Assignment 3: GraphRats

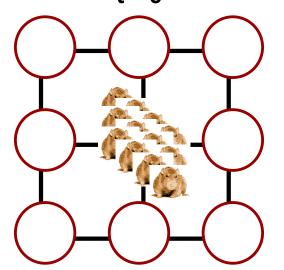


Topics

- Application
- Implementation Issues
- Optimizing for Parallel Performance
- Useful Advice

Basic Idea

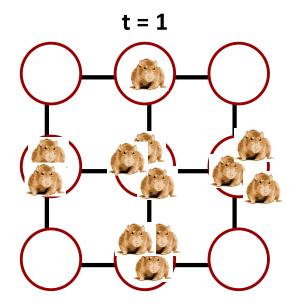
t = 0



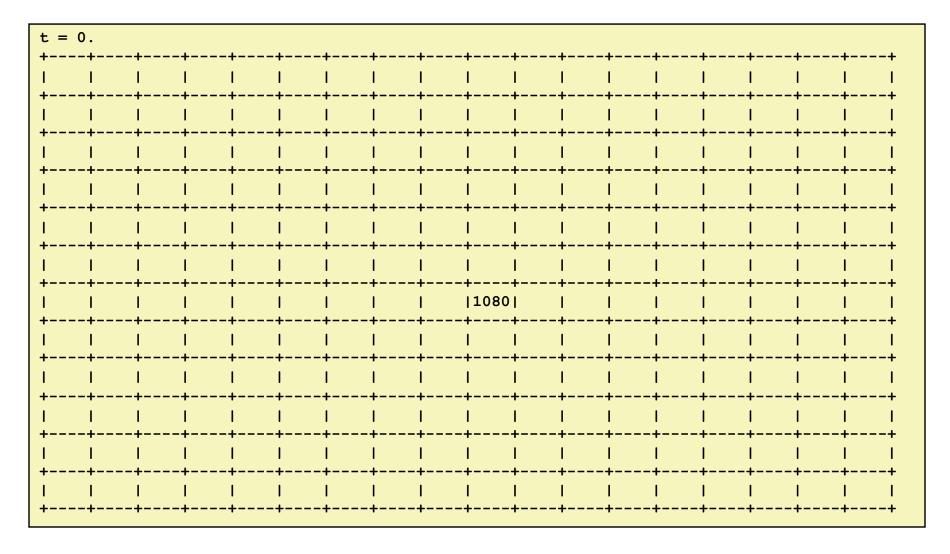
- Graph
 - WXH grid
- Initial State
 - Start with all *R* rats in center

Transitions

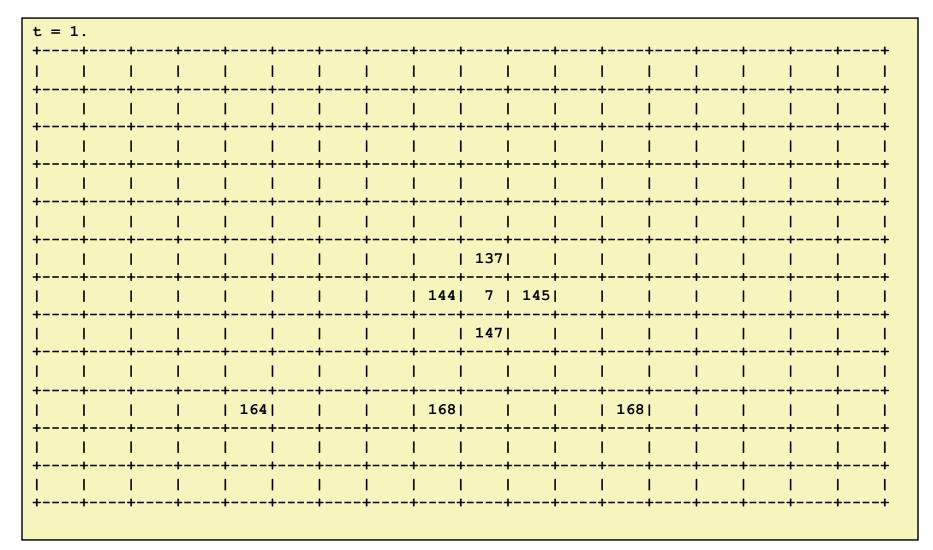
- Each rat decides where to move next
 - Don't like crowds
 - But also don't like to be alone
 - Weighted random choice



