# 15-213

*"The course that gives CMU its Zip!"*

# Machine-Level Programming II
# Control Flow
# Sept. 14, 2000

## Topics

- **Condition Codes**
  - Setting
  - Testing
- **Control Flow**
  - If-then-else
  - Varieties of Loops
  - Switch Statements



# Condition Codes

## Single Bit Registers

- `CF` Carry Flag
- `ZF` Zero Flag
- `SF` Sign Flag
- `OF` Overflow Flag

## Implicit Setting By Arithmetic Operations

`addl` *Src,Dest*

C analog: `t = a+b`

- **CF set if carry out from most significant bit**
  - Used to detect unsigned overflow
- **ZF set if** `t == 0`
- **SF set if** `t < 0`
- **OF set if two's complement overflow**

```
(a>0 && b>0 && t<0) || (a<0 && b<0 && t>0)
```

## *Not* Set by `leal` instruction



# Setting Condition Codes (cont.)

## Explicit Setting by Compare Instruction

`cmpl` *Src2,Src1*

- `cmpl b,a` **like computing** `a-b` **without setting destination**
- **CF set if carry out from most significant bit**
  - Used for unsigned comparisons
- **ZF set if** `a == b`
- **SF set if** `(a-b) < 0`
- **OF set if two's complement overflow**

```
(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)
```

## Explicit Setting by Test instruction

`testl` *Src2,Src1*

- **Sets condition codes based on value of** *Src1* `&` *Src2*
  - Useful to have one of the operands be a mask
- `testl b,a` **like computing** `a&b` **without setting destination**
- **ZF set when** `a&b == 0`
- **SF set when** `a&b < 0`



# Reading Condition Codes

## SetX Instructions

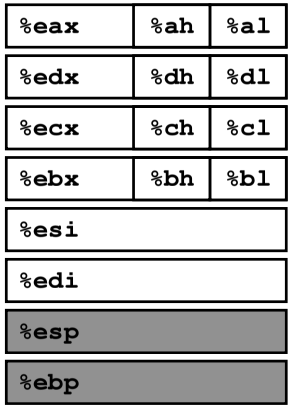
- **Set single byte based on combinations of condition codes**

| SetX | Condition | Description |
|---|---|---|
| `sete` | `ZF` | Equal / Zero |
| `setne` | `~ZF` | Not Equal / Not Zero |
| `sets` | `SF` | Negative |
| `setns` | `~SF` | Nonnegative |
| `setg` | `~(SF^OF)&~ZF` | Greater (Signed) |
| `setge` | `~(SF^OF)` | Greater or Equal (Signed) |
| `setl` | `(SF^OF)` | Less (Signed) |
| `setle` | `(SF^OF)\|ZF` | Less or Equal (Signed) |
| `seta` | `~CF&~ZF` | Above (unsigned) |
| `setb` | `CF` | Below (unsigned) |

# Reading Condition Codes (Cont.)

## SetX Instructions

- **Set single byte based on combinations of condition codes**
- **One of 8 addressable byte registers**
  - Embedded within first 4 integer registers
  - Does not alter remaining 3 bytes
  - Typically use `andl 0xFF,%eax` to finish job




```
int gt (int x, int y)
{
  return x > y;
}
```

Body

```
movl 12(%ebp),%eax   # eax = y
cmpl %eax,8(%ebp)    # Compare x : eax
setg %al             # al = x > y
andl $255,%eax       # Zero rest of %eax
```

Note inverted ordering!



# Jumping

## jX Instructions

- **Jump to different part of code depending on condition codes**

| jX | Condition | Description |
|---|---|---|
| jmp | 1 | Unconditional |
| je | ZF | Equal / Zero |
| jne | ~ZF | Not Equal / Not Zero |
| js | SF | Negative |
| jns | ~SF | Nonnegative |
| jg | ~(SF^OF)&~ZF | Greater (Signed) |
| jge | ~(SF^OF) | Greater or Equal (Signed) |
| jl | (SF^OF) | Less (Signed) |
| jle | (SF^OF)\|ZF | Less or Equal (Signed) |
| ja | ~CF&~ZF | Above (unsigned) |
| jb | CF | Below (unsigned) |



