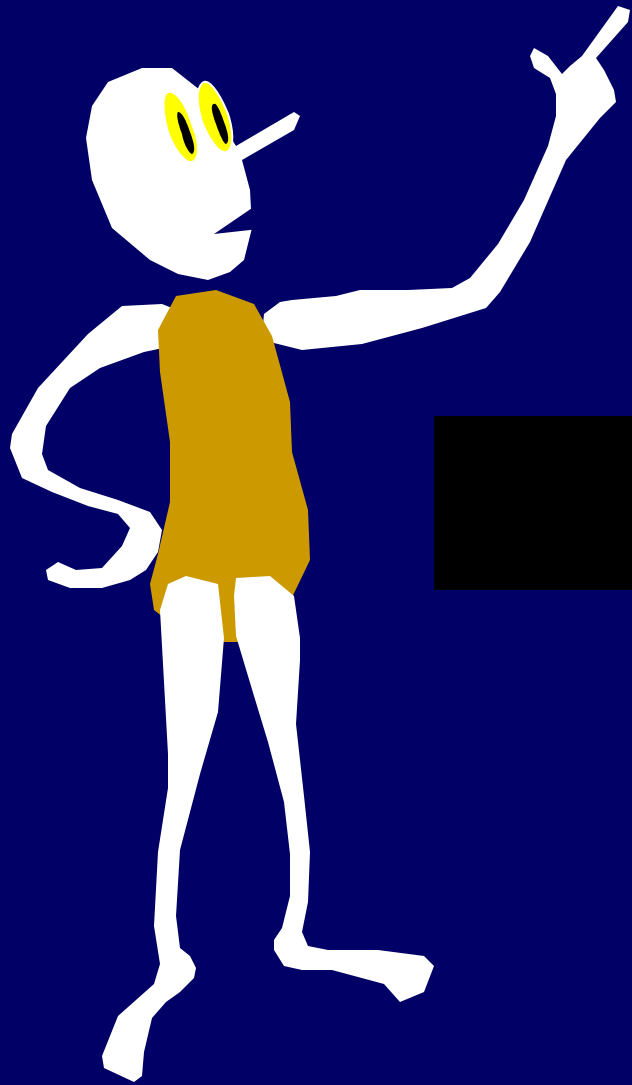


Counting III: Pascal's Triangle, Polynomials, and Vector Programs



$$1 + X^1 + X^2 + X^3 + \dots + X^{n-1} + X^n = \frac{X^{n+1} - 1}{X - 1}$$



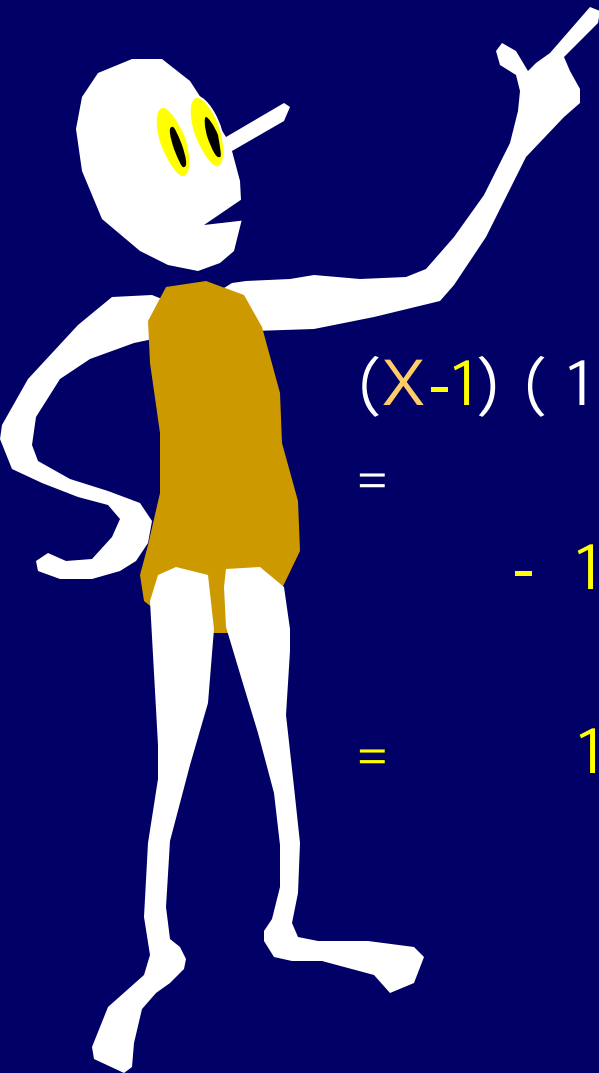
The Geometric Series

$$1 + X^1 + X^2 + X^3 + \dots + X^n + \dots = \frac{1}{1 - X}$$



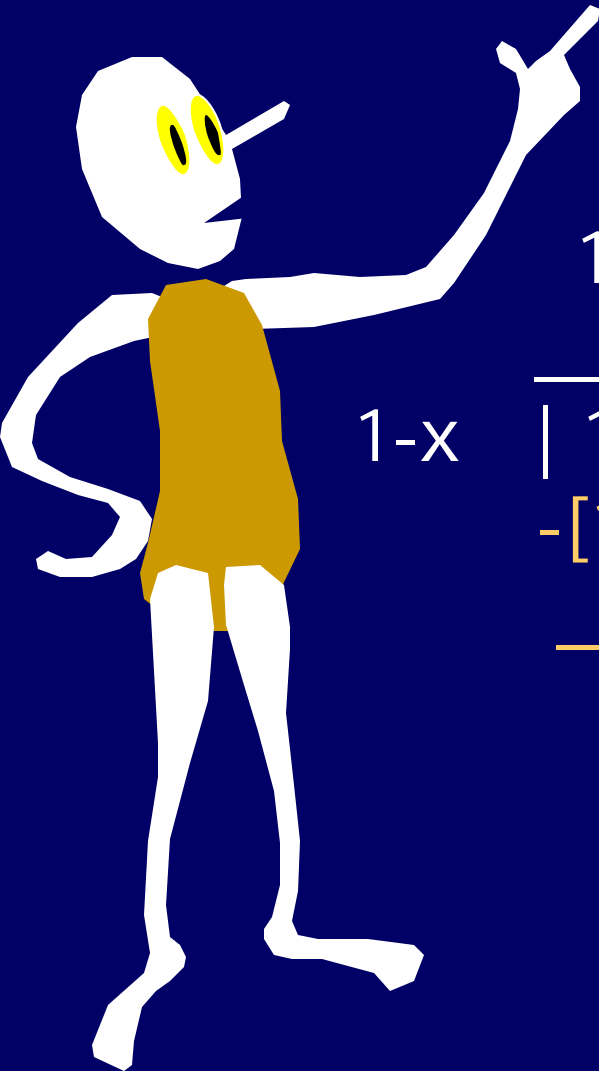
The Infinite Geometric Series

$$1 + X^1 + X^2 + X^3 + \dots + X^n + \dots = \frac{1}{1 - X}$$



$$\begin{aligned} & (X-1) (1 + X^1 + X^2 + X^3 + \dots + X^n + \dots) \\ &= X^1 + X^2 + X^3 + \dots + X^n + X^{n+1} + \dots \\ & \quad - 1 - X^1 - X^2 - X^3 - \dots - X^{n-1} - X^n - X^{n+1} - \dots \\ &= 1 \end{aligned}$$

