Bits, Bytes and Integers – Part 1

15-213/18-213/15-513: Introduction to Computer Systems 2nd Lecture, May 23, 2018

Instructors:

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Waitlist questions

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- 18-213: Zara Collier (zcollier@andrew.cmu.edu)
- 15-513: Amy Weis <u>alweis@andrew.cmu.edu</u>
- Please don't contact the instructors with waitlist questions.

Bootcamp

- Linux basics
- Git basics

Things like:

- How to ssh to the shark machines from windows or linux
- How to setup a directory on afs with the right permissions
- How to initialize a directory for git
- The basics of using git as you work on the assignment
- Basic linux tools like: tar, make, gcc, ...

First Assignment: Data Lab

- Datalab is out this afternoon
- Due: Thursday, 5/31 at 11:59pm
- Absolute last time to turn in: Saturday, 6/2 at 11:59pm
- Go to GitHub/Autolab soon and read the handout carefully
- Start early
- Don't be afraid to ask for help
 - Piazza
 - Office hours
- Based on lectures 2, 3 and 4

Today: Bits, Bytes, and Integers

Representing information as bits

- Bit-level manipulations
- Integers
 - Representation: unsigned and signed
 - Conversion, casting
 - Expanding, truncating
 - Addition, negation, multiplication, shifting
 - Summary
- Representations in memory, pointers, strings

Everything is bits

- Each bit is 0 or 1
- By encoding/interpreting sets of bits in various ways
 - Computers determine what to do (instructions)
 - ... and represent and manipulate numbers, sets, strings, etc...

Why bits? Electronic Implementation

- Easy to store with bistable elements
- Reliably transmitted on noisy and inaccurate wires



For example, can count in binary

Base 2 Number Representation

- Represent 15213₁₀ as 11101101101₂
- Represent 1.20₁₀ as 1.001100110011[0011]...,
- Represent 1.5213 X 10⁴ as 1.1101101101101₂ X 2¹³

Encoding Byte Values

Byte = 8 bits

- Binary 00000000, to 1111111,
- Decimal: 0₁₀ to 255₁₀
- Hexadecimal 00₁₆ to FF₁₆
 - Base 16 number representation
 - Use characters '0' to '9' and 'A' to 'F'
 - Write FA1D37B₁₆ in C as
 - 0xFA1D37B
 - Oxfa1d37b

He	^t De ^r	timal Binary
0	0	0000
1	1	0001
2	2	0010
З	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
Α	10	1010
В	11	1011
C	12	1100
D	13	1101
Ε	14	1110
F	15	1111

Example Data Representations

C Data Type	Typical 32-bit	Typical 64-bit	x86-64
char	1	1	1
short	2	2	2
int	4	4	4
long	4	8	8
float	4	4	4
double	8	8	8
pointer	4	8	8

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Boolean Algebra

Developed by George Boole in 19th Century

- Algebraic representation of logic
 - Encode "True" as 1 and "False" as 0

And

&

```
A&B = 1 when both A=1 and B=1
```

0 1

0

Or

Not

• ~A = 1 when A=0

0 0 0

Exclusive-Or (Xor)
• A^B = 1 when either A=1 or B=1, but not both

0

$$\begin{array}{c|ccc}
^{\ } 0 & 1 \\
\hline
0 & 0 & 1 \\
1 & 1 & 0
\end{array}$$

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General Boolean Algebras

Operate on Bit Vectors

Operations applied bitwise

	01101001	01101001		01101001		
<u>&</u>	01010101	01010101	^	01010101	~	01010101
	01000001	01111101		00111100		10101010

All of the Properties of Boolean Algebra Apply

Example: Representing & Manipulating Sets

Representation

- Width w bit vector represents subsets of {0, ..., w-1}
- a_i = 1 if j ∈ A
 - 01101001 { 0, 3, 5, 6 }
 - 7<u>65</u>4<u>3</u>210
 - 01010101 { 0, 2, 4, 6 }
 - 7<u>6</u>5<u>4</u>3<u>2</u>10

Operations

- & Intersection
 01000001
 { 0, 6 }
- Union
 01111101 { 0, 2, 3, 4, 5, 6 }
- Symmetric difference 00111100 { 2, 3, 4, 5 }
- Complement 10101010 { 1, 3, 5, 7 }

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Bit-Level Operations in C

■ Operations &, |, ~, ^ Available in C

- Apply to any "integral" data type
 - long, int, short, char, unsigned
- View arguments as bit vectors
- Arguments applied bit-wise

Examples (Char data type)

- ~0x41 →
- ► ~0x00 →
- 0x69 & 0x55 →
- 0x69 | 0x55 →

He	⁺ De ^r	cimal Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
Α	10	1010
В	11	1011
C	12	1100
D	13	1101
Е	14	1110
F	15	1111

Bit-Level Operations in C

■ Operations &, |, ~, ^ Available in C

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Examples (Char data type)

