

All Prefix Sums & Prescan

15-750

4/26/19

Let $\oplus$ be an associative binary op.

eg $+$, $\cdot$

Def All-Prefix-Sums

Input $[a_0, \dots, a_{n-1}]$

output $[a_0, a_0 \oplus a_1, a_0 \oplus a_1 \oplus a_2, \dots, a_0 \oplus \dots \oplus a_{n-1}]$

Def Prescan

output $[I, a_0, a_0 \oplus a_1, \dots, a_0 \oplus \dots \oplus a_{n-2}]$

Note $\text{PrefixSum}(A) = \text{Prescan}(A) \oplus A$

elementwise addition

Prescan Application

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Packing Out Memory

location	0	1	2	3	4	5	6	7
Contents	∅	*	∅	∅	*	*	∅	*

location	0	1	2	3	4	5	6	7
Contents	*	*	*	*	∅	∅	∅	∅

Use Prescan

Input (0, 1, 0, 0, 1, 1, 0, 1)

Prescan (0, 0, 1, 1, 1, 2, 3, 3)

Note Google: MapReduce not defined
for arrays!

Parallel Precan By Example

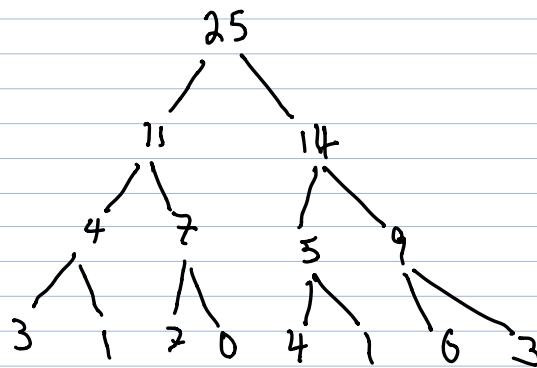
3

Input (3, 1, 7, 0, 4, 1, 6, 3) $\oplus$ addition

Precan Alg 1) Compute tree of partial Sums.

2) Set root to zero

3) Run Down!



Proc Down: I'll do

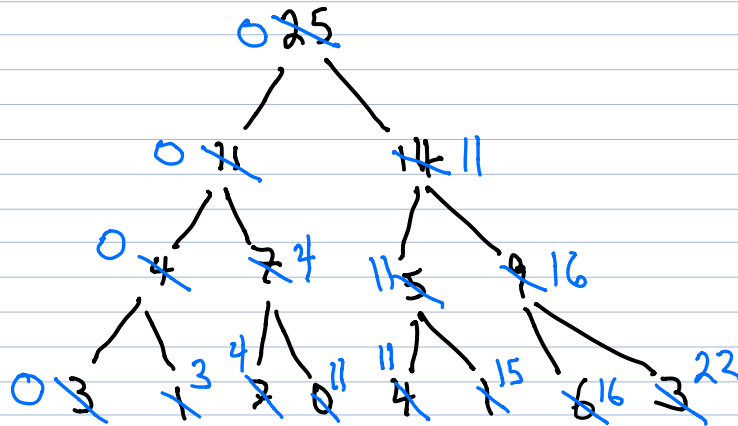
a) Right-Child $\leftarrow$ Parent $\oplus$ Left-child

b) Left-Child $\leftarrow$ Parent

c) Repeat on children

Down

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Preca (0, 3, 4, 11, 11, 15, 16, 22)

Input (3, 1, 7, 0, 4, 1, 6, 3)

