

Recitation 14: Proxy Lab Part 2

Instructor: TA(s)

Outline

- **Proxylab**
- **Threading**
- **Threads and Synchronization**

ProxyLab

- **ProxyLab is due in 1 week.**
 - No grace days
 - Make sure to submit well in advance of the deadline in case there are errors in your submission.
 - Build errors are a common source of failure

- **A proxy is a server process**
 - It is expected to be long-lived
 - To not leak resources
 - To be robust against user input

Proxies and Threads

■ Network connections can be handled concurrently

- Three approaches were discussed in lecture for doing so
- Your proxy should (eventually) use threads
- Threaded echo server is a good example of how to do this

■ Multi-threaded cache design

- Need to have multiple readers or one writer
- Be careful how you use mutexes – you do not want to serialize your readers
- Be careful how you maintain your object age

■ Tools

- Use Firefox's Network Monitor (Developer > Network) to see if all requests have been fulfilled

Join / Detach

- Does the following code terminate? Why or why not?

```
int main(int argc, char** argv)
{
...
    pthread_create(&tid, NULL, work, NULL);
    if (pthread_join(tid, NULL) != 0) printf("Done.\n");
...
void* work(void* a)
{
    pthread_detach(pthread_self());
    while(1);
}
```

Join / Detach cont.

- Does the following code terminate now? Why or why not?

```
int main(int argc, char** argv)
{
...
    pthread_create(&tid, NULL, work, NULL); sleep(1);
    if (pthread_join(tid, NULL) != 0) printf("Done.\n");
...
void* work(void* a)
{
    pthread_detach(pthread_self());
    while(1);
}
```

When should threads detach?

- In general, pthreads will wait to be reaped via `pthread_join`.
- When should this behavior be overridden?
- When termination status does not matter.
 - `pthread_join` provides a return value
- When result of thread is not needed.
 - When other threads do not depend on this thread having completed

Threads

- What is the range of value(s) that main will print?
- A programmer proposes removing `j` from thread and just directly accessing `count`. Does the answer change?

```
volatile int count = 0;

void* thread(void* v)
{
    int j = count;
    j = j + 1;
    count = j;
}

int main(int argc, char** argv)
{
    pthread_t tid[2];
    for(int i = 0; i < 2; i++)
        pthread_create(&tid[i], NULL,
                      thread, NULL);
    for (int i = 0; i < 2; i++)
        pthread_join(tid[i]);
    printf("%d\n", count);
    return 0;
}
```


Synchronization

- **Is not cheap**
 - 100s of cycles just to acquire without waiting
- **Is also not that expensive**
 - Recall your malloc target of 15000kops => ~100 cycles
- **May be necessary**
 - Correctness is always more important than performance

Which synchronization should I use?

- **Counting a shared resource, such as shared buffers**
 - Semaphore

- **Exclusive access to one or more variables**
 - Mutex

- **Most operations are reading, rarely writing / modifying**
 - RWLock

Threads Revisited

- Which lock type should be used?
- Where should it be acquired / released?

```
volatile int count = 0;

void* thread(void* v)
{
    int j = count;
    j = j + 1;
    count = j;
}

int main(int argc, char** argv)
{
    pthread_t tid[2];
    for(int i = 0; i < 2; i++)
        pthread_create(&tid[i], NULL,
                      thread, NULL);
    for (int i = 0; i < 2; i++)
        pthread_join(tid[i]);
    printf("%d\n", count);
    return 0;
}
```

Associating locks with data

- **Given the following key-value store**
 - Key and value have separate RWLocks: klock and vlock
 - When an entry is replaced, both locks are acquired.
- **Describe why the printf may not be accurate.**

```
typedef struct _data_t {
    int key;
    size_t value;
} data_t;

#define SIZE 10
data_t space[SIZE];
int search(int k)
{
    for(int j = 0; j < SIZE; j++)
        if (space[j].key == k) return j;
    return -1;
}
```

```
...
pthread_rwlock_rdlock(klock);
match = search(k);
pthread_rwlock_unlock(klock);

if (match != -1)
{
    pthread_rwlock_rdlock(vlock);
    printf("%zd\n", space[match]);
    pthread_rwlock_unlock(vlock);
}
```

Locks gone wrong

1. RWLocks are particularly susceptible to which issue:

- a. Starvation b. Livelock c. Deadlock

1. If some code acquires rwlocks as readers: LockA then LockB, while other readers go LockB then LockA. What, if any, order can a writer acquire both LockA and LockB?

No order is possible without a potential deadlock.

3. Design an approach to acquiring two semaphores that avoids deadlock and livelock, while allowing progress to other threads needing only one semaphore.

Client-to-Client Communication

- **Clients don't have to fetch content from servers**
 - Clients can communicate with each other
 - In a chat system, a server acts as a facilitator between clients
 - Clients could also send messages directly to each other, but this is more complicated (peer-to-peer networking)
- **Running the chat server**
 - `./chatserver <port>`
- **Running the client**
 - `telnet <hostname> <port>`
- **What race conditions could arise from having communication between multiple clients?**

Proxylab Reminders

- **Plan out your implementation**
 - “Weeks of programming can save you hours of planning”
 - – Anonymous
 - Arbitrarily using mutexes will not fix race conditions

- **Read the writeup**

- **Submit your code (days) early**
 - Test that the submission will build and run on Autolab

- **Final exam is only a few weeks away!**

Appendix

- **Calling `exit()` will terminate all threads**
- **Calling `pthread_join` on a detached thread is technically undefined behavior. Was defined as returning an error.**