Node Count Representation (18 x 12)



Simulation Example (18 x 12)



Note moves to nodes not connected by grid

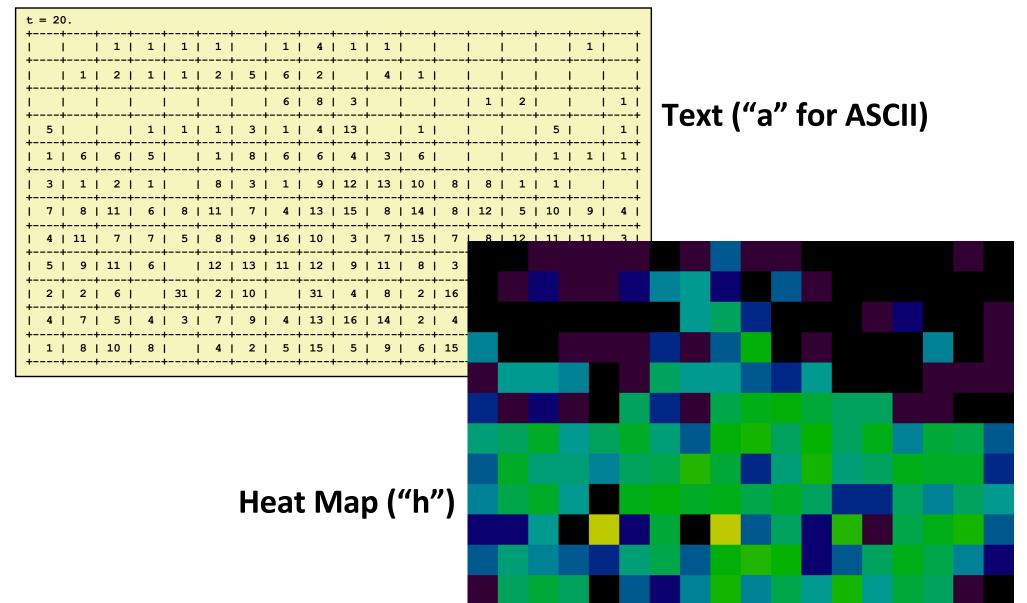
Explanation to follow

Simulation Example (18 x 12)

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-	1		10	-	-			-	15		•	-			10			

Rats dispersed across graph

Visualizations



Running it yourself

Demos

- 1: Text visualization, synchronous updates
- 2: Heap-map, synchronous updates

40



Count number of rats at current and adjacent locations

Adjacency structure represented as graph

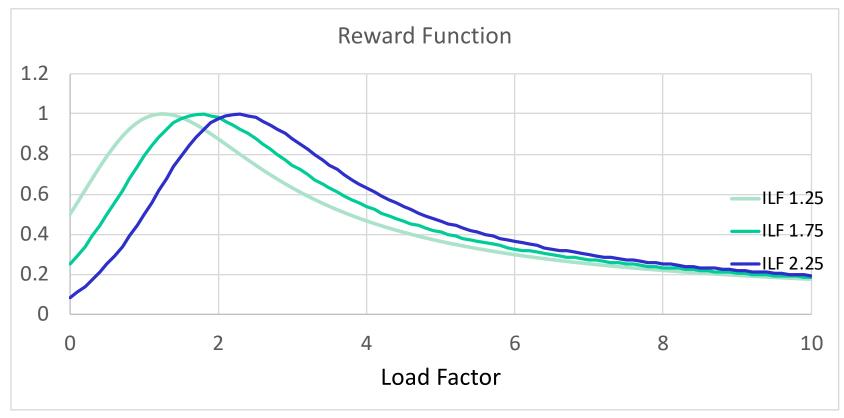
Compute reward value for each location

- Based on *load factor l* = count/average count
- *l* * Ideal load factor (ILF) (varying)
- α Fitting parameter (= 0.4)

