

15-251

**Cooking for
Computer Scientists**

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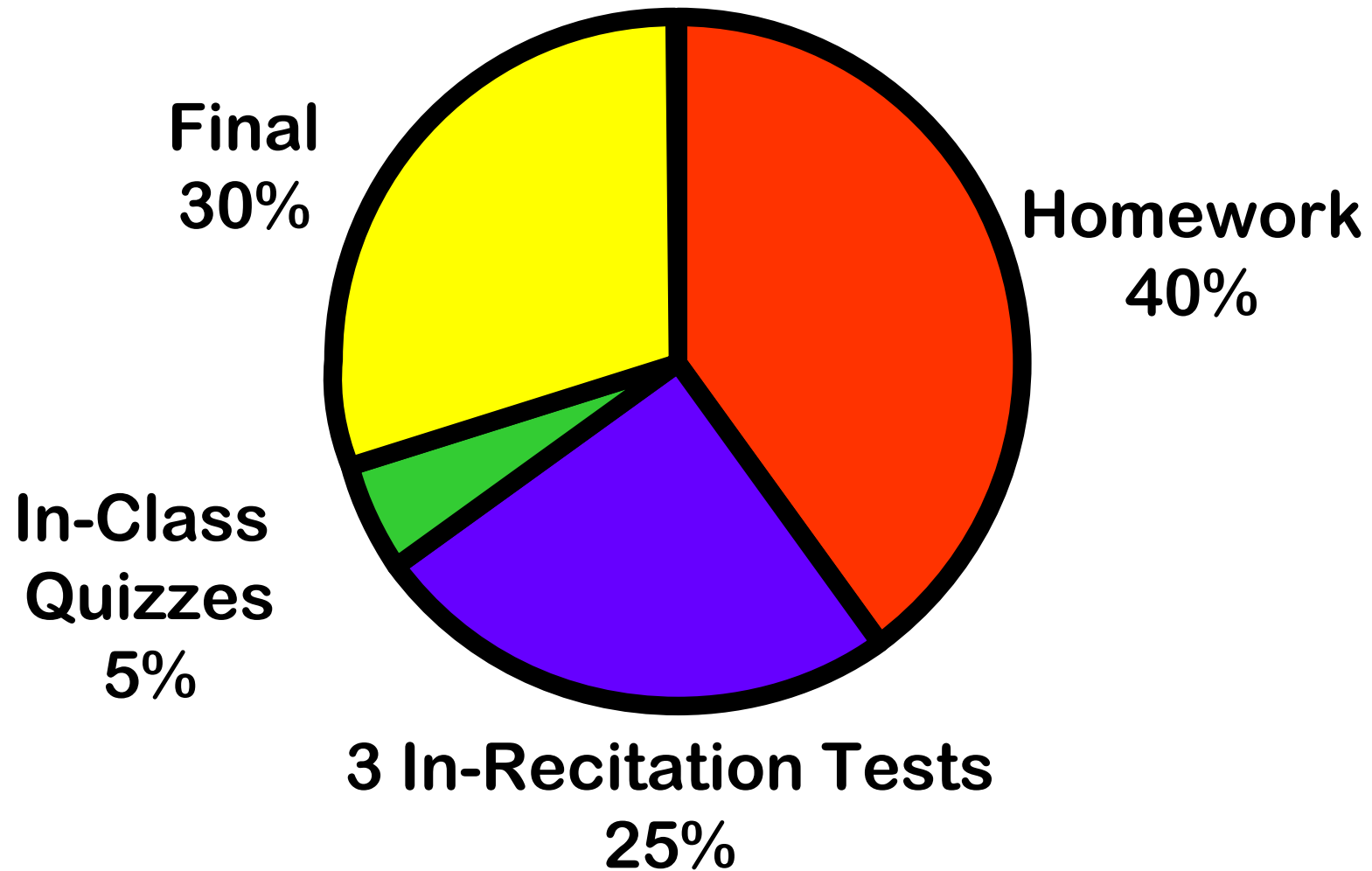
Brian Thompson

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Check this Website OFTEN!

Course Document



Course Document

Lowest homework
grade is dropped

Lowest test grade
is worth half

If Suzie gets 60,90,80 in her tests, how many
total test points will she have in her final grade?

$$(0.05)(60) + (0.10)(90) + (0.10)(80) = 20$$

Homework

Homework will go out every Tuesday
and is due the Tuesday after

Seven points per day late penalty

No homework will be accepted
more than seven days late

HOMEWORK #1 IS OUT TODAY!

Shared Secret

He was thoughtful and grave--but the
orders he gave

Were enough to bewilder a crew.

When he cried "Steer to starboard,
but keep her head larboard!"

What on earth was the helmsman to
do?

Shared Secret

H_a_h_d_g_v_-b_t
t_o

u_b_er_c.

