

15-441 Computer Networks

Lecture 1

Dave Eckhardt
Peter Steenkiste

(some slides courtesy of Hui Zhang)

Outline

- ▶ **First things first**
- ▶ **Administrative overview**
- ▶ **Course non-goals**
- ▶ **Course goals**
- ▶ **Key problems**
- ▶ **Network performance concepts**

First Things First

- ▶ **Please read Chapter 1 of the text**
 - ▶ To read ahead: most of Chapter 2, but ok to skip for now:
 - ▶ 2.5, Reliable Transmission
 - ▶ 2.8, Wireless
 - ▶ Also, don't kill yourself on 2.4, Error Detection
- ▶ **Please remind me to let you stretch**
 - I haven't taught an 80-minute class in 1.5 years

People

◆ Professors

- Peter Steenkiste (www.cs/~prs)
- Dave Eckhardt (www.cs/~davide)

◆ Teaching assistants

- Mike Cui
- Josh Hailpern
- David Murray
- [watch this space]

◆ Course secretary

- Barbara Grandillo, Wean Hall 8018

Information Sources

- **Watch the course web page**

- <http://www.cs.cmu.edu/~441>
- ***We expect you to read the syllabus!***
- Handouts, readings, ..

- **We expect you to read course bboards**

- Official announcements
 - academic.cs.15-441.announce
- Questions/answers
 - academic.cs.15-441

Information Sources

◆ Textbook

- Peterson and Davie, *Computer Networks: A Systems Approach*, 3rd Edition, Morgan Kaufmann, 2003

Information Sources

- ◆ **~30 lectures**
- ◆ **~3 paper homeworks**
- ◆ **1-2 lab homeworks**
 - Illustrate networking concepts
- ◆ **Mid-term and final exams**
- ◆ **1 programming assignment**
 - How to use a network
- ◆ **2 programming projects**
 - How to build a network

Grading

Homeworks	15%
Three projects	45%
Midterm exam	15%
Final exam	25%

- ▶ **Deadline means deadline**
 - Deadline is 11:59 pm on the specified date

Policy on Collaboration

- **Working together is important**
 - Discuss course material in general terms
 - Talk over tough debugging problems
- **Parts of the course must be done individually**
 - Homeworks, midterm, final, 1st programming assignment
- **Projects are done by two-student teams**
 - Learn how to collaborate
 - But each student must understand the entire project!
- **Web page has the details**

Course Non-goals

- ~~Learn how to configure a Cisco router~~
 - That requires a class all by itself
 - Cisco teaches those classes
 - Our perspective will be broader
- ~~Become “Internet Experts”~~
 - The Internet will be our frequent *motivating example*
 - Our perspective will be broader

Why not an “Internet class”?

- **Is there anything other than the Internet?**

Philosophy final exam question:

Define “Universe”. Give two examples.

Why not an “Internet class”?

- **Is there anything other than the Internet?**

Philosophy final exam question

Define “Universe”. Give two examples.

- **Yes – Internet in the 1800's!**

- **Yes – The secret network?**

- **[Yes – What's next?]**

Internet in the 1800's!?!?

- ▶ **Tom Standage, The Victorian Internet**
 - Telegraph!
 - Continent-spanning systems
 - Digital transmission of information
 - Crypto, code-breaking
 - Nerds
 - Attacks on the moral fiber of society
 - On-line dating (even an on-line wedding!)
- ▶ **Distributed message routing despite link outages**
 - Lines cut by armies in wartime!
- ▶ **Many problems; solutions eerily similar**

The Secret Network

- **With the Internet, who needs the phone system?**
 - It's a “new era”, etc.
- **One small detail...**
 - From inception, Internet has been a phone system application!
 - To connect two nodes, just ask your telco for a “circuit”...
 - ...somehow there's always copper/fiber waiting for you...
 - ...somehow when it breaks it gets fixed fast...
 - ...somehow your circuit can terminate anywhere...
 - Somehow?

