### 15-213-S09 Recitation #1

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## Little Endian to Big Endian

• 32-bit unsigned integer: 0x257950B2n =  $37 * 2^{24} + 121 * 2^{16} + 80 * 2^8 + 178$ 

• Little Endian: 31 24

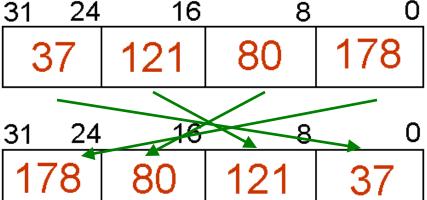
<u>31</u>	24	16	8	0
3	7	121	80	178

Big Endian:

31	24	16	8	0
17	'8	08	121	37

### Little Endian to Big Endian

· Little Endian:



```
• Big Endian: 178 80 121 37
```

```
unsigned int LE2BE(unsigned int x)

{
  unsigned int result =
    (x << 24) \mid ((x << 8) \& (0xFF << 16)) \mid ((x >> 8) \& (0xFF << 8)) \mid ((x >> 24) \& 0xFF);

return result;
}
```

### If-then-else

```
if (condition)
  expression1;
else
  expression2;
```

### If-then-else

```
if (condition)
  expression1;
else
  expression2;
```

```
Rewrite as:

( condition & expression1) |
(~condition & expression2)
```

### Datalab Example - Max

```
Return the max of 2 ints (x, y)
2 cases:
   Different signs – look at sign of x
   Same sign – look at sign of x-y
int max (int x, int y) {
   int sub = x - y;
   int signcomp = x \wedge y; // compare the signs
   // if signs are similar, take the sign of x-y, if not, take the sign of x.
   //This will be 1 if y is greater, 0 otherwise
   int mask = ((signcomp) \& x) | (\sim signcomp \& sub)) >> 31;
   return (mask & y) | (~mask & x);
```

### Datalab Example - Least Bit Pos

Return a mask of the least bit position of n

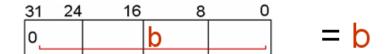
```
Take a binary number x (1011...010..0)
Flip the bits (0100...101...1)
Add 1 to it (0100...110...0)
Only 1's occupying the same position between each the new number and original are in the lowest bit position
int leastBitPos(int x) {
```

```
return x & (~x + 1);
```

### 2's complement: some examples

$$\frac{31 \quad 24}{1} \quad \frac{16}{b} \quad \frac{8}{1} \quad 0 = -2^{31} + b$$

• 
$$0x0 = 0$$



- 0x1 = 1
- $0x7FFFFFFFFF = 2^{31}-1$  // largest 32-bit int
- 0xFFFFFFF = -1
- 0xFFFFFFFE = -2
- $0x8000000000 = -2^{31}$  // smallest 32-bit int
- $0x80000001 = -2^{31} + 1$

### Simple 8-bit Integer Example

```
Some simple, positive examples:
```

```
0b0001 0110

= 128(0) + 64(0) + 32(0) + 16(1) + 8(0) + 4(1) + 2(1) + 1(0)

= 16 + 4 + 2

= 22

53

= 32 + 16 + 4 + 1

= 128(0) + 64(0) + 32(1) + 16(1) + 8(0) + 4(1) + 2(0) + 1(1)

= 0011 \ 0101
```

### Signed 8-bit Integer Examples

```
Using -x = (-x) + 1

-14

= -(14)

= -(8(1) + 4(1) + 2(1) + 1(0))

= -(0000 1110)

= -(0000 1110) + 1

= (1111 0001) + 1

= 1111 0010
```

Using negative sign bit (for numbers close to tMin)
-120
= -128(1) + 8(1)
= 1000 1000

