

Recitation 10: More Malloc Lab

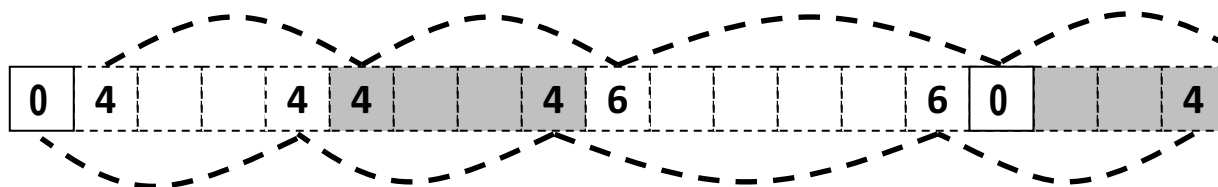
Instructor: TA(s)

Understanding Your Code

- Sketch out the heap
- Add Instrumentation
- Use tools

Sketch out the Heap

- Start with a heap, in this case implicit list



- Now try something, in this case, `extend_heap`

```

block_t *block = payload_to_header(bp);
write_header(block, size, false);
write_footer(block, size, false);
// Create new epilogue header
block_t *block_next = find_next(block);
write_header(block_next, 0, true);

```

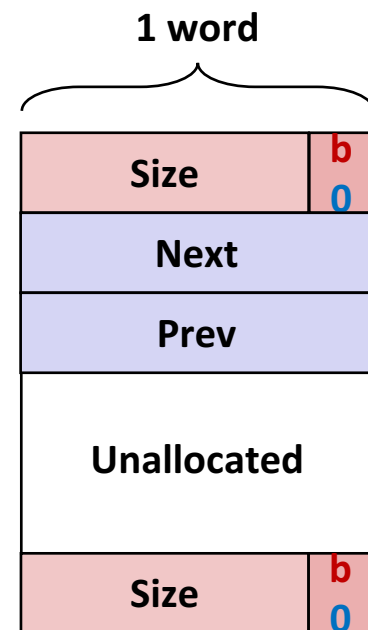
Sketch out the Heap

■ Here is a free block based on lectures 19 and 20

- Explicit pointers (will be well-defined see writeup and Piazza)
 - **This applies to ALL new fields you want inside your struct**
- Optional boundary tags

■ If you make changes to your design beyond this

- Draw it out.
- If you have bugs, pictures can help the staff help you
- Put a picture of your data structure into your file header (optional, but we will be impressed)



**Free
Block**

Common Problems

■ Throughput is very low

- Which operation is likely the most throughput intensive?
- Hint: It uses loops!
- Solution: ??

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- Hint: It extends the amount of memory that you have!
- Solution: ??

Common Problems

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Add Instrumentation

- **Remember that measurements inform insights.**
 - Add temporary code to understand aspects of malloc
 - Code can violate style rules or 128 byte limits, because it is temporary

- **Particularly important to develop insights into performance before making changes**
 - What is expensive throughput-wise?
 - How much might a change benefit utilization?

Add Instrumentation example

- Searching in `find_fit` is often the slowest step
- How efficient is your code? How might you know?
 - Compute the ratio of blocks viewed to calls

```
static block_t *find_fit(size_t asize)
{
    block_t *block; call_count++;
    for (block = heap_listp; get_size(block) > 0;
         block = find_next(block))
    {
        block_count++;
        if (!(get_alloc(block)) && (asize <= get_size(block)))
        {
            return block;
        }
    }
    return NULL; // no fit found
}
```

Add Instrumentation cont.

■ What size of requests?

- How many 8 bytes or less?
- How many 16 bytes or less?
- What other sizes?

■ What else could you measure? Why?

■ Remember that although the system's performance varies

- The mdriver's traces are deterministic
- Measured results should not change between runs

Use tools

■ Use `mm_checkheap()`

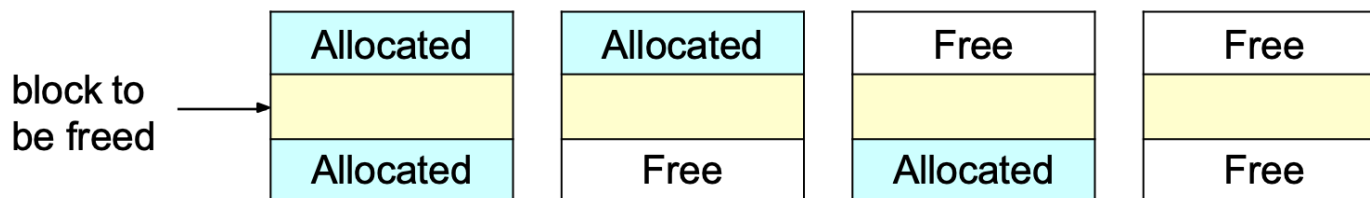
- Write it if you haven't done so already
- Add new invariants when you add new features
- Know how to use the heap checker.
 - Why do you need a heap checker? 2 reasons.

