

# 15-451 Algorithms, Spring 2017

## Recitation #6 Worksheet

---

### Dynamic Programming

**Off-Line Stock Market Problem** You're given a sequence of stock prices  $[p_1, p_2, \dots, p_n]$ . You want to find the maximum profit that you could have made on the stock in hindsight. In other words, you want to find  $i$  and  $j$  with  $1 \leq i \leq j \leq n$  such that  $p_j - p_i$  is maximal. Your algorithm should run in  $O(n)$  time.

**Longest Increasing Subsequence:** Given an array  $A$  of  $n$  integers like  $[7\ 2\ 5\ 3\ 4\ 6\ 9]$ , find the longest subsequence that's in increasing order (in this case, it would be  $2\ 3\ 4\ 6\ 9$ ). Give a dynamic-programming algorithm that runs in time  $O(n^2)$  to solve this problem.

1. To keep things simple, first let's say you just need to output the \*length\* of the longest-increasing subsequence. E.g., in the above case, the length is 5.

*Hint: suppose that for each  $i' < i$  you have computed the length of the LIS of  $A_{0\dots i'}$  that ends with  $A[i']$ . How would you use this to solve the corresponding problem for  $i$ ?*

2. Now extend your solution to actually find the LIS.

**Closest Depot in a Tree:** You're given a rooted tree  $T$  with  $n$  vertices. There are  $m \leq n$  special vertices called *depots*. You are to compute, for every node  $v$  of  $T$ , the distance from  $v$  to the nearest depot. The distance is the number of tree edges that must be traversed to get there. (If  $v$  is a depot the distance is 0.)

Your algorithm should work by doing one or two depth-first searches (DFS) of the tree starting from the root, and it should run in  $O(n)$  time.

**Making Change:** You are given denominations  $v_1, v_2, \dots, v_n$  (all integers) of the various kinds of currency you have. (Say  $v_1 = 1$ , so you can make change for any integer amount  $C \geq 1$ .) Given  $C$ , give a dynamic programming solution which makes change for  $C$  with the fewest bills possible.

*(Again, as a first stab, compute the number of bills required, and then extend the solution to output the number of bills of each denomination needed.)*

**Making Change (Part II):** Now suppose you have only one bill of each denomination  $i$ . Given  $C$ , give a dynamic programming solution which makes change for  $C$  using the fewest bills, using no more than one bill of each denomination  $i$  (or says this is not possible).

**Making Change (Part III):** Can you solve the problem if you have  $\ell_i$  bills of denomination  $i$ ?

**Balanced Partition.** You have a set of  $n$  integers each in the range  $0, \dots, K$ . In time  $O(n^2K)$ , partition these integers into two subsets such that you minimize  $|S_1 - S_2|$ , where  $S_1$  and  $S_2$  denote the sums of the elements in each of the two subsets.