ANITA'S SUPER AWESOME RECITATION SLIDES

15/18-213: Introduction to Computer Systems Assembly and GDB, 16 Sept 2013 Anita Zhang

MANAGEMENT AND STUFF

- o Bomb Lab due Tues, 24 Sept 2013, 11:59 PM
 - This is my favorite lab!
- o Buf Lab out Tues, 24 Sept 2013, 11:59 PM
 - One week long lab

WHAT'S ON THE MENU TODAY?

- Help (again)
- Books (again)
- Motivation
- Registers & Assembly
- Bomb Lab Overview
- GDB
- More Bomb Lab
- Must Know Unix Commands
- Walkthrough

HELPING US, HELPING YOU?

- o Email us: 15-213-staff@cs.cmu.edu
 - Please attach C files if you have a specific question
 - TAs + Professors → More coverage, fast replies
- All projects on Autolab: <u>autolab.cs.cmu.edu</u>
- o Office Hours: Sun-Thurs, 5:30PM − 8:30 PM
 - Wean 5207
- Peer Tutoring: Tues 8:30 11 PM
 - Mudge Reading Room

WHAT HAVE YOU READ?

- Randal E. Bryant and David R. O'Hallaron, Computer Systems: A Programmer's Perspective, Second Edition, Prentice Hall, 2011
- Brian W. Kernighan and Dennis M. Ritchie, The C Programming Language, Second Edition, Prentice Hall, 1988
- Koenig, Andrew. *C Traps and Pitfalls*. Reading, MA: Addison-Wesley, 1988
- Kernighan, Brian W., and Rob Pike. *The Practice of Programming*. Reading, MA: Addison-Wesley, 1999

WHY ARE WE DOING THIS AGAIN?



Insight for the Inquisitive

- Why are we not learning about the stack yet?
 - Because x86_64
- "Technology note"
 - x86(_64) only

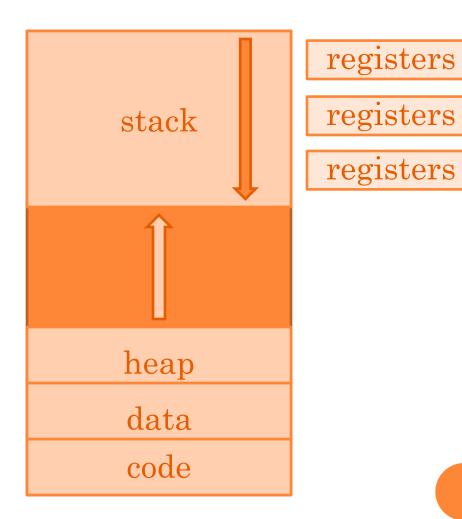
WHAT ARE REGISTERS?

Register

- Some place in hardware that stores bits
- It is NOT on the stack or in main memory

Important

 When moving data between registers and memory, only the DATA moves, not the register



REGISTERS AND ALL THEM BITS

%rax – 64 bits

%eax -32 bits

- Quad = 64 bits
- Doubleword = 32 bits
- \circ Word = 16 bits
- Byte = 8 bits

%ax – 16 bits

%ah 8 bits

%al 8 bits

These are all parts of the same register

WHAT WE'RE WORKING WITH

- x86_64 conventions on the next slide
- Specials
 - %eip instruction pointer
 - Points to the **NEXT** instruction to execute
 - %esp stack pointer
 - Points to top of the stack
 - %eax holds the return value
 - Also general purpose
- Conditional Flags
 - Sit in a special register of its own
 - Don't really need to worry about it

x86_64, LOTS of Registers!