$$1 + X^1 + X^2 + X^3 + \dots + X^n + \dots = \frac{1}{1 - X}$$



$$1 + X + \dots$$

1-x

$$\begin{array}{r} | 1 \\ -[1 - X] \\ \hline \end{array}$$

X

$$\begin{array}{r} -[X - X^2] \\ \hline \end{array}$$

X²

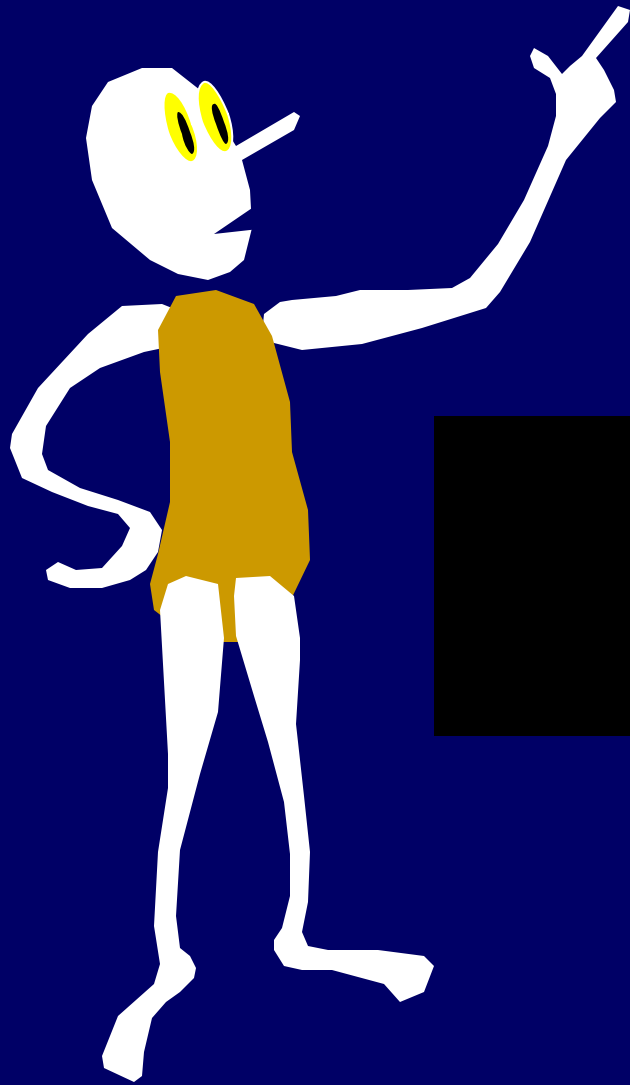
-.....

$$1 + aX^1 + a^2X^2 + a^3X^3 + \dots + a^nX^n + \dots = \frac{1}{1 - aX}$$



Geometric Series (Linear Form)

$$(1 + aX^1 + a^2X^2 + \dots + a^nX^n + \dots) (1 + bX^1 + b^2X^2 + \dots + b^nX^n + \dots) =$$



$$\frac{1}{(1 - aX)(1 - bX)}$$

Geometric Series
(Quadratic Form)

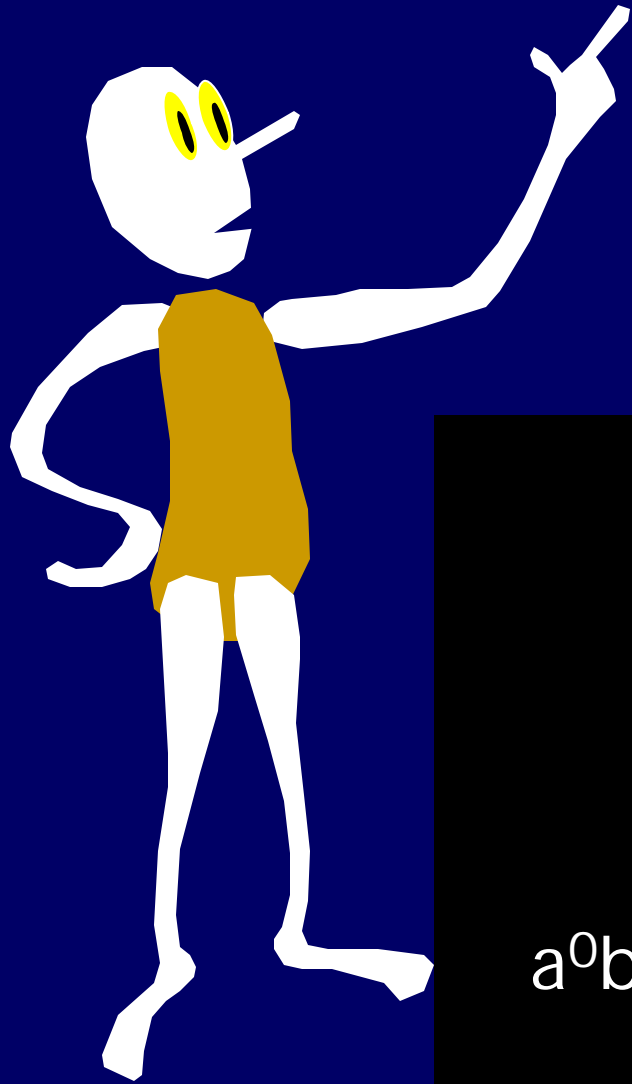
$$(1 + aX^1 + a^2X^2 + \dots + a^nX^n + \dots) (1 + bX^1 + b^2X^2 + \dots + b^nX^n + \dots) = 1 + c_1X^1 + \dots + c_k X^k + \dots$$



Suppose we multiply this out to get a single, infinite polynomial.

What is an expression for C_n ?

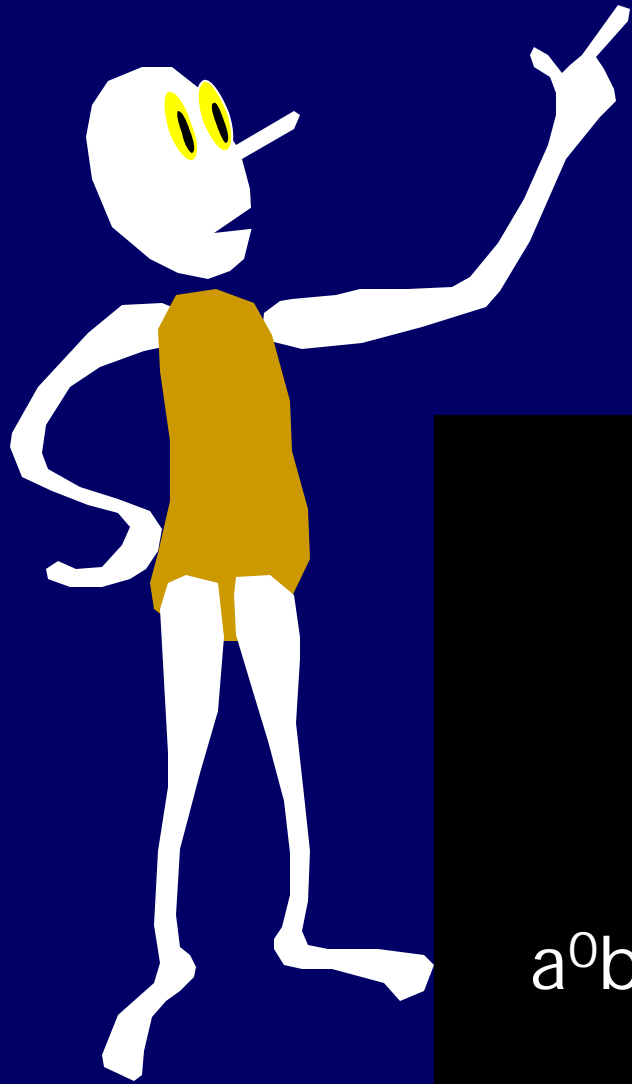
$$(1 + aX^1 + a^2X^2 + \dots + a^nX^n + \dots) (1 + bX^1 + b^2X^2 + \dots + b^nX^n + \dots) = 1 + c_1X^1 + \dots + c_k X^k + \dots$$



$$c_n =$$

$$a^0b^n + a^1b^{n-1} + \dots + a^i b^{n-i} \dots + a^{n-1}b^1 + a^n b^0$$

$$(1 + aX^1 + a^2X^2 + \dots + a^nX^n + \dots) (1 + bX^1 + b^2X^2 + \dots + b^nX^n + \dots) = 1 + c_1X^1 + \dots + c_k X^k + \dots$$

