## Midterm Exam

## 18-213/613 Midterm Exam (Fall 2021)

## Important notes:

- This exam contains 6 questions.
- You are not required to answer all of them. Please choose to answer questions within the constraints described below.
- There is no extra credit for answering additional questions.
- Should additional questions be answered, we will count the LOWER of the options. It is to your advantage to make choices.
- This exam is an individual effort.
- You are not permitted to help others, in any way, with this exam.
- You are not permitted to release or to discuss this exam with anyone, except the course staff, until given permission to do so by the instructors (which will not occur until all students have completed the exam. There may be exceptional cases that take it late).
- You are permitted to use only the official course textbook, the official course slides, and your own personal notes.
- A simple calculator is permitted, but won't prove to be helpful (we don't think).
- You have 90 minutes, from first exposure through submission to take this exam. Do not attempt to "peek", "check", or "test" the exam. This will start your clock.


## Answer EXACTLY ONE of these:

- Question 1: Integers
- Question 2: Floats
- Properties
- Special Values


## Answer EXACTLY ONE of these:

- Question 3: Assembly
- Basic control
- Switch
- Question 4: Calling Convention, Stack Discipline

Answer ***BOTH*** of these:

- Question 5: Data
- Structs
- Arrays
- Question 6: Caching and Memory Access
- Fully Associative Trace
- 2-Way Set Associative Trace
- Comparative Performance
- Memory Access Time


## Question 1: Integers

This question is based upon the following declaration on a machine using 5-bit two's complement arithmetic for signed integers.

Blank (A):

```
int x = -13;
unsigned uy = x;
```

Fill in the empty boxes in the table below.

- Show all digits for the "Binary" column, including any leading 0 s.
- You need not fill in entries marked with "-".
- TMax denotes the largest positive two's complement number
- TMin denotes the smallest negative two's complement number.
- Hint: Be careful with the promotion rules that $C$ uses for signed and unsigned ints, i.e. how the $C$ Language handles implicit casts between the types.

| Expression | Decimal Representation | Binary Representation |
| :---: | :---: | :---: |
| - | -3 | (A) |
| - | (B) | $O$ |
| $x$ | - | (C) |
| uy | (D) | - |
| $x-u y$ | - | (E) |
| TMax + 1 | (F) | - |
| TMin - 1 | - | (G) |
| TMin + 1 | (H) | - |
| TMin + TMin | - | (1) |

## 2

2.5 points

Blank (B): $\square$
2.5 points

Blank (C): $\square$
$4 \quad 2.5$ points

Blank (D):
$5 \quad 2.5$ points

Blank (E):
2.5 points

Blank (F):
$\square$
$\square$
$\square$
$\square$

## Question 2: Floats

## Part 1: Properties

- Consider the following 7-bit floating point representation based on the IEEE floating point format:
- The most significant bit is the sign bit
- The next $k=3$ bits are the exponent.
- The last $n=3$ bits are the significand.
- The bias is to balance the exponents in a way consistent with IEEE single and double precision floating point numbers, i.e. according to the formula and with the intuition we discussed in class.

Please answer the questions to the right.

## 11

2 points
2.1(A): What is the bias? (Decimal number)
$\square$

122 points
2.1(B): What is the actual exponent, e.g. what we called "E" in class, for denormalized numbers? (Give answer in decimal).
$\square$

## Part 2: Special values

This question is based upon the same number format as Part I.

Fill in the blank entries in the following table. Include nothing but 0 s and 1 s . Include no spaces.

| Description | Sign | Binary <br> Encoding |
| :--- | :--- | :--- |
| Zero | + | 0000000 |
| Smallest <br> Positive <br> (nonzero) | + | (A) |
| Largest <br> denormalized | - | (B) |
| Smallest <br> positive <br> normalized | + | (C) |

2.1(C): Consider any two adjacent denormalized floating point numbers.

What is the absolute value of their difference in base-2 binary? Fill in the blank, without any unnecessary trailing 0s.:
0. $\qquad$
2.1(D): Consider any two adjacent normalized numbers with a biased exponent field of exp=010.

Determine the absolute value of the difference in their base- 2 binary values and write it out in binary as $\mathbf{x} . \mathbf{y}$ without any unnecessary trailing Os and without any unnecessary leading 0s (include a single leading or trailing zero per field, as necessary, to avoid leaving either field entirely blank.):
$\qquad$
(x) .
$\qquad$ (y) $\qquad$
What is ( x )?
$\square$

## 15

2 points
2.1(E): Consider the scenario in question (D) above.
$\qquad$ (x) $\qquad$ .
(y) $\qquad$

What is (y)?
2.1(F) Consider any two adjacent normalized numbers with a biased exponent field of exp=011.