W_c "o bo
k_p h_a arb_ar

_ha_e_t_th_l_an_o
d

Collaboration + Cheating

You may NOT share written work

You may NOT use Google, or solutions to previous years' homework

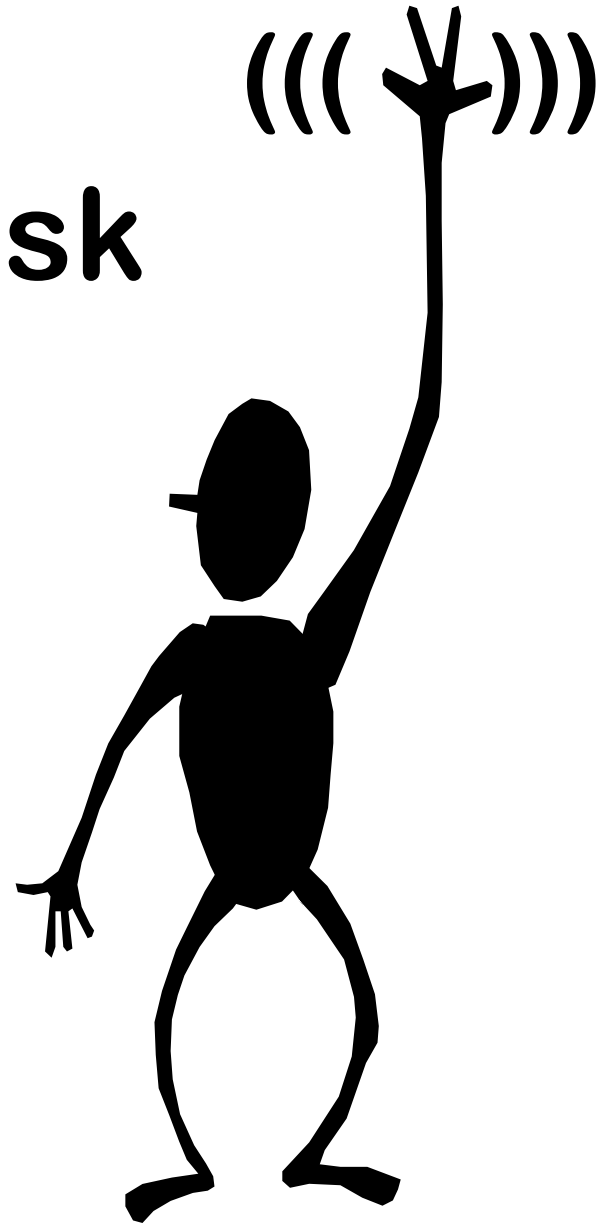
You MUST sign the class honor code

Textbook

There is NO textbook for this class

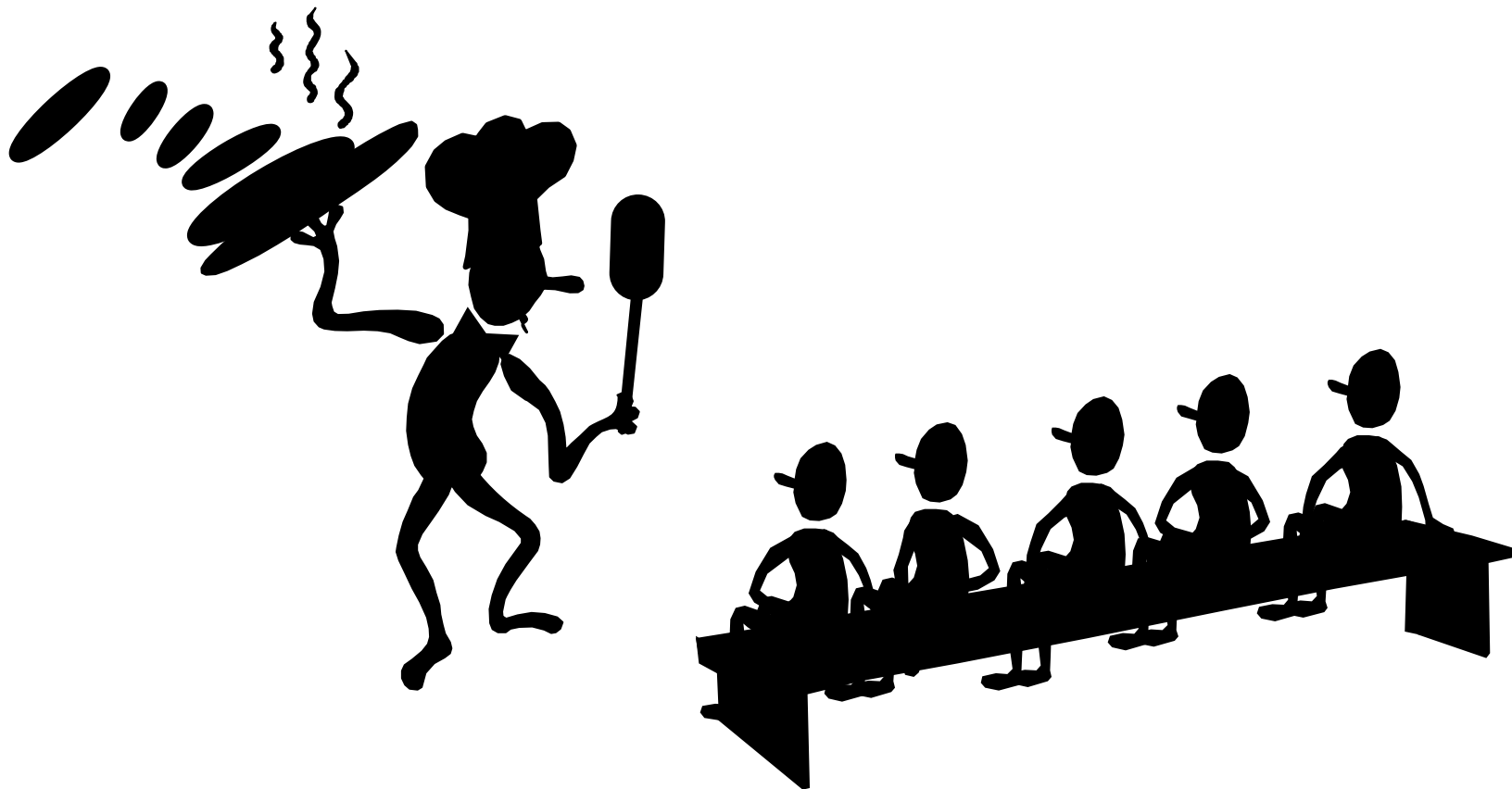
**Students from previous years tell us
that it's imperative to come to class**

**Feel free to ask
questions**



Pancakes With A Problem!

Lecture 1 (January 17, 2006)





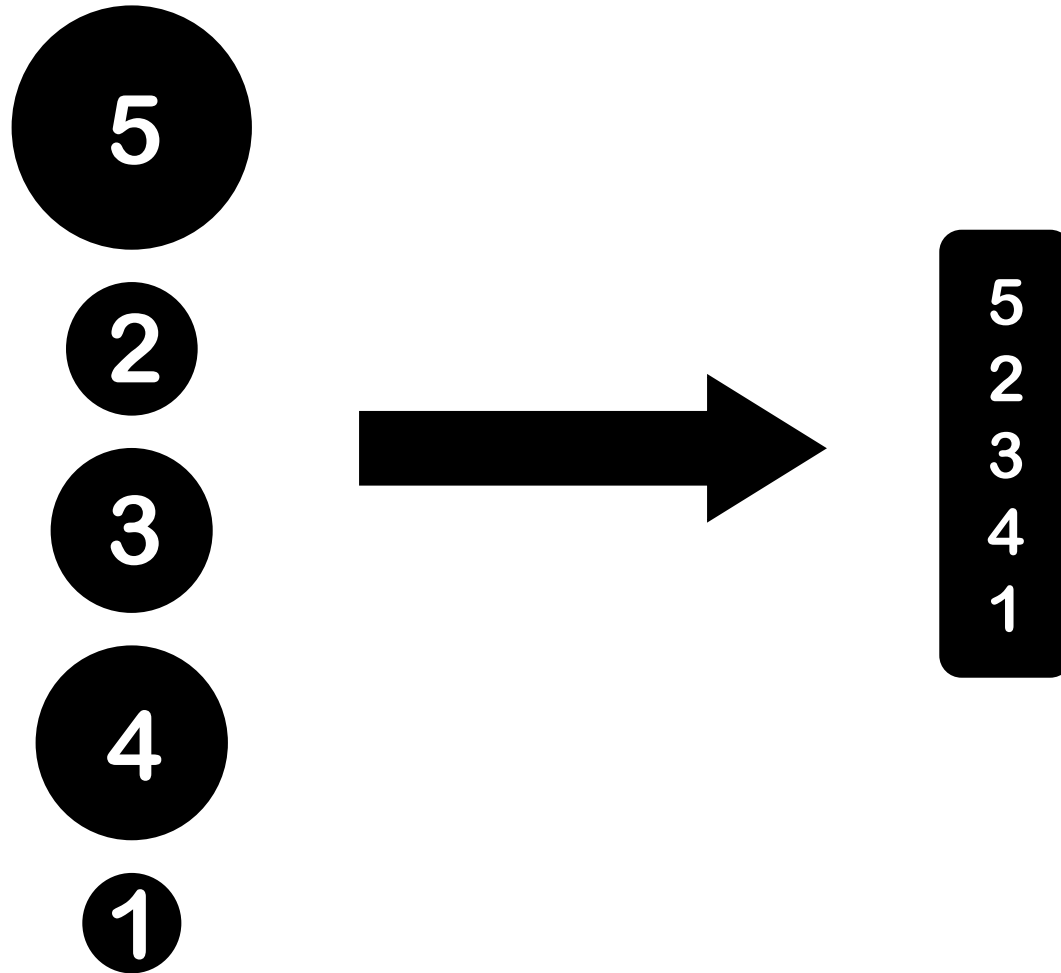
The chef at our place is sloppy: when he prepares pancakes, they come out all different sizes