Course Goals

- **Think “the network way”**
 - Distributed coordination is hard, let's go shopping
- **Learn how computer networks work**
 - Problems, approaches, protocols, software
- **Learn how to write network applications**
- **Hands-on understanding of network internals**
 - Build a simple network in software

Selected Key Problems

- ▶ **Two Generals**
- ▶ **Group Membership**
- ▶ **Scaling**

Two Generals

- **Problem (formulated by Jim Gray?)**
 - Two cooperating armies
 - Each size $2X$
 - Separated by...
 - One opposing army
 - Size $3X$
- **Idealized “combat” (think: Diplomacy, Risk, ...)**
 - $4X$ vs. $3X$: probable win
 - $2X$ vs. $3X$: certain loss

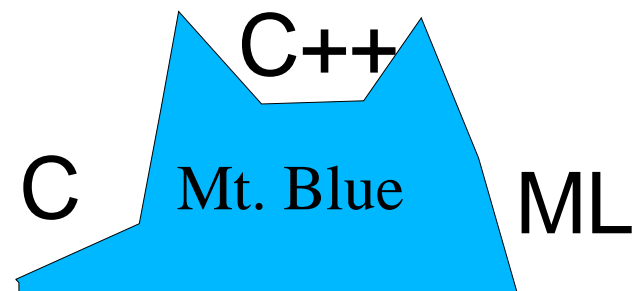
Two Generals

• Problem

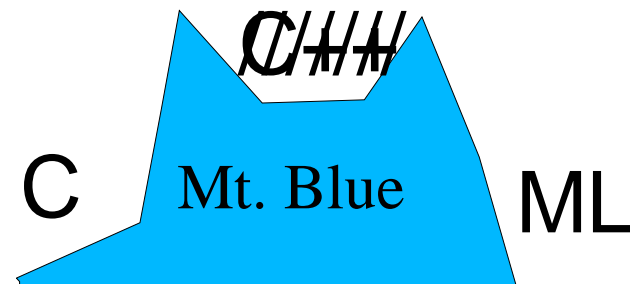
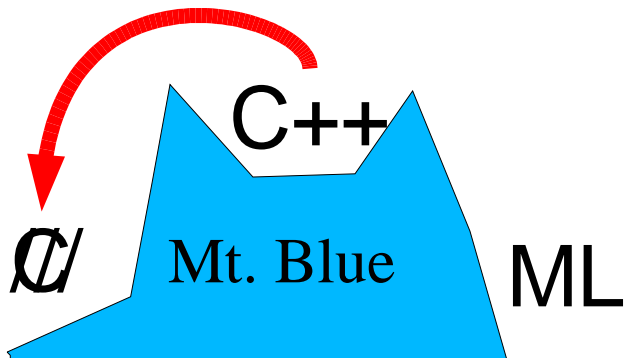
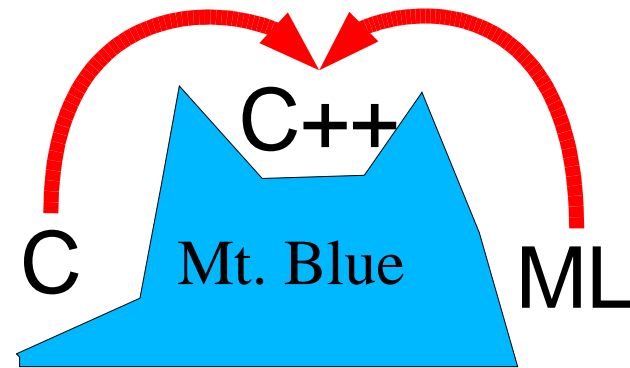
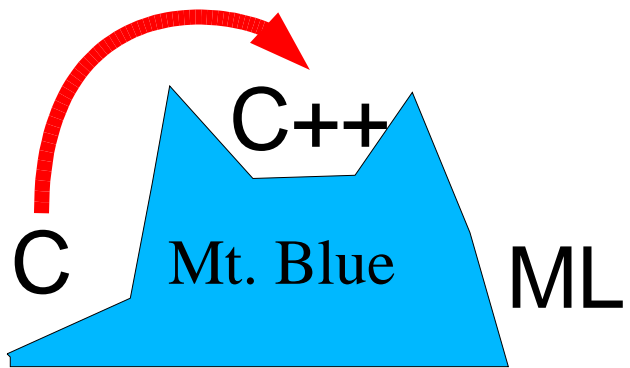
- Two cooperating armies
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• Idealized “combat”

- 4X vs. 3X: win
- 2X vs. 3X: lose



Two Generals



Two Generals

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 - Armies can communicate via messenger
- **Protocol 0**
 - C: “Attack at dawn!”
 - What if C's messenger is captured by C++?