Determine the absolute value of their difference in base-2 binary and write it out as $\mathbf{x}$. $\mathbf{y}$ without any unnecessary trailing 0s and without any unnecessary leading 0 s (include a single 0 per field as necessary to avoid leaving either field blank):
(x) $\qquad$ . $\qquad$ (y)
$\square$

171 point
2.1(E): Consider the scenario in question (D) above, what is $y$ ?
$\qquad$ (x) $\qquad$ .
(y)

What is (y)?
2.1(H) Which of the following explains the difference between your answers to (d), (e), and (f). Check all that apply:When the points on the number line are assigned to be closer in value, rounding error is reduced

When the points on the number line are assigned to be farther apart in value, a number line can cover a larger rangeIEEE wanted the number line to span a large range but to keep the rounding error approximately proportional to the magnitude of the number

$\square$
Denormalized numbers are relatively very small in magnitude and represent only a very small portion of the range, so it makes sense for them to be equidistant.

193 points
Part 2, Blank (A):

## 203 points

Part 2, Blank (C):
$\square$

## 21 <br> 3 points

Part 2, Blank (B):

Part 1: Control

Please consider the following assembly code and then answer the questions about it that follow:

Hint: We strongly suggest that, before answering the questions, you translate the code below into the $C$ Language and simplify it in writing.


## Part 2: Switch

Please consider the following assembly and memory dump:

Hint: Recall that the gdb command $\mathrm{x} / \mathrm{g}$ SOME_ADDRESS_EXPRESSION will examine an 8-byte word starting at the given address.

## 3.1(A): How many loops are there?

3.1(B)How would you describe the relationship among the loop(s). Choose one:

O One loop
$\bigcirc$ Nested
O Sequential
O Two or more of the above
O None of the above

## 242 points

3.1(C): If you had to choose one C Language loop construct to represent the loop(s) above, which of the following would you choose?
O While
O Do-While
$\bigcirc$ For

253 points
3.1(D): How many loop control variables are there, in total?

(Hint: A "loop control variable" is a variable that is evaluated as part of a loops test /and/ which is,or can be, changed within the loop's body or by its update (if a for loop).
(gdb) disassemble foo
Dump of assembler code for function foo:
0x0000000000400550<+0>: cmp
$\$ 0 \times 5$, \%esi
$0 \times 0000000000400553<+3>: j a$
0x40058b <foo 5 5 >
$0 \times 0000000000400555<+5>$ : mov
\%esi, \%eax
$0 \times 0000000000400557<+7>: j m p q$
*0x400630 (, \%rax, 8)
$0 \times 000000000040055 e<+14>$ : xchg
\%ax, \%ax
$0 \times 0000000000400560<+16>$ : add $\$ 0 x 2, \% e d i$
$0 \times 0000000000400563<+19>$ : mov
\%edi, \%eax
$0 \times 0000000000400565<+21>$ : mov $\$ 0 x 55555556$, \%edx
$0 x 000000000040056 a<+26>$ : sar
$\$ 0 x 1 f, \% e d i$
$0 \times 000000000040056 \mathrm{~d}<+29>:$ imul
\%edx
$0 \times 000000000040056 \mathrm{f}<+31>$ : sub \%edi, \%edx
$0 x 0000000000400571<+33>$ : mov
\%edx, \%eax
0x0000000000400573<+35>: retq
$0 \times 0000000000400574<+36>$ : nopl
$0 x 0$ (\%rax)
$0 \times 0000000000400578<+40>$ : add
\$0xa, \%edi
$0 \times 000000000040057 b<+43>: \quad$ lea
$0 x 0(, \% r d i, 4), \% e d x$
$0 \times 0000000000400582<+50>$ : mov
\%edx, \%eax
$0 \times 0000000000400584<+52>$ : retq
$0 \times 0000000000400585<+53>$ : nopl
(\%rax)
$0 \times 0000000000400588<+56>$ : and
$\$ 0 \times 1, \% e d i$
0x000000000040058b <+59>: lea
(\%rdi, \%rsi,1), \%edx
$0 \times 000000000040058$ e <+62>
mov
\%edx, \%eax
$0 \times 0000000000400590<+64>:$ retq