64 bits wide	32 bits wide	16 bits wide	8 bits wide	8 bits wide	Use
%rax	%eax	%ax	%ah	%al	Return Value
%rbx	%ebx	%bx	%bh	%bl	Callee Save
%rcx	%ecx	%cx	%ch	%cl	4 th Argument
%rdx	%edx	%dx	%dh	%dl	3 rd Argument
%rsi	%esi	%si		%sil	2 nd Argument
%rdi	%edi	%di		%dil	1st Argument
%rbp	%ebp	%bp		%bpl	Callee Save
%rsp	%esp	%sp		%spl	Stack Pointer
%r8	%r8d	%r8w		%r8b	5 th Argument
%r9	%r9d	%r9w		%r9b	6 th Argument
%r10	%r10d	%r10w		%r10b	Caller Save
%r11	%r11d	%r11w		%r11b	Caller Save
%r12	%r12d	%r12w		%r12b	Callee Save
%r13	%r13d	%r13w		%r12b	Callee Save
%r14	%r14d	%rw		%14b	Callee Save
%r15	%r15d	%r15w		%15b	Callee Save

SOME MORE DEFINITIONS

- Memory Addressing
 - How assemblers denote memory locations
 - Direct
 - Indirect
 - Relative
 - Absolute
 - o ...
 - Many different syntactical ways to represent the same address

REASONS WHY INTEL IS RIDICULOUS AND AWESOME

- Operations can take several forms:
 - Register-to-Register
 - Register-to-Memory / Memory-to-Register
 - Immediate-to-Register / Immediate-to-Memory
 - One address operations (push, pop)
 - Did I miss any?
- Fun fact: Why not memory to memory?

Representing Addresses

- x86(_64) Common Addressing
 - Offset(Base, Index, Scale)
 - $D(Rb, Ri, S) \rightarrow Mem[Rb + Ri*S + D]$
 - D can be any signed integer
 - Scale is 1, 2, 4, 8 (assume 1 if omitted)
 - Assume 0 for base if omitted

Representing Addresses

- Using parenthesis
 - Most of the time parenthesis means dereference
 - This is still only $x86(_64)$
- Examples of parenthesis usage:
 - (%eax)
 - Contents of memory at address stored, %eax
 - (%ebx, %ecx)
 - Contents of memory stored at address, %ebx + %ecx
 - (%ebx, %ecx, 8)
 - Contents of memory stored at address, %ebx + 8*%ecx
 - 4(%ebx, %ecx, 8)
 - Contents of memory stored at address, %ebx + 8*%ecx + 4

Representing Addresses

- Using parenthesis
 - Sometimes parenthesis are used just for addressing
 - This is still only $x86(_64)$
- Example
 - leal (%ebx, %ecx, 8), destination
 - Take only the values → %ebx + 8*%ecx
 - Does not dereference, uses the calculated value directly
 - jmpq *0x402660(,%rax,8)
 - The * does the dereference
- Examples of not using parenthesis
 - %eax
 - Use the value in %eax!
 - \$0x213
 - A constant value

REVIEW OF CONDITIONALS/ FLAGS

- Most operations will set conditional flags
 - Bit operations
 - Arithmetic
 - Comparisons...
- Core idea: For conditionals, look one instruction before it to see whether it is true or false
 - Will be explained

FLAGS WE (MIGHT) CARE ABOUT

- o Carry (CF)
 - Arithmetic carry/ borrow
- Parity (PF)
 - Odd or even number of bits set
- Zero (ZF)
 - Result was zero
- Sign (SF)
 - Most significant bit was set
- Overflow (OF)
 - Result does not fit into the location

PREP FOR ALL THE CHEAT SHEETS

- Warning: The following slides contain lots of assembly instructions.
 - All from CS:APP (our textbook BTW)
 - We're not going over every single one...
 - Use it as a reference for Bomb Lab
- Quick note on Intel vs. AT&T
 - This is AT&T syntax (also, Bomb Lab syntax)
 - Looks like: "src, dest"
 - Intel tends to follow "dest, src"
 - Check out their ISA sometime

ALL THE CHEAT SHEETS (MOVEMENT)

Instruction		Effect
movb	S, D	Move byte
movw	S, D	Move word
movl	S, D	Move doubleword
movsbw	S, D	Move byte to word (sign extended)
movsbl	S, D	Move byte to doubleword (sign extended)
movswl	S, D	Move word to doubleword (sign extended)
movzbw	S, D	Move byte to word (zero extended)
movzbl	S, D	Move byte to doubleword (zero extended)
movzwl	S, D	Move word to doubleword (zero extended)
pushl	S	Push double word (Mem[%esp] ← S; %esp = %esp 4)
popl	D	Pop double word (D \leftarrow Mem[%esp]; %esp = %esp + 4)