Reward(l) =
$$\frac{1}{1 + (\log_2 [1 + \alpha (l - l^*)])^2}$$

Reward Function

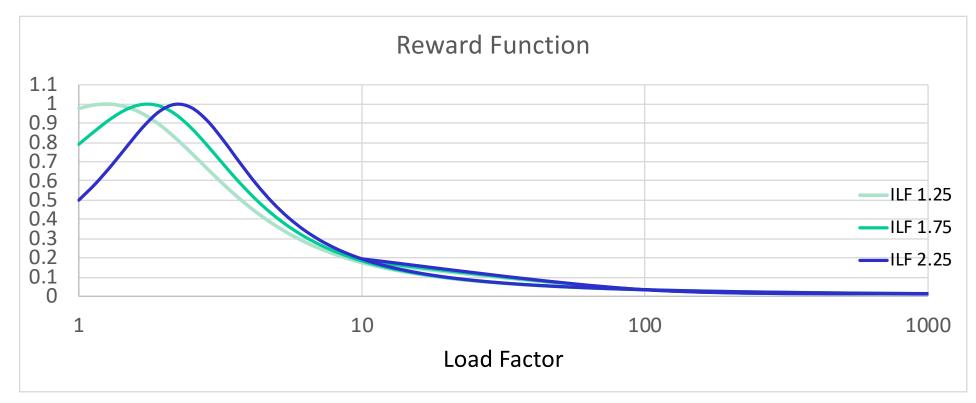
Reward(l) =
$$\frac{1}{1 + (\log_2 [1 + \alpha (l - l^*)])^2}$$



- Maximized at ILF
 - Just above average population
 - Drops for smaller loads (too few) and larger loads (too crowded)

Reward Function (cont.)

Reward(l) = $\frac{1}{1 + (\log_2 [1 + \alpha(l - l^*)])^2}$

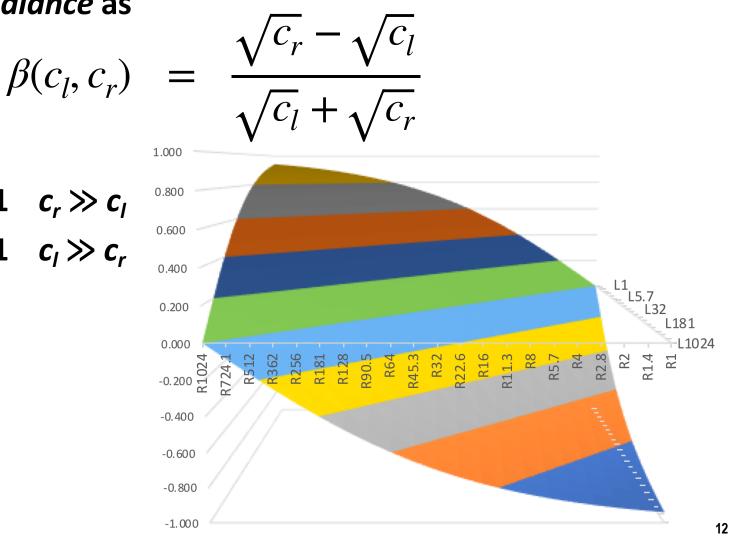


- Falls off gradually
 - *Reward*(1000) = 0.0132

Computing Ideal Load Factor (ILF)

- Suppose node has count c_l and neighbor has count c_r
- Compute *imbalance* as

- Maximum +1 $c_r \gg c_l$
- Minimum -1 $c_l \gg c_r$



Computing Ideal Load Factor (cont.)