# Conditional Branch Example

```
int max(int x, int y)
{
  if (x > y)
    return x;
  else
    return y;
}
```

```
_max:
        pushl %ebp              } Set
        movl %esp,%ebp          } Up

        movl 8(%ebp),%edx       }
        movl 12(%ebp),%eax      }
        cmpl %eax,%edx          } Body
        jle L9                  }
        movl %edx,%eax          }
L9:

        movl %ebp,%esp          }
        popl %ebp               } Finish
        ret                     }
```



# Conditional Branch Example (Cont.)

```
int goto_max(int x, int y)
{
  int rval = y;
  int ok = (x <= y);
  if (ok)
    goto done;
  rval = x;
done:
  return rval;
}
```

- **C allows "goto" as means of transferring control**
  - Closer to machine-level programming style
- **Generally considered bad coding style**

```
    movl 8(%ebp),%edx    # edx = x
    movl 12(%ebp),%eax   # eax = y
    cmpl %eax,%edx       # x : y
    jle L9               # if <= goto L9
    movl %edx,%eax       # eax = x   } Skipped when x ≤ y
L9:                      # Done:
```

# "Do-While" Loop Example

### C Code

```
int fact_do
  (int x)
{
  int result = 1;
  do {
    result *= x;
    x = x-1;
  } while (x > 1);
  return result;
}
```

### Goto Version

```
int fact_goto(int x)
{
  int result = 1;
loop:
  result *= x;
  x = x-1;
  if (x > 1)
    goto loop;
  return result;
}
```

- Use backward branch to continue looping
- Only take branch when "while" condition holds



# "Do-While" Loop Compilation

### Goto Version

```
int fact_goto
  (int x)
{
  int result = 1;
loop:
  result *= x;
  x = x-1;
  if (x > 1)
    goto loop;
  return result;
}
```

### Registers

```
%edx   x
%eax   result
```

### Assembly

```
_fact_goto:
   pushl %ebp          # Setup
   movl %esp,%ebp      # Setup
   movl $1,%eax        # eax = 1
   movl 8(%ebp),%edx   # edx = x

L11:
   imull %edx,%eax     # result *= x
   decl %edx           # x--
   cmpl $1,%edx        # Compare x : 1
   jg L11              # if > goto loop

   movl %ebp,%esp      # Finish
   popl %ebp           # Finish
   ret                 # Finish
```



# General "Do-While" Translation

### C Code

```
do
  Body
  while (Test);
```

### Goto Version

```
loop:
  Body
  if (Test)
    goto loop
```

- *Body* can be any C statement
  - Typically compound statement:

```
{
  Statement1;
  Statement2;
    ...
  Statementn;
}
```

- *Test* is expression returning integer
  - = 0 interpreted as false
  - ≠0 interpreted as true



# "While" Loop Example #1

### C Code

```
int fact_while
  (int x)
{
  int result = 1;
  while (x > 1) {
    result *= x;
    x = x-1;
  };
  return result;
}
```

### First Goto Version

```
int fact_while_goto
  (int x)
{
  int result = 1;
loop:
  if (!(x > 1))
    goto done;
  result *= x;
  x = x-1;
  goto loop;
done:
  return result;
}
```

- Is this code equivalent to the do-while version?
- Must jump out of loop if test fails

# Actual "While" Loop Translation

### C Code

```
int fact_while(int x)
{
  int result = 1;
  while (x > 1) {
    result *= x;
    x = x-1;
  };
  return result;
}
```

- Uses same inner loop as do-while version
- Guards loop entry with extra test

### Second Goto Version

```
int fact_while_goto2
  (int x)
{
  int result = 1;
  if (!(x > 1))
    goto done;
loop:
  result *= x;
  x = x-1;
  if (x > 1)
    goto loop;
done:
  return result;
}
```



# General "While" Translation

### C Code

```
while (Test)
  Body
```

### Do-While Version

```
  if (!Test)
    goto done;
  do
    Body
    while(Test);
done:
```

### Goto Version

```
  if (!Test)
    goto done;
loop:
  Body
  if (Test)
    goto loop;
done:
```



# "While" Loop Example #2

```
/* Compute x raised to nonnegative power p */
int ipwr_while(int x, unsigned p)
{
  int result = 1;
  while (p) {
    if (p & 0x1)
      result *= x;
    x = x*x;
    p = p>>1;
  }
  return result;
}
```

