Andrew login ID:_________________________
Full Name:__________________________

CS 15-213, Fall 2001

Final Exam

December 13, 2001

Instructions:

- Make sure that your exam is not missing any sheets, then write your full name and Andrew login ID on the front.

- Write your answers in the space provided below the problem. If you make a mess, clearly indicate your final answer.

- The exam has a maximum score of 120 points.

- This exam is OPEN BOOK. You may use any books or notes you like. You may use a calculator, but no laptops or other wireless devices. Good luck!

<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>1 (20):</td>
<td></td>
</tr>
<tr>
<td>2 (10):</td>
<td></td>
</tr>
<tr>
<td>3 (10):</td>
<td></td>
</tr>
<tr>
<td>4 (8):</td>
<td></td>
</tr>
<tr>
<td>5 (12):</td>
<td></td>
</tr>
<tr>
<td>6 (6):</td>
<td></td>
</tr>
<tr>
<td>7 (14):</td>
<td></td>
</tr>
<tr>
<td>8 (10):</td>
<td></td>
</tr>
<tr>
<td>9 (16):</td>
<td></td>
</tr>
<tr>
<td>10 (14):</td>
<td></td>
</tr>
</tbody>
</table>

TOTAL (120):
Problem 1. (20 points):
We are running programs on a machine with the following characteristics:

- Values of type `int` are 32 bits. They are represented in two’s complement, and they are right shifted arithmetically. Values of type `unsigned` are 32 bits.

- Values of type `float` are represented using the 32-bit IEEE floating point format, while values of type `double` use the 64-bit IEEE floating point format.

We generate arbitrary values `x`, `y`, and `z`, and convert them to other forms as follows:

```c
/* Create some arbitrary values */
int x = random();
int y = random();
int z = random();
/* Convert to other forms */
unsigned ux = (unsigned) x;
unsigned uy = (unsigned) y;
double dx = (double) x;
double dy = (double) y;
double dz = (double) z;
```

For each of the following C expressions, you are to indicate whether or not the expression always yields 1. If so, circle “Y”. If not, circle “N”. You will be graded on each problem as follows:

- If you circle no value, you get 0 points.
- If you circle the right value, you get 2 points.
- If you circle the wrong value, you get -1 points (so don’t just guess wildly).

<table>
<thead>
<tr>
<th>Expression</th>
<th>Always True?</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>(x&lt;y) == (-x&gt;-y)</code></td>
<td>Y</td>
</tr>
<tr>
<td><code>((x+y)&lt;&lt;4) + y-x == 17*y+15*x</code></td>
<td>Y</td>
</tr>
<tr>
<td><code>~x+~y+1 == ~(x+y)</code></td>
<td>Y</td>
</tr>
<tr>
<td><code>ux-uy == -(y-x)</code></td>
<td>Y</td>
</tr>
<tr>
<td>`(x &gt;= 0)</td>
<td>(x &lt; ux)`</td>
</tr>
<tr>
<td><code>((x &gt;&gt; 1) &lt;&lt; 1) &lt;= x</code></td>
<td>Y</td>
</tr>
<tr>
<td><code>(double)(float) x == (double) x</code></td>
<td>Y</td>
</tr>
<tr>
<td><code>dx + dy == (double) (y+x)</code></td>
<td>Y</td>
</tr>
<tr>
<td><code>dx + dy + dz == dz + dy + dx</code></td>
<td>Y</td>
</tr>
<tr>
<td><code>dx * dy * dz == dz * dy * dx</code></td>
<td>Y</td>
</tr>
</tbody>
</table>
Problem 2. (10 points):
A C function `looper` and the assembly code it compiles to on an IA-32 machine running Linux/GAS is shown below:

```
looper:
    pushl %ebp
    movl %esp,%ebp
    pushl %esi
    pushl %ebx
    movl 8(%ebp),%ebx
    movl 12(%ebp),%esi
    xorl %edx,%edx
    xorl %ecx,%ecx
    cmpl %ebx,%edx
    jge .L25
.L27:
    movl (%esi,%ecx,4),%eax
    cmpl %edx,%eax
    jle .L28
    movl %eax,%edx
.L28:
    incl %edx
    incl %ecx
    cmpl %ebx,%ecx
    jl .L25
.L25:
    movl %edx,%eax
    poapl %ebx
    poapl %esi
    movl %ebp,%esp
    poapl %ebp
    ret
```

Based on the assembly code, fill in the blanks in the C source code.