For node *u* with population *p*(*u*)

$$\hat{\beta}(u) = Avg_{(u,v)\in E} \left[\beta(p(u), p(v))\right]$$

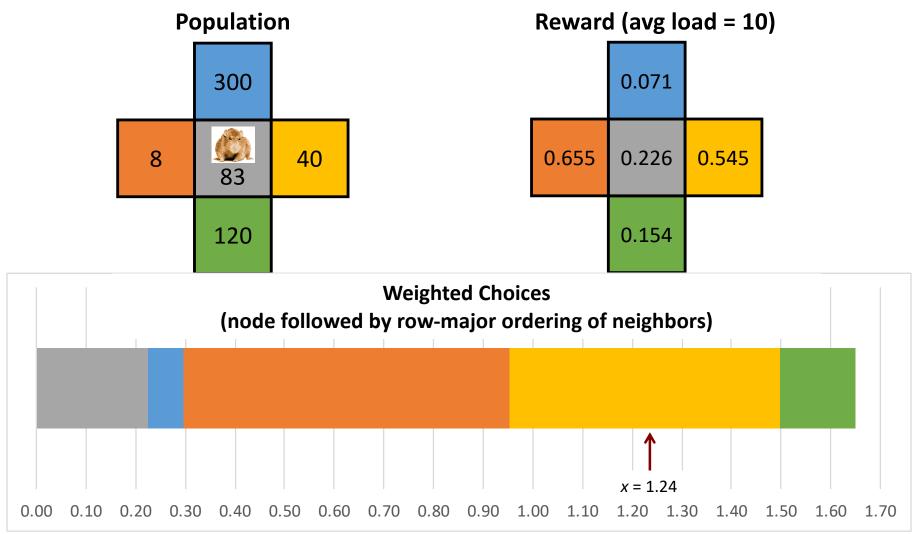
Define ILF as

$$l^*(u) = 1.75 + 0.5 \cdot \hat{\beta}(u)$$

Minimum 1.25

- When adjacent nodes much less crowded
- Maximum 2.25
 - When adjacent nodes much more crowded
- Changes as rats move around

Selecting Next Move



- Choose random number between 0 and sum of rewards
- Move according to interval hit

Update Models

Synchronous

- Demo 3
- Compute next positions for all rats, and then move them
- Causes oscillations/instabilities

Rat-order

- Demo 4
- For each rat, compute its next position and then move it
- Smooth transitions, but costly

Batch

- Demo 5
- For each batch of B rats, compute next moves and then move them
- B = 0.02 * *R*
- Smooth enough, with better performance possibilities

What We Provide

Python version of simulator

- Demos 1–2
- Very slow

C version of simulator

- Fast sequential implementation
- Demos 3–5: 36X32 grid, 11,520 rats
- Demos 6–11: 180X160 grid, 1,008,000 rats
 - That's what we'll be using for benchmarks
- Generate visualizations by piping C simulator output into Python simulator
 - Operating in visualization mode
 - See Makefile for examples

Correctness

Simulator is Deterministic

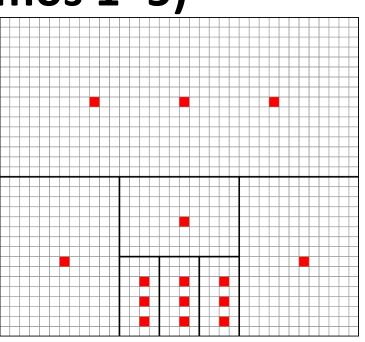
- Global random seed
- Random seeds for each rat
- Process rats in fixed order

