

Grothendieck's Inequality:

Given numbers $a_{ij} \in \mathbb{R}$ (both positive or negative) want to figure out how.

OPT :=
$$\max_{\substack{x_i \in \{\pm 1\} \\ y_j \in \{\pm 1\}}} \sum a_{ij} x_i y_j$$

relates to

SDP :=
$$\max_{\substack{\|u_i\|_2 = 1 \\ \|v_j\|_2 = 1}} \sum a_{ij} \langle u_i, v_j \rangle \quad \leftarrow \text{the SDP relaxation.}$$

Fact: $SDP \geq OPT$ (it is a relaxation)

Thm!: $OPT \geq 0.56 \cdot SDP$ [Krivine, can do better...]

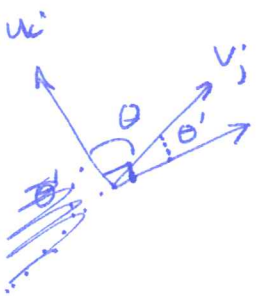
Pf: Suppose we try the simplest idea. These vectors u_i, v_j are on unit ball. So take the random hyperplane, and set $\hat{x}_i = 1$ if $\langle \hat{g}, u_i \rangle \geq 0$
 $= -1$ if $\langle \hat{g}, u_i \rangle < 0$ etc

$$\Rightarrow \hat{x}_i = \text{sign}(\langle \hat{g}, u_i \rangle), \hat{y}_j = \text{sign}(\langle \hat{g}, v_j \rangle).$$

then
$$E[\hat{x}_i \hat{y}_j] = -1 \cdot \Pr(i, j \text{ separated}) + 1 \cdot \Pr(i, j \text{ not separated})$$

$$= -1 \cdot \frac{\theta}{\pi} + 1 \cdot (1 - \frac{\theta}{\pi}) = 1 - \frac{2\theta}{\pi} = \frac{2}{\pi} (\frac{\pi}{2} - \theta)$$

$$= \frac{2}{\pi} \theta'$$



$$SDP_{ij} = \langle u_i, v_j \rangle = \cos \theta = \sin \theta'$$

~~$$\Rightarrow E[A_{ij}] = \sum a_{ij} \hat{x}_i \hat{y}_j = SDP_{ij}$$~~

Here's the problem: for Max Cut we were happy to say that $E[A_{ij}]$ was bigger than SDP_{ij} . But now we have weights that are negative, so cannot do this!

Krivine's Trick:

Lemma: Given $u_1, \dots, u_n, v_1, \dots, v_m$, can get \tilde{u}_i, \tilde{v}_j such that

$$\frac{\pi}{2} E[\underbrace{\text{sign}(\langle \tilde{u}_i, g \rangle)}_{\hat{x}_i} \cdot \underbrace{\text{sign}(\langle \tilde{v}_j, g \rangle)}_{\hat{y}_j}] \stackrel{\text{equality!!}}{=} c \cdot \langle u_i, v_j \rangle$$

N.b. LHS = $\frac{\pi}{2} \cdot A_{ij}$ $\Rightarrow E[A_{ij}] = E[\sum_{ij} a_{ij} \hat{x}_i \hat{y}_j] = \sum_{ij} a_{ij} \cdot \frac{c}{(\pi/2)} \langle u_i, v_j \rangle = \frac{c}{(\pi/2)} \cdot \text{SDP}$.

Also: LHS, by calculation above = $\sin^{-1}(\langle \tilde{u}_i, \tilde{v}_j \rangle)$.

So we want to prove that ~~sin~~ $\langle \tilde{u}_i, \tilde{v}_j \rangle = \sin(c \langle u_i, v_j \rangle)$.

Pf: $\sin(c \cdot \langle u_i, v_j \rangle) = \sum_{k=0}^{\infty} (-1)^k \cdot \frac{c^{2k+1}}{(2k+1)!} (u_i \cdot v_j)^{2k+1}$
 $= \sum_{k=0}^{\infty} (-1)^k \cdot \frac{c^{2k+1}}{(2k+1)!} (u_i^{\otimes 2k+1} \cdot v_j^{\otimes 2k+1})$

existence

So define ~~\tilde{u}_i~~ to be a vector whose k^{th} set of coordinates

are $(-1)^k \cdot \sqrt{\frac{c^{2k+1}}{(2k+1)!}} \cdot u_i^{\otimes 2k+1}$

\tilde{v}_j has $\sqrt{\frac{c^{2k+1}}{(2k+1)!}} v_j^{\otimes 2k+1}$

$(u \cdot v)^2 \text{ etc.} = u^{\otimes 2} \cdot v^{\otimes 2}$

then by defⁿ $\langle \tilde{u}_i, \tilde{v}_j \rangle = \sin(c \langle u_i, v_j \rangle)$.

What are lengths of \tilde{u}_i ? $\|\tilde{u}_i\|^2 = \sum_{k=0}^{\infty} \frac{c^{2k+1}}{(2k+1)!} \|u_i\|^{(2k+1) \cdot 2} = \sinh(c \|u_i\|^2) = \sinh(c)$

So choose $c = \sinh^{-1}(1) = \ln(1 + \sqrt{2}) = 1$.

similarly $\|\tilde{v}_j\|^2 = \sinh(c) = 1$.

Finally: this was an infinite construction. But this shows \exists vectors that achieve this.

How do you find these vectors? write an SDP for vectors \tilde{u}_i, \tilde{v}_j st. (3)

$$\langle \tilde{u}_i, \tilde{v}_j \rangle = \cos(\angle c \langle u_i, v_j \rangle) \quad \forall i, j$$

↑ inputs!

this SDP gives the solution you want.

Further work: this gave the Grothendieck constant to be ≈ 0.56 . ← Krivine's bound.

[Alon Makarychev², Naor]. Given a general graph, the gap between OPT and SDP

$$\text{OPT} = \max_{\{x_i\}} \sum_{ij \in E} a_{ij} \langle x_i, x_j \rangle \quad \text{vs} \quad \text{SDP} = \max_{\|v_i\|=1} \sum_{ij \in E} a_{ij} \langle v_i, v_j \rangle$$

↑
same set of vars

is at most $O(\log \chi(\bar{G})) \leftarrow K(G)$

↑ the Lovasz theta fn on the complement of G .

For bipartite graphs (which are perfect)

$$\chi(\bar{G}) = \chi(G), \text{ get constant factor } \leftarrow \text{we gave explicit bound.}$$

↑ is true in general

(*) They also show that the Krivine bound is not tight, can beat it by a little bit.

Open Question: What is true bound for K_G ? Current bounds:-

$$\Omega(\log w(G)) \leq K_G \leq O(\log \chi(\bar{G}))$$