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- **Protocol 1**
 - C: “Attack at dawn! Ok?”
 - ML: “Ok!”
 - What if ML's messenger is captured?
- **Seemingly-trivial coordination is *impossible!***

Group Membership

- ▶ **Group of nodes on a network**
 - Require distributed election of a “leader”
- ▶ **Sample solution**
 - “Distributed election” algorithm chooses among group members
 - If a node enters or leaves during election, re-start algorithm
- ▶ **Results**
 - Works great for 10 nodes
 - Fails horribly for 1,000,000 nodes
 - If inter-node-join time approximates election time...
 - Election process never completes

Group Membership

- **Problem: “group membership” is undefined**
 - By the time you can compute it, it's changed
- **Lots of algorithms will run into trouble**
 - “To acquire a node number, add one to the largest current node number” – oops!
- **Key network functions must face this environment**
 - Routing, naming

Scaling

- **“DOD Standard Internet Protocol”**
 - RFC 760, 1980: Addresses are fixed length of four octets (32 bits). An address begins with a one octet network number, followed by a three octet local address. This three octet field is called the "rest" field.
 - Result: 254 networks (surely enough!)

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- ▶ **Subsequently revised to Class A/B/C networks**
 - ~16k “Class B” networks of ~64k hosts (CMU)
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- ▶ **Then “subnets”, then “CIDR”**
- ▶ **“Surely enough” evaporates pretty fast!**

Network Performance Concepts

➤ **Throughput**

- “How many things per unit time?”
- Mb/s = megabits per second
- KB/s = kilobytes per second

➤ **Latency**

- “How long until my message arrives?”
- ms = millisecond (10^{-6}), μ s = microsecond (10^{-9})

▪ **Reciprocal “in theory”**

- bits/second = $(1/(\text{seconds/bit}))$

▪ **Relationship much more complex**

Hen Performance

➤ Old riddle

- “If a hen and a half lays an egg and a half in a day and a half, how long does it take to get a dozen eggs?”

➤ Egg Latency

- How long does it take for one hen to lay one egg?

➤ Henhouse throughput (eggs per day)

- Increases with number of hens
- Does not mean you can build henhouse, get first egg in 1 hour
- What is minimum time to 12 eggs?

Latency and Throughput

- ▶ **Radio a message to your friend Mike**
 - 1-megabyte photo
 - 1-megabit radio link
 - How long?

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Two Problems

• **Small problem**

- Mega != Mega
- Computer people: Megabyte (MB) = 2^{20} bytes
- Network people: Megabit (Mb) = 10^6 bits
- It's 8.4 seconds, not 8

Two Problems

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- Computer people: MegaByte (MB) = 2^{20} bytes
- Network people: Megabit (Mb) = 10^6 bits
- It's 8.4 seconds, not 8

• Big problem

- I forgot to tell you... Mike lives on the Moon
 - (Extra credit: What is Mike's last name?)
- It takes radio waves 1.3 seconds to get there