- ~0x41 → 0xBE
 - $\sim 0100\ 0001_2 \rightarrow 1011\ 1110_2$
- $\sim 0x00 \rightarrow 0xFF$
 - $\sim 0000\ 0000_2 \rightarrow 1111\ 1111_2$
- 0x69 & 0x55 → 0x41
 - $0110\ 1001_2$ & $0101\ 0101_2 \rightarrow 0100\ 0001_2$
- $0x69 \mid 0x55 \rightarrow 0x7D$
 - $0110\ 1001_2 \mid 0101\ 0101_2 \rightarrow 0111\ 1101_2$



Contrast: Logic Operations in C

Contrast to Bit-Level Operators

- Logic Operations
 - View 0 as "Fals
 - Anything nonzo
 - Alway
 - Early Watch out for && vs. & (and || vs. |)...
- Example one of the more common oopsies in
 - IOx41 → C programming
 - !0x00 →
 - ‼0x41→ 0x01
 - 0x69 && 0x55 → 0x01
 - 0x69 || 0x55 → 0x01
 - p && *p (avoids null pointer access)

Shift Operations

Left Shift: x << y</p>

- Shift bit-vector x left y positions
 - Throw away extra bits on left
 - Fill with 0's on right

Right Shift: x >> y

- Shift bit-vector x right y positions
 - Throw away extra bits on right
- Logical shift
 - Fill with 0's on left
- Arithmetic shift
 - Replicate most significant bit on left

Undefined Behavior

Shift amount < 0 or ≥ word size

${\bf Argument } {\bf x}$	01100010
<< 3	00010000
Log. >> 2	<i>00</i> 011000
Arith. >> 2	<i>00</i> 011000

Argument x	10100010
<< 3	00010000
Log. >> 2	<i>00</i> 101000
Arith. >> 2	11101000

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Encoding Integers



C short 2 bytes long

	Decimal	Hex	Binary	
x	15213	3B 6D	00111011 01101101	
У	-15213	C4 93	11000100 10010011	

Sign Bit

- For 2's complement, most significant bit indicates sign
 - 0 for nonnegative
 - 1 for negative

Two-complement: Simple Example

$$10 = \begin{bmatrix} -16 & 8 & 4 & 2 & 1 \\ 0 & 1 & 0 & 1 & 0 \end{bmatrix}$$

$$8+2 = 10$$

$$-10 = \begin{bmatrix} -16 & 8 & 4 & 2 & 1 \\ 1 & 0 & 1 & 1 & 0 \end{bmatrix}$$

$$-16+4+2 = -10$$

Two-complement Encoding Example (Cont.)

x =	15213:	00111011		01101101	
у =	-15213:	1100	0100	100	10011
Weight	15213			-1521	L3
1	1	1		1	1
2	0	о		1	2
4	1	4		0	0
8	1	8		0	0
16	0	0		1	16
32	1	32		0	О
64	1	64		0	0
128	0	0		1	128
256	1	256		0	0
512	1	512		0	0
1024	0	0		1	1024
2048	1	2048		0	0
4096	1	4096		0	0
8192	1	8192		0	0
16384	0	0		1	16384
-32768	0	0		1	-32768
Sum		15213			-15213

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Numeric Ranges

- Unsigned Values
 - *UMin* = 0
 - 000...0
 - $UMax = 2^w 1$

111...1

Two's Complement Values

- $TMin = -2^{w-1}$ 100...0
- $TMax = 2^{w-1} 1$

011...1

Minus 1

111...1

Values for *W* = 16

	Decimal	Hex	Binary
UMax	65535	FF FF	11111111 11111111
TMax	32767	7F FF	01111111 11111111
TMin	-32768	80 00	1000000 00000000
-1	-1	FF FF	11111111 11111111
0	0	00 00	0000000 00000000

Values for Different Word Sizes

	W					
	8	16	32	64		
UMax	255	65,535	4,294,967,295	18,446,744,073,709,551,615		
ТМах	127	32,767	2,147,483,647	9,223,372,036,854,775,807		
TMin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808		

Observations

- *|TMin| = TMax + 1*
 - Asymmetric range
- *UMax*= 2 * *TMax* + 1

C Programming

- #include <limits.h>
- Declares constants, e.g.,
 - ULONG_MAX
 - LONG_MAX
 - LONG_MIN
- Values platform specific

Unsigned & Signed Numeric Values

X	B2U(<i>X</i>)	B2T(<i>X</i>)
0000	0	0
0001	1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	-8
1001	9	-7
1010	10	-6
1011	11	-5
1100	12	-4
1101	13	-3
1110	14	-2
1111	15	-1

Equivalence

Same encodings for nonnegative values

Uniqueness

- Every bit pattern represents unique integer value
- Each representable integer has unique bit encoding

■ ⇒ Can Invert Mappings

- $U2B(x) = B2U^{-1}(x)$
 - Bit pattern for unsigned integer
- $T2B(x) = B2T^{-1}(x)$
 - Bit pattern for two's comp integer

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Integers

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Mapping Between Signed & Unsigned



Mappings between unsigned and two's complement numbers: Keep bit representations and reinterpret

