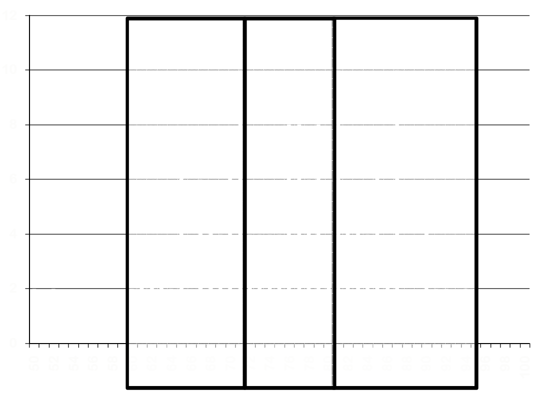
15-251

Some ~~Great~~ Theoretical Ideas ~~in~~ for Computer Science



Internet Memes **Luis von Ahn**

Collaborative Filtering Digg Network Dynamics

15-396 A TTh 3:00-4:20pm

Social Network Theory Web Spam

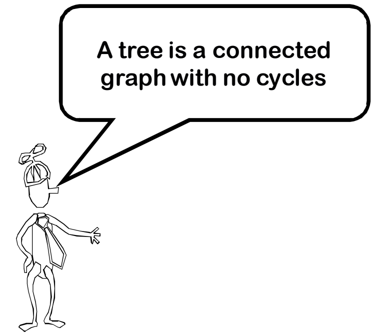
PageRank Recommender Systems

Science of the Web

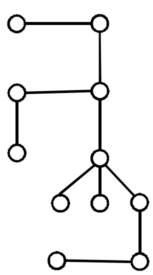
Graphs

Lecture 20 (March 27, 2008)

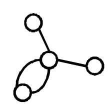
What's a tree?



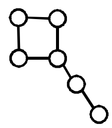
Tree



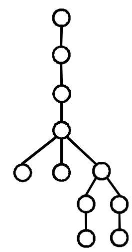
Not Tree



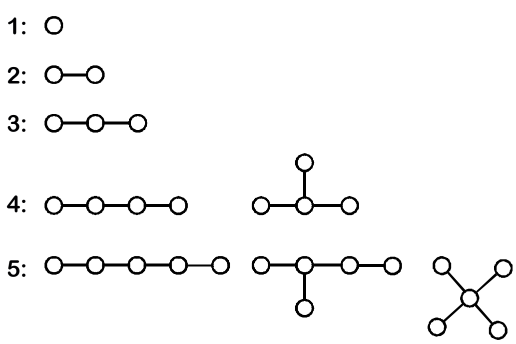
Not Tree



Tree



How Many n-Node Trees?



Notation

In this lecture:

n will denote the number of nodes in a graph

e will denote the number of edges in a graph

Theorem: Let G be a graph with n nodes and e edges

The following are equivalent:

1. G is a tree (connected, acyclic)
2. Every two nodes of G are joined by a unique path
3. G is connected and n = e + 1
4. G is acyclic and n = e + 1
5. G is acyclic and if any two non-adjacent points are joined by a line, the resulting graph has exactly one cycle

To prove this, it suffices to show

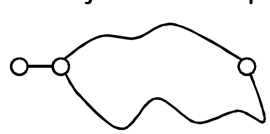
$$1 \Rightarrow 2 \Rightarrow 3 \Rightarrow 4 \Rightarrow 5 \Rightarrow 1$$

$1 \Rightarrow 2$

1. G is a tree (connected, acyclic)
2. Every two nodes of G are joined by a unique path

Proof: (by contradiction)

Assume G is a tree that has two nodes connected by two different paths:



Then there exists a cycle!

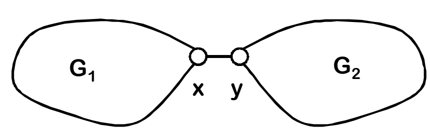
$2 \Rightarrow 3$

2. Every two nodes of G are joined by a unique path
3. G is connected and n = e + 1

Proof: (by induction)

Assume true for every graph with < n nodes

Let G have n nodes and let x and y be adjacent



Let n_1, e_1 be number of nodes and edges in G_1

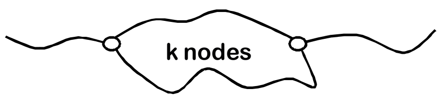
Then $n = n_1 + n_2 = e_1 + e_2 + 2 = e + 1$