$$T(n) = O(\log n)$$

$$P(n) = O(n)$$

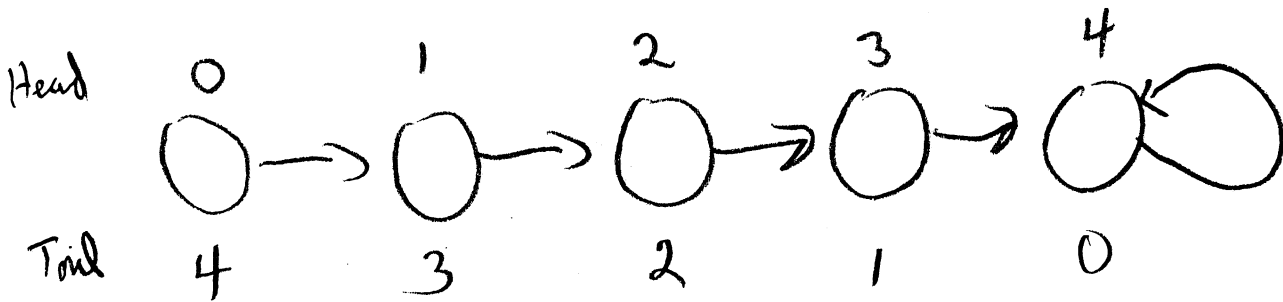
$$\Rightarrow P.T = O(n \cdot \log n)$$

List Ranking

input: linked-list

output: A mark on each node s.t.

mark = distance from head or
= distance to tail



Assume!

- 1) pointers are in consecutive memory
- 2) we know location of head & tail
- 3) pointers in arbitrary order.

Wyllie's Alg

Init rank(!) := 1 ; rank(tail) = 0

Init while succ(head) ≠ nil do

if succ(!) ≠ nil do

rank(!) := rank(!) ⊕ rank(succ(!))

succ(!) := succ(succ(!))