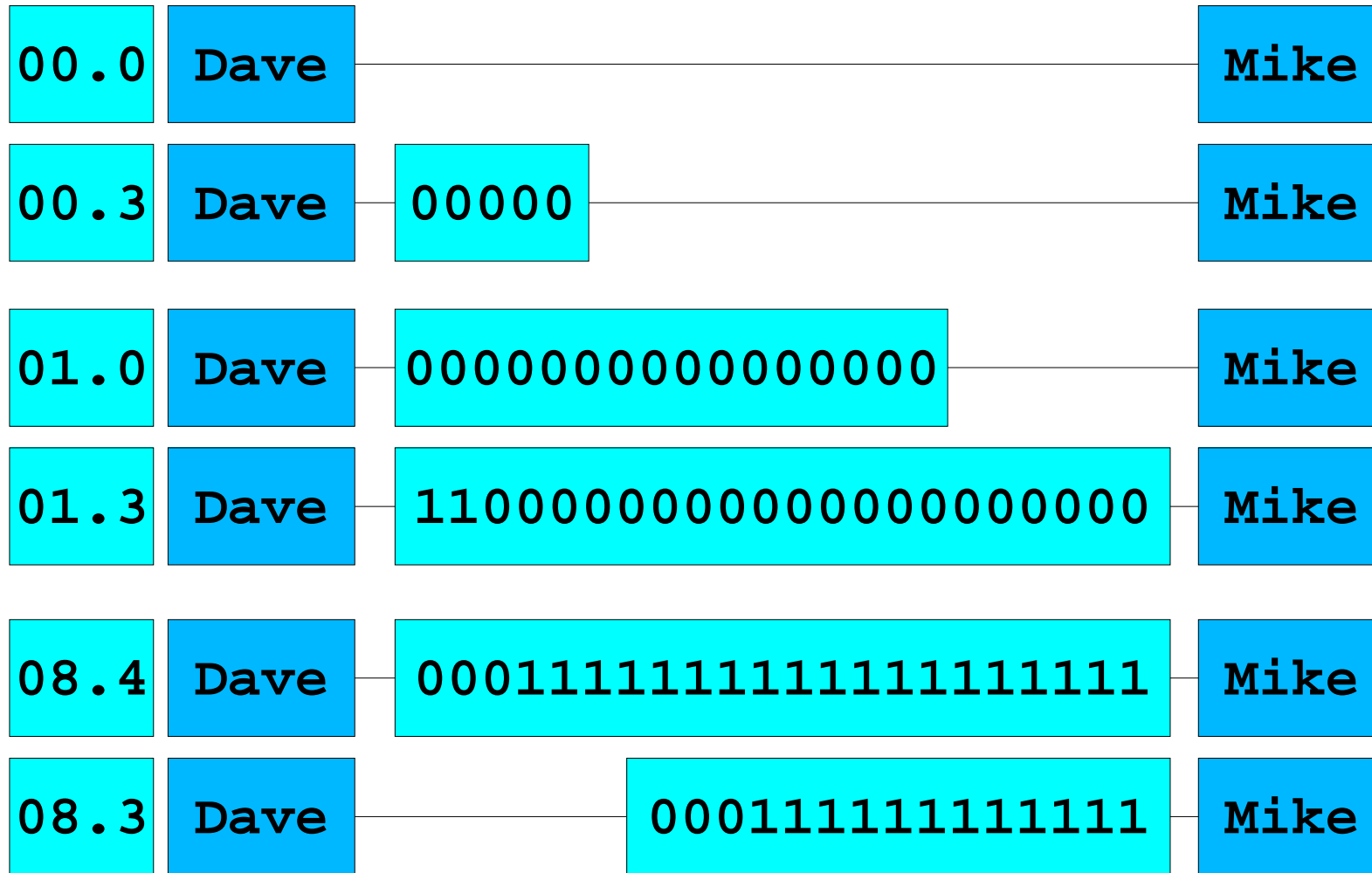
Message Latency

- **Message latency = sum of**
 - Propagation delay (distance/lightspeed)
 - Transmission time (size/throughput)
 - Queue delay (ignore for now)
- **Message to Mike**
 - Propagation delay is 1.3 seconds (one-way)
 - Also known as “link delay”
 - Transmission time is 8.4 seconds, total is 9.7 (121% of 8)
- **By the way: RTT (round-trip-time)**
 - Time to send a 0-bit message there and back: 2.6 seconds

Message Latency

- **Propagation delay vs. transmission time**
 - May vary widely
 - Earth-to-Moon is 1.3 seconds ($\ll 8.4$)
 - Delay is a minor compared to transmission time
 - Can transmit part of message, receive back status
 - “Got that part ok” or “Oops, send it again”