When the waiter delivers them to a customer, he rearranges them (so that smallest is on top, and so on, down to the largest at the bottom)

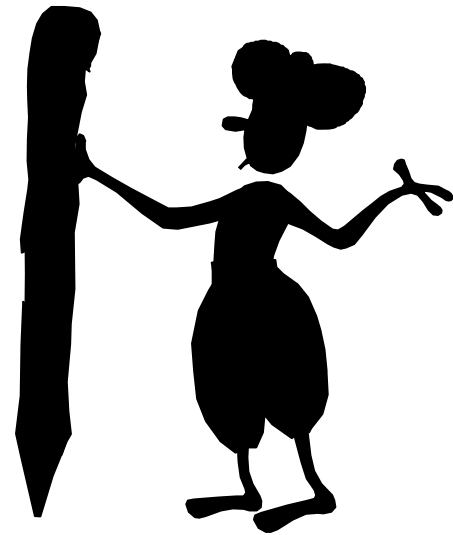
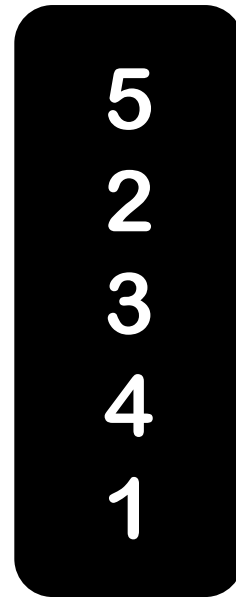
He does this by grabbing several from the top and flipping them over, repeating this (varying the number he flips) as many times as necessary



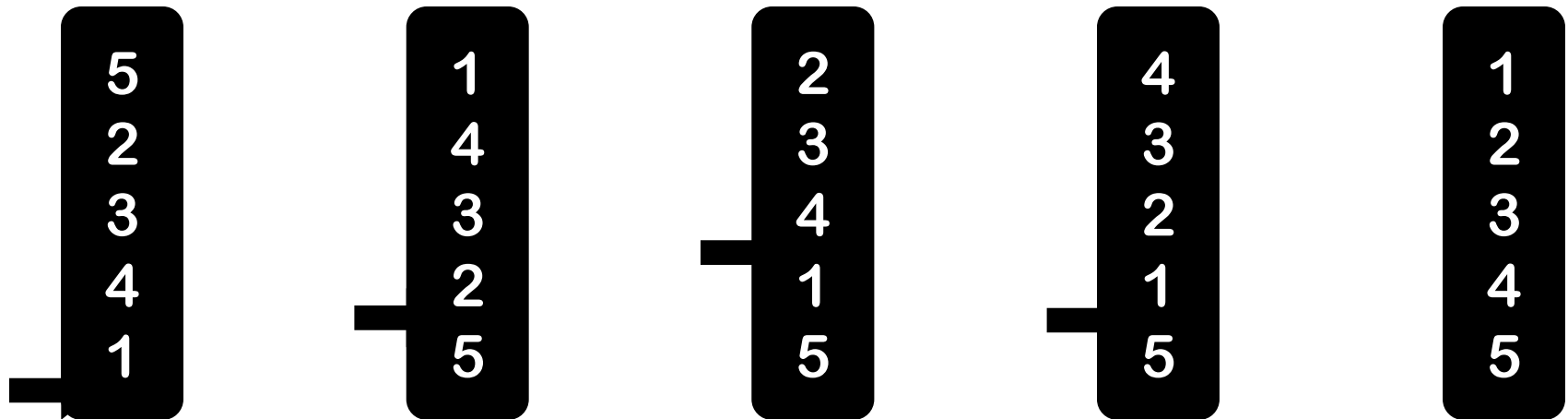
Developing A Notation: Turning pancakes into numbers



How do we sort this stack?
How many flips do we need?



4 Flips Are Sufficient



Best Way to Sort

X = Smallest number of
flips required to sort:

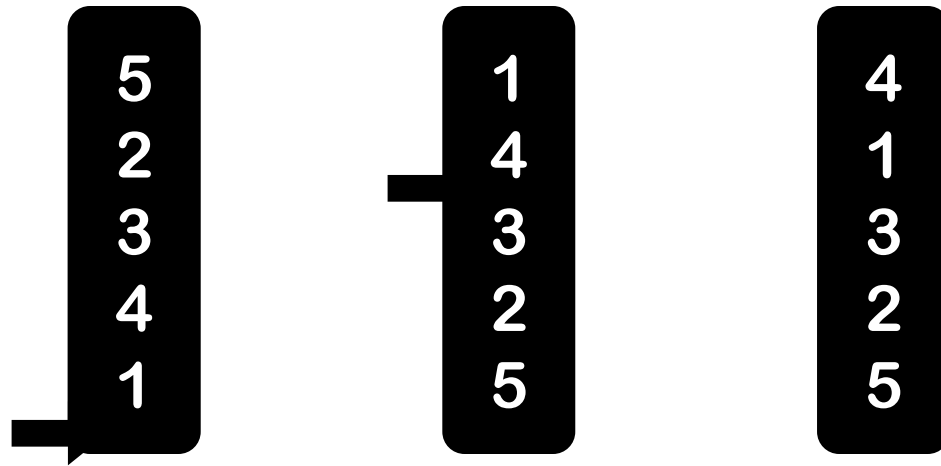
5
2
3
4
1

Lower
Bound

$? \leq X \leq ?$

Upper
Bound

Four Flips Are Necessary



Flip 1 has to put 5 on bottom

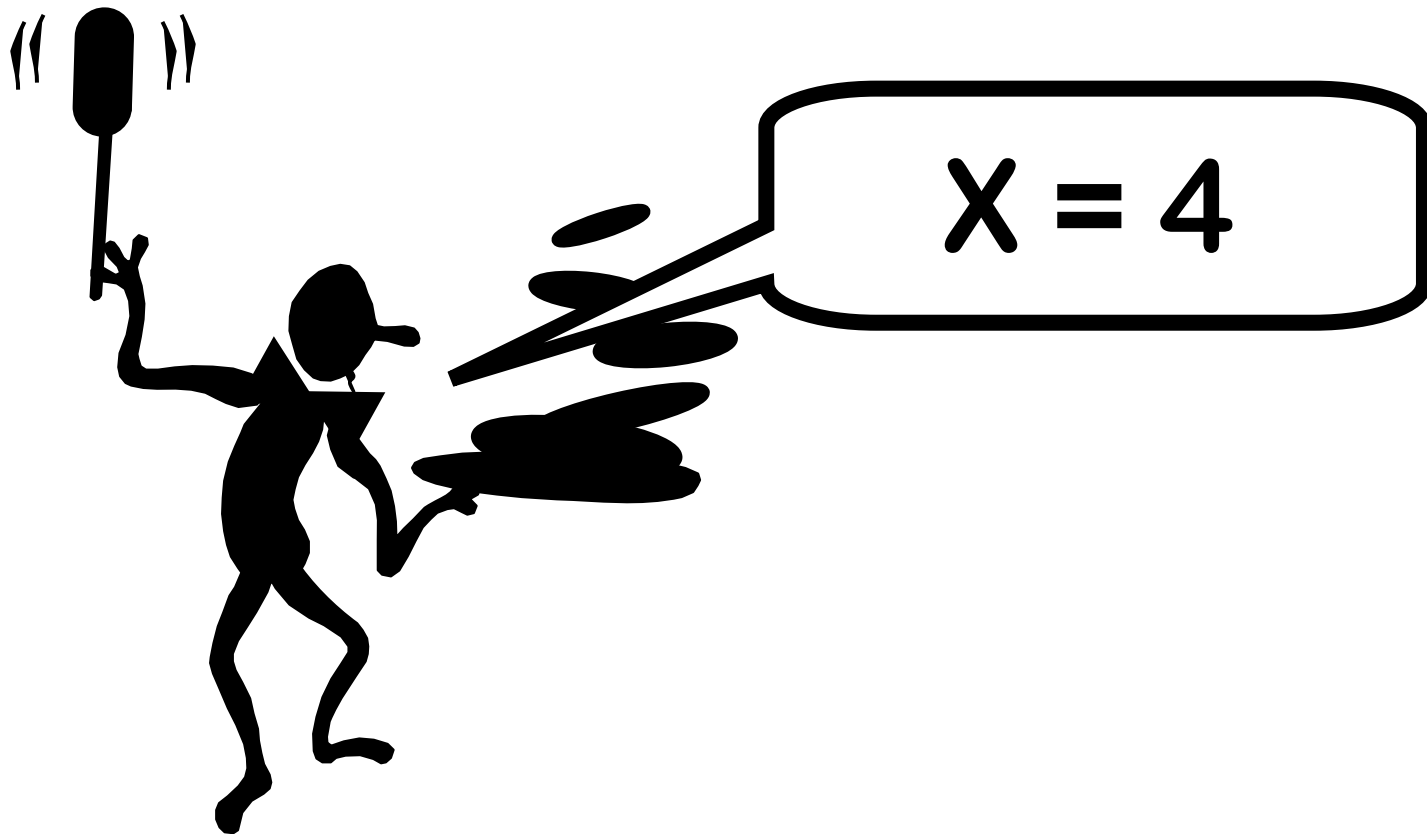
Flip 2 must bring 4 to top (if it didn't, we would spend more than 3)