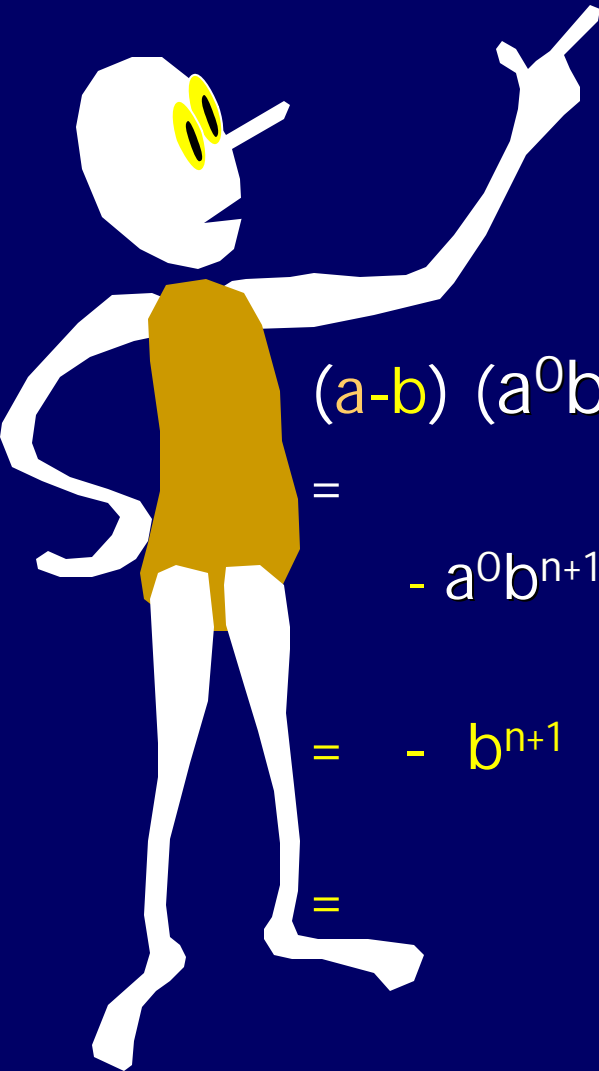


If $a = b$ then

$$c_n = (n+1)(a^n)$$

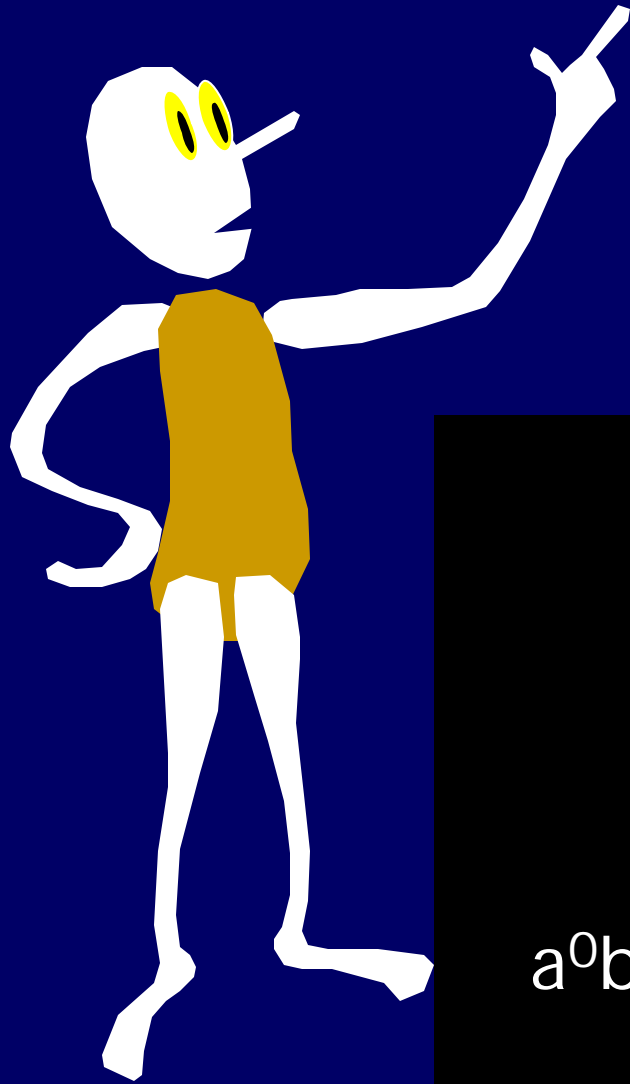
$$a^0b^n + a^1b^{n-1} + \dots + a^i b^{n-i} \dots + a^{n-1}b^1 + a^n b^0$$

$$a^0b^n + a^1b^{n-1} + \dots + a^ib^{n-i} + \dots + a^{n-1}b^1 + a^nb^0 = \frac{a^{n+1} - b^{n+1}}{a - b}$$



$$\begin{aligned}
 & (a-b) (a^0b^n + a^1b^{n-1} + \dots + a^ib^{n-i} + \dots + a^{n-1}b^1 + a^nb^0) \\
 = & \quad a^1b^n + \dots + a^{i+1}b^{n-i} + \dots + a^nb^1 + a^{n+1}b^0 \\
 & - a^0b^{n+1} - a^1b^n - \dots - a^{i+1}b^{n-i} - \dots - a^{n-1}b^2 - a^nb^1 \\
 = & \quad - b^{n+1} + a^{n+1} \\
 = & \quad a^{n+1} - b^{n+1}
 \end{aligned}$$

$$(1 + aX^1 + a^2X^2 + \dots + a^nX^n + \dots) (1 + bX^1 + b^2X^2 + \dots + b^nX^n + \dots) = 1 + c_1X^1 + \dots + c_k X^k + \dots$$



if $a \neq b$ then

$$C_n = \frac{a^{n+1} - b^{n+1}}{a - b}$$

$$a^0b^n + a^1b^{n-1} + \dots + a^i b^{n-i} \dots + a^{n-1}b^1 + a^n b^0$$

$$(1 + aX^1 + a^2X^2 + \dots + a^nX^n + \dots) (1 + bX^1 + b^2X^2 + \dots + b^nX^n + \dots) =$$