fi

processors = n

Goal

Time = $O(\lg n)$

$O(\lg n)$

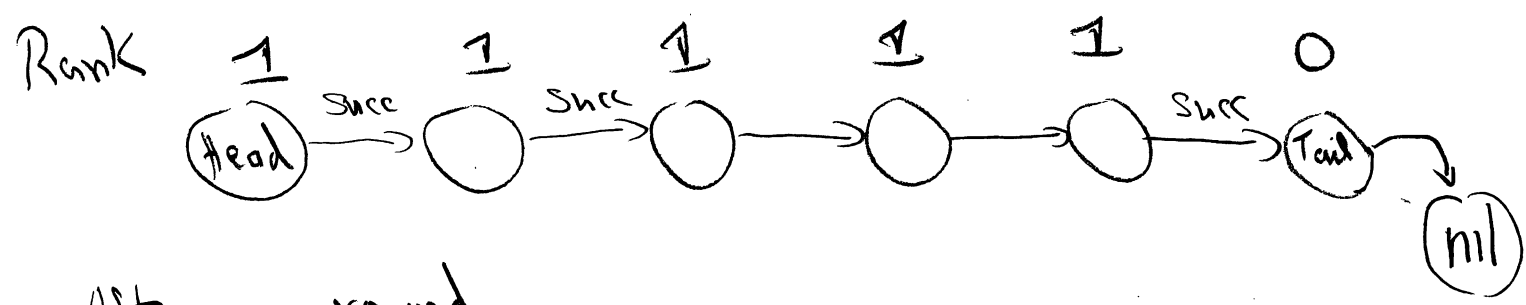
Work = $O(n \lg n)$

$O(n)$

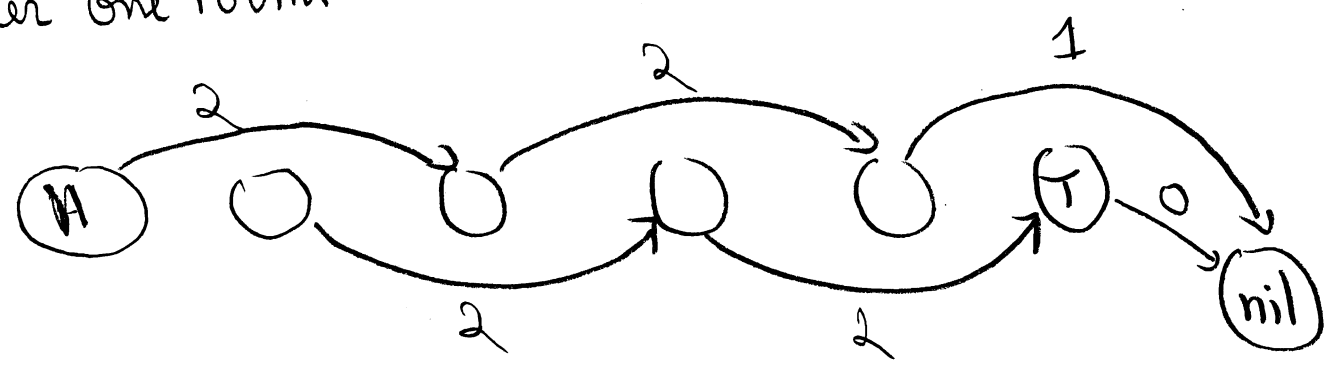
CREW model

Wyllie's Alg

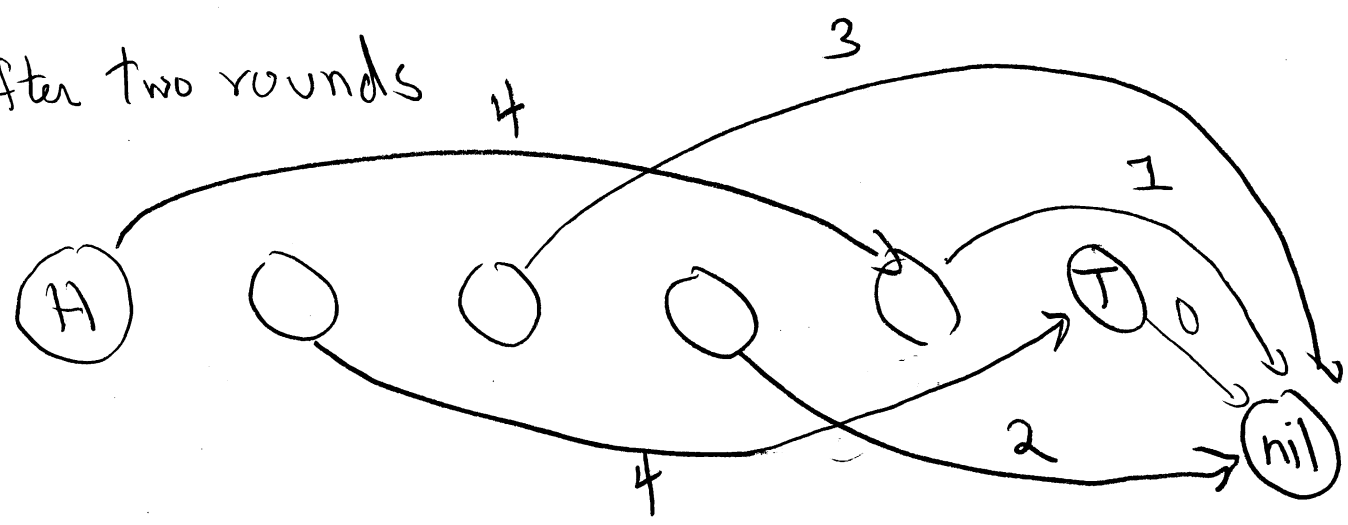
Initial



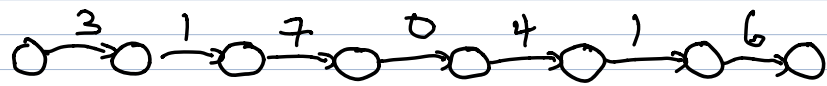
After one round



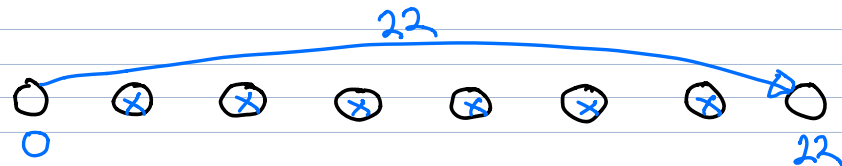
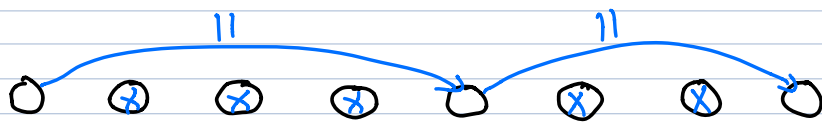
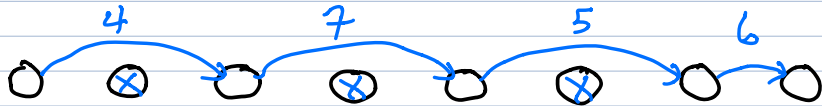
After two rounds