### Other Signed 8-bit Integer Examples

$$-10 =$$

$$-100 =$$

### Other Signed 8-bit Integer Examples

$$12 = 0000 \ 1100$$

$$63 = 0011 \ 1111$$

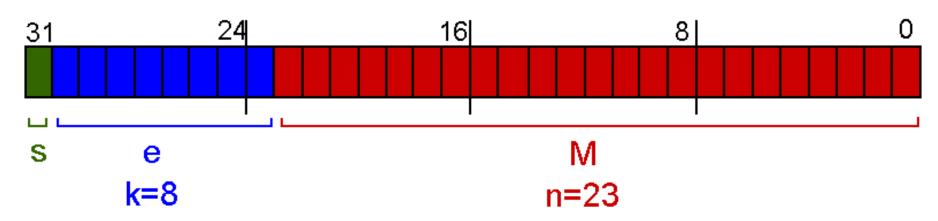
$$72 = 0100 \ 1000$$

$$-10 = \sim (0000\ 1010) + 1 = 1111\ 0101 + 1 = 1111\ 0110$$

$$-100 = -128 + 28 = 1001 1100$$

# Floating point

### Single precision:

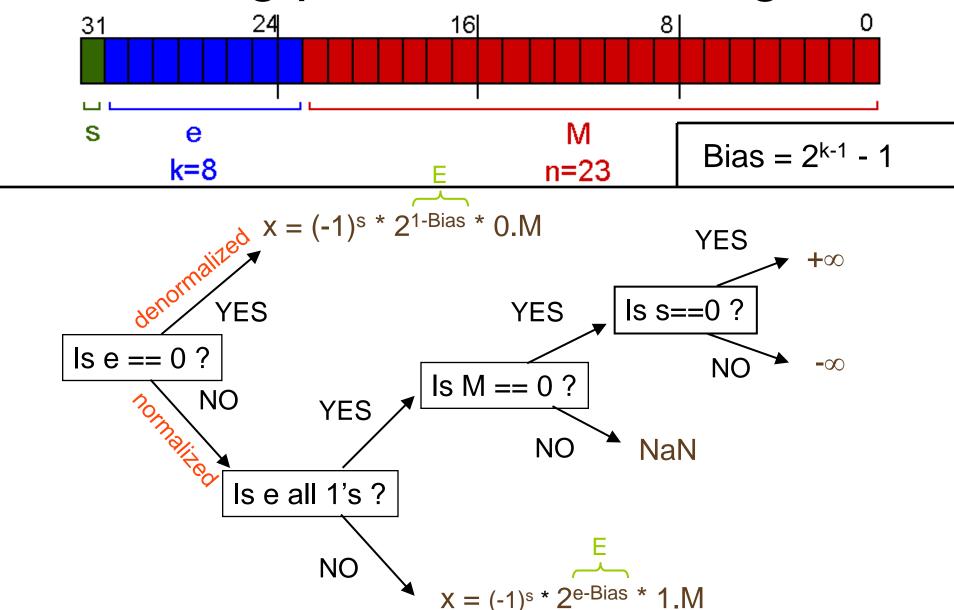


Note: exponent boundary is NOT aligned with byte boundary e.g. 0xFF7FFFF has lowest exponent bit zero (is normalized v.)

#### Double precision:

$$k = 11, n = 52$$

# Floating point decision diagram

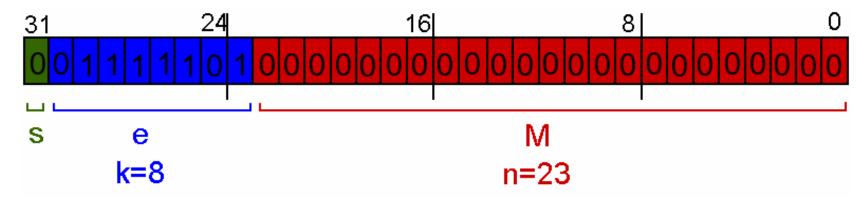


### Example

$$1/4 = 1.0 * 2^{-2}$$

$$s=0$$
,  $e = -2 + Bias = 125$ ,  $M = 0$ 

representation = 3E800000



# 8-bit Floating Point Example

```
8-bit floating point (1 sign, 3 exponent, 4 fraction)
  bias = 2^{(3 - 1)} - 1 = 3
encode 13/16
sign bit = 0 (13/16 >= 0)
13/16 = 13 \times 2^{-4}
  = 1101 \times 2^{(-4)}
  = 1.1010 \times 2^{(-1)} (remove the 1 because it's normalized)
exponent = -1 + bias = 2 = 010
result = 0.010 1010
```

# 8-bit Floating Point Example

```
8-bit floating point (1 sign, 3 exponent, 4 fraction)
  bias = 2^{(3 - 1)} - 1 = 3
encode -3/32
sign bit = 1(-3/32 < 0)
3/32 = 3 \times 2^{-5}
  = 0b11 \times 2^{(-5)}
  = 1.1 \times 2^{(-4)} Note: Lowest possible exponent is -2
                         Must denormalize
  = .011 \times 2^{(-2)}
exponent = 0 (denormalized)
result = 1 000 0110
```

# Other 8-bit Floating Point Examples

$$10/4 =$$

$$-10 =$$

$$1/64 =$$

# Other 8-bit Floating Point Examples

$$10/4 = 0\ 100\ 0100$$

$$-3 = 1 100 1000$$

$$-10 = 1 \ 110 \ 0100$$

$$1/64 = 0\ 000\ 0001$$

## Good coding style

- Consistent indentation
- Avoid long sequences of commands without a comment
- Each source file should have an appropriate header
- Have a brief comment at the beginning of each function

### Version Control - Do It

Version Control allows you to maintain old versions of your source code, potentially saving you if you were to do something like rm -rf \*

#### Good SVN tutorial:

http://artis.imag.fr/~Xavier.Decoret/resources/svn/index.html

#### Note:

repoDir=`pwd` #of working dir
Path to repo is svn+ssh://unix.andrew.cmu.edu/\${repoDir}

#### Example:

svn+ssh://unix.andrew.cmu.edu/afs/andrew.cmu.edu/usr17/agartrel/svn/15-410