ALL THE CHEAT SHEETS (BIT OPS)

Instruct	ion	Effect
LEAL	S, D	D ← &S (Load address of source into destination)
INC	D	D ← D + 1
DEC	D	D ← D − 1
NEG	D	$D \leftarrow -D$
NOT	D	D ← ~D
ADD	S, D	$D \leftarrow D + S$
SUB	S, D	$D \leftarrow D - S$
IMUL	S, D	$D \leftarrow D * S$
XOR	S, D	$D \leftarrow D \land S$
OR	S, D	$D \leftarrow D \mid S$
AND	S, D	$D \leftarrow D \& S$
SAL	k, D	D ← D << k
SHL	k, D	D ← D << k
SAR	k, D	D ← D >> k (arithmetic shift)
SHR	k, D	$D \leftarrow D >> k \text{ (logical shift)}$

ALL THE CHEAT SHEETS (SPECIALS)

Instructi	on	Effect
imull	S	R[%edx]:R[%eax] ← S * R[%eax] Signed multiply of %eax by S Result stored in %edx:%eax
mull	S	R[%edx]:R[%eax] ← S * R[%eax] Unsigned multiply of %eax by S Result stored in %edx:%eax
cltd		R[%edx]:R[%eax] ← SignExtend(R[%eax]) Sign extend %eax into %edx
idivl	S	R[%edx] ← R[%edx]:R[%eax] mod S; R[%eax] ← R[%edx]:R[%eax] ÷ S Signed divide of %eax by S Quotient stored in %eax Remainder stored in %edx
divl	S	$R[\%edx] \leftarrow R[\%edx]:R[\%eax] \mod S;$ $R[\%eax] \leftarrow R[\%edx]:R[\%eax] \div S$ Unsigned divide of %eax by S Quotient stored in %eax Remainder stored in %edx

ALL THE CHEAT SHEETS (COMPARISONS)

Instruction		Effect
cmpb	S2, S1	Compare byte S1 and S2, Sets conditional flags based on $S1 - S2$.
cmpw	S2, S1	Compare word S1 and S2, Sets conditional flags based on $S1 - S2$.
cmpl	S2, S1	Compare double word S1 and S2, Sets conditional flags based on $S1-S2$.
testb	S2, S1	Compare byte S1 and S2, Sets conditional flags based on S1 & S2.
testw	S2, S1	Compare word S1 and S2, Sets conditional flags based on S1 & S2.
testl	S2, S1	Compare double word S1 and S2, Sets conditional flags based on S1 & S2.

ALL THE CHEAT SHEETS (SET)

Instruction		Effect
sete/ setz	D	D ← ZF ("set if equal to 0")
setne/ setnz	D	D ← ~ZF (set if not equal to 0)
sets	D	$D \leftarrow SF$ (set if negative)
setns	D	D ← ~SF (set if nonnegative)
setg/ setnle	D	$D \leftarrow \sim (SF \land OF) \& \sim ZF \text{ (set if greater (signed >))}$
setge/ setnl	D	D ← ~(SF ^ OF) (set if greater or equal (signed >=))
setl/ setnge	D	D ← SF ^ OF (set if less than (signed <))
setle/ setng	D	$D \leftarrow (SF \land OF) \mid ZF \text{ (set if less than or equal (signed <=))}$
seta/ setnbe	D	$D \leftarrow \sim CF \& \sim ZF \text{ (set if above (unsigned >))}$
setae/ setnb	D	D ← ~CF (set if above or equal (unsigned >=))
setb/ setnae	D	$D \leftarrow CF$ (set if below (unsigned <))
setbe/ setna	D	D ← CF ZF (set if below or equal (unsigned <=))

ALL THE CHEAT SHEETS (JUMP)