Message Latency

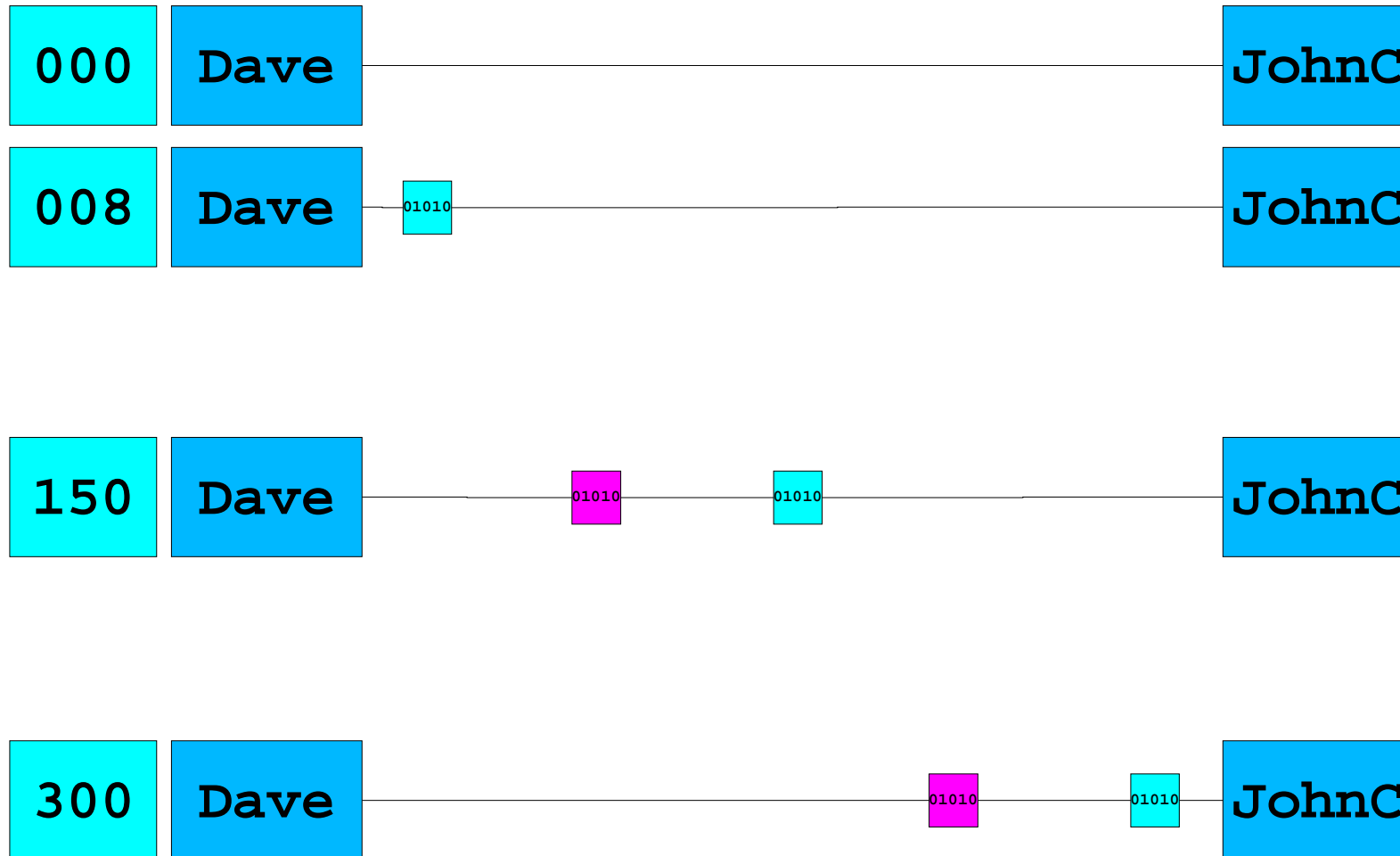


Message Latency

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 - “Got that part ok” or “Oops, send it again”
- Earth-to-Mars is 300-1225 seconds ($\gg 8.4$)
 - Delay vastly exceeds transmission time
 - Link holds *multiple entire messages*

Message Latency



Latency and Throughput

- ▶ **Bandwidth-delay product**
 - megabits/second X link-delay
 - This many bits are always “in flight” / “queued in link”
- ▶ **What if Mike says “Stop!! My buffer is full!”?**
 - One b-d product of bits are “in flight” to him already
 - You will queue another b-d product before you hear his alert!
- ▶ **Message throughput (\neq link throughput)**
 - How many messages per second can you send to Mike?
 - Depends on b-d product vs. message size
 - Depends on message protocol (= waiting protocol) you use

Latency and Throughput

- **See text for more-dignified treatment**
 - No hens, no Loonies, no Martians
- **Things to watch out for**
 - Is “delay” one-way or round-trip?
 - Mega vs. mega, kilo vs. kilo
 - Do we mean link latency or message latency?
 - Do we mean link throughput or message throughput?

Things Which Aren't Throughput

• **Bandwidth**

- Properly, measured in Hertz
- Difference between max & min frequency of transmission band
- Routinely abused by CS people to mean “throughput”

• **Goodput**

- Used to mean “productive throughput”
- Ignore “waste” if part of a message is transmitted multiple times

Back to the Internet

- **Another reason the Internet isn't perfect**
 - Fatally overoptimized for single-planet case
 - Will work to low-Earth orbit
 - Efficiency problems talking to the Moon
 - Forget about Mars
- **InterPlanetary Internet**
 - <http://www.ipnsig.org/>

Summary

- ◆ **First things first**
 - Read Chapter 1
 - Study socket-programming example as a refresher
- ◆ **Project 1 (individual) out Wednesday**
- ◆ **Course non-goals, goals**
 - “Networking perspective”, Internet as running example
- ◆ **Key problems**
 - Distributed coordination; scaling
- ◆ **Network performance concepts**
 - Throughput vs. latency, ...