■ Use `gdb`

- You can call `print` or `mm_checkheap` whenever you want in `gdb`. No need to add a while lot of `printf`'s.
- Offers useful information whenever you crash, like `backtrace`.
- Write helper functions to print out free lists that are **ONLY** called from `GDB`

Write your own traces!

- Write short traces that test simple sequences of malloc and free
- Write a trace that simply tests all 4 coalescing cases



- Read the README file in the traces directory to see how trace files need to be written

mdriver-emulate

- **Testing for 64-bit address space**
- **Use correctly sized masks, constants, and other variables**
- **Be careful about subtraction between size types (may result in underflow/overflow)**
- **Reinitialize your pointers in mm_init**

Garbled Bytes

■ Malloc library returns a block

- mdriver writes bytes into payload (using memcpy)
- mdriver will check that those bytes are still present
- If malloc library has overwritten any bytes, then report garbled bytes
 - Also checks for other kinds of bugs

■ Now what?

■ The mm_checkheap call is catching it right?

■ If not, we want to find the garbled address and watch it

Garbled Bytes and gdb

- Get out a laptop
- Login to shark machine
- `wget http://www.cs.cmu.edu/~213/activities/rec11b.tar`
- `tar xf rec11b.tar`

- **mm.c is a fake explicit list implementation.**
 - Source code is based on mm.c starter code
 - A few lines of code are added that vaguely resembles what an explicit list implementation could have.

GDB Exercise

■ `gdb --args ./mdriver -c ./traces/syn-array-short.rep -D`

```
(gdb) r
```

```
// Sample output follows
```

```
Throughput targets: min=6528, max=11750, benchmark=13056
```

```
Malloc size 9904 on address 0x800000010.
```

```
Malloc size 50084 on address 0x8000026d0.
```

```
ERROR [trace ././traces/syn-array-short.rep, line 7]:  
block 0 has 8 garbled bytes, starting at byte 0
```

```
...
```

```
ERROR [trace ././traces/syn-array-short.rep, line 7]:  
block 0 has 8 garbled bytes, starting at byte 0
```

```
...
```

```
Terminated with 14 errors
```

```
[Inferior 1 (process 8456) exited normally]
```

```
(gdb)
```


GDB Exercise cont.

■ What is the first address that was garbled?

- Use gdb watch to find out when / what garbled it.

```
(gdb) watch * 0x800000010
```

```
(gdb) run
```

```
// Keep continuing through the breaks:
```

```
// mm_init()
```

```
// 4 x memcpy
```

```
Hardware watchpoint 1: *0x800000010
```

We just broke in
after overwriting



```
Old value = -7350814
```

```
New value = 9928
```

```
mm_malloc (size=50084) at mm.c:214
```

Second Exercise

Well fine, the bug from the first exercise was very artificial.
No one just sets bytes to 0 for no reason.

Try this more plausible exercise:

```
$ gdb --args ./mdriver-2 -c traces/syn-array-short.rep
```

What error was printed to the console?

The function that prints the error is named `malloc_error`. Add a breakpoint for it if you want.

Second Exercise

The library must've written the header and footer for the out-of-bounds payload at some point. Add a watchpoint for either address, or both.

```
(gdb) watch *0x8000036c8
```

```
(gdb) run
```

...So, the writes occurred in `place`. Is the `place` function wrong, or was it just given a bad argument?

Hint: the bug is found in at basically the same place as last recitation's bug.

It's caused by a careless typo, like nearly all others bugs.

Tips for using our tools

- Run `mdriver` with the `-D` option to detect garbled bytes as early as possible. Run it with `-V` to find out which trace caused the error.
- Note that sometimes, you get the error within the first few allocations. If so, you could set a breakpoint for `mm_malloc / mm_free` and step through every line.
- Print out local variables and convince yourself that they have the right values.
- For `mdriver-emulate`, you can still read memory from the simulated 64-bit address space using `mem_read(address, 8)` instead of `x / gx`.

MallocLab

- **Due next Tuesday**
- **7% of final grade (+ 4% for checkpoint)**
 - Style matters! Don't let all of your hard work get wasted.
 - There are many different implementations and TAs will need to know the details behind your implementation.
- **Read the writeup. It even has a list of tips on how to improve memory utilization.**
- **Read the malloc roadmap posted on Piazza**
- **Rubber duck method**
 - If you explain to a rubber duck what your function does step-by-step, while occasionally stopping to explain why you need each of those steps, you'd may very well find the bug in the middle of your explanation.
 - Remember the “debug thought process” slide from Recitation 9?

Style

■ Well organized code is easier to debug and easier to grade!

- Modularity: Helper functions to respect the list interface.
- Documentation:
 - File Header: Describes all implementation details, including block structures.
- Code Structure:
 - Minimal-to-no pointer arithmetic.
 - Loops instead of conditionals, where appropriate.