

15-414 — Bug Catching — Fall 2006

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Assignment 2

Due date: Tuesday, November 14, 2006

1 CTL

This is a problem given in Huth and Ryan's book. Consider the model \mathcal{M} in Figure 1. Check whether $\mathcal{M}, s_0 \models \phi$ and $\mathcal{M}, s_2 \models \phi$ hold for the CTL formula ϕ :

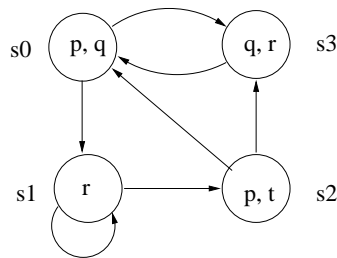


Figure 1: A model with four states

- (a) $\mathbf{AF}q$
- (b) $\mathbf{AG}(\mathbf{EF}(p \vee r))$
- (c) $\mathbf{EX}(\mathbf{EX}r)$
- (d) $\mathbf{AG}(\mathbf{AF}q)$
- (e) Give the SMV code to check the above properties with the initial state as s_2 . Also double check your answers to the previous parts by running a model checker on the SMV code.

2 More CTL

Consider the model \mathcal{M} in Figure 2. Check whether $\mathcal{M}, s_0 \models \phi$ hold for the CTL formula ϕ :

- (a) $\mathbf{AG}(\mathbf{AF}c)$
- (b) $\mathbf{AG}(\mathbf{AF}(\mathbf{AX}c))$
- (c) $\mathbf{AG}(\mathbf{AF}b)$
- (d) $\mathbf{AG}(\mathbf{AX}c \vee \mathbf{AX}(\mathbf{AX}(\neg c)))$
- (e) $\mathbf{EF}(\mathbf{AG}c)$

You may want to double check your answers by writing the above model in SMV. We do not need to the SMV code.

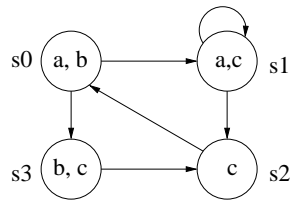


Figure 2: A model with four states

3 Writing specifications in CTL

Express the following properties in CTL. You can use atomic propositions such as *req*, *ack*, *restart* to write the properties.

- A request is eventually acknowledged.
- A request is acknowledged in atmost three cycles (transitions).
- Infinitely many requests occur during the operation of the system.
- From any state it is possible to get to a restart state.
- From any state it is possible to get to a restart state in exactly two steps.
- After p , q is never true.

4 SMV to state machines

In this problem we will use the INIT, TRANS construct from SMV language. Consider the SMV program in Figure 3.

```

MODULE main

VAR b1: boolean;
VAR b2: boolean;
VAR b3: boolean;

INIT (b1 & b2 & b3)

TRANS
  (b1 & !b2 & !b3 & next(b1) & !next(b2) & !next(b3)) |
  (b1 & b2 & !b3 & !next(b1) & !next(b2)) |
  (b1 & b2 & b3 & next(b1) & next(b3)) |
  (b1 & !b2 & next(b1) & !next(b2) & next(b3)) |
  (!b1 & !b2 & !next(b1) & !next(b2)) |
  (b1 & b3 & next(b1) & !next(b2) & !next(b3)) |
  (!b1 & b2 & !b3 & !next(b1) & next(b2) & !next(b3)) |
  (!b1 & b2 & b3 & !next(b1) & next(b2) & !next(b3))

```

Figure 3: A SMV program.

- (a) Draw the state transition graph represented by the above SMV program? Mark the initial states.

Answer whether the program satisfies the following specifications using the state transition graph.

- (b) **AG (b1)**.
(c) **AF (!b2)**.
(d) **AG(EF (b2))**.
(e) **EF(!b1 & b2)**.

Check your answers by using NuSMV or Cadence SMV. The above program is present in a SMV file available separately.