Notes:

- You may only use the C variable names `n`, `a`, `i` and `x`, not register names.
- Use array notation in showing accesses or updates to elements of `a`.
Problem 3. (10 points):
Consider the following incomplete definition of a C struct along with the incomplete code for a function `func` given below.

```c
typedef struct node {
    ______________ x;
    ______________ y;
    struct node *next;
    struct node *prev;
} node_t;

node_t n;
void func() {
    node_t *m;
    m = ______________________;
    m->y /= 16;
    return;
}
```

When this C code was compiled on an IA-32 machine running Linux, the following assembly code was generated for function `func`.

```assembly
func:
    pushl %ebp
    movl n+12,%eax
    movl 16(%eax),%eax
    movl %esp,%ebp
    movl %ebp,%esp
    shrw $4,8(%eax)
    popl %ebp
    ret
```

Given these code fragments, fill in the blanks in the C code given above. Note that there is a unique answer.

The types must be chosen from the following table, assuming the sizes and alignment given.

<table>
<thead>
<tr>
<th>Type</th>
<th>Size (bytes)</th>
<th>Alignment (bytes)</th>
</tr>
</thead>
<tbody>
<tr>
<td>char</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>short</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>unsigned short</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>int</td>
<td>4</td>
<td>4</td>
</tr>
<tr>
<td>unsigned int</td>
<td>4</td>
<td>4</td>
</tr>
<tr>
<td>double</td>
<td>8</td>
<td>4</td>
</tr>
</tbody>
</table>
Problem 4. (8 points):
Consider the source code below, where M and N are constants declared with \#define.

```c
int array1[M][N];
int array2[N][M];

void copy(int i, int j)
{
    array1[i][j] = array2[j][i];
}
```

Suppose the above code generates the following assembly code:

```assembly
copy:
    pushl %ebp
    movl %esp,%ebp
    pushl %ebx
    movl 8(%ebp),%ecx
    movl 12(%ebp),%eax
    leal 0(,%eax,4),%ebx
    leal 0(,%ecx,8),%edx
    subl %ecx,%edx
    addl %ebx,%eax
    sall $2,%eax
    movl array2(%eax,%ecx,4),%eax
    movl %eax,array1(%ebx,%edx,4)
    popl %ebx
    movl %ebp,%esp
    popl %ebp
    ret
```

What are the values of M and N?

M =

N =
The following problem will test your understanding of the runtime stack. You are given the following declarations on an x86 architecture:

```c
struct file_spec {
    int fs_tag, parent_dir, size;
};

struct l_node {
    int tag;
    struct l_node *next;
};

struct l_node get_list_head()
{
    /* some irrelevant code */
}

void make_alias(..., ..., ..., ..., ...)
{
    /* some more irrelevant code */
}

void save_file(struct file_spec *file, int len, char *descriptor)
{
    int result = 0;
    struct l_node root = get_list_head();

    make_alias(..., ..., ..., ..., ...);

    /* yet more irrelevant code */
}
```

On the next page, you have the diagram of the stack immediately before the call assembly instruction for `make_alias()` in `save_file()` is executed. Argument `descriptor` given to `save_file()` is stored at `0xbffff5c0`. You can make the following assumptions:

- The function `make_alias()` takes exactly five arguments.
- The allocation order of local variables on the stack corresponds to their definition order in the source code.
- The compiler does not insert any additional unused space on the stack apart from unused space required for alignment restrictions of variables.
- No registers (apart from `%ebp`) are being saved on the stack.
Feel free to make comments or notes in the third column of the table - they will not be graded.

<table>
<thead>
<tr>
<th>Address</th>
<th>Numeric Value</th>
<th>Comments/Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0xbffff5c0</td>
<td>0xbffff7a0</td>
<td></td>
</tr>
<tr>
<td>0xbffff5bc</td>
<td>0x000feedb</td>
<td></td>
</tr>
<tr>
<td>0xbffff5b8</td>
<td>0xbffff780</td>
<td></td>
</tr>
<tr>
<td>0xbffff5b4</td>
<td>0x080459ec</td>
<td></td>
</tr>
<tr>
<td>0xbffff5b0</td>
<td>0xbffff630</td>
<td></td>
</tr>
<tr>
<td>0xbffff5ac</td>
<td>0x00000000</td>
<td></td>
</tr>
<tr>
<td>0xbffff5a8</td>
<td>0x83045c30</td>
<td></td>
</tr>
<tr>
<td>0xbffff5a4</td>
<td>0x00000045</td>
<td></td>
</tr>
<tr>
<td>0xbffff5a0</td>
<td>0xbffff5ac</td>
<td></td>
</tr>
<tr>
<td>0xbffff59c</td>
<td>0x83045c30</td>
<td></td>
</tr>
<tr>
<td>0xbffff598</td>
<td>0x00000045</td>
<td></td>
</tr>
<tr>
<td>0xbffff594</td>
<td>0xbffff780</td>
<td></td>
</tr>
<tr>
<td>0xbffff590</td>
<td>0x000feedb</td>
<td></td>
</tr>
<tr>
<td>0xbffff58c</td>
<td>0xbffff7a0</td>
<td></td>
</tr>
</tbody>
</table>
Problem 5. (12 points):