$$= \frac{1}{(1 - aX)(1 - bX)}$$

$$= \sum_{n=0..1} \dots$$

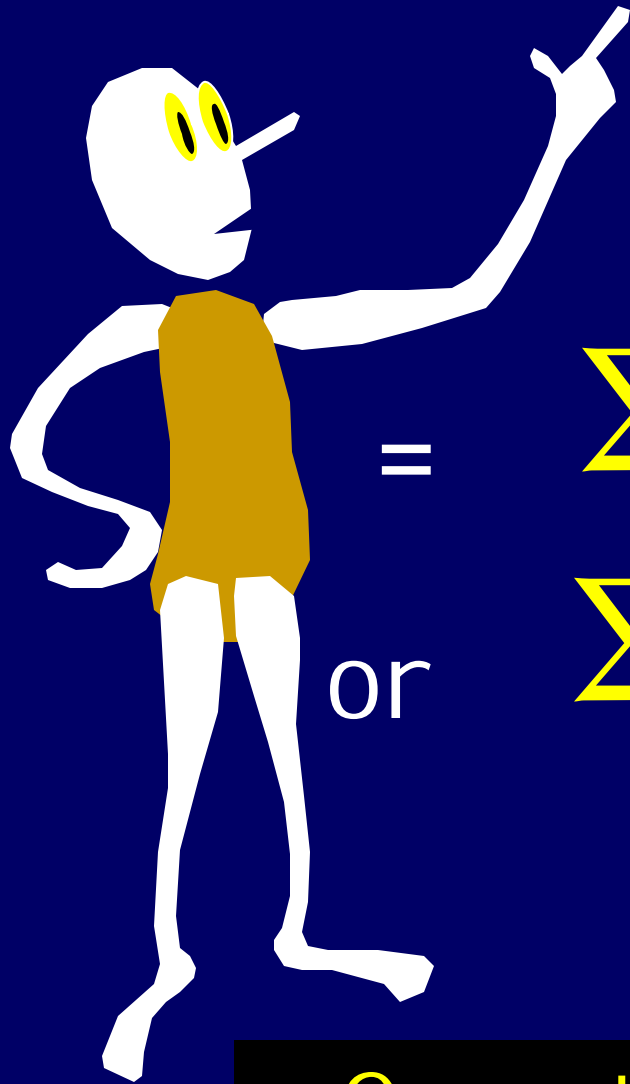
$$\frac{a^{n+1} - b^{n+1}}{a - b} X^n$$

or $\sum_{n=0..1} \dots$

$$(n+1)a^n X^n$$

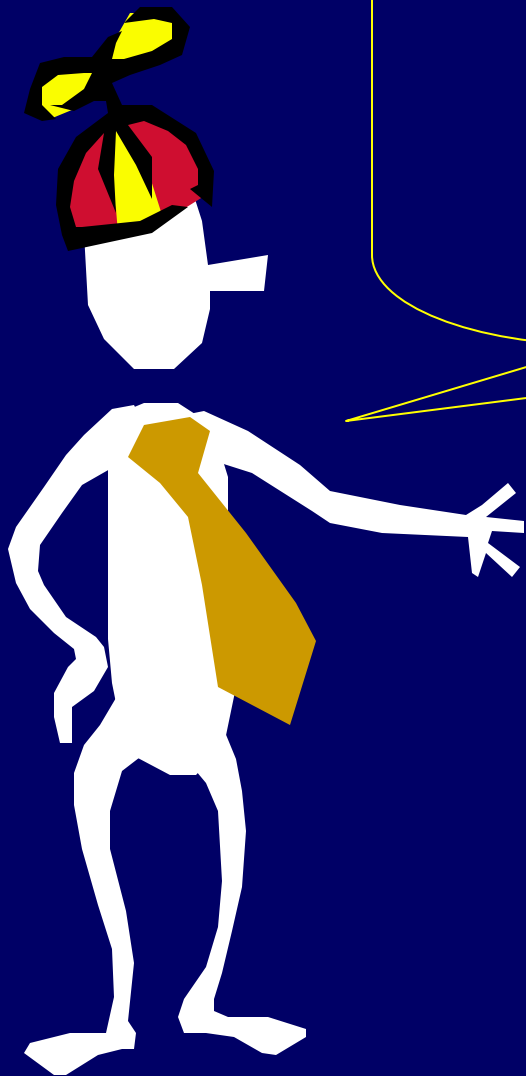
when a=b

Geometric Series (Quadratic Form)



Previously, we saw that

Polynomials Count!

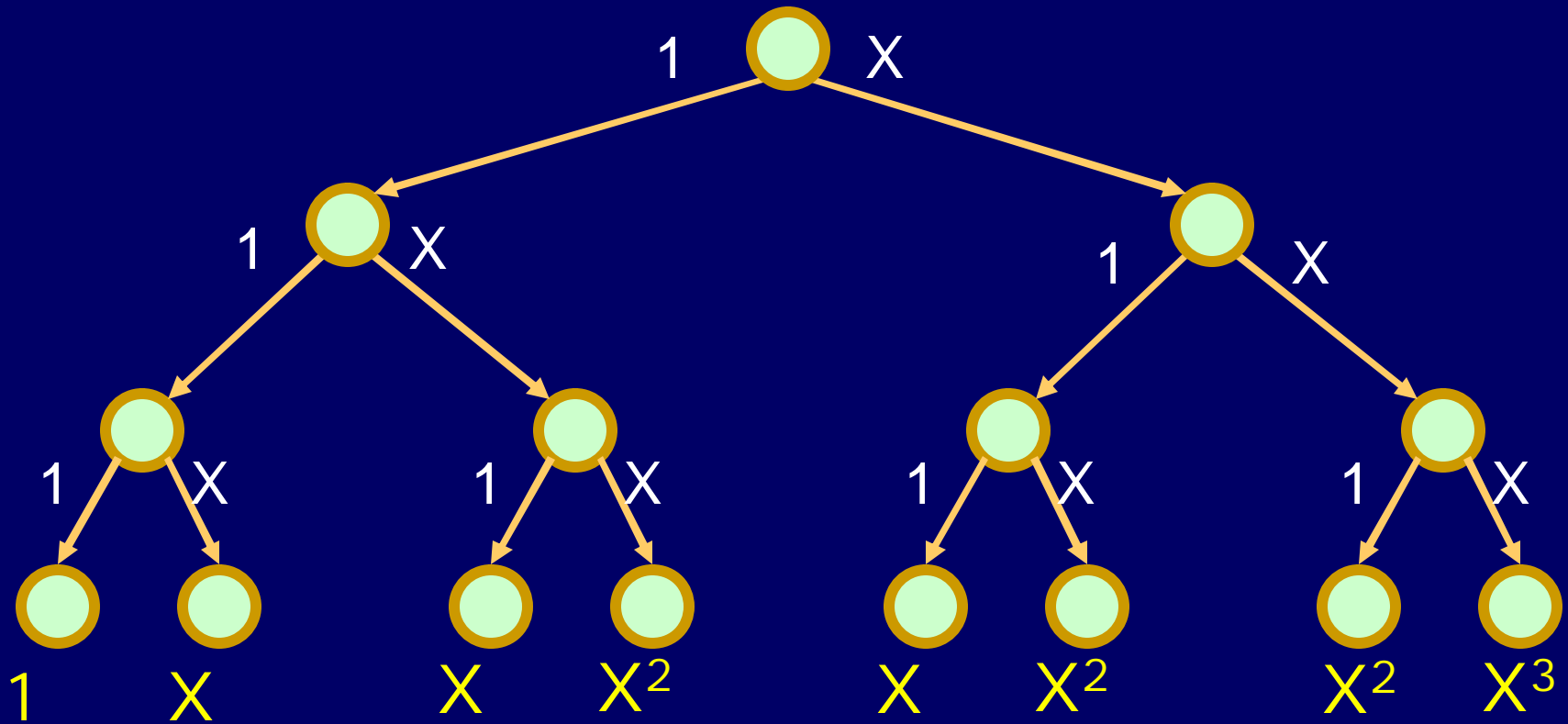




What is the
coefficient of
 BA^3N^2 in the
expansion of
 $(B + A + N)^6$?

The number of
ways to rearrange
the letters in the
word BANANA.

Choice tree for terms of $(1+X)^3$



Combine like terms to get $1 + 3X + 3X^2 + X^3$

The Binomial Formula

$$(1 + X)^n = \binom{n}{0} + \binom{n}{1}X + \binom{n}{2}X^2 + \dots + \binom{n}{k}X^k + \dots + \binom{n}{n}X^n$$

Binomial Coefficients

binomial
expression

The Binomial Formula

$$(1 + x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$

One polynomial, two representations

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$

“Product form” or
“Generating form”

“Additive form” or
“Expanded form”

Power Series Representation

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$

“Closed form” or
“Generating form”

$$= \sum_{k=0}^{\infty} \binom{n}{k} \cdot x^k$$

Since $\binom{n}{k} = 0$ if $k > n$

“Power series” (“Taylor series”) expansion

By playing these two representations against each other we obtain a new representation of a previous insight:

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$

Let $x=1$.

$$2^n = \underbrace{\sum_{k=0}^n \binom{n}{k}}$$

The number of
subsets of an
 n -element set

By varying x , we can discover new identities

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$

Let $x = -1$.

$$0 = \sum_{k=0}^n \binom{n}{k} \cdot (-1)^k$$

Equivalently,

$$\sum_{k \text{ even}}^n \binom{n}{k} = \sum_{k \text{ odd}}^n \binom{n}{k} = 2^{n-1}$$

The number of even-sized subsets of an n element set is the same as the number of odd-sized subsets.