Mapping Signed ↔ Unsigned

Bits	Signed		Unsigned
0000	0		0
0001	1		1
0010	2		2
0011	3		3
0100	4		4
0101	5		5
0110	6		6
0111	7	│ ╡╶ <u>╷</u> ┲┘╡── │	7
1000	-8	•	8
1001	-7		9
1010	-6		10
1011	-5		11
1100	-4		12
1101	-3		13
1110	-2		14
1111	-1		15

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Mapping Signed ↔ Unsigned



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Relation between Signed & Unsigned



Conversion Visualized

2's Comp. \rightarrow Unsigned



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Signed vs. Unsigned in C

Constants

- By default are considered to be signed integers
- Unsigned if have "U" as suffix

OU, 4294967259U

Casting

Explicit casting between signed & unsigned same as U2T and T2U

int tx, ty; unsigned ux, uy; tx = (int) ux; uy = (unsigned) ty;

Implicit casting also occurs via assignments and procedure calls

tx = ux; int fun(unsigned u); uy = ty; uy = fun(tx);

Casting Surprises

Expression Evaluation

- If there is a mix of unsigned and signed in single expression, signed values implicitly cast to unsigned
- Including comparison operations <, >, ==, <=, >=
- Examples for W = 32: TMIN = -2,147,483,648, TMAX = 2,147,483,647

Constant ₁ Co		Со	nstant ₂ Relation Evaluation
0	0U	==	unsigned
-1	0	<	signed
-1	0U	>	unsigned
214	4748364	7-21	47483647-1 > signed
214	4748364	7U	-2147483647
-1	-2	>	signed
(un	signed)-	1-2	> unsigned
21	4748364	7	2147483648 ⊌ unsigned
21	4748364	7	(int) 2147483648 signed

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Unsigned vs. Signed: Easy to Make Mistakes

```
Can be very subtle
```

```
#define DELTA sizeof(int)
int i;
for (i = CNT; i-DELTA >= 0; i-= DELTA)
. . .
```

Summary

Casting Signed ↔ Unsigned: Basic Rules

- Bit pattern is maintained
- But reinterpreted
- Can have unexpected effects: adding or subtracting 2^w
- Expression containing signed and unsigned int
 - int is cast to unsigned!!

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Sign Extension

Task:

- Given *w*-bit signed integer *x*
- Convert it to w+k-bit integer with same value

Rule:

Make k copies of sign bit:



Sign Extension: Simple Example



Larger Sign Extension Example

short int x = 15213; int ix = (int) x; short int y = -15213; int iy = (int) y;

	Decimal	Hex	Binary
x	15213	3B 6D	00111011 01101101
ix	15213	00 00 3B 6D	00000000 00000000 00111011 01101101
У	-15213	C4 93	11000100 10010011
iy	-15213	FF FF C4 93	11111111 1111111 11000100 10010011

Converting from smaller to larger integer data type

C automatically performs sign extension

Truncation

Task:

- Given k+w-bit signed or unsigned integer X
- Convert it to w-bit integer X' with same value for "small enough" X

Rule:

- Drop top *k* bits:
- $X' = x_{w-1}, x_{w-2}, \dots, x_0$



Truncation: Simple Example



Summary: Expanding, Truncating: Basic Rules

Expanding (e.g., short int to int)

- Unsigned: zeros added
- Signed: sign extension
- Both yield expected result

Truncating (e.g., unsigned to unsigned short)

- Unsigned/signed: bits are truncated
- Result reinterpreted
- Unsigned: mod operation
- Signed: similar to mod
- For small numbers yields expected behavior

Fake real world example

- Acme, Inc. has developed a state of the art voltmeter they are connecting to a pc. It is precise to the millivolt and does not drain the unit under test.
- Your job is to develop the driver software.



printf("%d\n", getValue());

Fake real world example

- Acme, Inc. has developed a state of the art voltmeter they are connecting to a pc. It is precise to the millivolt and does not drain the unit under test.
- Your job is to develop the driver software.



Lets run some tests

printf("%d\n", getValue());

Lets run some tests

int x=getValue(); printf("%d %08x\n",x, x);

- **50652** 0000c5dc
- 1500 000005dc
- 9692 000025dc
- 26076 000065dc
- **17884** 000045dc
- **42460** 0000a5dc
- **34268** 000085dc
- **50652** 0000c5dc



Only care about least significant 12 bits

```
int x=getValue();
x=(x & 0x0fff);
printf("%d\n",x);
```



Only care about least significant 12 bits

```
int x=getValue();
x=x(&0x0fff);
printf("%d\n",x);
```



Must sign extend

```
int x=getValue();
x=(x&0x007ff) | (x&0x0800?0xfffff000:0);
printf("%d\n",x);
```



There is a better way.

Because you graduated from 213

```
int x=getValue();
x=(x&0x007ff) | (x&0x0800?0xfffff000:0);
printf("%d\n",x);
```



Lets be really thorough

```
int x=getValue();
x=(x&0x00fff)|(x&0x0800?0xf
printf(``%d\n",x);
```



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