$$4 \leq X \leq 4$$

Lower
Bound

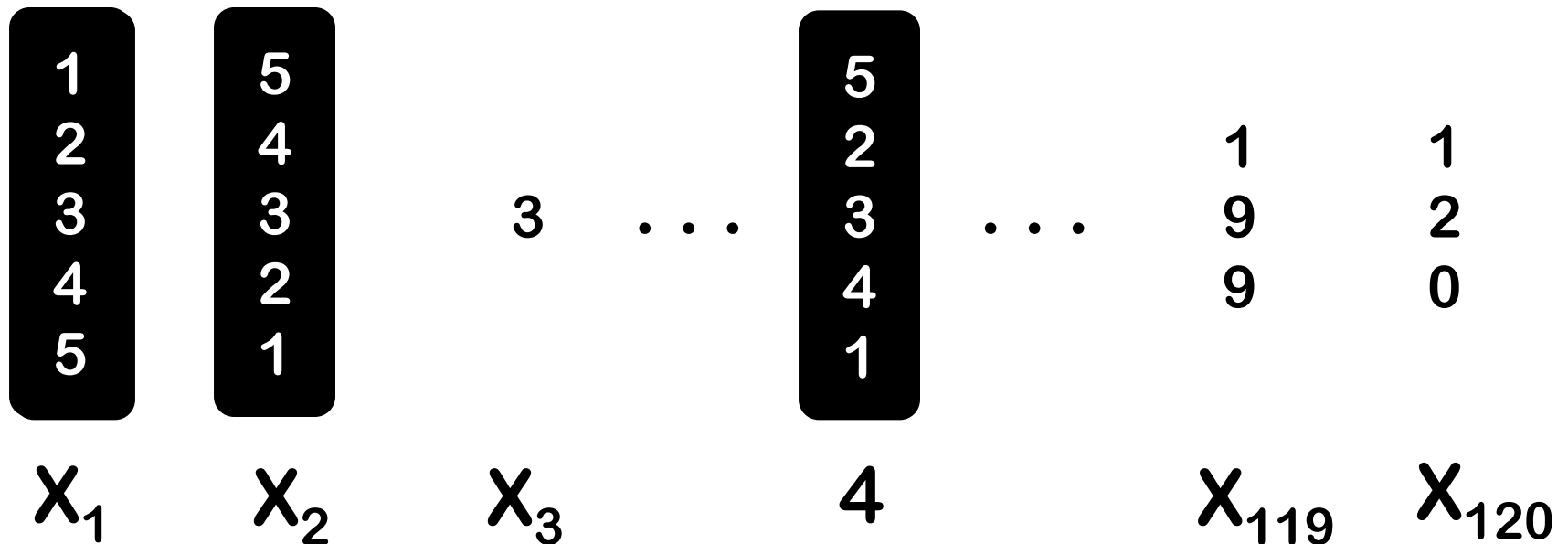


Upper
Bound

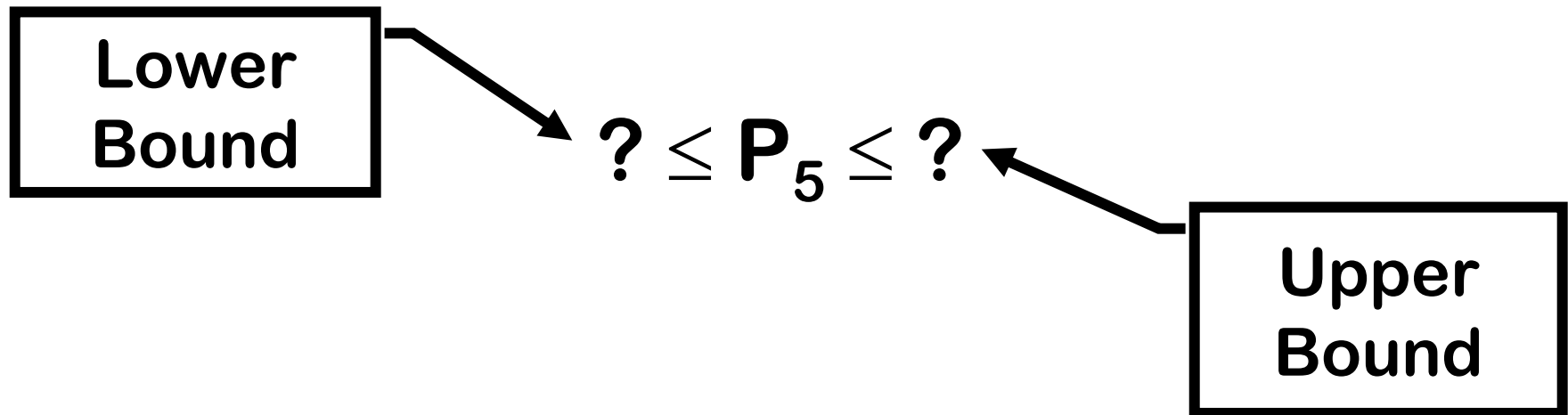


5th Pancake Number

P_5 = Number of flips required to sort the worst case stack of 5 pancakes



5th Pancake Number



$P_n = \text{MAX over } s \in \text{stacks of } n \text{ pancakes}$
of MIN # of flips to sort s

$P_n = \text{The number of flips required to sort}$
the worst-case stack of n pancakes

What is P_n for small n ?



Can you do
 $n = 0, 1, 2, 3$?