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$

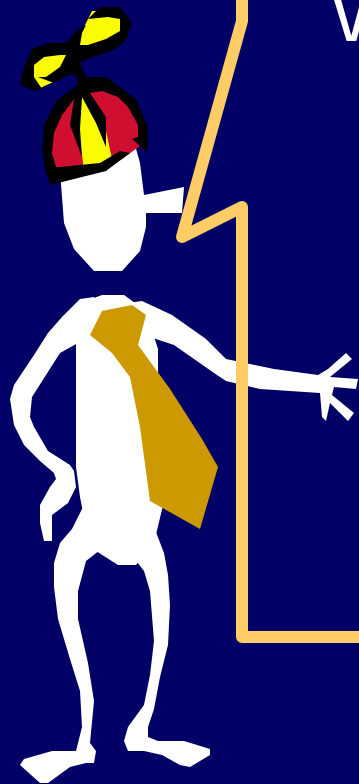
Let $x = -1$.

$$0 = \sum_{k=0}^n \binom{n}{k} \cdot (-1)^k$$

Equivalently,

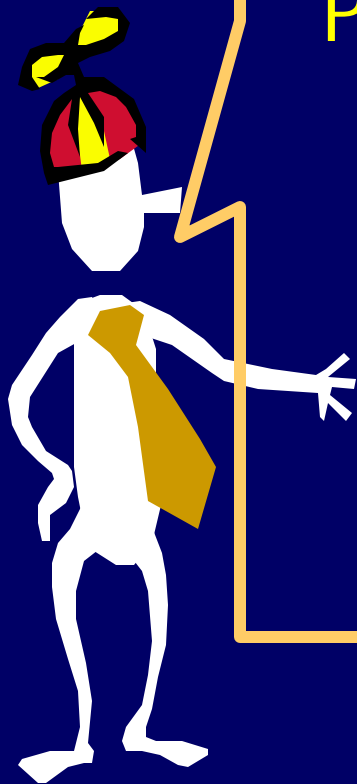
$$\sum_{k \text{ even}}^n \binom{n}{k} = \sum_{k \text{ odd}}^n \binom{n}{k} = 2^{n-1}$$

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$



We could discover new identities by substituting in different numbers for X . One cool idea is to try complex roots of unity, however, the lecture is going in another direction.

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$



Proofs that work by manipulating algebraic forms are called "algebraic" arguments. Proofs that build a 1-1 onto correspondence are called "combinatorial" arguments.

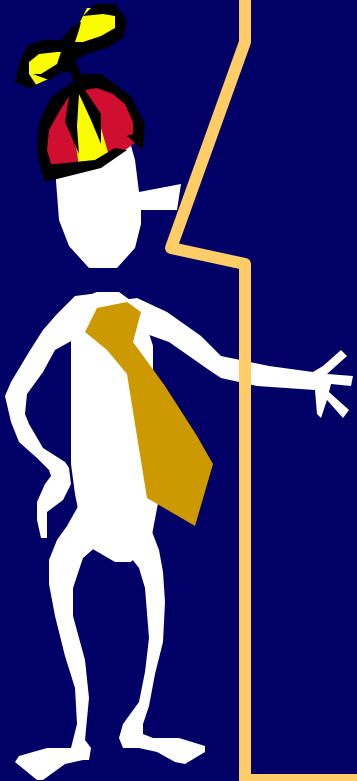
$$\sum_{k \text{ even}}^n \binom{n}{k} = \sum_{k \text{ odd}}^n \binom{n}{k} = 2^{n-1}$$

Let O_n be the set of binary strings of length n with an **odd** number of ones.

Let E_n be the set of binary strings of length n with an **even** number of ones.

We gave an algebraic proof that

$$|O_n| = |E_n|$$



A Combinatorial Proof

Let O_n be the set of binary strings of length n with an **odd** number of ones.

Let E_n be the set of binary strings of length n with an **even** number of ones.

A combinatorial proof must construct a **one-to-one correspondence** between O_n and E_n

An attempt at a correspondence

Let f_n be the function that takes an n -bit string and flips all its bits.

f_n is clearly a one-to-one and onto function

for odd n . E.g. in f_7 we have

0010011 \rightarrow 1101100

1001101 \rightarrow 0110010

...but do even n work? In f_6 we have

110011 \rightarrow 001100

101010 \rightarrow 010101

Uh oh. Complementing maps evens to evens!

A correspondence that works for all n

Let f_n be the function that takes an n -bit string and flips only *the first bit*.

For example,

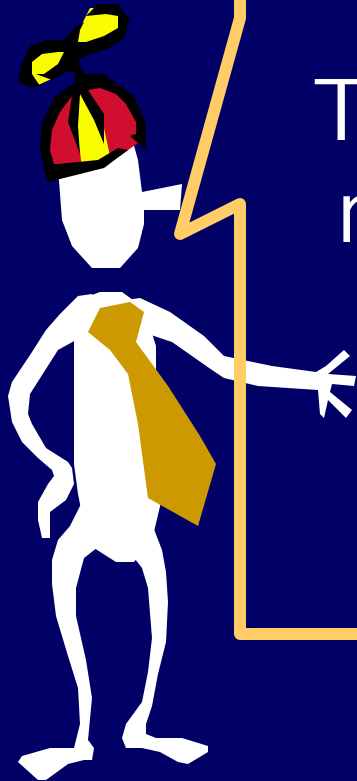
0010011 \rightarrow 1010011

1001101 \rightarrow 0001101

110011 \rightarrow 010011

101010 \rightarrow 001010

$$(1+x)^n = \sum_{k=0}^n \binom{n}{k} \cdot x^k$$



The binomial coefficients have so many representations that many fundamental mathematical identities emerge...

The Binomial Formula

$$(1+X)^0 = 1$$

$$(1+X)^1 = 1 + 1X$$

$$(1+X)^2 = 1 + 2X + 1X^2$$

$$(1+X)^3 = 1 + 3X + 3X^2 + 1X^3$$

$$(1+X)^4 = 1 + 4X + 6X^2 + 4X^3 + 1X^4$$

Pascal's Triangle:

k^{th} row are the coefficients of $(1+X)^k$

$$(1+X)^0 = 1$$

$$(1+X)^1 = 1 + 1X$$

$$(1+X)^2 = 1 + 2X + 1X^2$$

$$(1+X)^3 = 1 + 3X + 3X^2 + 1X^3$$

$$(1+X)^4 = 1 + 4X + 6X^2 + 4X^3 + 1X^4$$

k^{th} Row Of Pascal's Triangle:

$$\binom{n}{0}, \binom{n}{1}, \binom{n}{2}, \dots, \binom{n}{k}, \dots, \binom{n}{n}$$

$$(1+X)^0 = 1$$

$$(1+X)^1 = 1 + 1X$$

$$(1+X)^2 = 1 + 2X + 1X^2$$

$$(1+X)^3 = 1 + 3X + 3X^2 + 1X^3$$

$$(1+X)^4 = 1 + 4X + 6X^2 + 4X^3 + 1X^4$$

Inductive definition of kth entry of nth row:

$$\text{Pascal}(n,0) = \text{Pascal}(n,n) = 1;$$

$$\text{Pascal}(n,k) = \text{Pascal}(n-1,k-1) + \text{Pascal}(n-1,k)$$

$$(1+X)^0 = 1$$

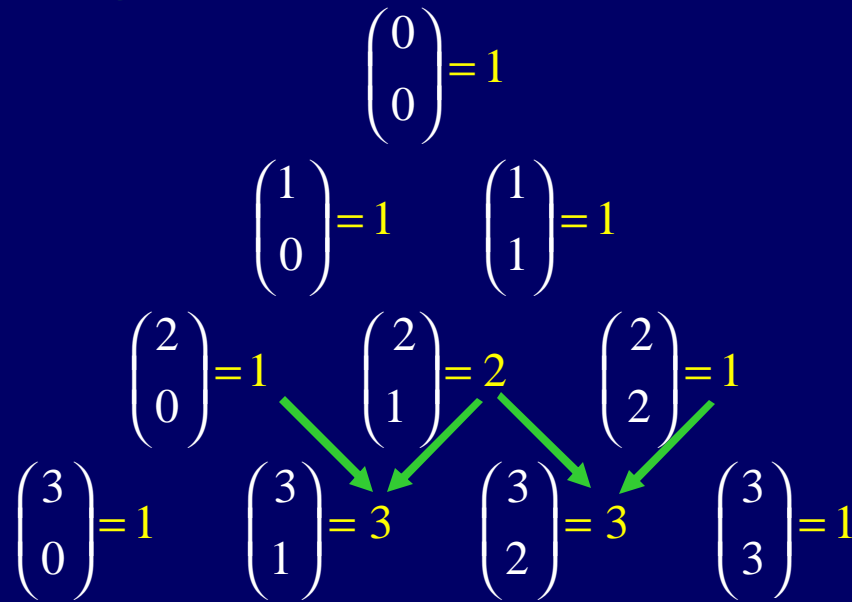
$$(1+X)^1 = 1 + 1X$$

$$(1+X)^2 = 1 + 2X + 1X^2$$

$$(1+X)^3 = 1 + 3X + 3X^2 + 1X^3$$

$$(1+X)^4 = 1 + 4X + 6X^2 + 4X^3 + 1X^4$$

"Pascal's Triangle"