### Algorithm

- Exploit property that $p = p_0 + 2p_1 + 4p_2 + \ldots 2^{n-1}p_{n-1}$
- Gives: $x^p = z_0 \cdot z_1^{\,2} \cdot (z_2^{\,2})^2 \cdot \ldots \cdot (\ldots((z_{n-1}^{\,2})^2)\ldots)^2$ ($n$ times)
  - $z_i = 1$ when $p_I = 0$
  - $z_i = x$ when $p_I = 1$
- Complexity $O(\log p)$

**Example**

$3^{10}$

$= 3^2 * 3^8$

$= 3^2 * ((3^2)^2)^2$



# `ipwr` Computation

```
int ipwr(int x, unsigned p)
{
  int result = 1;
  while (p) {
    if (p & 0x1)
      result *= x;
    x = x*x;
    p = p>>1;
  }
  return result;
}
```

| result | x | p |
|---|---|---|
| 1 | 3 | 10 |
| 1 | 9 | 5 |
| 9 | 81 | 2 |
| 9 | 6561 | 1 |
| 531441 | 43046721 | 0 |

# "While" → "Do-While " → "Goto "

```
  int result = 1;
  while (p) {
    if (p & 0x1)
      result *= x;
    x = x*x;
    p = p>>1;
  }
```

```
  int result = 1;
  if (!p) goto done;
  do {
    if (p & 0x1)
      result *= x;
    x = x*x;
    p = p>>1;
  } while (p);
done:
```

```
  int result = 1;
  if (!p)
    goto done;
loop:
  if (!(p & 0x1))
    goto skip;
  result *= x;
skip:
  x = x*x;
  p = p>>1;
  if (p)
    goto loop;
done:
```

- Also converted conditional update into test and branch around update code



# Example #2 Compilation

## Goto Version

```
  int result = 1;
  if (!p)
    goto done;
loop:
  if (!(p & 0x1))
    goto skip;
    result *= x;
skip:
    x = x*x;
    p = p>>1;
  if (p)
    goto loop;
done:
```

## Registers

| | |
|---|---|
| %ecx | x |
| %edx | p |
| %eax | result |

```
   pushl %ebp          # Setup
   movl %esp,%ebp      # Setup
   movl $1,%eax        # eax = 1
   movl 8(%ebp),%ecx   # ecx = x
   movl 12(%ebp),%edx  # edx = p
   testl %edx,%edx     # Test p
   je L36              # If 0, goto done
L37:                 # Loop:
   testb $1,%dl        # Test p & 0x1
   je L38              # If 0, goto skip
   imull %ecx,%eax     # result *= x
L38:                 # Skip:
   imull %ecx,%ecx     # x *= x
   shrl $1,%edx        # p >>= 1
   jne L37             # if p goto Loop
L36:                 # Done:
   movl %ebp,%esp      # Finish
   popl %ebp           # Finish
   ret                 # Finish
```



# "For" Loop Example

```
  int result;
  for (result = 1;
       p != 0;
       p = p>>1) {
    if (p & 0x1)
      result *= x;
    x = x*x;
  }
```

## General Form

```
for (Init; Test; Update)
    Body
```

## *Init*

```
result = 1
```

## *Test*

```
p != 0
```

## *Update*

```
p = p >> 1
```

## *Body*

```
{
  if (p & 0x1)
    result *= x;
  x = x*x;
}
```



# "For"→ "While"

## For Version

```
for (Init; Test; Update)
    Body
```

## While Version

```
Init;
while (Test) {
    Body
    Update ;
}
```

## Do-While Version

```
  Init;
  if (!Test)
    goto done;
  do {
    Body
    Update ;
  } while (Test)
done:
```

## Goto Version

```
  Init;
  if (!Test)
    goto done;
loop:
  Body
  Update ;
  if (Test)
    goto loop;
done:
```

# "For" Loop Compilation

## Goto Version

```
  Init;
  if (!Test)
    goto done;
loop:
  Body
  Update ;
  if (Test)
    goto loop;
done:
```

```
  result = 1;
  if (p == 0)
    goto done;
loop:
  if (p & 0x1)
    result *= x;
  x = x*x;
  p = p >> 1;
  if (p != 0)
    goto loop;
done:
```