Initial Values of P_n

n	0	1	2	3
P_n	0	0	1	3

$$P_3 = 3$$

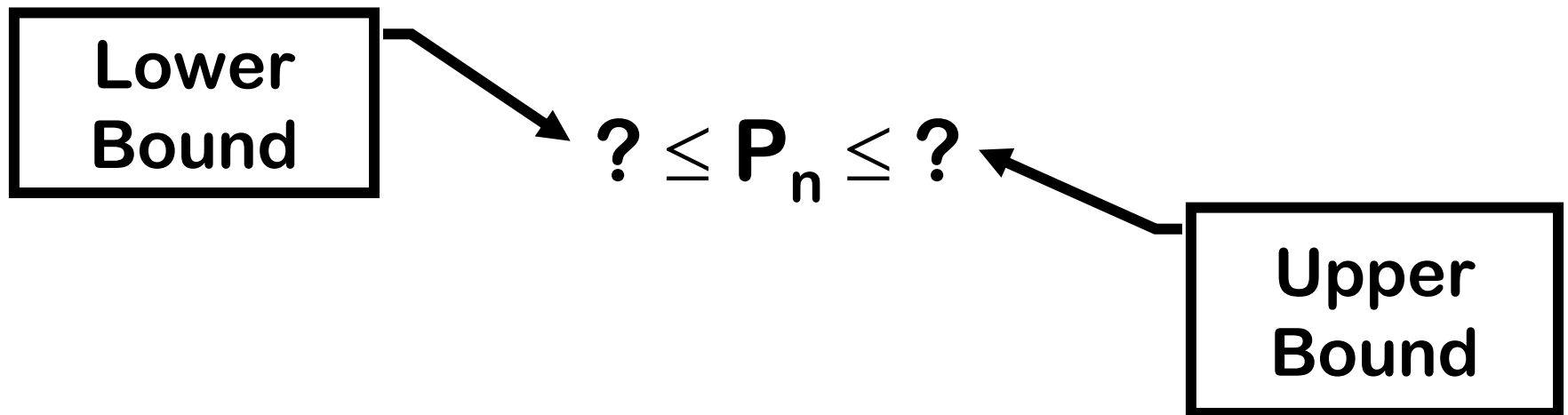
1
3
2

requires 3 Flips, hence $P_3 \geq 3$

ANY stack of 3 can be done by getting the big one to the bottom (≤ 2 flips), and then using ≤ 1 flips to handle the top two

n^{th} Pancake Number

P_n = Number of flips required to sort the worst case stack of n pancakes

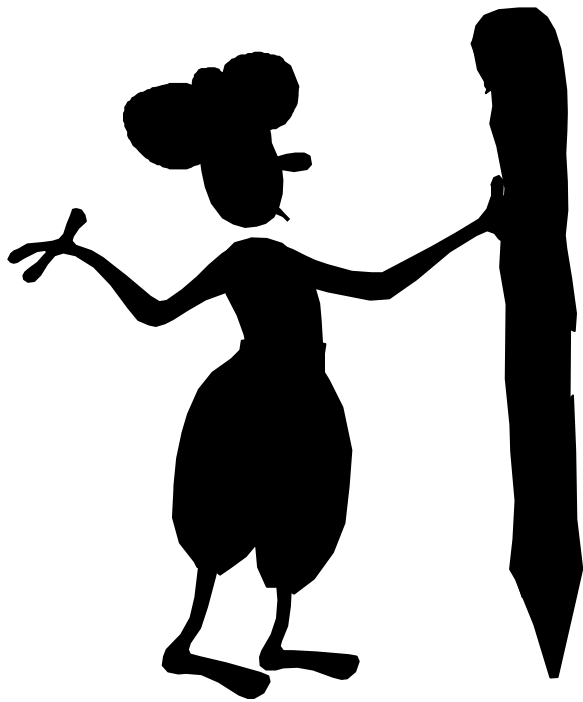


Bracketing:

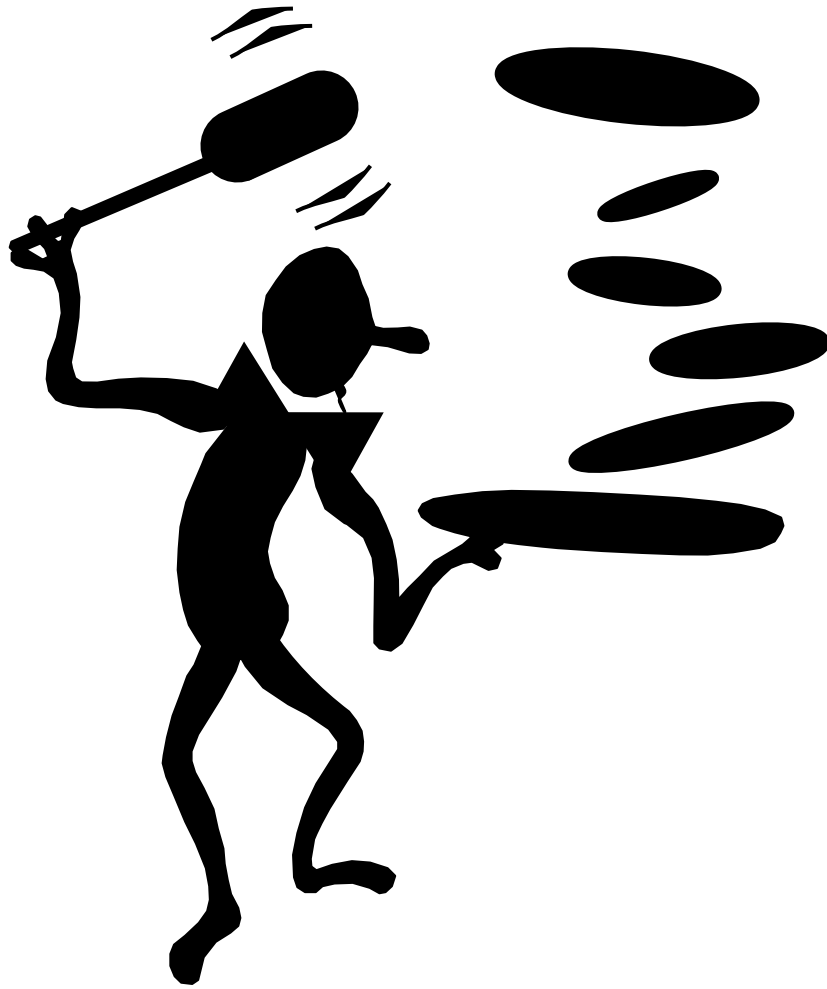
What are the best lower and upper bounds that I can prove?

$$? \leq P_n \leq ?$$

Try to find upper and
lower bounds on P_n ,
for $n > 3$



Bring-to-top Method



**Bring biggest to top
Place it on bottom**

**Bring next largest to top
Place second from
bottom**

And so on...