Al-Karaji, Baghdad 953-1029

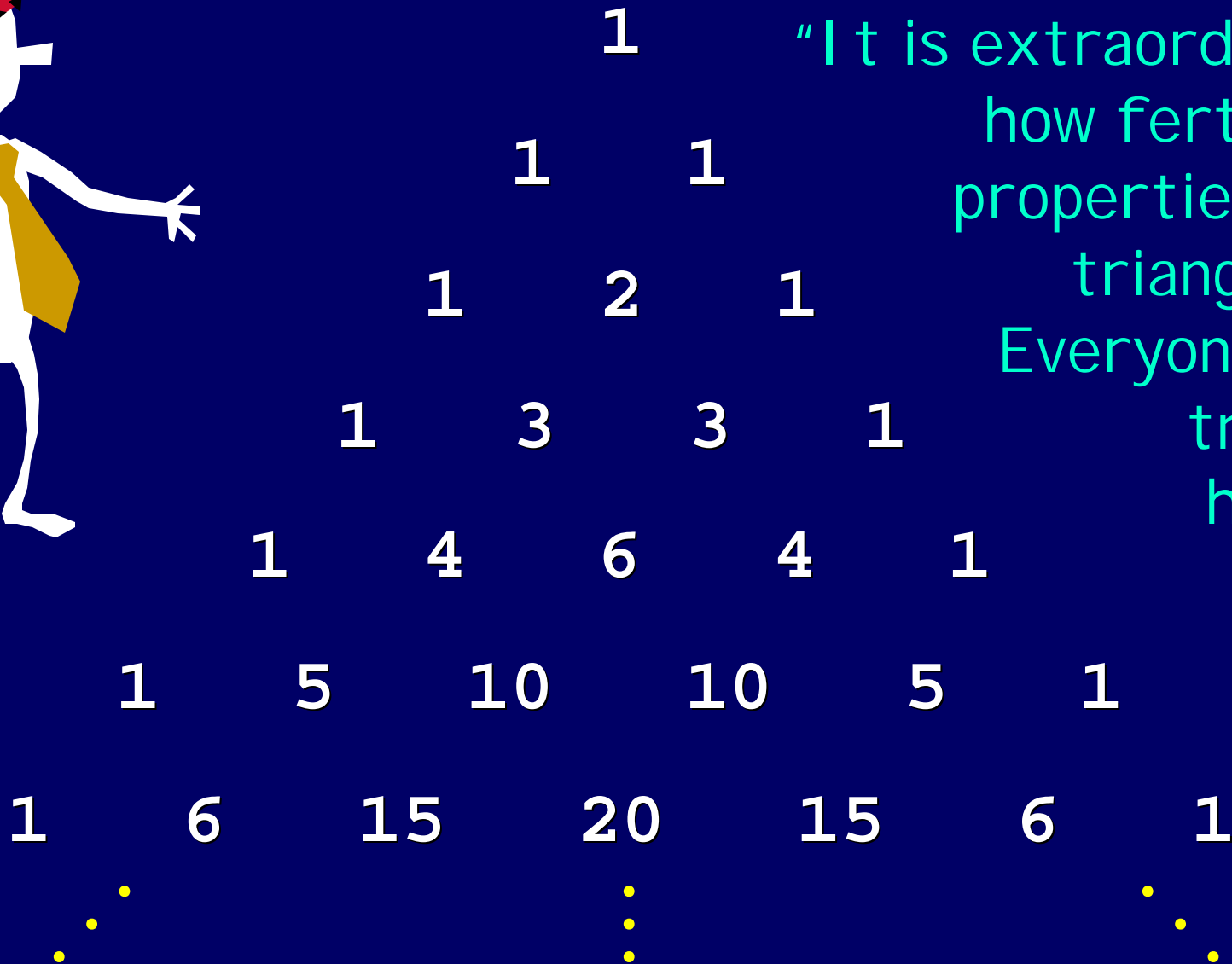
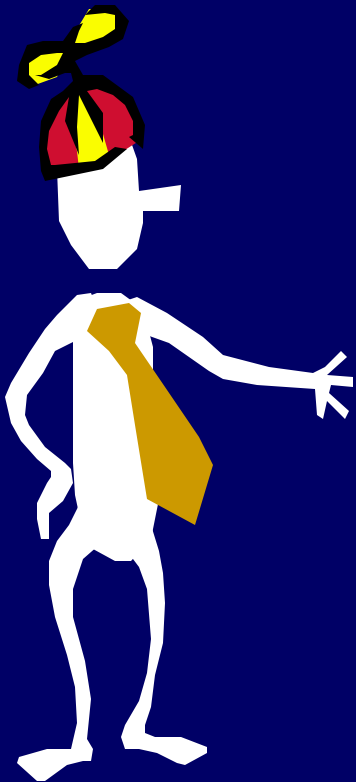
Chu Shin-Chieh 1303