262 points
What is the output of the code shown?
$\square$

2 points
3.2(A): Blank (A): 0x

28 2 points
3.2(B): Blank (B): 0x
$\square$

292 points
3.2(C): Blank (C): 0x
$\square$

## 30

2 points

## 3.2(D): Blank (D): 0x

$\square$

312 points
3.2(E): Blank (E): 0x
$\square$

End of assembler dump.
(gdb) disassemble 0x400550 Dump of
assembler code for function foo:

$$
0 \times 0000000000400550<+0>: c m p
$$

$$
\$ 0 \times 5, \% \text { esi } 0 \times 0000000000400553<+3>:
$$

ja 0x40058b <foo+59>

$$
0 \times 0000000000400555<+5>: \text { mov }
$$

\%esi,\%eax 0x0000000000400557<+7>:
jmpq *0x400630(,\%rax, 8)

$$
0 \times 000000000040055 \mathrm{e}<+14>: \text { xchg }
$$

$$
\text { \%ax, \%ax } 0 x 0000000000400560<+16>:
$$

$$
\text { add } \$ 0 x 2, \% e d i \quad 0 x 0000000000400563
$$

$$
<+19>: \text { mov \%edi, \%eax }
$$

$$
0 \times 0000000000400565<+21>: \text { mov }
$$

$$
\$ 0 \times 55555556, \% e d x \quad 0 x 000000000040056 a
$$

$$
<+26>: \text { sar } \$ 0 x 1 f, \% e d i
$$

$$
0 x 000000000040056 \mathrm{~d}<+29>\text { : imul \%edx }
$$

$$
0 \times 000000000040056 \mathrm{f}<+31>\text { : sub }
$$

$$
\% e d i, \% e d x 0 x 0000000000400571<+33>:
$$

$$
\text { mov \%edx, \%eax } 0 x 0000000000400573
$$

$$
<+35>: \text { retq } 0 x 0000000000400574
$$

$$
<+36>: \text { nopl 0x0 (\%rax) }
$$

$$
0 \times 0000000000400578<+40>: \text { add }
$$

$$
\$ 0 x a, \% e d i \quad 0 x 000000000040057 b<+43>:
$$

$$
\text { lea } 0 x 0(, \% r d i, 4), \% e d x
$$

$$
0 \times 0000000000400582<+50>: \mathrm{mov}
$$

\%edx, \%eax 0x0000000000400584<+52>:

$$
\text { retq } 0 \times 0000000000400585<+53>\text { : nopl }
$$

$$
\text { (\%rax) } 0 \times 0000000000400588<+56>: \text { and }
$$

$$
\$ 0 x 1, \% e d i \quad 0 x 000000000040058 b<+59>:
$$

lea (\%rdi,\%rsi,1), \%edx

$$
0 x 000000000040058 \mathrm{e}<+62>: \mathrm{mov}
$$

$$
\text { \%edx, \%eax } 0 x 0000000000400590<+64>:
$$

retq End of assembler dump.

Please fill in the switch jump table corresponding to the gdb dump above. Do not include any leading zeros and note that the answer should be in hexadecimal without the leading $0 x$, as it is given.
(gdb) $x / 6 g$ 0x400630
0x400630: 0x $\qquad$ (A) $\qquad$
0 x $\qquad$ (B) $\qquad$
0×400640:
$0 x$ $\qquad$ (C) $\qquad$
$0 x$ $\qquad$ (D) $\qquad$
0x400650:
0 x $\qquad$ (E) $\qquad$
$0 x$ $\qquad$ (F) $\qquad$

Question 4: Stack Use and Calling

Convention

## Calling Convention and Stack Discipline

The following stack and register dump is from a Linux x86-64 machine like the shark hosts. It is taken immediately AFTER a function has been called, right before the first instruction within that function has been executed. The original function was written in the C Language.

```
(gdb) info registers
rax 0x6 6
rbx 0x0 0
rcx 0x4 4
rdx 0x9 9
rsi 0x8 8
rdi 0x6 6
rbp
0x7fffffffe0b0
rsp
0x7fffffffe0b0
r8
0x7ffff7dd5060
r9
0x7fffffffe528
r10 0x4 4
r11 0x0 0
r12 0x4004404195392
r13
0x7fffffffe1d0
r14 0x0 0
r15 0x0 0
rip 0x40053d
0x40053d <add+16>
(gdb) x/10xg 0x7ffffffffe0a8
0x7fffffffe0a8:
0x00007fffff7a44900
0x00007ffffffffe0f0
0x7fffffffe0b8:
0x000000000004005e9
0x00007ffffffffe1d8
0x7fffffffe0c8:
0x0000000700000000
4.1(A) 1st argument:
\(\square\)
\(\square\)

354 points 4.1(C) 3rd argument:
4.1(D): Return address: 0x
\(\square\)
37
 4 points Number of arguments:
\(\square\)

385 points 4.1(E) C Language data type for 3rd argument:
\(\bigcirc\) int
○ float
O long
\(\bigcirc\) double
〇 Unknowable
O None of the above

Please fill in the following, or indicate that the value is not knowable from the provided trace:

\section*{Question 5: Data}

\section*{Part 1: Structs}

Consider the following struct as compiled on a system using "natural alignment", i.e. the size of a data type is also its alignment requirement, and where chars are 1 byte, shorts are 2 bytes, ints are 4 bytes, and longs are 8 bytes, and then answer the questions that follow:
```