Upper Bound On P_n :

Bring-to-top Method For n Pancakes

If $n=1$, no work required — we are done!
Otherwise, flip pancake n to top and
then flip it to position n

Now use:

Bring To Top Method
For $n-1$ Pancakes

Total Cost: at most $2(n-1) = 2n - 2$ flips

Better Upper Bound On P_n :

Bring-to-top Method For n Pancakes

If $n=2$, at most one flip and we are done!
Otherwise, flip pancake n to top and
then flip it to position n

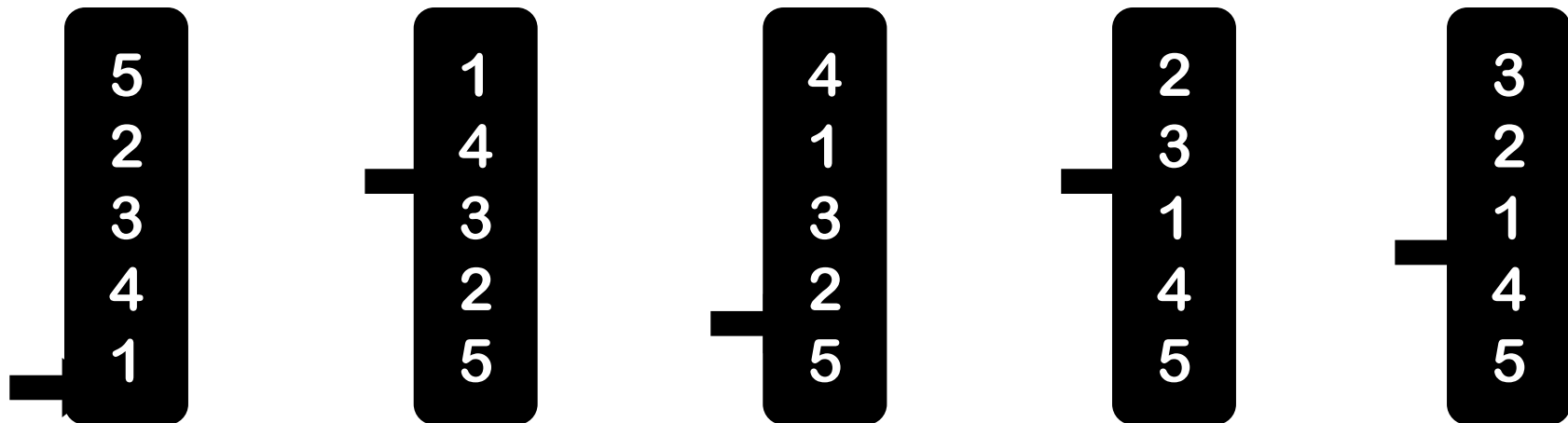
Now use:

Bring To Top Method
For $n-1$ Pancakes

Total Cost: at most $2(n-2) + 1 = 2n - 3$ flips

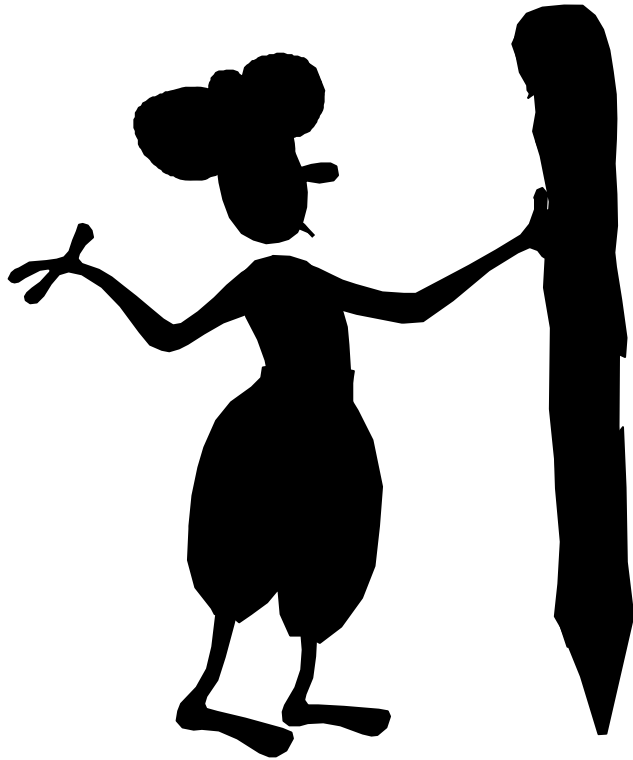
$$? \leq P_n \leq 2n-3$$

Bring-to-top not always optimal for a particular stack



Bring-to-top takes 5 flips,
but we can do in 4 flips

$$? \leq P_n \leq 2n-3$$



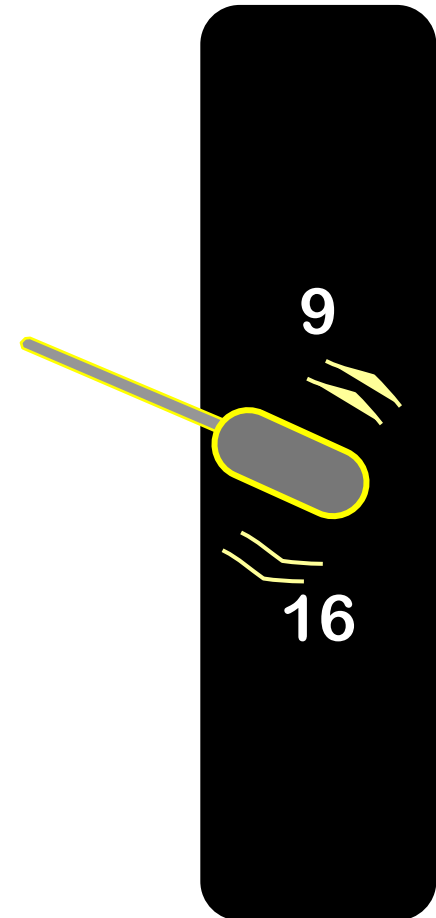
What other
bounds can you
prove on P_n ?

Breaking Apart Argument

Suppose a stack S has a pair of adjacent pancakes that will not be adjacent in the sorted stack

Any sequence of flips that sorts stack S must have one flip that inserts the spatula between that pair and breaks them apart

Furthermore, this is true of the “pair” formed by the bottom pancake of S and the plate



S

$$n \leq P_n$$

Suppose n is even

S contains n pairs that will need to be broken apart during any sequence that sorts it

2
4
6
8
.
.
.
 n
1
3
5
.
.
.
 $n-1$

Detail: This construction only works when $n > 2$

2
1

S

$$n \leq P_n$$

Suppose n is odd

S contains n pairs that will need to be broken apart during any sequence that sorts it

1
3
5
7
.
.
.
n
2
4
6
.
.
.
n-1

Detail: This
construction only
works when $n > 3$

1
3
2

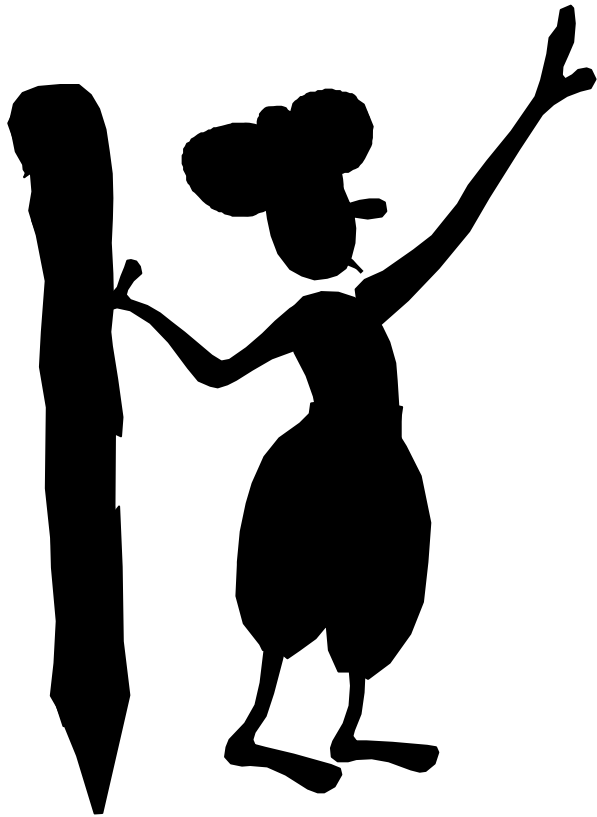
$$n \leq P_n \leq 2n - 3 \text{ for } n > 3$$



Bring-to-top is
within a factor
of 2 of optimal!