You Must Preserve Exact Same Behavior

- Python simulator generates same result as C simulator
- Use regress.py to check
 - Only checks small cases
 - Useful sanity check
- Benchmark program compares your results to reference solution
 - Handles full-sized graphs

Fractal Graph fracZ (Demos 1–5)

Rats spread quickly within region More slowly across regions Hub nodes tend to have high counts



Base grid

W X H nodes, each with nearest neighbor connectivity

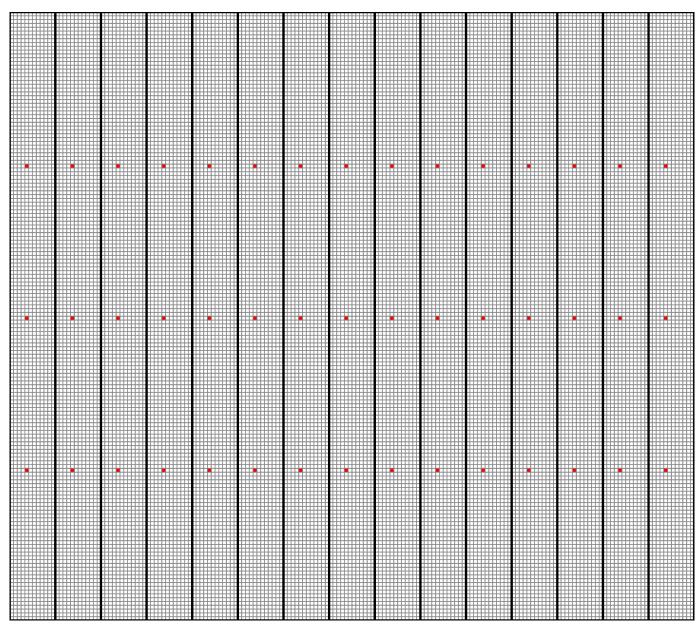
Regions

- Recursively partition rectangles into two or three subrectangles
- Leaves of tree form regions

Hubs

- Connect to every other node in region
- Each region has one or three hubs

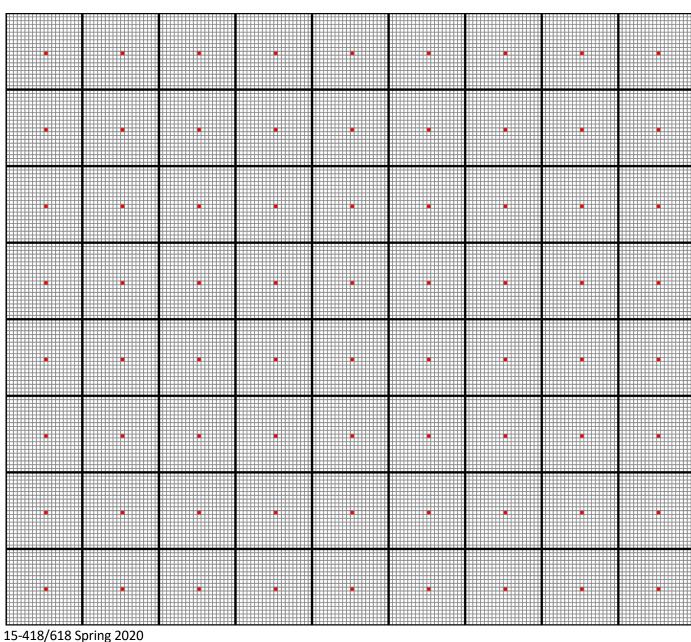
Benchmark Graph UniA (Demo 7)



