

# Recitation 13

## Hashing

### 13.1 Announcements

- *DPLab* has been released, and is due **Monday, April 18**. It's worth 140 points.
- *PASLLab* will be released on Monday, April 18.

## 13.2 Removing Duplicates

Removing duplicates is a crucial substep of many interesting algorithms. For example, in BFS, consider the step where we construct a new frontier. One viable method would be to generate the sequence of all out-neighbors, and then remove duplicates:

$$F' = \text{removeDuplicates } \langle v : u \in F, v \in N_G^+(u) \rangle$$

So, how fast is it to remove duplicates? Can we do it in parallel?

### 13.2.1 Sequential

Before we think about parallelism, we should acquaint ourselves with a good sequential algorithm solving the same problem. This way, we know what to shoot for in terms of work bounds, since we want our parallel algorithm to be asymptotically work-efficient.

**Task 13.1.** *Describe a sequential algorithm which performs expected  $O(n)$  work to remove duplicates from a sequence of length  $n$ . Also argue that  $\Omega(n)$  work is necessary in order to solve this problem, and conclude that your algorithm is asymptotically optimal.*

*Hint: try hashing elements one at a time.*

### 13.2.2 Parallel

**Task 13.2.** *Implement a function*

```
val removeDuplicates : ( $\alpha \times \text{int} \rightarrow \text{int}$ )  $\rightarrow \alpha \text{ Seq.t} \rightarrow \alpha \text{ Seq.t}$ 
```

*where  $(\text{removeDuplicates } h \ S)$  returns a sequence of all unique elements of  $S$ , given that  $h(e, m)$  hashes the element  $e$  to a uniform random integer in the range  $[0, m)$  (thus the probability of collision for any two distinct elements is  $1/m$ ).*

*Hint: as a first attempt, try simultaneously hashing as many elements as possible all at the same time. What do you do when elements collide?*