15-213/18-243 Introduction to Computer Systems Spring 2009 Recitation 7

Your Name Here

Overview

- System Calls
 - fork()
 - execve()
 - wait()
 - exit()
- **■** Signals
- Shell Lab
 - Race Conditions
 - I/O Redirection

- fork()
 - Splits the execution of one process into two processes which execute concurrently.
 - Returns the PID of the child process in the parent.
 - Returns 0 in the child.
 - The parent's address space is COPIED into the child's address space.
 - Writes to variables, arrays, etc. are not visible in the parent
 - File descriptors (stdin, stdout, stderr) are SHARED.

- execve(char* path, char** argv, char** envp)
 - Loads the specified executable file and executes it in the current process.
 - argv is the argument vector passed to main().
 - envp is the environment variable list.
 - This is a common form of the exec system call others exist (execv, execl, etc.).

- exit(int status)
 - Causes the current process to terminate and sets the exit status.
 - Exit status 0 usually denotes success, any other value denotes failure.
 - Wait can retrieve the status set by exit().

- wait(int* child_status)
 - Waits (blocks) until ANY child of the current process terminates.
 - Stores the exit status of the child in *child_status and then returns to the parent.

- waitpid(int child_pid, int* child_status, int options)
 - Waits for a specific child to exit, or all children if child_pid = -1.
 - Blocking and other behavior can be configured via the options variable (a bit vector)
 - WNOHANG don't block if no children have exited, just return 0
 - WUNTRACED also return if the child has only stopped (received the SIGSTOP signal), not exited

Overview

- Signals are "sent" to a process when an exceptional event occurs.
- Process "receives" signals by jumping to a signal handler function.
- Signals can be received (and thus signal handlers may be called) at any time during process execution.

- Useful Signals
 - SIGINT sent when the interrupt sequence (usually Ctrl-C) is entered.
 - SIGSTOP sent when the stop sequence (usually Ctrl-Z) is entered.
 - SIGCHLD sent when any child process exits
 - SIGSEGV sent when a process makes an improper request to read or write memory (SEGmentation Violation).

- Signals are not queued!
 - If the same signal is sent more than once to a process before the process receives it, the signal handler function for the process will only be called once.
 - For example, if two children exit before the signal handler is called in the parent, the handler for SIGCHLD will only be called once.

- Blocking Signals
 - Signals can be blocked (rather, deferred) so that the program does not run a signal handler during a "critical section" of code.
 - Use the sigprocmask() API to block and unblock signals.
 - Signals sent to a process while blocked will be received when (if) the process unblocks them.
 - Again, if multiple signals of the same type are sent to a process while the signal is blocked, the signal will only be received ONCE when the signal is unblocked.

Overview

- Your goal is to create a functioning Unix shell with foreground and background processes.
- Exercises the process control system calls as well as signal handling.
- You must write a main command line evaluation function which parses the command line and handles the creation of child processes.
- You must also write signal handlers to detect and clean up terminated child processes.
- Beware of race conditions!

- Race Conditions
 - A race condition occurs when the state of a program is *unexpectedly nondeterministic*.
 - Nondeterminism occurs when multiple concurrent threads of execution operate upon a shared set of data, and the final value of the data changes depending on the order in which the threads are scheduled.
 - If a piece of code unexpectedly produces nondeterministic results, that code is said to have a race condition.

- Race Conditions and Signals
 - Recall that, in the absence of blocking, signals can be received by a program at any time.
 - Suppose both the main thread of execution and a signal handler modify shared state.
 - If the main thread is modifying this state and the signal handler is called during this operation, the signal handler sees *inconsistent state*.
 - If the signal handler then decides to use the shared state, the program may begin to behave in unexpected ways (if you are lucky, it will crash).

Race Conditions and Signals: A Trivial Example

```
int i = 0:
void sigalrm handler() {
  i++;
int main(...) {
  int temp;
  /* set up signal handling */
  alarm(1);
  temp = i;
  temp = temp + 1;
  i = temp;
  pause();
  printf("%d", i);
  return 0;
```

What is the output?

Depends on when the signal handler is called!

If the alarm expires during the execution of these three lines, the program outputs 1 instead of 2!

Race Conditions and Signals: A Trivial Example

```
int i = 0;
void sigalrm handler() {
  i++;
int main(...) {
  int temp;
  /* set up signal handling */
  alarm(1);
  sigprocmask(SIG BLOCK, SIGALRM);
  temp = i;
  temp = temp + 1;
  i = temp;
  sigprocmask(SIG UNBLOCK, SIGALRM);
  pause();
  printf("%d", i);
  return 0;
```

How do we make this program deterministic (i.e. make it output 2 on every execution)?

Block the alarm signal during the critical section!

Critical Section

- I/O Redirection
 - You may need to set the stdin or stdout of a child process to file descriptors referencing files on disk.
 - Use dup2 (old_fd, new_fd) to do this.
 - For example, this redirects stdout to a file:

```
int main(int argc, char** argv) {
  int fd = open("/tmp/foo", O_TRUNC);
  dup2(fd, STDOUT_FILENO);
  printf("Hello World!");
  return 0;
}
```

/tmp/foo will contain the text "Hello World!"

Be careful where you call dup2 in your shell!

- There are lots more hints in the shell lab handout, be sure to read these!
- Be sure to review the lecture slides if you get stuck.
- Come to office hours or email the staff list if you have any questions.

Questions?