

**15-451/651 Algorithms, Fall 2017**  
**Recitation #12 Worksheet**

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## Multiplicative Weights

1. In lecture we saw that the simple procedure that multiplied the weight of each expert by  $\frac{1}{2}$  whenever the expert made a mistake, resulted in

$$m = \text{\#mistakes of algorithm} \leq 2.41(M + \log_2 n),$$

where  $M = \text{\#mistakes made by the best expert}$  and  $n = \text{\# of experts}$ . If we change the weight by  $2/3$  at each time, how does this analysis change?

2. In the lecture: in order to get a better mistake bound of  $(1 + \epsilon)M + O(\frac{\ln n}{\epsilon})$ , we used randomization. Let us now show that you cannot get better than a factor of 2 if you don't use randomness.

*There are two experts. One always predicts 0. The other always predicts 1. Fix any deterministic algorithm  $A$  for prediction. Here is one sequence of days: each day, the actual outcome is the opposite of what the algorithm predicts.*

After  $T$  days, the algorithm would have made  $T$  mistakes. Show that the better of the two experts makes at most  $T/2$  mistakes. Hence infer that  $m \geq 2M$ .

## Finding misplaced keys (Online Algorithms)

You are at the origin of  $\mathbb{Z}$  and have misplaced your keys at some integer  $n$ , the value of  $n$  is unknown to you. If you are at position  $i$ , in a single step you can move to  $i + 1$  or  $i - 1$ . To check if your keys are at some position  $i$ , you must visit  $i$ . Your goal is to find your keys while minimizing the number of steps you take. If you knew what  $n$  was, you would simply walk there in  $n$  steps, but since you don't know the value of  $n$ , you are satisfied with taking  $cn$  steps, where  $c$  is some small constant.

1. A natural strategy to try is to visit nodes in the following order:  $1, -1, 2, -2, 3, -3, \dots$ . Show that this is not a good strategy, that is, it has a competitive ratio of  $\Omega(n)$ .
2. Consider this alternate strategy. You visit nodes in the order  $1, -1, 2, -2, 4, -4, 8, -8, \dots$ . Show that this strategy has a competitive ratio of at most 13.
3. As a slight optimization over part (b), visit nodes in order  $1, -2, 4, -8, 16, -32, \dots$ . Show that this strategy has a competitive ratio of at most 9.