>0</td>
<td>-1</td>
<td>-2</td>
<td>-3</td>
<td>-4</td>
</tr>
<tr>
<td>A</td>
<td>c</td>
<td>-2</td>
<td>-1</td>
<td>0</td>
<td>3</td>
<td>2</td>
<td>1</td>
</tr>
<tr>
<td>a</td>
<td>-3</td>
<td>0</td>
<td>-1</td>
<td>2</td>
<td>5</td>
<td>4</td>
<td>3</td>
</tr>
<tr>
<td>t</td>
<td>-4</td>
<td>-1</td>
<td>2</td>
<td>1</td>
<td>4</td>
<td>6</td>
<td>5</td>
</tr>
</tbody>
</table>

3 blanks inserted = -3
1 t-c match = 1
3 perfect matches = 6
TOTAL = 4

Optimizations

Space efficiency:
The divide-and-conquer technique still works.

The Ukkonen/Myers algorithm:
A variant works, but O(dn) time is no longer
guaranteed since the distance from the diagonal
cannot in general be directly bounded by the score.

Bounds, however, can be given in terms of relative
weights of matrix elements and the technique
works reasonably well in practice.

Real Problem: solves global-alignment problem when
biologists care about the local-alignment problem.

Gap Penalties

Problem with technique so far: Longer indels
(insertions or deletions) should not be weighted as
the sum of single indels.

Gap: indel of k characters is a gap of length k

Gap score: let x_i be the score of a gap of length k

What is a good gap scoring function?
Can the dynamic programming approach be
extended?
**Possible Gap Scores**

- $x_k = \alpha + \beta k$

**Waterman-Smith-Beyer Algorithm**

\[
\begin{align*}
W_{00} &= 0 \\
W_{i0} &= x_i \\
W_{0j} &= x_j \\
W_{ij} &= \max \left\{ \frac{W_{i-1,j-1} + \sigma(a_i, b_j)}{\max_{k=1}^{l=1} \{W_{i-k,j-k} + x_k\}} \right\}
\]

Every cell in the matrix has to calculate its score based on all previous elements in its row and column.

**Affine Gap Model: $x_k = \alpha + \beta k$**

**Algorithm:** (Gotoh '82) for each cell of the $n \times m$ matrix keep three values, $E$, $F$ and $W$.

- $E$ = optimal alignment of form $A$, $B$: __
- $F$ = optimal alignment of form $A$: __ $B$
- $W$ = optimal alignment

Constant time per cell.

Total time: $O(nm)$
### Other Gap Models

<table>
<thead>
<tr>
<th>Function Form</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>General</td>
<td>$O(nm^2 + n^2m)$</td>
</tr>
<tr>
<td>Affine</td>
<td>$O(nm)$</td>
</tr>
<tr>
<td>Logarithmic</td>
<td>$O(nm)$</td>
</tr>
<tr>
<td>Concave Downwards</td>
<td>$O(nm \log n)$</td>
</tr>
<tr>
<td>Piecewise linear</td>
<td>$O(lnm)$</td>
</tr>
</tbody>
</table>

#### Local Alignment

**Algorithm:** (Smith and Waterman '81)

\[
\begin{align*}
W_{0i} &= 0 \\
W_{0j} &= 0 \\
W_{ij} &= \max \left\{ \begin{array}{l}
W_{i-1,j-1} + \sigma(a_i, b_j) \\
\max_{k=1}^{\max \{i, j\}} \{ W_{i-1,k-1} + \sigma(a_i, x_k) \} \\
\max_{k=1}^{\max \{i, j\}} \{ W_{i,k-1} + \sigma(x_k, b_j) \} \end{array} \right. \\
W_{ij} &= \max \left\{ \begin{array}{l}
W_{i-1,j-1} + \sigma(a_i, b_j) \\
\max_{k=1}^{\max \{i, j\}} \{ W_{i-1,k-1} + \sigma(a_i, x_k) \} \\
\max_{k=1}^{\max \{i, j\}} \{ W_{i,k-1} + \sigma(x_k, b_j) \} \end{array} \right. \\
\end{align*}
\]

With Gap Scores

Without Gap Scores

Only real difference from before is the 0 in the max. We might want **global** or **local maximums** of $W_{ij}$

### Example

The algorithm finds 3 "local maximum" corresponding to 3 hits. Although the three matches shown are fully on diagonals, this is not always the case.
Database Search

Basic model:
1. User selects a database and submits a source sequence S, typically via the web (or email).
2. The remote computer compares S to each target T in its database. The runtime depends on the length of S and the size of the database.
3. The remote computer returns a ranked list of the "best" hits found in the database, usually based on local alignment.
Example of BLAST

Algorithms in the “real world”

Dynamic programming is too expensive even with optimizations.
Heuristics are used that approximate the dynamic programming solution. Dynamic program is often used at end to give final score.
Main two programs in practice:
- FASTA (1985)
- BLAST (Basic Local Alignment Search Tool) (1990)
Lipman involved in both, Myers involved in BLAST. There are many variants of both.

