RECITATION 7: SIGNALS AND I/O

15-213-M12

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Today:

- Shell Lab
- Midterm Exam
- Signals
- Process Control
- Unix I/O

Shell Lab

- Due Tuesday
- Testing
 - Test Suite
 - Race conditions

Midterm Exam

- Local students: In class next Thursday (7/12)
- Distance students: At a test center (or self-proctored)
 between 7/12 and 7/14
- Coverage:
 - Integers / Bitwise operations
 - Floating Point
 - Structures / Alignment
 - Assembly / Stack Discipline
 - Caching
 - Signals
 - Unix I/O

Signals

- Communication from the OS to the program
- Received on context switch
- Responses program dependent, changed with signal() or sigaction()
 - Ignore
 - Terminate
 - Suspend (or resume on SIGCONT)
 - Execute signal handler
- Default action varies by signal
- SIGSTOP, SIGKILL cannot be changed
 - Reliable way to get rid of an unresponsive process

Signals – Blocking

- Signals can be "blocked" by a process
- Blocked signals are not ignored, but are not received until the signal is unblocked
- Blocked (or not-yet-received) signals are not queued handler executes once no matter how many times the signal was sent

Important Signals

SIGINT

- By default, terminates the process
- Can be handled to clean up resources
- Ctrl+C sends SIGINT to terminal's active process group

SIGTSTP

- By default, suspends process
- Ctrl+Z sends SIGTSTP to terminal's active process group

SIGKILL

- Always terminates process
- OS sends SIGTERM, then later SIGKILL to all processes on shutdown
- "kill -9"

SIGCHLD

- Received by parent process when child terminates or suspends
- Ignored by default

Process Control

- Process IDs
- fork()
 - Creates a new process
 - Called once, returns twice
 - Returns 0 to child, child's PID to parent
- Process Groups
 - Inherited on fork()
 - Can be changed with setpgid()
- execve()
 - Current process starts executing a different executable file
 - Uses current address space, system resources
 - Called once, never returns

Process Control

- kill()
 - Sends a signal to a process
 - Must be owned by the same user
- wait(), waitpid()
 - Cleans up terminated process resources
 - By default, waits for process to terminate
 - Can set WNOHANG option
 - status argument to waitpid() gives calling process information about the terminated process
 - See "man waitpid" for details
- Terminated processes must be "reaped" by wait
 - init (PID 1) reaps processes if they are reparented to it (if their parent dies)

Unix I/O

File Descriptors

- Small integer identifying a file in use by the process
- Each process has a table mapping file descriptors to OS-level file structures (stores mode, offset, etc)
- File structures shared per process, descriptors not
- Returned by open(), dup(), etc
- Duplication (dup, dup2) only duplicates file descriptor, not file structure
- Pass file descriptors to read(), write(), etc
- Close files when done leaking file descriptors over the course of a long running program (like a shell!) is bad
- Automatically cleaned up on exit from program