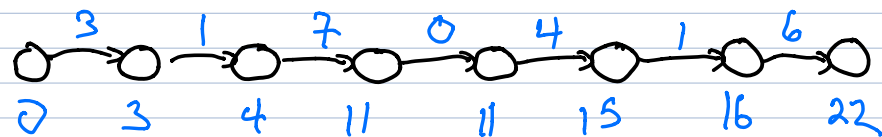
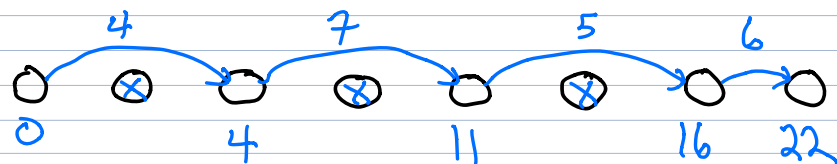
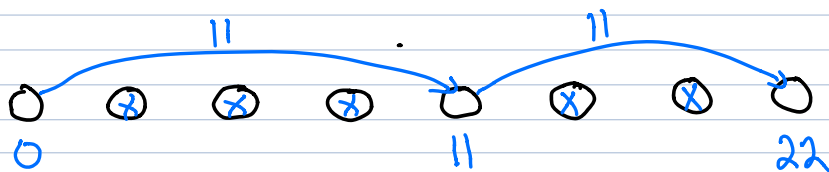
Prescan on a linked list



Recursively splice out every other element!



Fill in missing node values.

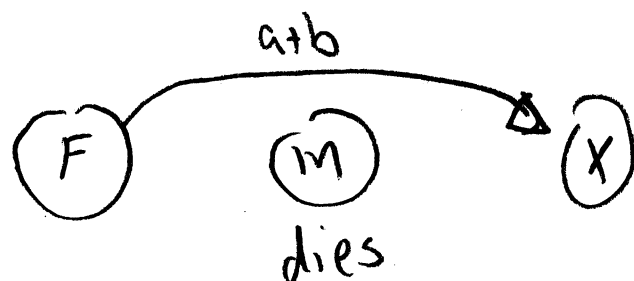


Random-Mate

1) Contraction Phase

1) Each live node randomly picks a sex

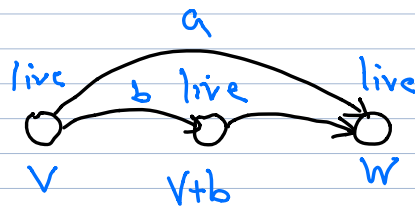
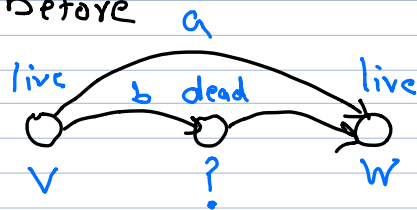
2) If $(F) \xrightarrow{a} (M) \xrightarrow{b} (X)$ then



3) Stop when head points to nil.
(only head is live)

In the expansion phase we run contraction phase "backward".

eg before



Random Mate Timing Analysis

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Thm The contraction phase for Random Mate terminates with high prob.

Pf Let $P_i^k \equiv$ Event that node i is still live after k rounds.

Note If node i not head then $\text{Prob}[P_i^1] = 3/4$

Since each round is independent

$$\text{Prob}[P_i^k] = (3/4)^k \text{ for node } i \text{ not head.}$$

$$\text{Set } k = \lfloor \log_{4/3} n \rfloor$$

$$\text{Prob}[P_i^k] = (3/4)^k \leq \frac{1}{(4/3)^{\lfloor \log_{4/3} n \rfloor}} = \frac{1}{n^c}$$

Let $P^k \equiv$ Event that some non-head node is still live.