$3 \Rightarrow 4$

3. G is connected and n = e + 1
4. G is acyclic and n = e + 1

Proof: (by contradiction)

Assume G is connected with n = e + 1, and G has a cycle containing k nodes



Note that the cycle has k nodes and k edges

Start adding nodes and edges until you cover the whole graph

Number of edges in the graph will be at least n

Corollary: Every nontrivial tree has at least two endpoints (points of degree 1)

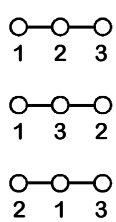
Proof (by contradiction):

Assume all but one of the points in the tree have degree at least 2

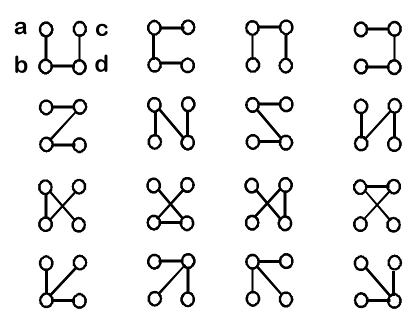
In any graph, sum of the degrees = 2e

Then the total number of edges in the tree is at least (2n-1)/2 = n - 1/2 > n - 1

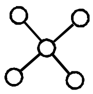
How many labeled trees are there with three nodes?



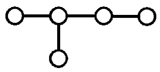
How many labeled trees are there with four nodes?



How many labeled trees are there with five nodes?

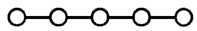


5
labelings



5 x 4 x 3

labelings



5!/ 2

labelings

125 labeled trees

How many labeled trees are there with n nodes?

3 labeled trees with 3 nodes

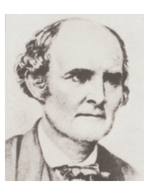
16 labeled trees with 4 nodes

125 labeled trees with 5 nodes

n^{n-2} labeled trees with n nodes

Cayley's Formula

The number of labeled trees on n nodes is n^{n-2}



The proof will use the correspondence principle

Each labeled tree on n nodes

corresponds to

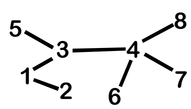
A sequence in $\{1,2,\ldots,n\}^{n-2}$ (that is, n-2 numbers, each in the range [1..n])

How to make a sequence from a tree?

Loop through i from 1 to n-2

- Let L be the degree-1 node with the lowest label
- Define the i^{th} element of the sequence as the label of the node adjacent to L
- Delete the node L from the tree

Example:



1 3 3 4 4 4

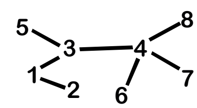
How to reconstruct the unique tree from a sequence S:

Let I = {1, 2, 3, …, n}

Loop until S is empty

- Let i = smallest # in I but not in S
- Let s = first label in sequence S
  - Add edge {i, s} to the tree
  - Delete i from I
  - Delete s from S

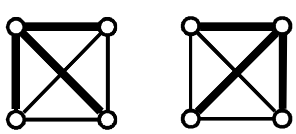
Add edge {a,b}, where I = {a,b}



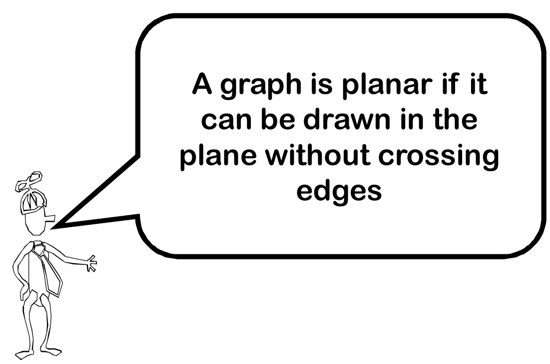
1 3 3 4 4 4

Spanning Trees

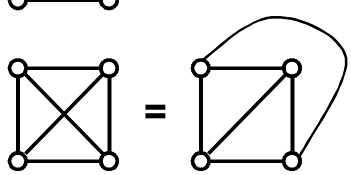
A spanning tree of a graph G is a tree that touches every node of G and uses only edges from G



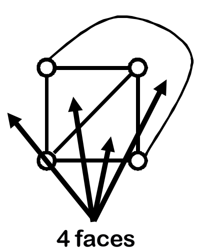
Every connected graph has a spanning tree