**Init**

```
result = 1
```

**Test**

```
p != 0
```

**Update**

```
p = p >> 1
```

**Body**

```
{
  if (p & 0x1)
    result *= x;
  x = x*x;
}
```



# Switch Statements

```
typedef enum
 {ADD, MULT, MINUS, DIV, MOD, BAD}
    op_type;

char unparse_symbol(op_type op)
{
  switch (op) {
  case ADD :
    return '+';
  case MULT:
    return '*';
  case MINUS:
    return '-';
  case DIV:
    return '/';
  case MOD:
    return '%';
  case BAD:
    return '?';
  }
}
```

## Implementation Options

- **Series of conditionals**
  - Good if few cases
  - Slow if many
- **Jump Table**
  - Lookup branch target
  - Avoids conditionals
  - Possible when cases are small integer constants
- **GCC**
  - Picks one based on case structure
- **Bug in example code**
  - No default given



# Jump Table Structure

## Switch Form

```
switch(op) {
  case 0:
    Block 0
  case 1:
    Block 1
    • • •
  case n–1:
    Block n–1
}
```

## Jump Table

```
jtab:  Targ0
       Targ1
       Targ2
         •
         •
         •
       Targn–1
```

## Jump Targets

```
Targ0:   Code Block 0
Targ1:   Code Block 1
Targ2:   Code Block 2
           •
           •
           •
Targn–1: Code Block n–1
```

## Approx. Translation

```
target = JTab[op];
goto *target;
```



# Switch Statement Example

## Branching Possibilities

```
typedef enum
  {ADD, MULT, MINUS, DIV, MOD,
BAD}
    op_type;

char unparse_symbol(op_type op)
{
  switch (op) {
    • • •
  }
}
```

## Enumerated Values

```
ADD   0
MULT  1
MINUS 2
DIV   3
MOD   4
BAD   5
```

## Setup:

```
unparse_symbol:
 pushl %ebp              # Setup
 movl %esp,%ebp          # Setup
 movl 8(%ebp),%eax       # eax = op
 cmpl $5,%eax            # Compare op : 5
 ja .L49                 # If > goto done
 jmp *.L57(,%eax,4)      # goto Table[op]
```

# Assembly Setup Explanation

## Symbolic Labels

- Labels of form `.LXX` translated into addresses by assembler

## Table Structure

- Each target requires 4 bytes
- Base address at `.L57`

## Jumping

`jmp .L49`

- Jump target is denoted by label `.L49`

`jmp *.L57(,%eax,4)`

- Start of jump table denoted by label `.L57`
- Register `%eax` holds `op`
- Must scale by factor of 4 to get offset into table
- Fetch target from effective Address `.L57 + op*4`



# Jump Table

## Table Contents

```
.section .rodata
   .align 4
.L57:
  .long .L51 #Op = 0
  .long .L52 #Op = 1
  .long .L53 #Op = 2
  .long .L54 #Op = 3
  .long .L55 #Op = 4
  .long .L56 #Op = 5
```

## Enumerated Values

```
ADD    0
MULT   1
MINUS  2
DIV    3
MOD    4
BAD    5
```

## Targets & Completion

```
.L51:
    movl $43,%eax # '+'
    jmp .L49
.L52:
    movl $42,%eax # '*'
    jmp .L49
.L53:
    movl $45,%eax # '-'
    jmp .L49
.L54:
    movl $47,%eax # '/'
    jmp .L49
.L55:
    movl $37,%eax # '%'
    jmp .L49
.L56:
    movl $63,%eax # '?'
    # Fall Through to .L49
```



# Switch Statement Completion

```
.L49:                   # Done:
    movl %ebp,%esp      # Finish
    popl %ebp           # Finish
    ret                 # Finish
```

## Puzzle

- What value returned when `op` is invalid?