Assume that node 0 is the head.

$$P^k = P_1^k \cup \dots \cup P_n^k$$

$$\text{Prob}[P^k] = \text{Prob}[P_1^k \cup \dots \cup P_n^k]$$

$$\leq \text{Prob}[P_1^k] + \dots + \text{Prob}[P_n^k] \leq \frac{n}{n^c} = \frac{1}{n^{c-1}}$$

□

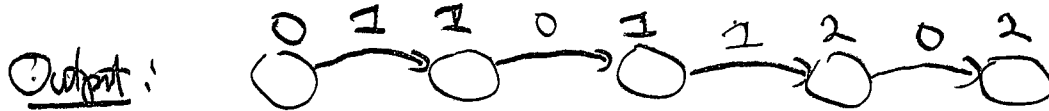
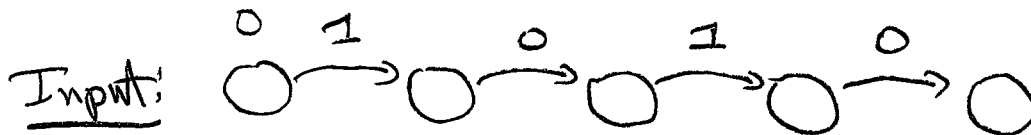
Euler Tours

Prob: Preorder number a tree

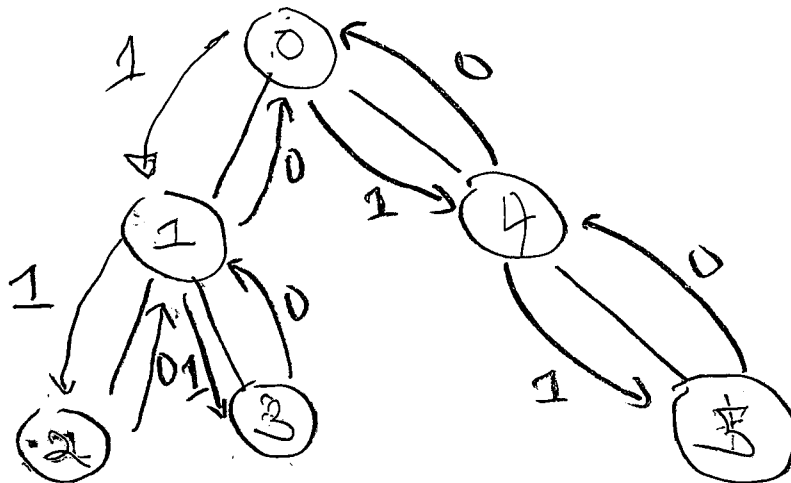
Input: A tree stored using pointers

Output: Preorder numbering

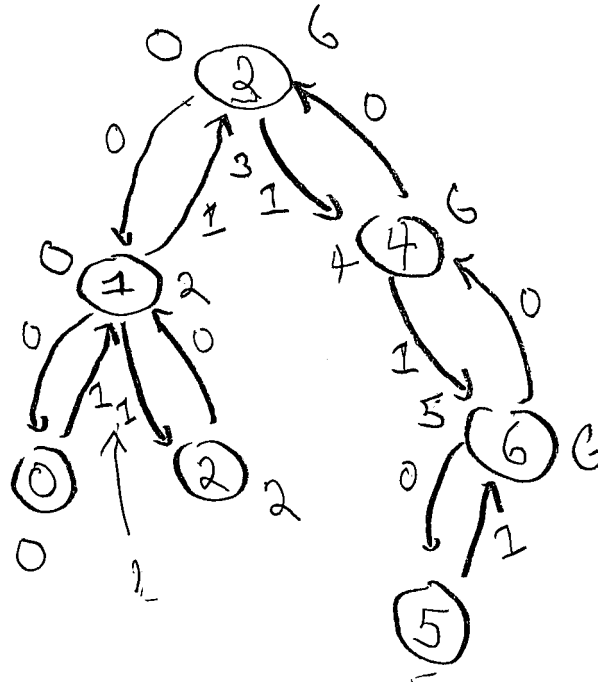
Weighted List Ranking



Preorder



Use First-Visit value.

Preorder1 $\equiv$ weight of down edge0 $\equiv$ weight of up edgeInorder

left-edge

0/1

right-edge

1/0

Use "Middle-Visit" Value

Postorder?

Size of Subtree?