Instructio	ns	Effect
jmp	Label	Jump to label
jmp	*Operand	Jump to specified locations
je/ jz	Label	Jump if equal/zero (ZF)
jne/ jnz	Label	Jump if not equal/ nonzero (~ZF)
js	Label	Jump if negative (SF)
jns	Label	Jump if nonnegative (~SF)
jg/ jnle	Label	Jump if greater (signed) (~(SF $^{\circ}$ OF) & ~ZF)
jge/jnl	Label	Jump if greater or equal (signed) (~(SF ^ OF))
jl/ jnge	Label	Jump if less (signed) (SF ^ OF)
jle/ jng	Label	Jump if less or equal (signed) ((SF ^ OF) ZF)
ja/ jnbe	Label	Jump if above (unsigned) (~CF & ~ZF)
jae/ jnb	Label	Jump if above or equal (unsigned) (~CF)
jb/ jnae	Label	Jump if below (unsigned) (CF)
jbe/ jna	label	Jump if below or equal (unsigned) (CF ZF)

ALL THE CHEAT SHEETS (CMOVE)

Instruction		Effect
cmove/ cmovz	S, R	S ← R if Equal/zero (ZF)
cmovne/ cmovnz	S, R	S ← R if Not equal/ not zero (~ZF)
cmovs	S, R	$S \leftarrow R$ if Negative (SF)
cmovns	S, R	S ← R if Nonnegative (~SF)
cmovg/ cmovnle	S, R	$S \leftarrow R$ if Greater (signed >) (~(SF ^ OF) & ~ZF)
cmovge/ cmovnl	S, R	S ← R if Greater or equal (signed >=) (~(SF ^ OF))
cmovl/ cmovnge	S, R	S ← R if Less (signed <) (SF ^ OF)
cmovle/ cmovg	S, R	S ← R if Less or equal (signed <=) ((SF ^ OF) ZF)
cmova/ cmovnbe	S, R	$S \leftarrow R$ if Above (unsigned >) (~CF & ~ZF)
cmovae/ cmovnb	S, R	S ← R if Above or equal (unsigned >=) (~CF)
cmovb/ cmovnae	S, R	$S \leftarrow R$ if Below (unsigned <) (CF)
cmovbe/ cmovna	S, R	S ← R if Below or equal (unsigned <=) (CF SF)

ALL THE CHEAT SHEETS (CALLING)

Instruction		Effect
call	Label	Push return and jump to label
call	*operand	Push return and jump to specified location
leave		Prepare stack for return. Set stack pointer to %ebp and pop top stack into %ebp. In assembly: mov %ebp, %esp pop %ebp
ret		Pop return address from stack and jump there

JUMPS, IN DEPTH

test %al,%al 4011ed jne



jump to 4011ed

• The test instruction is usually followed by jump if equal/ not equal

cmpl \$0x5,0x14(%rsp)4011d0 jg



if ((%al & %al) != 0) if (0x14(%rsp) > \$0x5) jump to 4011d0

> • For conditional jumps, it is usually the second argument greater/less than first argument

JE, JNE, JLE, JGE, ETC

- Jump if equal == Jump if zero
 - If the previous result was 0, jump
- Jump if not equal == Jump if not zero
 - If the previous result was not 0, jump
- Don't worry about the conditional flags
 - Just remember "if second argument greater/less than first argument"

DR. EVIL AND BOMBLAB

- 6 stages, each asking for input
 - Wrong input → bomb explodes (lose 1/2 point)
 - Score rounds up, so first explosion is free
 - Each stage may have multiple answers
- You get:
 - Bomb executable
 - Partial source of Dr. Evil mocking you
- Speed up next phase traversal with a text file
 - Place answers on each line
 - Run with bomb as ./bomb <solution file>

HOW IT WORKS

- "But how do I find the solutions if I don't have C code to work from?"
 - Read a lot of bomb disassembly
 - All of the phases are just loops and patterns
 - GDB
- If you're not working on a shark machine, your bomb won't work.
 - Will get an "illegal host" error

WORKING THROUGH THIS THING

- Read the disassembly
 - phase_1, phase_2, phase_3...
 - explode_bomb
 - Possible to reason through solutions without using GDB

• GNU Debugger

- Step through each instruction, examine registers...
- Set up breakpoints
- Make sure to run "kill" when you hit the explode_bomb breakpoint
 - You're screwed once you hit here, so why not exit?

BUT I DON'T KNOW HOW TO GDB??