The Precious Mirror of the Four Elements

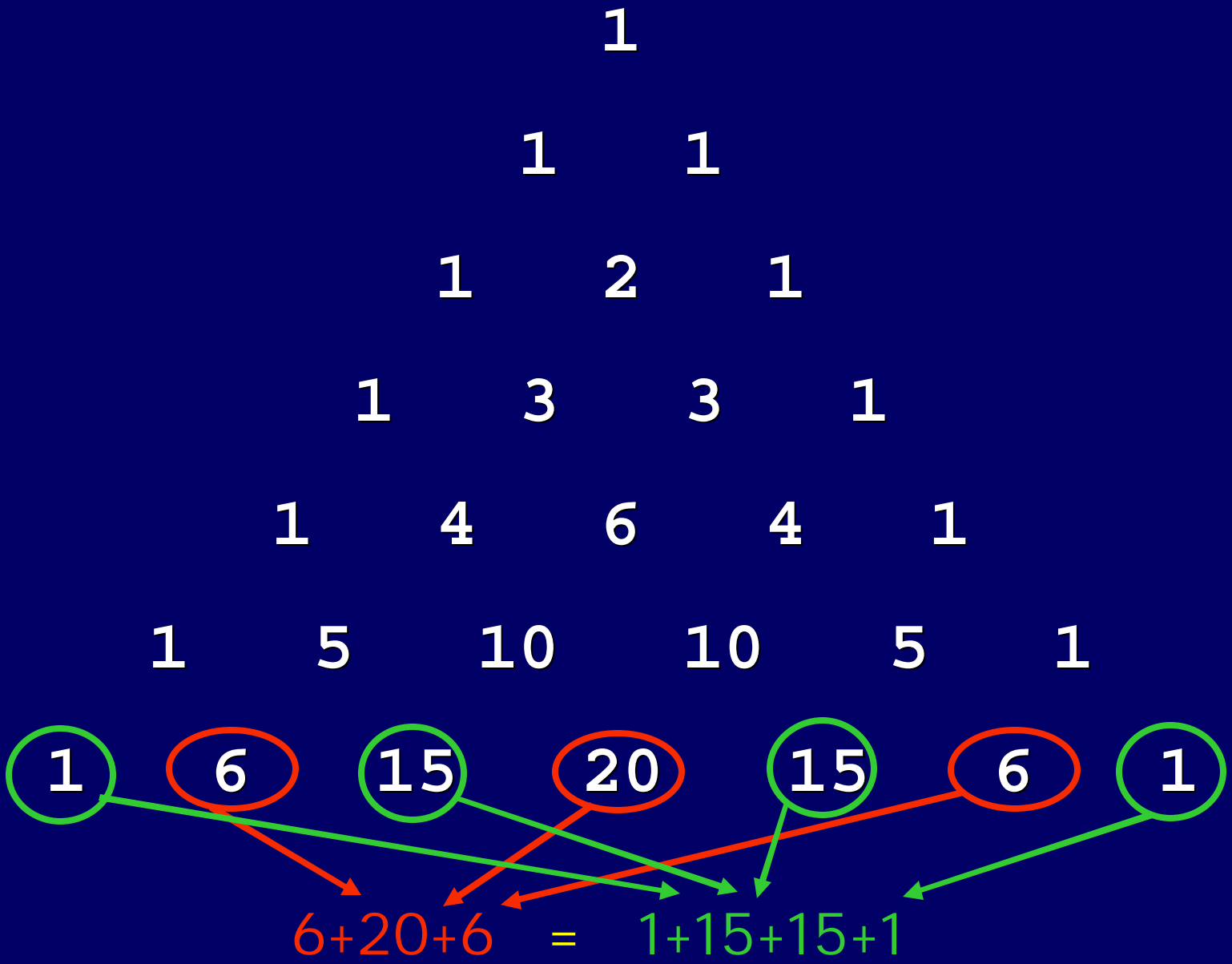
... Known in Europe by 1529

Blaise Pascal 1654

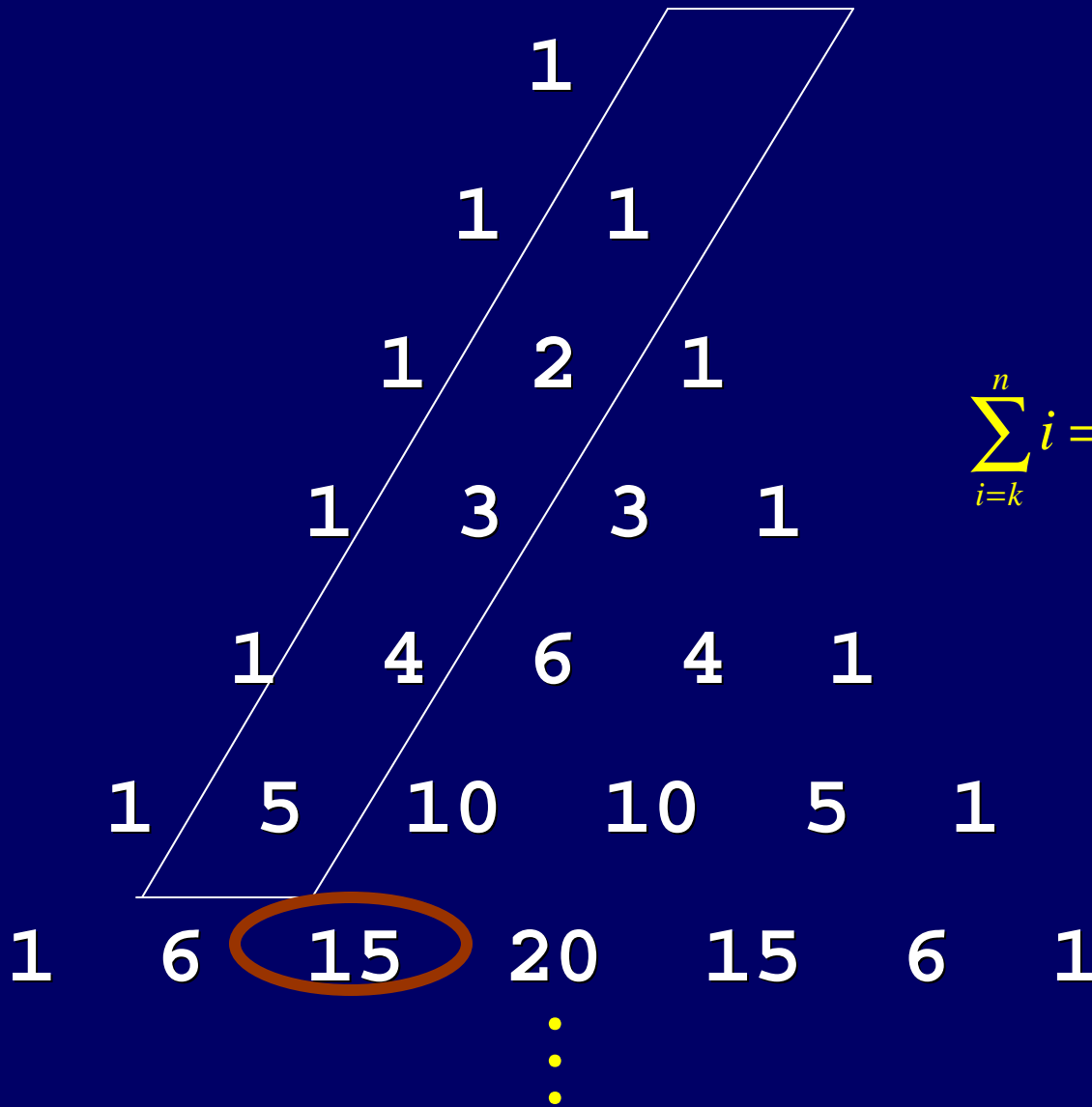
Pascal's Triangle



"It is extraordinary how fertile in properties the triangle is. Everyone can try his hand."