struct {
char c;
short s;
long l;
int i;
} initial;

```

39 2 points
5.1(A): How many bytes of alignment does the struct as a whole require?

\section*{40} 2 points
5.1(B): How many bytes of padding does the compiler add before the first (char c) field?
\(\square\)
5.1(C): How many bytes of padding does the compiler add after the last (int i) field?
\(\square\)
5.1(D): How many bytes of alignment does the compiler add between fields, e.g. neither at the beginning nor at the end?
\(\square\)

Please answer the questions to the right.
5.1(E): How many bytes can be saved in a single instance of the struct by reorganizing the fields?
\(\square\)

1 point
5.1(F): Given the reorganized struct you contemplated for (E) above, how many bytes would be saved across an array of four (4) such structs as compared to an array of four (4) of the original structs?

\section*{45}

2 points
5.2(A): In total, how many bytes are allocated, directly and/or indirectly, to array1? If you don't have enough information to answer or if the answer isn't knowable, write "-1".
\(\square\)
\(46 \quad 2\) points
5.2(B): What is the minimum number of bytes allocated directly to array2?
5.2(C): In total, how many bytes are allocated, directly and/or indirectly, to array2? If you don't have enough information to answer or if the answer isn't knowable, write "-1".
5.2(D): Consider the addresses of array1[1][1] and array \(1[3][2]\). What is the absolute difference as measured in bytes? If you don't have enough information to answer or if the answer isn't knowable, write "-1".

\section*{492 points}
5.2(E): Consider the addresses of array2[1][1] and array2[3][2]. What is the absolute difference as measured in bytes? If you don't have enough information to answer or if the answer isn't knowable, write "-1".
\(\square\)

\section*{\(50 \quad 2\) points}
5.2(F): If the entirety of array1 is initialized, is the value of array1[1][6], knowable? Yes or No


Yes
\(\bigcirc \mathrm{No}\)

\section*{512 points}
5.2(G): If the entirety of array2, including the indirect components, is initialized, is the value of array2[1][6], knowable? Yes or No
\(\bigcirc \mathrm{Yes}\)
\(\bigcirc \mathrm{No}\)

\section*{Question 6: Caching and Memory Access}

This question tests your understanding of cache behavior, asks you to simulate and describe the behavior of the same memory access trace on two different cache configurations, asks you some questions about the performance, and then asks you about the impact of caching upon memory access time.

Given the following information, please fill in the table below. If no set bits are decoded, fill in 0 for the set number.

The cache configuration for Part-1 is described as follows:
- 2-way set-associative ( \(\mathrm{E}=2\) )
- Address with = 6 bits
- Block size \(=8\) bytes
- 32byte total cache size
\begin{tabular}{|l|l|l|l|l|l}
\hline Time & \begin{tabular}{l} 
Mem \\
Addr \\
(Hex)
\end{tabular} & \begin{tabular}{l} 
Set \\
(Decimal)
\end{tabular} & \begin{tabular}{l} 
Tag \\
(Binary)
\end{tabular} & \begin{tabular}{l} 
Hit/Miss \\
(H/M)
\end{tabular} & \begin{tabular}{l} 
TyI \\
Mis \\
(Cl \\
Co \\
Ca \\
N//
\end{tabular} \\
\hline 0 & \(0 \times 1 \mathrm{~A}\) & (A) & (B) & (C) & (D) \\
\hline 2 & \(0 \times 2 \mathrm{~A}\) & & & & (E) \\
\hline 3 & \(0 \times 05\) & & & & \\
\hline 4 & \(0 \times 0 \mathrm{~A}\) & (F) & (G) & (H) & (I) \\
\hline 5 & \(0 \times 23\) & & & & \\
\hline 6 & \(0 \times 16\) & & & & (J) \\
\hline 6 & \(0 \times 00\) & & & & (K) \\
\hline
\end{tabular}

