A machine so simple that you can understand it in less than one minute

The machine accepts a string if the process ends in a double circle
The machine accepts a string if the process ends in a double circle state.

The alphabet of a finite automaton is the set where the symbols come from: \{0,1\}

The language of a finite automaton is the set of strings that it accepts.

The Language of Machine M

L(M) = \{ w | w has an even number of 1s\}
**Notation**

An alphabet $\Sigma$ is a finite set (e.g., $\Sigma = \{0, 1\}$)

A string over $\Sigma$ is a finite-length sequence of elements of $\Sigma$

For $x$ a string, $|x|$ is the length of $x$

The unique string of length 0 will be denoted by $\varepsilon$ and will be called the empty or null string

A language over $\Sigma$ is a set of strings over $\Sigma$

A finite automaton is a 5-tuple $M = (Q, \Sigma, \delta, q_0, F)$

- $Q$ is the set of states
- $\Sigma$ is the alphabet
- $\delta : Q \times \Sigma \rightarrow Q$ is the transition function
- $q_0 \in Q$ is the start state
- $F \subseteq Q$ is the set of accept states

$L(M) =$ the language of machine $M$

$= \text{set of all strings machine } M \text{ accepts}$

**Example Automaton**

$M = (Q, \Sigma, \delta, q_0, F)$ where $Q = \{q_0, q_1, q_2, q_3\}$

$\Sigma = \{0, 1\}$

$\delta : Q \times \Sigma \rightarrow Q$ transition function

$q_0 \in Q$ is start state

$F = \{q_0, q_1\} \subseteq Q$ accept states

Build an automaton that accepts all and only those strings that contain 001
Build an automaton that accepts all strings whose length is divisible by 2 but not 3

A language is regular if it is recognized by a deterministic finite automaton

L = \{ w \mid w \text{ contains 001} \} is regular
L = \{ w \mid w \text{ has an even number of 1s} \} is regular

**Union Theorem**

Given two languages, \( L_1 \) and \( L_2 \), define the union of \( L_1 \) and \( L_2 \) as
\[
L_1 \cup L_2 = \{ w \mid w \in L_1 \text{ or } w \in L_2 \}
\]

Theorem: The union of two regular languages is also a regular language

**Proof Sketch:** Let \( M_1 = (Q_1, \Sigma, \delta_1, q_0^1, F_1) \) be finite automaton for \( L_1 \) and \( M_2 = (Q_2, \Sigma, \delta_2, q_0^2, F_2) \) be finite automaton for \( L_2 \). We want to construct a finite automaton \( M = (Q, \Sigma, \delta, q_0, F) \) that recognizes \( L = L_1 \cup L_2 \)
Idea: Run both $M_1$ and $M_2$ at the same time!

$Q$ = pairs of states, one from $M_1$ and one from $M_2$
$= \{ (q_1, q_2) | q_1 \in Q_1 \text{ and } q_2 \in Q_2 \}$
$= Q_1 \times Q_2$

Theorem: The union of two regular languages is also a regular language

Automaton for Union

Automaton for Intersection
Theorem: The union of two regular languages is also a regular language

Corollary: Any finite language is regular

The Regular Operations

Union: $A \cup B = \{ w \mid w \in A \text{ or } w \in B \}$

Intersection: $A \cap B = \{ w \mid w \in A \text{ and } w \in B \}$

Reverse: $A^R = \{ w_1 \ldots w_k \mid w_k \ldots w_1 \in A \}$

Negation: $\neg A = \{ w \mid w \notin A \}$

Concatenation: $A \cdot B = \{ vw \mid v \in A \text{ and } w \in B \}$

Star: $A^* = \{ w_1 \ldots w_k \mid k \geq 0 \text{ and each } w_i \in A \}$

Regular Languages Are Closed Under The Regular Operations

We have seen part of the proof for Union. The proof for intersection is very similar. The proof for negation is easy.

The “Grep” Problem

Input: Text $T$ of length $t$, string $S$ of length $n$

Problem: Does string $S$ appear inside text $T$?

Naive method:

$\boxed{a_1, a_2, a_3, a_4, a_5, \ldots, a_t}$

Cost: Roughly $nt$ comparisons
Automata Solution

Build a machine $M$ that accepts any string with $S$ as a consecutive substring
Feed the text to $M$
Cost: $t$ comparisons + time to build $M$
As luck would have it, the Knuth, Morris, Pratt algorithm builds $M$ quickly

Real-life Uses of DFAs

- Grep
- Coke Machines
- Thermostats (fridge)
- Elevators
- Train Track Switches
- Lexical Analyzers for Parsers

Are all languages regular?

Consider the language $L = \{ a^n b^n | n > 0 \}$
i.e., a bunch of a’s followed by an equal number of b’s
No finite automaton accepts this language
Can you prove this?
a^n b^n is not regular. No machine has enough states to keep track of the number of a's it might encounter.

That is a fairly weak argument

Consider the following example...

$L = \text{strings where the \# of occurrences of the pattern } ab \text{ is equal to the number of occurrences of the pattern } ba$

Can't be regular. No machine has enough states to keep track of the number of occurrences of ab

$M$ accepts only the strings with an equal number of ab's and ba's!
Theorem: $L = \{a^nb^n \mid n > 0 \}$ is not regular
Proof (by contradiction):
Assume that $L$ is regular
Then there exists a machine $M$ with $k$ states that accepts $L$
For each $0 \leq i \leq k$, let $S_i$ be the state $M$ is in after reading $a^i$
$\exists i, j \leq k$ such that $S_i = S_j$, but $i \neq j$
$M$ will do the same thing on $a^ib^i$ and $a^ib^i$
But a valid $M$ must reject $a^ib^i$ and accept $a^ib^i$