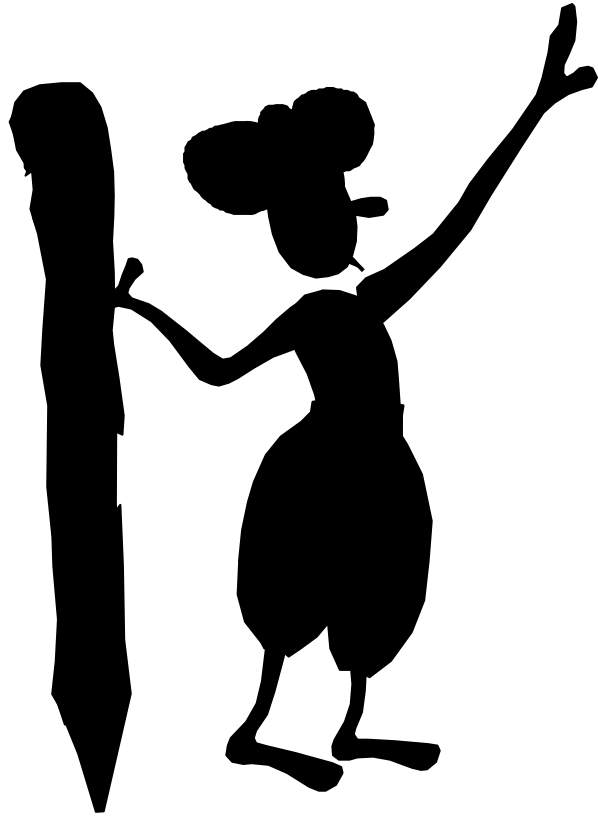
From ANY stack to sorted stack in $\leq P_n$

From sorted stack to ANY stack in $\leq P_n$?



Reverse the sequences
we use to sort

Hence, from ANY stack
to ANY stack in $\leq 2P_n$



Can you find
a faster way
than $2P_n$ flips
to go from
ANY to ANY?

ANY Stack S to ANY stack T in $\leq P_n$

S: 4,3,5,1,2

T: 5,2,4,3,1

1,2,3,4,5

3,5,1,2,4

“new T” 

Rename the pancakes in S to be 1,2,3,...,n

**Rewrite T using the new naming scheme
that you used for S**

**The sequence of flips that brings the sorted
stack to the “new T” will bring S to T**

The Known Pancake Numbers

n	P_n
1	0
2	1
3	3
4	4
5	5
6	7
7	8
8	9
9	10
10	11
11	13
12	14
13	15

P_{14} is Unknown

$1 \cdot 2 \cdot 3 \cdot 4 \cdot \dots \cdot 13 \cdot 14 = 14!$ orderings of 14 pancakes

$$14! = 87,178,291,200$$

Is This Really Computer Science?



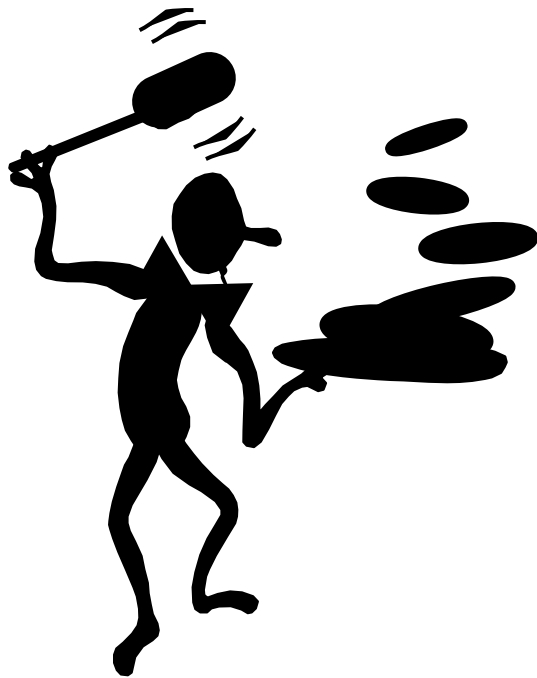


Sorting By Prefix Reversal

Posed in *Amer. Math. Monthly* 82 (1) (1975),
“Harry Dweighter” a.k.a. Jacob Goodman

$$(17/16)n \leq P_n \leq (5n+5)/3$$

William Gates and Christos Papadimitriou.
Bounds For Sorting By Prefix Reversal.
Discrete Mathematics, vol 27, pp 47-57, 1979.



$$(15/14)n \leq P_n \leq (5n+5)/3$$

H. Heydari and H. I. Sudborough. On the Diameter of the Pancake Network. *Journal of Algorithms*, vol 25, pp 67-94, 1997.



**How many different
stacks of n pancakes
are there?**

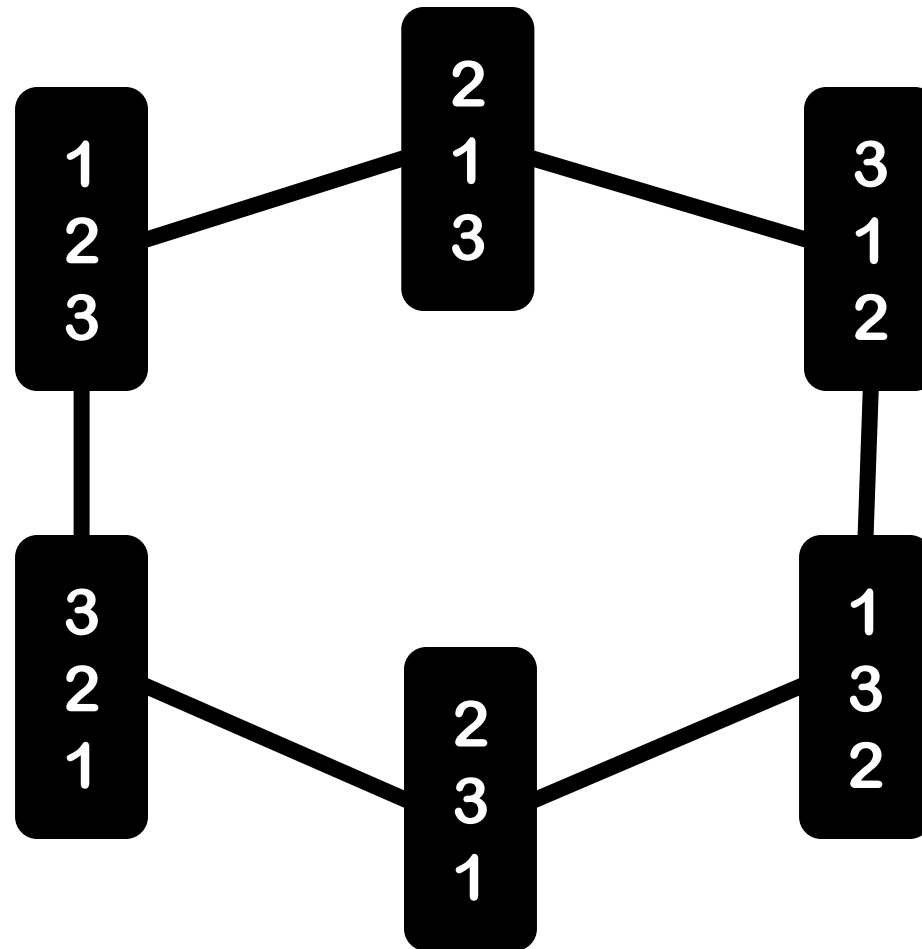
$$n! = 1 \times 2 \times 3 \times \dots \times n$$