Examples of Planar Graphs



http://www.planarity.net

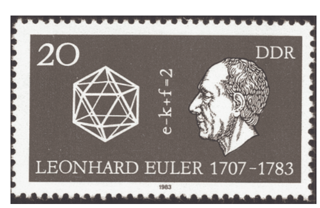


Faces

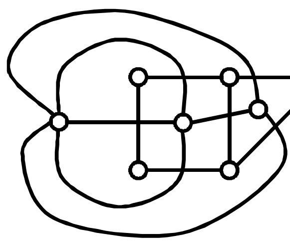
A planar graph splits the plane into disjoint faces

Euler's Formula

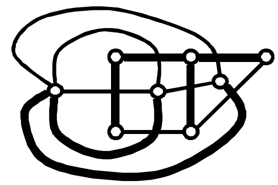
If G is a connected planar graph with n vertices, e edges and f faces, then $n - e + f = 2$



Rather than using induction, we'll use the important notion of the dual graph



Dual = put a node in every face, and an edge for each edge joining two adjacent faces



Let G* be the dual graph of G

Let T be a spanning tree of G

Let T* be the graph where there is an edge in dual graph for each edge in G – T

Then T* is a spanning tree for G*

$n = e_T + 1$ $\quad$ $n + f = e_T + e_{T^*} + 2$

$f = e_{T^*} + 1$ $\quad$ $= e + 2$

Corollary: Let G be a simple planar graph with n > 2 vertices. Then:

1. G has a vertex of degree at most 5
2. G has at most 3n – 6 edges

Proof of 1:

In any graph, (sum of degrees) = 2e

Assume all vertices have degree ≥ 6

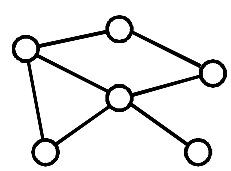
Then $e \geq 3n$

Furthermore, since G is simple, $3f \leq 2e$

So $3n + 3f \leq 3e \Rightarrow 3(n-e+f) \leq 0$, contradict.

Graph Coloring

A coloring of a graph is an assignment of a color to each vertex such that no neighboring vertices have the same color



Graph Coloring

Arises surprisingly often in CS

Register allocation: assign temporary variables to registers for scheduling instructions. Variables that interfere, or are simultaneously active, cannot be assigned to the same register

Theorem: Every planar graph can be 6-colored

Proof Sketch (by induction):

Assume every planar graph with less than n vertices can be 6-colored

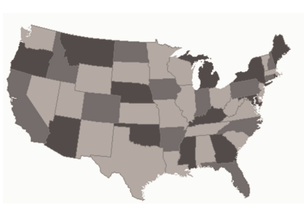
Assume G has n vertices

Since G is planar, it has some node v with degree at most 5

Remove v and color by Induction Hypothesis

Not too difficult to give an inductive proof of 5-colorability, using same fact that some vertex has degree ≤ 5

4-color theorem remains challenging!



Implementing Graphs

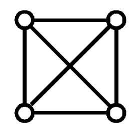
Adjacency Matrix

Suppose we have a graph G with n vertices. The adjacency matrix is the n x n matrix $A=[a_{ij}]$ with:

$a_{ij} = 1$ if (i,j) is an edge
$a_{ij} = 0$ if (i,j) is not an edge

Good for dense graphs!

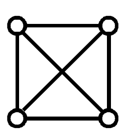
Example



$$A = \begin{pmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

Counting Paths

The number of paths of length k from node i to node j is the entry in position (i,j) in the matrix A^k



$$A^2 = \begin{pmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}\begin{pmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}$$

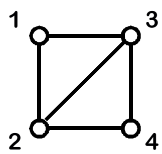
$$= \begin{pmatrix} 3 & 2 & 2 & 2 \\ 2 & 3 & 2 & 2 \\ 2 & 2 & 3 & 2 \\ 2 & 2 & 2 & 3 \end{pmatrix}$$

Adjacency List

Suppose we have a graph G with n vertices. The adjacency list is the list that contains all the nodes that each node is adjacent to

Good for sparse graphs!

Example



1: 2,3
2: 1,3,4
3: 1,2,4
4: 2,3

Here's What You Need to Know…

Trees
- Counting Trees
- Different Characterizations

Planar Graphs
- Definition
- Euler's Theorem
- Coloring Planar Graphs

Adjacency Matrix and List
- Definition
- Useful for counting