

Lecture Notes on First-Order Modal Logic

15-816: Modal Logic
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Lecture 12
February 25, 2010

1 Introduction to This Lecture

In this lecture, we will introduce first-order modal logic and start considering its relationship to classical first-order logic.

2 Quantified Modal Logic

In this section, we define the syntax and semantics of quantified modal logic [Car46, Kri63]. An excellent source on first-order modal logic, its various variations and pitfalls is the book by Fitting and Mendelsohn [FM99].

We fix a set Σ of function symbols and predicate symbols, with arities associated (number of arguments) and a set of logical variables. Terms are defined as in first-order logic. The syntax of classical first-order modal logic is defined as follows:

Definition 1 (First-order modal formulas) *The set $\text{Fml}_{\text{FOML}}(\Sigma)$ of formulas of classical quantified modal logic a.k.a. first-order modal logic is the smallest set with:*

- If $p \in \Sigma$ is a predicate symbol of arity n and $\theta_1, \dots, \theta_n$ are terms then $p(\theta_1, \dots, \theta_n) \in \text{Fml}_{\text{FOML}}(\Sigma)$.
- If $\phi, \psi \in \text{Fml}_{\text{FOML}}(\Sigma)$, then $\neg\phi, (\phi \wedge \psi), (\phi \vee \psi), (\phi \rightarrow \psi) \in \text{Fml}_{\text{FOML}}(\Sigma)$.
- If $\phi \in \text{Fml}_{\text{FOML}}(\Sigma)$ and x is a logical variable, then $(\forall x \phi) \in \text{Fml}_{\text{FOML}}(\Sigma)$ and $(\exists x \phi) \in \text{Fml}_{\text{FOML}}(\Sigma)$.

- If $\phi \in \text{Fml}_{\text{FOML}}(\Sigma)$ and $x \in V$, then $(\Box\phi), (\Diamond\phi) \in \text{Fml}_{\text{FOML}}(\Sigma)$.

There are several variations for the definition of semantics for quantified modal logic. Here is one variant:

Definition 2 (Kripke structure) A Kripke structure $K = (W, \rho, M)$ consists of Kripke frame (W, ρ) and a mapping M that assigns first-order structures $M(s)$ to each world s such that, for each $s, t \in W$ with $s \rho t$, the structure $M(s)$ is a substructure of $M(t)$, i.e.:

- the universe of $M(s)$ is a subset of the universe of $M(t)$ (monotonicity), and
- the structures $M(s)$ and $M(t)$ agree on the interpretation of all function symbols on the (smaller) universe of $M(s)$.

Another common case in the semantics is that of *constant domain*, where all worlds in a Kripke structure are required to share the same universe.

Definition 3 (Interpretation of quantified modal formulas) Given a Kripke structure $K = (W, \rho, M)$, the interpretation \models of modal formulas in a world s is defined as

1. $K, s \models p(\theta_1, \dots, \theta_n)$ iff $M(s) \models p(\theta_1, \dots, \theta_n)$.
2. $K, s \models \phi \wedge \psi$ iff $K, s \models \phi$ and $K, s \models \psi$.
3. $K, s \models \phi \vee \psi$ iff $K, s \models \phi$ or $K, s \models \psi$.
4. $K, s \models \neg\phi$ iff it is not the case that $K, s \models \phi$.
5. $K, s \models \forall x \phi(x)$ iff $K, s \models \phi(d)$ for all d in the universe of s
6. $K, s \models \exists x \phi(x)$ iff $K, s \models \phi(d)$ for some d in the universe of s
7. $K, s \models \Box\phi$ iff $K, t \models \phi$ for all worlds t with $s \rho t$.
8. $K, s \models \Diamond\phi$ iff $K, t \models \phi$ for some world t with $s \rho t$.

When K is clear from the context, we sometimes abbreviate $K, s \models \phi$ by $s \models \phi$.

In constant domain semantics, quantifiers refer to the same set of objects in all worlds. In varying domain semantics, quantifiers may possibly refer to a different set of objects, depending on the world.

Exercises

Exercise 1 Recall Definition 3 of interpretation of quantified modal formulas. The definition is imprecise at some points. What is the problem, why is it a problem, and what can be done to fix it?

References

- [Car46] Rudolf Carnap. Modalities and quantification. *J. Symb. Log.*, 11(2):33–64, 1946.
- [FM99] Melvin Fitting and Richard L. Mendelsohn. *First-Order Modal Logic*. Kluwer, Norwell, MA, USA, 1999.
- [Kri63] Saul A. Kripke. Semantical considerations on modal logic. *Acta Philosophica Fennica*, 16:83–94, 1963.