\section*{Part 2: Fully-Associative Cache}

Given the following information, please fill in the table below. If no set bits are decoded, fill in 0 for the set \#.
- Fully associative (All cache lines in same set)
- Address with \(=6\) bits
- 3 tag bits
- 32byte total cache size

54 1 point Blank (C) ○ (H) it O (M)iss

\section*{57 \\ 1 point} Blank (F)
\(\square\)
\begin{tabular}{|l|l|l|l|l|l}
\hline Time & \begin{tabular}{l} 
Mem \\
\\
Addr
\end{tabular} & \begin{tabular}{l} 
Set \\
(Decimal)
\end{tabular} & \begin{tabular}{l} 
Tag \\
(Binary)
\end{tabular} & \begin{tabular}{l} 
Hit/Miss \\
\((H / M)\)
\end{tabular} & \begin{tabular}{l} 
Tyı \\
Mi:
\end{tabular} \\
& & & & \begin{tabular}{l} 
(Cı \\
Co
\end{tabular}
\end{tabular}
\begin{tabular}{|l|l|l|l|l|l}
\hline & & & & & \begin{tabular}{l} 
Ca \\
\(\mathrm{N} / 4\)
\end{tabular} \\
\hline 0 & \(0 \times 1 \mathrm{~A}\) & (A) & (B) & (C) & (D) \\
\hline 2 & \(0 \times 2 \mathrm{~A}\) & & & & \\
\hline 3 & \(0 \times 05\) & & & & \((\mathrm{E})\) \\
\hline 4 & \(0 \times 0 \mathrm{~A}\) & (F) & (G) & (H) & \((\mathrm{I})\) \\
\hline 5 & \(0 \times 23\) & & & & \((\mathrm{~J})\) \\
\hline 6 & \(0 \times 16\) & & & & \\
\hline 6 & \(0 \times 00\) & & & & \((\mathrm{~K})\) \\
\hline
\end{tabular}

Part 3: Comparison
Please answer the question to the right.

\section*{Part 4: Memory Access}

Consider a memory system with the following properties:
- Level 1 cache: SRAM, 10nS access tie
- Main memory: DRAM, 100nS access time.
- Cache hit rate: \(95 \%\)

Please answer the questions to the right.

Blank (H)
O (H) it
O (M)iss

60
1 point
Blank (I)
O Cold
Conflict
O Capacity
\(\bigcirc N / A\)

\section*{611 point}

Blank (J)
\(\bigcirc\) Cold
Conflict
O Capacity
○ \(N / A\)

\section*{62} Blank (K)
\(\bigcirc\) Cold
Conflict
- Capacity

○ \(N / A\)

\section*{63} 1 point

Blank (A)

Blank (B): \(\square\)

651 point
Blank (C)
(H)it
\((M) i s s\)

661 point
Blank (D)
O Cold
O Conflict
O Capacity
O N/A

67
1 point
Blank (E)
Cold
Conflict
Capacity
N/A

681 point
Blank (F)

O (M)iss

71 1 point Blank (I)
Cold
Conflict
Capacity
N/A

721 point Blank (J)
Cold

Conflict
O Capacity
\(\bigcirc\) N/A
73 \begin{tabular}{ll}
1 point \\
Blank (K) \\
Oold \\
Conflict \\
Capacity \\
N/A
\end{tabular}
6.3: Did either cache configuration perform better for the given traces than the other? If so, how do you know
O They performed equally well for the given trace
\(\bigcirc\)
It isn't possible to know, given the traces provided
○
The cache configuration in Part 1 had fewer hits than the cache configuration Part 2, so the cache configuration in Part 2 performed better.

〇 The cache configuration in Part 1 had fewer misses than the cache configuration in Part 2, so the cache configuration in Part 1 performed better.

O None of the above

\section*{75}

1 point
6.4(A): What is the cache miss rate?

Fill in the blank: \(\qquad\) \%.

\section*{76 \\ 1 point}
6.4(B): What is the cache miss penalty (in nS )? Fill in the blank: \(\qquad\) nS.
\(\square\)

\section*{77}

1 point
6.4(C): What is the average access time to the nearest 0.01 nS ?
Fill in the blank: \(\qquad\) nS.

780 points
Feel free to provide us any feedback, comments, or notes here. For example, if you made any assumptions, etc. If you do, after the dust has settled (grades are back), please ping one of us and let us know that we should take a look. Remember -- grades can be adjusted at any time. And, we are humans, just like you. We're happy to discuss anything with you. Thanks!```

