

Shell Lab + How to C

Agenda

- **Problem 5 and 9**
- C Stuff from Cachelab.
- Processes and Signals.
- Shellab hints and tips.

Problem 5

- Initial stack frame

VALUE	ADDRESS
UNK	0XFFFF1008
18	0XFFFF1004
213	0XFFFF1000
Return address	0XFFFF0FFC

- 1st call stack frame

VALUE	ADDRESS
UNK	0XFFFF0FF8
UNK	0XFFFF0FF4
UNK	0XFFFF0FF0
15	0XFFFF0FEC
18	0XFFFF0FE8
0X080483B7	0XFFFF0FE4

Problem 5

- 2nd stack frame

VALUE	ADDRESS
0XFFFF0FF8	0XFFFF0FE0
UNK	0XFFFF0FDC
UNK	0XFFFF0FD8
3	0XFFFF0FD4
15	0XFFFF0FD0
0X080483B7	0XFFFF0FCC

- 3rd stack frame

VALUE	ADDRESS
0XFFFF0FE0	0XFFFF0FC8
UNK	0XFFFF0FC4
UNK	0XFFFF0FC0
0	0XFFFF0FBC
3	0XFFFF0FB8

Problem 9

- Part A:
 - Accessing A row wise for 1 miss ... 15 hits.
 - Accessing B column wise for 1 miss ... 1 hit.
 - Total is roughly $\frac{1}{2}$
 - The idea here is to look at access patterns and not individual memory accesses.
 - Need to move to a higher level.

Problem 9

- Part 9
 - Access A row wise.
 - Access B row wise.
 - Both are 1 miss ... 15 hits.
 - Approx miss rate $1/16$.

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C Stuff from Cachelab

- Always check for NULL return from a malloc.
- If your code does not look like this:

```
char *foo = malloc(200);  
if (!foo) {  
    error("malloc failed");  
}
```

- Then you are doing something terribly wrong.
- This should be second nature.

C Stuff from Cachelab

- Sanity checking arguments.
- If you get a value from outside of your program, and you expect it to have a particular value, you *must* check to make sure it has that value.

C Stuff from Cachelab

- 80 column limit:
 - `grep '\{81,\}' *.c`
- Indentation, convert all tabs to spaces.
 - `:set expandtab` and `:retab` in VIM
 - `M-x untabify` in EMACS

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Processes

- Four basic process control functions
 - fork()
 - exec()
 - exit()
 - wait()
- Standard on all Unix systems

Processes

- `fork()`
 - Creates a process
 - Parent and child are exactly alike
 - Equal but separate
 - Execution (%eip)
 - Registers
 - Memory
 - File **descriptors**, the files themselves are shared.

Processes

- `exec()`
 - Replaces process context
 - How programs are run
 - Replace memory image with a new program
 - Set up stack with arguments
 - Start execution at the entry point (main)
 - A family of functions (man 3 exec)

Processes

- `exit()`
 - Terminates a process
 - OS frees resources used by the process
 - Tiny leftover data
 - Exit status for the parent
 - Must be freed
 - Which brings us to ...

Processes

- `wait()`
 - Waits for a child to change state
 - If a child terminates, the parent “reaps” the child, freeing all resources and getting the exit status
 - Lots of details (man 2 wait)

Signals

- Interprocess communication/notification
- **Asynchronous** with normal execution
- Come in many types (man 7 signal)
- Sent in various ways
 - ^C, ^Z, ^\
 - kill (which we will demonstrate)

Signals

- Disposition
 - Ignore
 - Catch and run a signal handler
 - Terminate
 - man sigaction
- Blocking
 - man sigprocmask
- Waiting
 - man sigsuspend (don't need this, but cool)

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Shellab

- Read the code we've given you.
 - There is a lot of stuff you don't need to write for yourself.
 - It's a good example of the kind of code we expect from you.

Shellab

- If you find yourself using `sleep()` as a way of avoiding race conditions, you are doing it VERY wrong. We will dock performance points for this.
- You should only use it for performance to avoid your code having to execute useless instructions. Your code should still work if we remove calls to sleep.

Shellab

- Hazards
 - Race conditions
 - Hard to debug so start early.
 - Reaping zombies
 - Race conditions
 - Fiddle with signals
 - Waiting for foreground job
 - One of the only places where sleep is acceptable (though you don't NEED it)