- Here have a cheat sheet
 - http://csapp.cs.cmu.edu/public/docs/gdbnotes-x86-64.pdf
 - Everything you need to use GDB to solve bomb lab
- The Internet has a great range of commands you might find useful

GDB'S MOST USEFUL

- o run/ run <arguments>
 - Runs the program up till the next breakpoint.
- o disassemble/ disas
 - Shows the current function with an arrow to the next
 - WARNING: shortcut "disa" disables all breakpoints
- o step/ stepi/ nexti
 - stepi steps to the next line of Assembly.
 - nexti does the same but doesn't stop in function calls.
 - stepi *n* or nexti *n* steps through *n* lines.

GDB'S MOST USEFUL

o break <location>

- Sets breakpoint. Location can be function name or address.
- Stop at an instruction address with break *address
- You have to reset your break points when you restart GDB!

o x <address/register>

- Dereference the address or value in the register and print the contents to the console
- Give it a format to print out to, ie. "x/s" prints as string

o p <address/register/variable>

- Print the contents of the register, or the variable, or the address to the console
- Give it a format to print out to, ie. "p/s" prints as string

GETTING STARTED

- Download and untar ON A SHARK MACHINE
 - tar xvf *labhandout.tar*
- o shark> objdump -d bomb > filename
 - Outputs the whole bomb assembly code to a filename
- o shark> objdump -t bomb > filename
 - Contains locations of globals, variables, etc
- o shark> strings bomb > filename
 - All printable strings used in your bomb
- o shark> gdb bomb
 - Prepares to run the bomb in gdb

SPEED UP THE WAIT

- When you have solutions, put it into a text file
 - Separate each solution with a newline
 - Your bomb will auto-advance completed phases with pre-filled solutions
- Then when you run gdb next time:
 - (gdb)> run solution_file

BOMB LAB SPECIFICS

- o int sscanf (const char *s, const char *format, ...);
 - S
 - Source string to retrieve data from
 - format
 - Formatting string used to get values from the source string
 - ...
 - Depending the format string, one location (address) per formatter used to hold values extracted from source string

SSCANF EXAMPLE

```
#include <stdio.h>
int main () {
  char sentence[]="Rudolph is 12 years old";
  char str[20];
  int i;
  sscanf (sentence, "%s %*s %d", str, &i);
  printf ("%s -> %d\n", str, i);
  return 0;
```

Outputs: Rudolph -> 12

RELEVANCE TO BOMB LAB

- Why do we care about sscanf?
 - Mainly used to read in arguments
 - Keep track of which locations the read in values will be stored
 - Important for knowing where arguments will be stored
 - And how they will be used
 - They will usually be store in memory/ on the stack

More Bomb Lab Specifics

Jump tables

- In memory, you can think of it as an "array" of locations to jump to
- Using assembly it is possible to index into the "array"
- Each entry of the array will hold addresses of instructions

JUMP TABLES

- The tip-off is something like this:
 - jmpq *0x400600(,%rax,8)
 - Empty base means implied 0
 - o %rax is the "index"
 - 8 is the "scale"
 - In a jump tables, 64-bit machine addresses are 8 bytes
 - * indicates a dereference (as in regular C)
 - o Like leal: does not do a dereference even with parenthesis
 - Put it all together: "Jump to the address stored in the address 0x400600 + %rax*8"
- Using GDB (example output): x/8g 0x400600

0x400600: 0x00000000004004d1 0x00000000004004c8

0x400610: 0x00000000004004c8 0x00000000004004be

0x400620: 0x00000000004004c1 0x00000000004004d7

0x400630: 0x00000000004004c8 0x00000000004004be

TOP UNIX COMMANDS

- o man to read manual pages
- cd to change directories
- 1s to list contents of the current directory
- o ls −l to list with more info, including permissions
- scp to send files between your computer and others
- ssh to log into time shares
- tar to tar (-cvf) and untar (-xvf) (-z for optional gzip)
- chmod (ugo)+/-(rwx) to change permission bits
- Helpful hints
 - Tab auto-completes.
 - An up arrow scrolls up through your last few commands.

DEMO TIME



CREDITS & QUESTIONS

- StackOverflow on Assembly Projects
- P. 274 of CS:APP x86_64 Registers
- P. 171 221 of CS:APP Assembly Instructions
- o CPlusPlus Reference on sscanf