## Answer

- Register `%eax` set to `op` at beginning of procedure
- This becomes the returned value

## Advantage of Jump Table

- Can do *k*-way branch in *O*(1) operations



# Object Code

## Setup

- Label `.L49` becomes address `0x804875c`
- Label `.L57` becomes address `0x8048bc0`

```
08048718 <unparse_symbol>:
8048718:55              pushl   %ebp
8048719:89 e5           movl    %esp,%ebp
804871b:8b 45 08        movl    0x8(%ebp),%eax
804871e:83 f8 05        cmpl    $0x5,%eax
8048721:77 39           ja      804875c <unparse_symbol+0x44>
8048723:ff 24 85 c0 8b jmp     *0x8048bc0(,%eax,4)
```

# Object Code (cont.)

## Jump Table

- Doesn't show up in disassembled code
- Can inspect using GDB

```
gdb code-examples
(gdb) x/6xw 0x8048bc0
```

  - Examine 6 hexadecimal format "words" (4-bytes each)
  - Use command "`help x`" to get format documentation

```
0x8048bc0 <_fini+32>:
  0x08048730
  0x08048737
  0x08048740
  0x08048747
  0x08048750
  0x08048757
```



# Extracting Jump Table from Binary

## Jump Table Stored in Read Only Data Segment (.rodata)

- Various fixed values needed by your code

## Can examine with objdump

```
objdump code-examples -s --section=.rodata
```

- Show everything in indicated segment.

## Hard to read

- Jump table entries shown with reversed byte ordering

```
Contents of section .rodata:
 8048bc0 30870408 37870408 40870408 47870408  0...7...@...G...
 8048bd0 50870408 57870408 46616374 28256429  P...W...Fact(%d)
 8048be0 203d2025 6c640a00 43686172 203d2025   = %ld..Char = %
 …
```

- E.g., `30870408` really means `0x08048730`



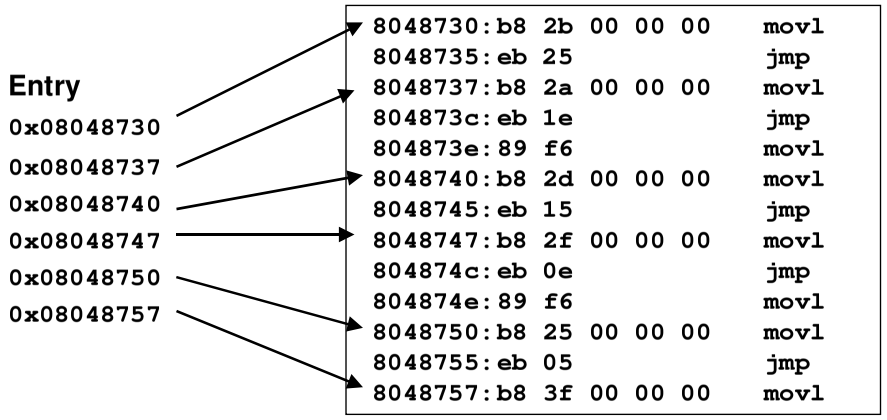
# Disassembled Targets

- No-operations (`movl %esi,%esi`) inserted to align target addresses

```
8048730:b8 2b 00 00 00   movl   $0x2b,%eax
8048735:eb 25            jmp    804875c <unparse_symbol+0x44>
8048737:b8 2a 00 00 00   movl   $0x2a,%eax
804873c:eb 1e            jmp    804875c <unparse_symbol+0x44>
804873e:89 f6            movl   %esi,%esi
8048740:b8 2d 00 00 00   movl   $0x2d,%eax
8048745:eb 15            jmp    804875c <unparse_symbol+0x44>
8048747:b8 2f 00 00 00   movl   $0x2f,%eax
804874c:eb 0e            jmp    804875c <unparse_symbol+0x44>
804874e:89 f6            movl   %esi,%esi
8048750:b8 25 00 00 00   movl   $0x25,%eax
8048755:eb 05            jmp    804875c <unparse_symbol+0x44>
8048757:b8 3f 00 00 00   movl   $0x3f,%eax
```



# Matching Disassembled Targets



# Summarizing

## C Control

- if-then-else
- do-while
- while
- switch

## Assembler Control

- jump
- Conditional jump

## Compiler

- Must generate assembly code to implement more complex control

## Standard Techniques

- All loops converted to do-while form
- Large switch statements use jump tables

## Conditions in CISC

- CISC machines generally have condition code registers

## Conditions in RISC

- Use general registers to store condition information
- Special comparison instructions
- E.g., on Alpha:

```
cmple $16,1,$1
```

  - Sets register $1 to 1 when Register $16 <= 1