Nodes	28,800
Edges	286,780
Regions	15
Hubs	45
Max Degree	1,919

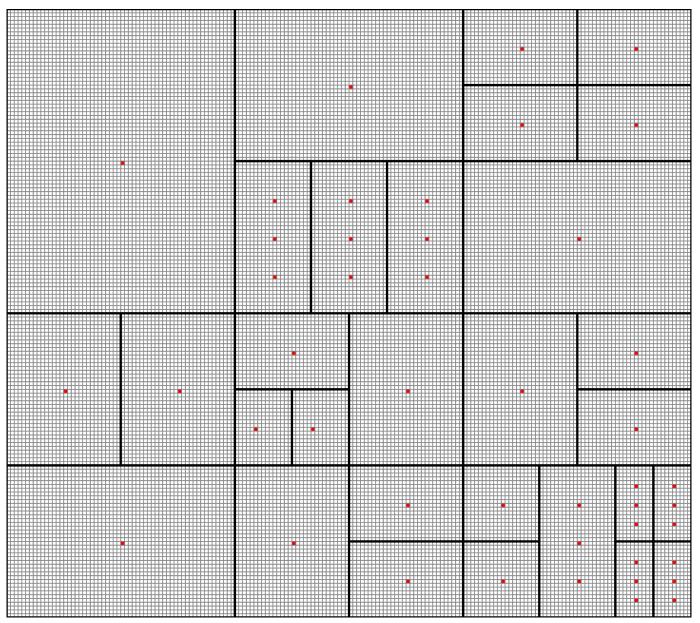
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Benchmark Graph UniB (Demo 8)



Nodes	28,800
Edges	171,400
Regions	72
Hubs	72
Max Degree	399

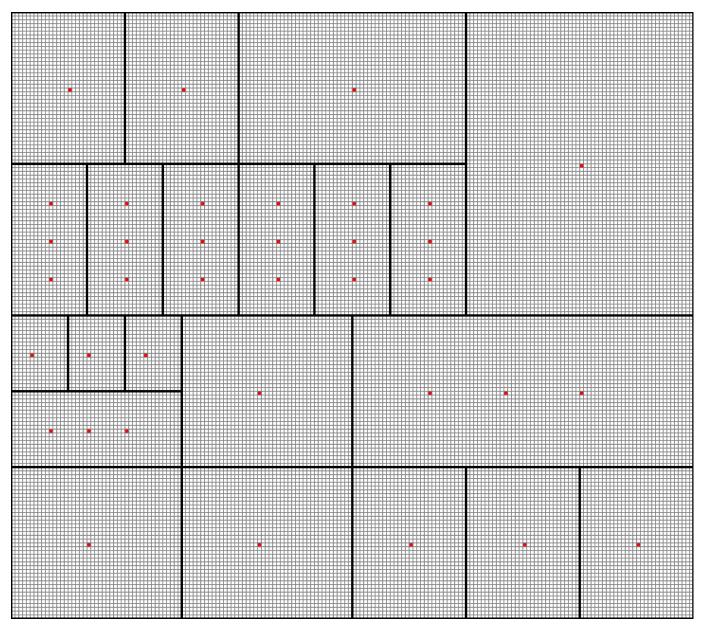
Benchmark Graph FracC (Demos 9–10)



Nodes	28,800
Edges	187,612
Regions	30
Hubs	46
Max Degree	4,899

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Benchmark Graph FracD



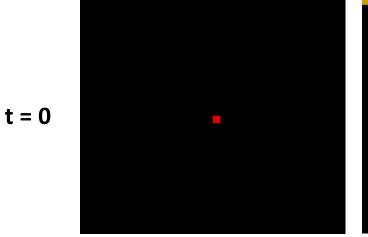
Nodes	28,800			
Edges	208,902			
Regions	21			
Hubs	37			
Max Degree	4,899			

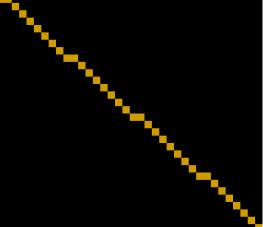
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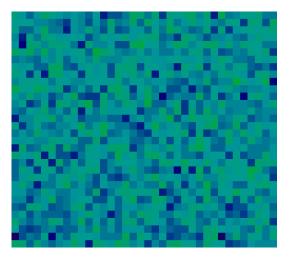
Initial States (fracZ)