A. Give the current value of the frame pointer (machine register %ebp).

B. The declaration of `make_alias()` is missing the types of its parameters. Give the types in the order they would appear in the source code. The names of the parameters do not matter.

C. List the arguments passed to `make_alias()` in `save_file()`, in the order they would appear in the source code.

```
make_alias( , , , , , , , );
```
Problem 6. (6 points):
The following table gives the parameters for a number of different caches, where $m$ is the number of physical address bits, $C$ is the cache size (number of data bytes), $B$ is the block size in bytes, and $E$ is the number of lines per set. For each cache, determine the number of cache sets ($S$), tag bits ($t$), set index bits ($s$), and block offset bits ($b$).

<table>
<thead>
<tr>
<th>Cache</th>
<th>$m$</th>
<th>$C$</th>
<th>$B$</th>
<th>$E$</th>
<th>$S$</th>
<th>$t$</th>
<th>$s$</th>
<th>$b$</th>
</tr>
</thead>
<tbody>
<tr>
<td>1.</td>
<td>32</td>
<td>1024</td>
<td>4</td>
<td>4</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>2.</td>
<td>32</td>
<td>1024</td>
<td>4</td>
<td>256</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>3.</td>
<td>32</td>
<td>1024</td>
<td>8</td>
<td>1</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>4.</td>
<td>32</td>
<td>1024</td>
<td>8</td>
<td>128</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>5.</td>
<td>32</td>
<td>1024</td>
<td>32</td>
<td>1</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>6.</td>
<td>32</td>
<td>1024</td>
<td>32</td>
<td>4</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Problem 7. (14 points):
Consider a direct mapped cache of size 64K with block size of 16 bytes. Furthermore, the cache is write-back and write-allocate. You will calculate the miss rate for the following code using this cache. Remember that sizeof(int) == 4. Assume that the cache starts empty and that local variables and computations take place completely within the registers and do not spill onto the stack.

A. Now consider the following code to copy one matrix to another. Assume that the src matrix starts at address 0 and that the dest matrix follows immediately follows it.

```c
void copy_matrix(int dest[ROWS][COLS], int src[ROWS][COLS])
{
    int i, j;

    for (i=0; i<ROWS; i++) {
        for (j=0; j<COLS; j++) {
            dest[i][j] = src[i][j];
        }
    }
}
```

1. What is the cache miss rate if ROWS = 128 and COLS = 128?
   Miss rate = __________%  

2. What is the cache miss rate if ROWS = 128 and COLS = 192?
   Miss rate = __________%  

3. What is the cache miss rate if ROWS = 128 and COLS = 256?
   Miss rate = __________%  

B. Now consider the following two implementations of a horizontal flip and copy of the matrix. Again assume that the src matrix starts at address 0 and that the dest matrix follows immediately follows it.

```c
void copy_n_flip_matrix1(int dest[ROWS][COLS], int src[ROWS][COLS])
{
    int i, j;

    for (i=0; i<ROWS; i++) {
        for (j=0; j<COLS; j++) {
            dest[i][COLS - 1 - j] = src[i][j];
        }
    }
}
```

1. What is the cache miss rate if ROWS = 128 and COLS = 128?
   Miss rate = __________%

2. What is the cache miss rate if ROWS = 128 and COLS = 192?
   Miss rate = __________%

```c
void copy_n_flip_matrix2(int dest[ROWS][COLS], int src[ROWS][COLS])
{
    int i, j;

    for (j=0; j<COLS; j++) {
        for (i=0; i<ROWS; i++) {
            dest[i][COLS - 1 - j] = src[i][j];
        }
    }
}
```