FASTA and BLAST

The idea of both algorithms is to find approximate matches by composing smaller exact matches.

Both algorithms loop over each string T in the database and find the "heuristically" best match for the search string along with its score (assuming the score is above some threshold).
The matches across the T in the database are returned in rank order (highest-score first).

FASTA: Step 1

Break source S into k-tuples (adjacent sequence of k characters) using a "sliding window". Typically k=1 or 2 for proteins, and k=4-6 for DNA.

Create a table that maps each k-tuple value found to all the start positions with that value.
**FASTA: Step 2**

Linearly search \( T \) for each \( k \)-tuple belonging to \( S \) and bucket the hits (called "hot spots") by diagonal.

At \( T \), there are two hits in the table giving positions \( i_1 \) and \( i_2 \) in \( S \), which are in diagonals \( j - i_1 \) and \( j - i_2 \).

**FASTA: Step 3**

Join hits along diagonals into runs by

1. Each hit gets a positive score and each space between hits a negative score that increases with distance.
2. Finding runs that have maximal score (i.e. the sum of the scores cannot be increase by extending the run). There might be more than one such run in a diagonal.

**FASTA: Step 4**

Select top 10 scores from step 3, and re-score them using a substitution matrix (e.g. BLOSUM62). The best score is called \( \text{init}_1 \). Any scores below a threshold is thrown out. This leaves between 0 and 10 runs.

**FASTA: Step 5 (method 1)**

Each diagonal run \( k \) can be described by the start location in the matrix, \( (i_k, j_k) \), and its length, \( d_k \). We say that run \( l \) dominates run \( k \) if \( i_l \geq i_k + d \) and \( j_l \geq j_k + d \).

For all pairs of runs \( l, k \) such that \( l \) dominates \( k \), score the gap from the end of run \( k \) to the start of run \( l \).

I does not dominate \( k \)

I dominates \( k \)
FASTA: Step 6
Create a graph in which
1. each run is a vertex weighted by its score
2. each pair (l,k) with l dominating k is an edge
weighted by the gap score
Find the heaviest weight path in the graph, where
the vertex weight is included in the path length.

FASTA: Step 7
After the path is found, dynamic programming is
used to give a more accurate score to the path.
We now have a global alignment problem (since we
know the start and end of each string), so we can
use the Myers/Ukkonen algorithm.

FASTA: Summary
1. Break source S into k-tuples (adjacent sequence of k
characters). Typically k=2 for proteins.
For each T in the database
2. Find all occurrences of tuples of S in T.
3. Join matches along diagonals if they are nearby
4. Re-score top 10 matches using, e.g. Blosum matrix.
Method 1: (init.)
5. With top 10 scored matches, make a weighted graph
representing scores and gap scores between matches.
6. Find heaviest path in the graph.
7. re-score the path using dynamic programming
Method 2: (opt)
5. With top match, score band of width w around the diagonal
Rank order the matches found in steps 2-7.

FASTA: Step 5 (method 2)
Select best score for a diagonal (this was init.) Use
dynamic programming to score a band of width w
(typically 16 or 32 for proteins) around the best
diagonal
BLAST

1. Break source $S$ into $k$-tuples (adjacent sequence of $k$ characters). Typically $k=3$ for proteins.
2. For each $k$-tuple $w$ (word), find all possible $k$-tuples that score better than threshold $t$ when compared to $w$ (using e.g. BLOSUM matrix).
   This gives an expanded set $S_k$ of $k$-tuples.

For each $T$ in the database

3. Find all occurrences (hits) of $S_k$ in $T$.
4. Extend each hit along the diagonal to find a locally maximum score, and keep if above a threshold $s$. This is called a high-scoring pair (HSP). This extension takes 90% of the time.
   Optimization: only do this if two hits are found nearby on the diagonal.

The initial BLAST did not deal with indels (gaps), but a similar method as in FASTA (i.e. based on a graph) can be, and is now, used.

The BLAST thresholds are set based on statistical analysis to make sure that few false positives are found, while not having many false negatives.

Note that the main difference from FASTA is the use of a substitution matrix in the first stage, thus allowing a larger $k$ for the same accuracy.

A finite-state-machine is used to find $k$-tuples in $T$, and runs in $O(|T|)$ time independently of $k$.

The BLAST page