Center (r) ~Demo 7 Diagonal (d) ~Demo 9

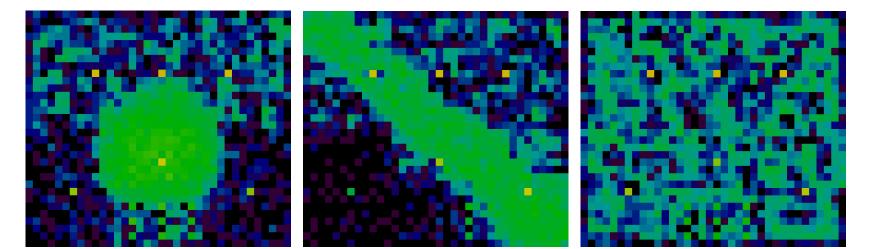
Random (u) ~Demo 10





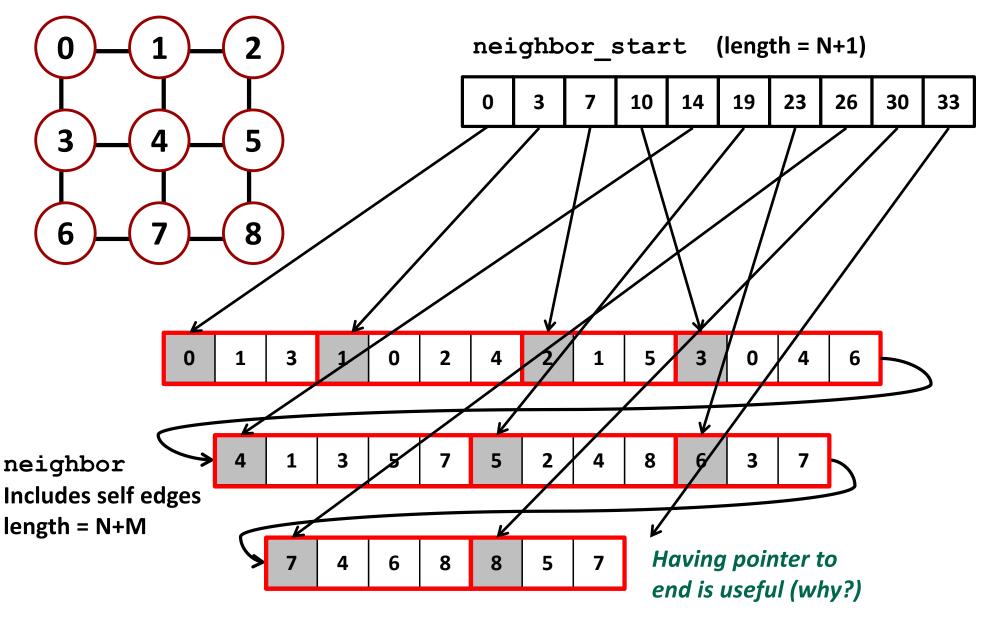


t = 20



Graph Representation

N node, M edges



- From sim.c
- Compute reward value for node

```
/* Compute weight for node nid */
static inline double compute_weight(state_t *s, int nid)
{
    int count = s->rat_count[nid];
    double ilf = neighbor_ilf(s, nid);
    return mweight((double) count/s->load_factor, ilf);
}
```

- Simulation state stored in state_t struct
- Reward function computed by mweight

- From sim.c
- Compute reward value for all nodes

```
static inline void compute_all_weights(state_t *s) {
    graph_t *g = s->g;
    double *node_weight = s->node_weight;
    int nid;
    for (nid = 0; nid < g->nnode; nid++)
        node_weight[nid] = compute_weight(s, nid);
}
```