Pancake Network: Definition For $n!$ Nodes

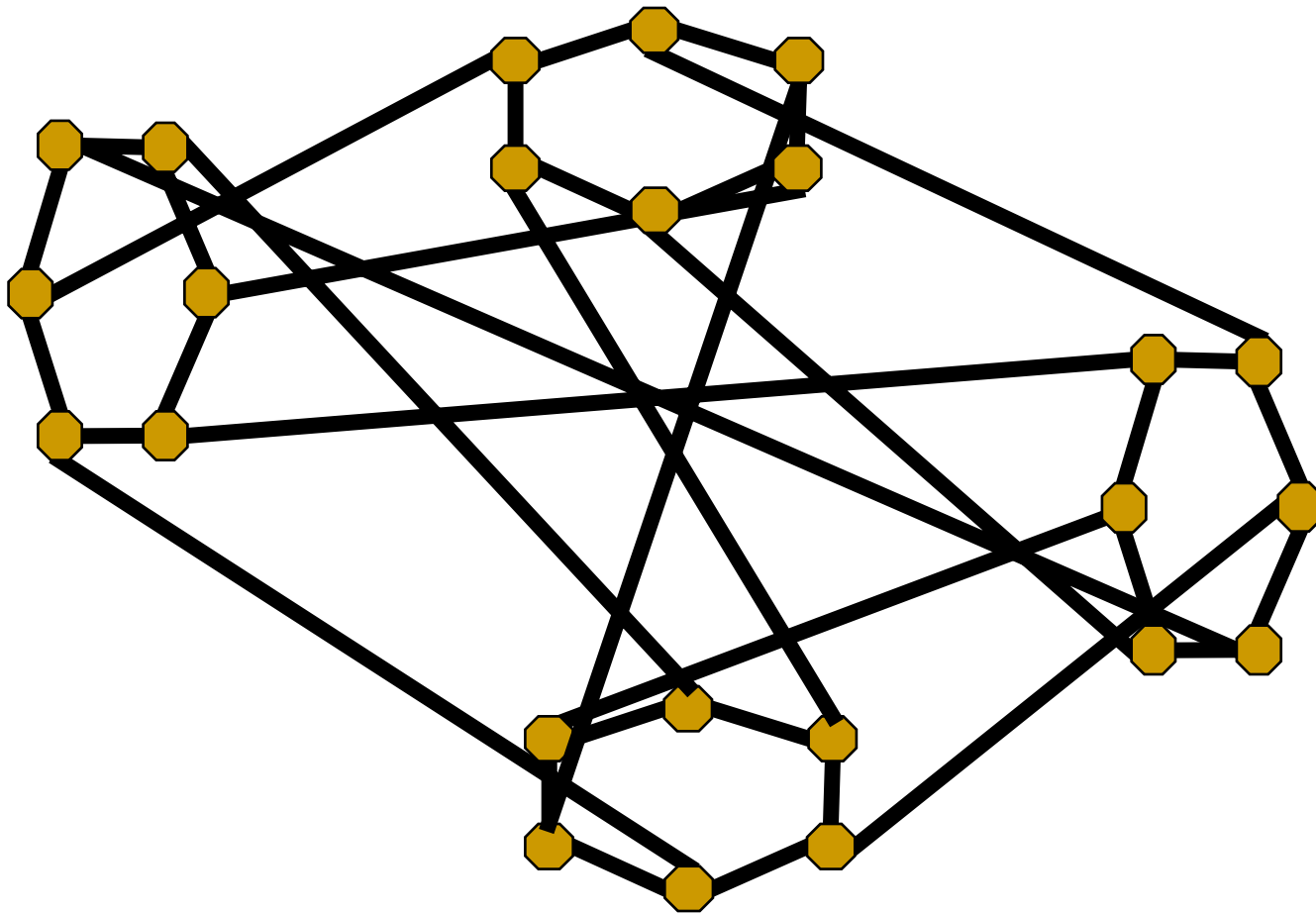
For each node, assign it the name of one of the $n!$ stacks of n pancakes

Put a wire between two nodes if they are one flip apart

Network For $n = 3$

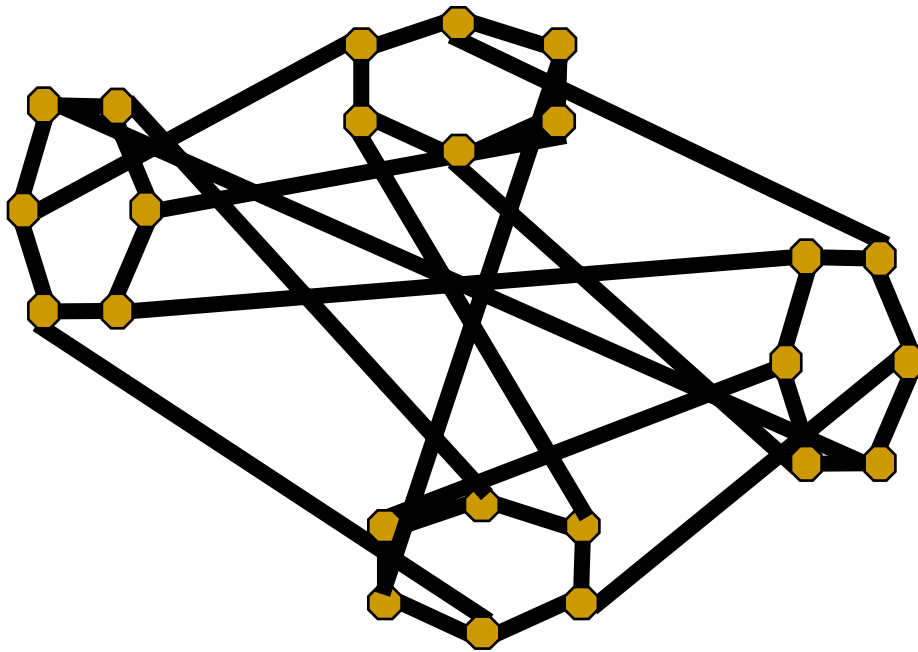


Network For $n=4$



Pancake Network: Message Routing Delay

What is the maximum distance between two nodes in the pancake network?



P_n

Pancake Network: Reliability

If up to $n-2$ nodes get hit by lightning, the network remains connected, even though each node is connected to only $n-1$ others

The Pancake Network is optimally reliable for its number of edges and nodes

Mutation Distance

Head Cabbage
(*Brassica oleracea capitata*)



Turnip
(*Brassica rapa*)



One “Simple” Problem



A host of
problems and
applications at
the frontiers of
science

High Level Point

Computer Science is not merely about computers and programming, it is about mathematically modeling our world, and about finding better and better ways to solve problems

Today's lecture is a microcosm of this exercise





Study Bee

Definitions of:

n^{th} pancake number

lower bound

upper bound

Proof of:

ANY to ANY in $\leq P_n$

Important Technique:

Bracketing