1. What is the cache miss rate if ROWS = 128 and COLS = 128?
   Miss rate = __________%

2. What is the cache miss rate if ROWS = 192 and COLS = 128?
   Miss rate = __________%
Problem 8. (10 points):

Consider the following function for computing the product of an array of \( n \) integers. We have unrolled the loop by a factor of 3.

```c
int aprod(int a[], int n)
{
    int i, x, y, z;
    int r = 1;
    for (i = 0; i < n-2; i+= 3) {
        x = a[i]; y = a[i+1]; z = a[i+2];
        r = r * x * y * z; // Product computation
    }
    for (; i < n; i++)
        r *= a[i];
    return r;
}
```

For the line labeled Product computation, we can use parentheses to create 5 different associations of the computation, as follows:

\[
\begin{align*}
  r &= ((r \times x) \times y) \times z; & \text{A1} \\
  r &= (r \times (x \times y)) \times z; & \text{A2} \\
  r &= r \times ((x \times y) \times z); & \text{A3} \\
  r &= r \times (x \times (y \times z)); & \text{A4} \\
  r &= (r \times x) \times (y \times z); & \text{A5}
\end{align*}
\]

We express the performance of the function in terms of the number of cycles per element (CPE). As described in the book, this measure assumes the run time, measured in clock cycles, for an array of length \( n \) is a function of the form \( Cn + K \), where \( C \) is the CPE.

We measured the 5 versions of the function on an Intel Pentium III. Recall that the integer multiplication operation on this machine has a latency of 4 cycles and an issue time of 1 cycle.
The following table shows some values of the CPE, and other values missing. The measured CPE values are those that were actually observed. “Theoretical CPE” means that performance that would be achieved if the only limiting factor were the latency and issue time of the integer multiplier.

<table>
<thead>
<tr>
<th>Version</th>
<th>Measured CPE</th>
<th>Theoretical CPE</th>
</tr>
</thead>
<tbody>
<tr>
<td>A1</td>
<td>4.00</td>
<td></td>
</tr>
<tr>
<td>A2</td>
<td>2.67</td>
<td></td>
</tr>
<tr>
<td>A3</td>
<td></td>
<td>$4/3 = 1.33$</td>
</tr>
<tr>
<td>A4</td>
<td>1.67</td>
<td></td>
</tr>
<tr>
<td>A5</td>
<td></td>
<td>$8/3 = 2.67$</td>
</tr>
</tbody>
</table>

Fill in the missing entries. For the missing values of the measured CPE, you can use the values from other versions that would have the same computational behavior. For the values of the theoretical CPE, you can determine the number of cycles that would be required for an iteration considering only the latency and issue time of the multiplier, and then divide by 3.
Problem 9. (16 points):
This problem tests your understanding of exceptional control flow in C programs. Assume we are running code on a Unix machine. The following problems all concern the value of the variable `counter`.

Part I (6 points)

```c
int counter = 0;

int main()
{
    int i;

    for (i = 0; i < 2; i ++){
        fork();
        counter ++;
        printf("counter = %d\n", counter);
    }

    printf("counter = %d\n", counter);
    return 0;
}
```

A. How many times would the value of `counter` be printed: ____________

B. What is the value of `counter` printed in the first line? ____________

C. What is the value of `counter` printed in the last line? ____________
Part II (6 points)

```c
pid_t pid;
int counter = 0;

void handler1(int sig)
{
    counter ++;
    printf("counter = %d\n", counter);
    fflush(stdout); /* Flushes the printed string to stdout */
    kill(pid, SIGUSR1);
}

void handler2(int sig)
{
    counter += 3;
    printf("counter = %d\n", counter);
    exit(0);
}

main()
{
    signal(SIGUSR1, handler1);
    if ((pid = fork()) == 0) {
        signal(SIGUSR1, handler2);
        kill(getppid(), SIGUSR1);
        while(1) {};
    } else {
        pid_t p; int status;
        if ((p = wait(&status)) > 0) {
            counter += 2;
            printf("counter = %d\n", counter);
        }
    }
}
```

What is the output of this program?
Part III (4 points)

```c
int counter = 0;

void handler(int sig)
{
    counter ++;
}

int main()
{
    int i;
    signal(SIGCHLD, handler);
    for (i = 0; i < 5; i ++){
        if (fork() == 0){
            exit(0);
        }
    }
    /* wait for all children to die */
    while (wait(NULL) != -1);
    printf("counter = %d\n", counter);
    return 0;
}
```

A. Does the program output the same value of `counter` every time we run it?    Yes   No

B. If the answer to A is Yes, indicate the value of the `counter` variable. Otherwise, list all possible values of the `counter` variable.