Simulation state stored in state_t struct

- From sim.c
- Compute sum of reward values for node
- Store cumulative value for each edge
- Store total sum for later reuse

```
static inline void find_all_sums(state_t *s) {
    graph_t *g = s->g;
    int nid, eid;
    for (nid = 0; nid < g->nnode; nid++) {
        double sum = 0.0;
        for (eid = g->neighbor_start[nid];
            eid < g->neighbor_start[nid+1];
            eid++) {
            sum += s->node_weight[g->neighbor[eid]];
            s->neighbor_accum_weight[eid] = sum;
        }
        s->sum_weight[nid] = sum;
    }
}
```

Compute next move for rat

```
static inline int fast next random move(state t *s, int r) {
    int nid = s->rat position[r];
    graph t *g = s - >g;
    random t *seedp = &s->rat seed[r];
    double tsum = s->sum weight[nid];
    double val = next random float(seedp, tsum);
    int estart = g->neighbor start[nid];
    int elen = g->neighbor start[nid+1] - estart;
    /* Find location by binary search */
    int offset = locate value(val,
                               &s->neighbor accum weight[estart],
                              elen);
    return g->neighbor[estart + offset];
}
```

Instrumented Code

- From sim.c
- Wrap major sections with instrumentation macros

```
static inline void find all sums(state t *s) {
    graph t *g = s \rightarrow g;
    START ACTIVITY (ACTIVITY SUMS);
    int nid, eid;
    for (nid = 0; nid < g->nnode; nid++) {
        double sum = 0.0;
        for (eid = g->neighbor start[nid];
             eid < g->neighbor start[nid+1];
             eid++) {
            sum += s->node weight[g->neighbor[eid]];
            s->neighbor_accum_weight[eid] = sum;
        s->sum weight[nid] = sum;
    FINISH ACTIVITY (ACTIVITY SUMS);
```

Running Instrumented Code

Demo 11

```
./crun-seq -q data/q-180x160-fracC.qph
   -r data/r-180x160-r35.rats -u b -n 50 -I -q
50 steps, 1008000 rats, 13.770 seconds
        228 ms 1.6 %
                          startup
      10677 ms 76.3 %
                          compute weights
        750 ms 5.4 %
                          compute sums
                          find moves
       2340 ms 16.7 %
         2 ms 0.0 %
                          unknown
      13998 ms
                100.0 %
                          elapsed
```

- Shows breakdown of where time spent
- See speedups of different parts of code
- Can instrument both your code & reference version

Finding Parallelism

Sequential constraints

- Must complete time steps sequentially
- Must complete each batch before starting next
 - ILF values and weights then need to be recomputed

Sources of parallelism

- Over nodes
 - Computing ILFs and reward functions
 - Computing cumulative sums
- Over rats (within a batch)
 - Computing next moves
 - Updating node counts

Performance Measurements

Nanoseconds per move (NPM)

- R rats running for S steps
- Requires time T seconds
- NPM = 10⁹ * T / (R * S)
- Reference solution:
 - Average 290 NPM for 1 thread
 - Average 44.6 NPM for 12 threads
 - Speedups:
 - UniA 6.78
 - UniB 6.43
 - FracC 6.45
 - FracD 6.55
- Maybe you can do better!

Performance Targets

Benchmarks

- 4 combinations of graph/initial state
- Each counts 16 points

Target performance

- T = measured time
- T_r = time for reference solution
- T_r / T = How well you reach reference solution performance
 - Full credit when ≥ 0.95
 - Partial when ≥ 0.60

Machines

Latedays cluster

- 16 worker nodes + 1 head node
- Each is 12-core Xeon processor (dual socket with 6 cores each)
- You submit jobs to batch queue
- Assigned single processor for entire run
- Python script provided

Code Development

- OK to do code development and testing on other machines
- But, they have different performance characteristics
- Max threads on GHC cluster = 8
- Code should run for any number of threads (up to machine limit)

Some Logos



GraphChi: Going small with GraphLab



