

## Lecture 3

### Local Optimizations

- I Basic blocks/Flow graphs
- II Abstraction 1: DAG
- III Abstraction 2: Value numbering

### I. Basic Blocks & Flow Graphs

- **What is**
  - a basic block?
  - a flow graph?
- **How do we restructure a sequential list of instructions into a flow graph of basic blocks?**
  - Muchnick pp. 174-177
- **Reachability of basic blocks**

```
if x {                                bfls r1, L1
    ...
    return;                            ...
                                         ret
                                         jmp L2
} else {                               L1: ...
    ...
}
```

## II. Local Optimizations

- Common subexpression elimination
  - array expressions
  - field access in records
  - access to parameters

## Graph Abstractions

- Example 1: an expression

$$a + a * (b - c) + (b - c) * d$$

- Muchnick Section 4.9.3

## How well do DAGs hold up across statements?

- **Example 2**

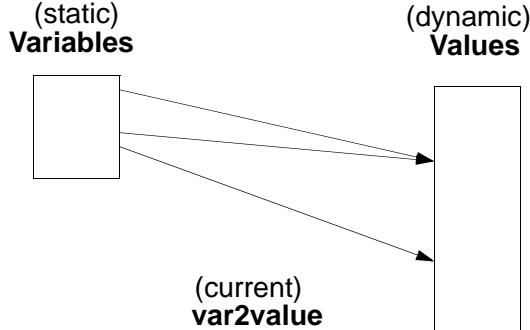
```
a = b+c;  
b = a-d;  
c = b+c;  
d = a-d;
```

## Critique of DAGs

- **Cause of problems**
  - Assignment statements
  - Value of variable depends on TIME
- **How to fix problem?**
  - build graph in order of execution
  - attach variable name to latest value
- **Final graph created is not very interesting**
  - Key: variable->value mapping across time
  - loses appeal of abstraction

### III. Value Numbering: Another Abstraction

- John Cocke & Jack Schwartz in unpublished book: "Programming Languages and their Compilers", (1970) (*Muchnick Section 12.4*)
- More explicit with respect to VALUES, and TIME



- each value has its own “number”
  - common subexpression means same value number
- var2value: current map of variable to value
  - used to determine the value number of current expression

$r1 + r2 \Rightarrow \text{var2value}(r1) + \text{var2value}(r2)$

### Algorithm

Data structure:

```
VALUES = Table of
    expression
    var      (temporary holding variable)
```

For each instruction ( $dst = op\ src1\ src2$ ) in execution order

```
IF [OP var2value(src1) var2value(src2)] is in VALUES
    v = the index of expression
    Replace instruction with CPY dst = VALUES[v].var
ELSE
    Add
        expression =[OP var2value(src1) var2value(src2)]
        var       =dst
        to VALUES
        v = index of new entry

set_var2value (dst, v)
```

## More Details

- What are the initial values of the variables?
  - values at beginning of the basic block
- Possible implementations:
  - Initialization: create “initial values” for all variables
  - Or dynamically create them as they are used
- Implementation of VALUES and var2value: hash tables

## Example

```
Assign: a->r1,b->r2,c->r3,d->r4
a = b+c;      ADD  t1 = r2,r3
              CPY   r1 = t1
b = a-d;      SUB  t2 = r1,r4
              CPY   r2 = t2
c = b+c;      ADD  t3 = r2,r3
              CPY   r3 = t3
d = a-d;      SUB  t4 = r1,r4
              CPY   r4 = t4
```

## Conclusions

- Comparisons of two abstractions
  - DAGs
  - Value numbering
- Value numbering
  - VALUE: distinguish between variables and VALUES
  - TIME
    - Interpretation of instructions in order of execution
    - Keep dynamic state information

## Question

- How do you extend value numbering to constant folding?

```
a = 1  
b = 2  
c = a+b
```