Answer: `counter` = ______________________
Problem 10. (14 points):
Consider an allocator with the following specification:

- Uses a single explicit free list.
- All memory blocks have a size that is a multiple of 8 bytes and is at least 16 bytes.
- All headers, footers, and pointers are 4 bytes in size
- Headers consist of a size in the upper 29 bits, a bit indicating if the block is allocated in the lowest bit, and a bit indicating if the previous block is allocated in the second lowest bit.
- Allocated blocks consist of a header and a payload (no footer)
- Free blocks consist of a header, two pointers for the next and previous free blocks in the free list, and a footer at the end of the block.
- All freed blocks are immediately coalesced.
- The heap starts with 0 bytes, never shrinks, and only grows large enough to satisfy memory requests.
- The heap contains only allocated blocks and free blocks. There is no space used for other data or special blocks to mark the beginning and end of the heap.
- When a block is split, the lower part of the block becomes the allocated part and the upper part becomes the new free block.
- Any newly created free block (whether it come from a call to free, the upper part of a split block, or the coalescing of several free blocks) is inserted at the beginning of the free list.
- All searches for free blocks start at the head of the list and walk through the list in order.
- If a request can be fulfilled by using a free block, that free block is used. Otherwise the heap is extended only enough to fulfill the request. If there is a free block at the end of the heap, this can be used along with the new heap space to fulfill the request.
A. Simulating Malloc (10 points)

Below you are given a series of memory requests. You are asked to show what the heap looks like after each request is completed using a first fit and a best fit placement policy. The heap is represented as a series of boxes, where each box is a single block on the heap, and the bottom of the heap is the left most box. In each block, you should write the total size (including headers and footers) of the block in bytes and either ‘f’ or ‘a’ to mark it as free or allocated, respectively. For example, the following heap contains an allocated block of size 16, followed by a free block of size 32.

Assume that the heap is empty before each of the sequences is run. You do not necessarily have to use all the boxes provided for the heap. Some of the boxes are already filled in to help you.

<table>
<thead>
<tr>
<th>First Fit</th>
<th>Best Fit</th>
</tr>
</thead>
<tbody>
<tr>
<td>ptr1 = malloc(32);</td>
<td>24a</td>
</tr>
<tr>
<td>ptr2 = malloc(16);</td>
<td></td>
</tr>
<tr>
<td>ptr3 = malloc(16);</td>
<td></td>
</tr>
<tr>
<td>ptr4 = malloc(40);</td>
<td></td>
</tr>
<tr>
<td>free(ptr3);</td>
<td></td>
</tr>
<tr>
<td>free(ptr1);</td>
<td></td>
</tr>
<tr>
<td>ptr5 = malloc(16);</td>
<td>24a</td>
</tr>
<tr>
<td>free(ptr4);</td>
<td></td>
</tr>
<tr>
<td>ptr6 = malloc(48);</td>
<td>24a</td>
</tr>
<tr>
<td>free(ptr2);</td>
<td></td>
</tr>
</tbody>
</table>
B. Code for Malloc (4 points)

For this part, you are asked to complete some small functions which are used to setup blocks. Each function will be missing a line of code and you are given three choices for this line of code. Circle the choice that completes the function correctly.

**Function 1**

```c
/* Input: void *block: a pointer to a block
   * unsigned long size: the size of the block,
   * char alloc: the lower order bit indicates if this block is
     allocated
   * char palloc: the lower order bit indicates if the previous
     block is allocated
   * Actions: This function will construct a header from the last 3
     parameters and place it in the header of the block pointed
     to by the first parameter.
   */
void make_header(void *block, unsigned long size, char alloc, char palloc)
{
    long header;
    __________;
    *(long *)block = header;
}
```

A. header = (size >> 3) | ((alloc & 0x1) << 31) | ((palloc & 0x1) << 30);
B. header = (size & ~0x7) | (alloc & 0x1) | ((palloc & 0x1) << 1);
C. header = size | alloc | palloc;
Function 2

/* Input: void *block: a pointer to a block
 * char palloc: the low order bit indicates if the previous
 * block is allocated
 * Actions: Sets just the bit in the header indicating if the previous
 * block is allocated. Does nothing to the rest of the header.
 */

void set_palloc_in_header(void *block, char palloc)
{
    long curr = *(long *)block;
    ________;
}

A. *(long *)block = (curr & ~0x2) | ((palloc & 0x1) << 1);
B. *(long *)block = (curr & ~(0x1 << 30)) | ((palloc & 0x1) << 30);
C. *(long *)block = (curr | (0x1 << 30)) & ~((palloc & 0x1) << 30);