

# Recitation 3

## Scan

### 3.1 Announcements

- *SkylineLab* has been released, and is due **Friday afternoon**. It's worth 125 points.
- *BignumLab* will be released on Friday.

## 3.2 What is scan?

In the SEQUENCE library, there is a symmetry among certain aggregation functions:

Sequential	Parallel
iterate	reduce
iteratePrefixes	scan
iteratePrefixesIncl	scanIncl

We can see this symmetry in their types...

iterate	$(\beta * \alpha \rightarrow \beta) \rightarrow \beta \rightarrow \alpha \text{ seq} \rightarrow \beta$
reduce	$(\alpha * \alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \alpha \text{ seq} \rightarrow \alpha$
iteratePrefixes	$(\beta * \alpha \rightarrow \beta) \rightarrow \beta \rightarrow \alpha \text{ seq} \rightarrow \beta \text{ seq} * \beta$
scan	$(\alpha * \alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \alpha \text{ seq} \rightarrow \alpha \text{ seq} * \alpha$
iteratePrefixesIncl	$(\beta * \alpha \rightarrow \beta) \rightarrow \beta \rightarrow \alpha \text{ seq} \rightarrow \beta \text{ seq}$
scanIncl	$(\alpha * \alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \alpha \text{ seq} \rightarrow \alpha \text{ seq}$

...as well as their output behavior: each of the parallel functions has output identical to its sequential analog under the condition that the first two arguments are an *associative function* and a corresponding *identity*, respectively.

**Definition 3.1.** A function  $f$  is associative if for every  $x, y, z$ ,

$$f(f(x, y), z) = f(x, f(y, z)).$$

**Definition 3.2.** A value  $b$  is an identity of a binary function  $f$  if for every  $x$ ,

$$f(b, x) = x = f(x, b).$$

So, for now, you can think of `scan` as a magical function which performs iteration of an associative function in parallel. If the function is constant-time, then an application of `scan` has linear work and logarithmic span.

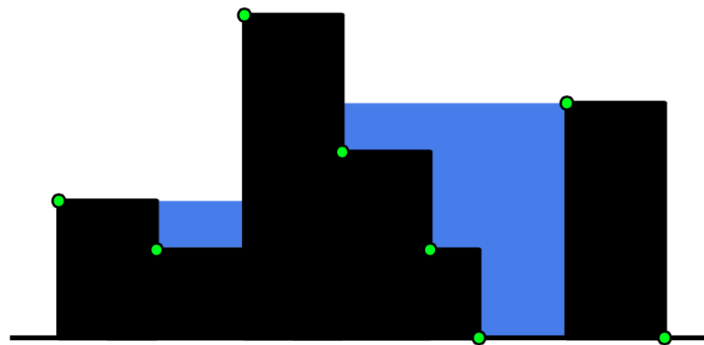
**Remark 3.3.** In reality, we can relax the constraint on the identity. It only needs to be an identity for the values encountered during the execution of the `reduce`, `scan`, or `scanIncl`. For example, if  $S$  is a sequence of non-negative integers, then  $(\text{scan Int.max } 0 S)$  will still be logically equivalent to  $(\text{iteratePrefixes Int.max } 0 S)$ , despite the fact that, in general,  $0$  is not an identity for `Int.max`.

### 3.3 Skyline-Fill

For this example, we'll use the same conventions given in *SkylineLab*:

- Skylines are sequences of points  $(x, y)$  sorted by  $x$ -coordinate,
- all  $x$ -coordinates are unique and non-negative, and
- all  $y$ -coordinates (heights) are non-negative.

Imagine pouring water on a skyline. How much water can it hold?



**Task 3.4.** Implement the function

```
val fill : (int * int) Seq.t → int
```

where  $(fill\ S)$  returns the area of water which can fill the skyline  $S$ . Your implementation should have  $O(|S|)$  work and  $O(\log |S|)$  span.

### 3.4 A Group at Dinner

A group of  $n$  friends sit around a circular table at a restaurant. Some of them know what they want to order; some of them don't. The ones who don't know what to order decide to pick the same thing as the person on their left.

**Task 3.5.** *Implement the function*

```
val groupOrder : (int → α option) → int → α Seq.t
```

where  $(\text{groupOrder } f \ n)$  returns the sequence of orders of a group of  $n$  people.  $f(i)$  is either the preferred order of the  $i^{\text{th}}$  person, or `NONE` if they don't know what they want. Assume the people are labeled  $0$  to  $n - 1$  counter-clockwise, and that at least one person originally knows what they want to order. Your implementation should have  $O(n)$  work and  $O(\log n)$  span.

## 3.5 Bonus Exercises

**Exercise 3.6.** Implement `parenMatch` (from the previous recitation) using `scan` such that it has linear work and logarithmic span. Try adapting the iterative approach.

**Exercise 3.7.** Implement `parenDist` (from `ParenLab`) using `scan` such that it has linear work and logarithmic span.

**Exercise 3.8.** Did you know that you can calculate the first  $n$  Fibonacci numbers in  $O(n)$  time and  $O(\log n)$  span? We claim that if we extend the Fibonacci sequence as so...

$$\begin{aligned} F_{-1} &= 1 \\ F_0 &= 0 \\ F_1 &= 1 \\ &\vdots \\ F_n &= F_{n-1} + F_{n-2} \end{aligned}$$

...that the following holds for  $n \geq 0$  (easily provable via induction):

$$\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^n = \begin{pmatrix} F_{n+1} & F_n \\ F_n & F_{n-1} \end{pmatrix}$$

Using this fact, implement a function

```
val fibs : int → int Seq.t
```

which returns the first  $n$  Fibonacci numbers in  $O(n)$  work and  $O(\log n)$  span. Use `scan` to compute prefixes of matrix multiplications.

**Exercise 3.9.** Implement a function

```
val parenPairs : Paren.t Seq.t → (int * int) Seq.t
```

where `(parenPairs S)` returns a sequence of index-pairs where each pair contains the index of a parenthesis as well as its matching partner. For example the input `(( ) ) ( )` should yield some permutation of  $\{(0, 3), (4, 5), (1, 2)\}$ . Your implementation should have  $O(|S| \log |S|)$  work and  $O(\log^2 |S|)$  span.

*Hint:* try marking each parenthesis with how many other parentheses are enclosing it. You might also need to `sort` at some point...