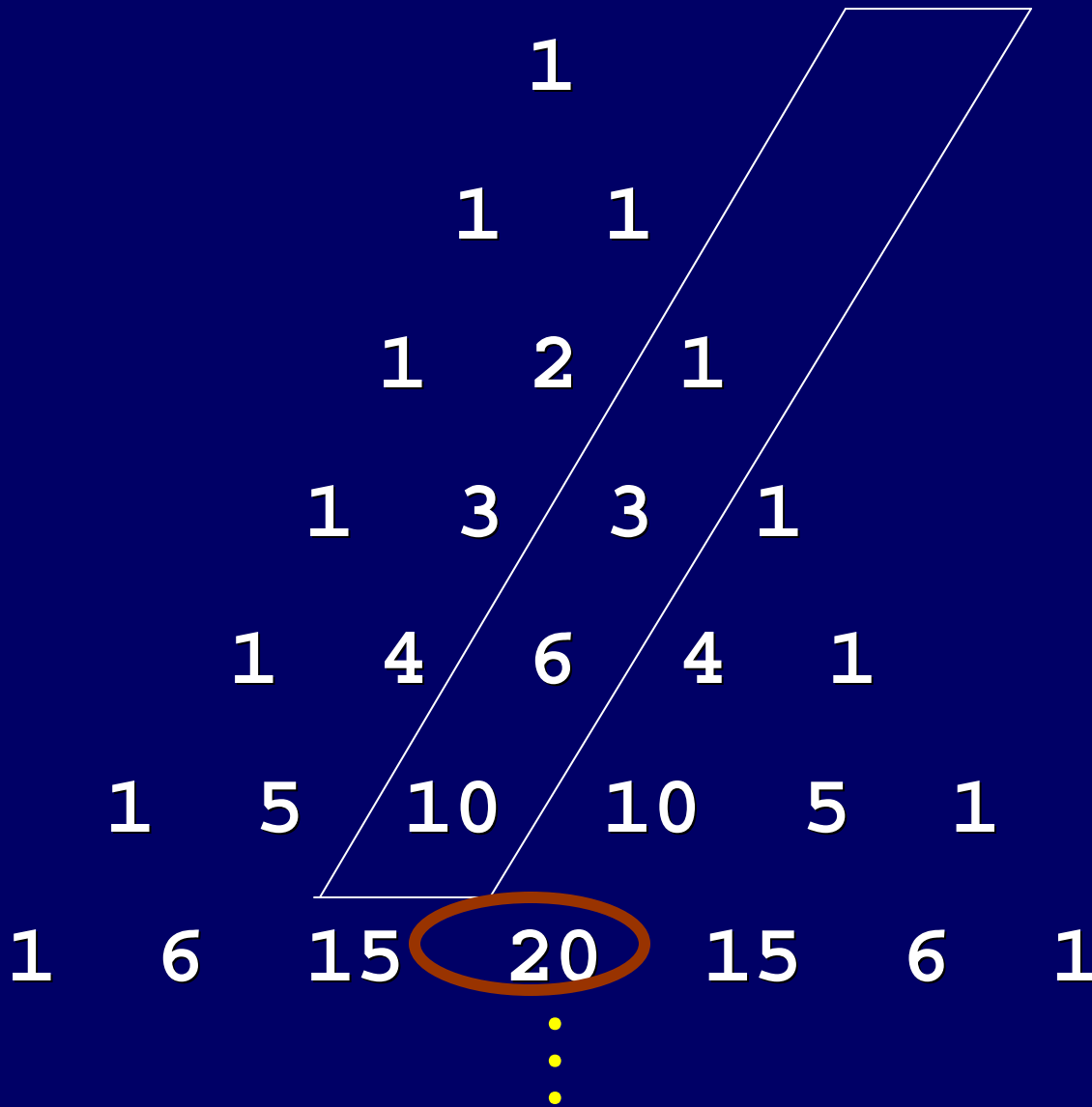


Summing on 1st Avenue



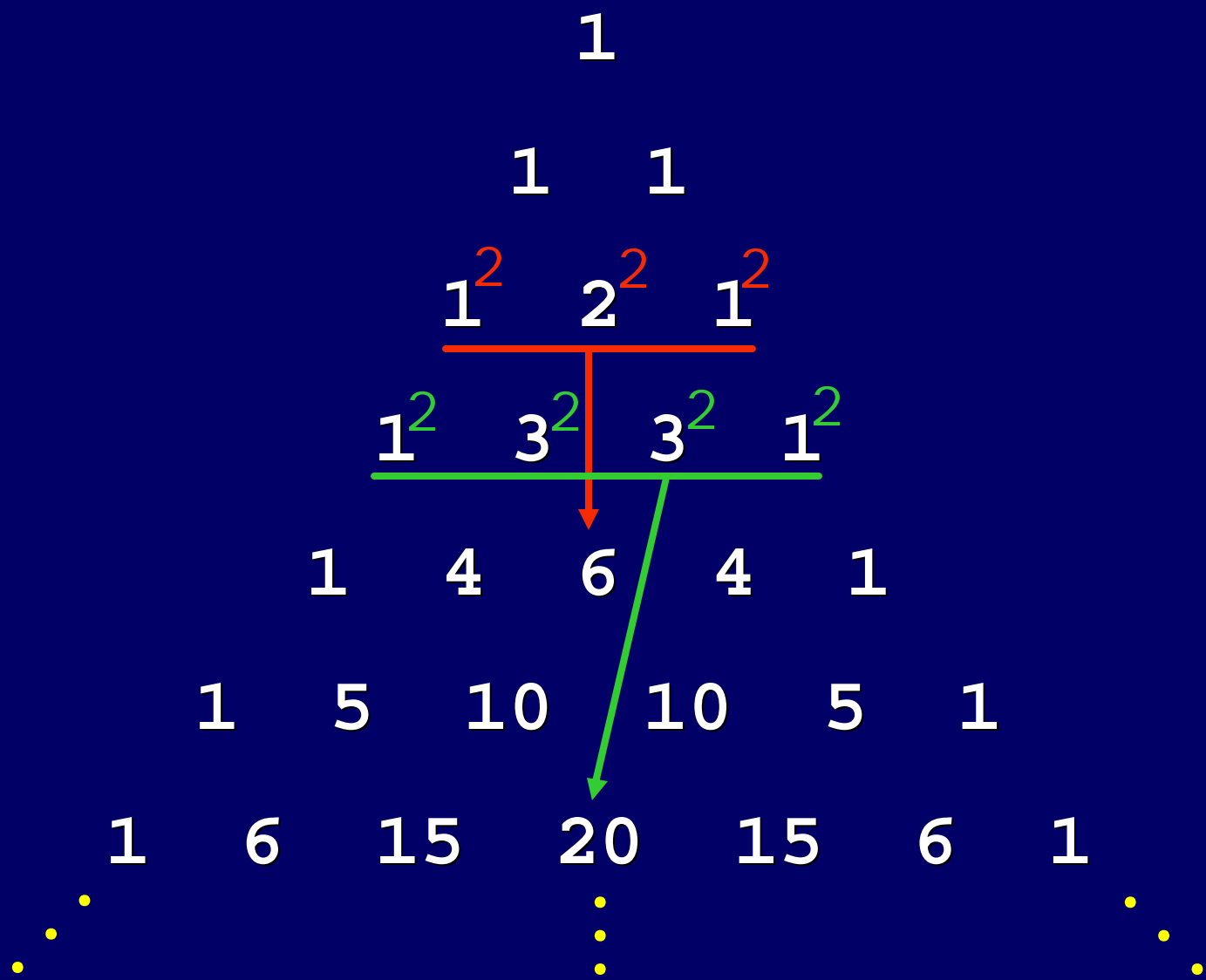
$$\sum_{i=k}^n i = \sum_{i=k}^n \binom{i}{1} = \binom{n+1}{2} = \frac{n \cdot (n+1)}{2}$$

Summing on k^{th} Avenue



$$\sum_{i=k}^n \binom{i}{k} = \binom{n+1}{k+1}$$





Al-Karaji Squares

$$1 = 0$$

$$1 + 1 = 1$$

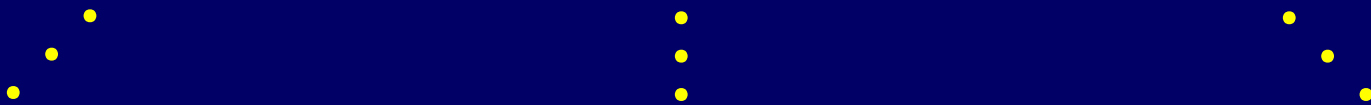
$$1 + 2 + 2 * 1 = 4$$

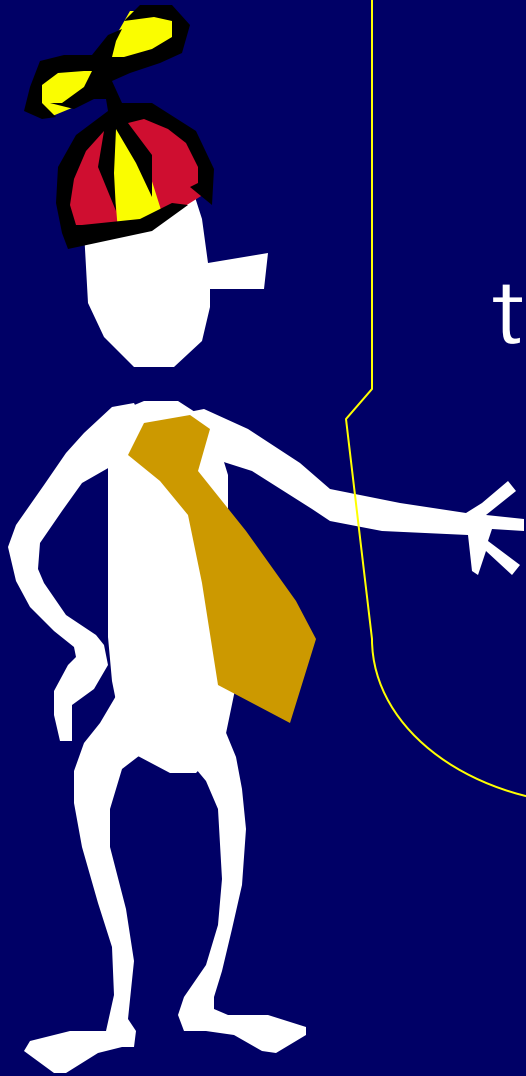
$$1 + 3 + 2 * 3 + 1 = 9$$

$$1 + 4 + 2 * 6 + 4 + 1 = 16$$

$$1 + 5 + 2 * 10 + 10 + 5 + 1 = 25$$

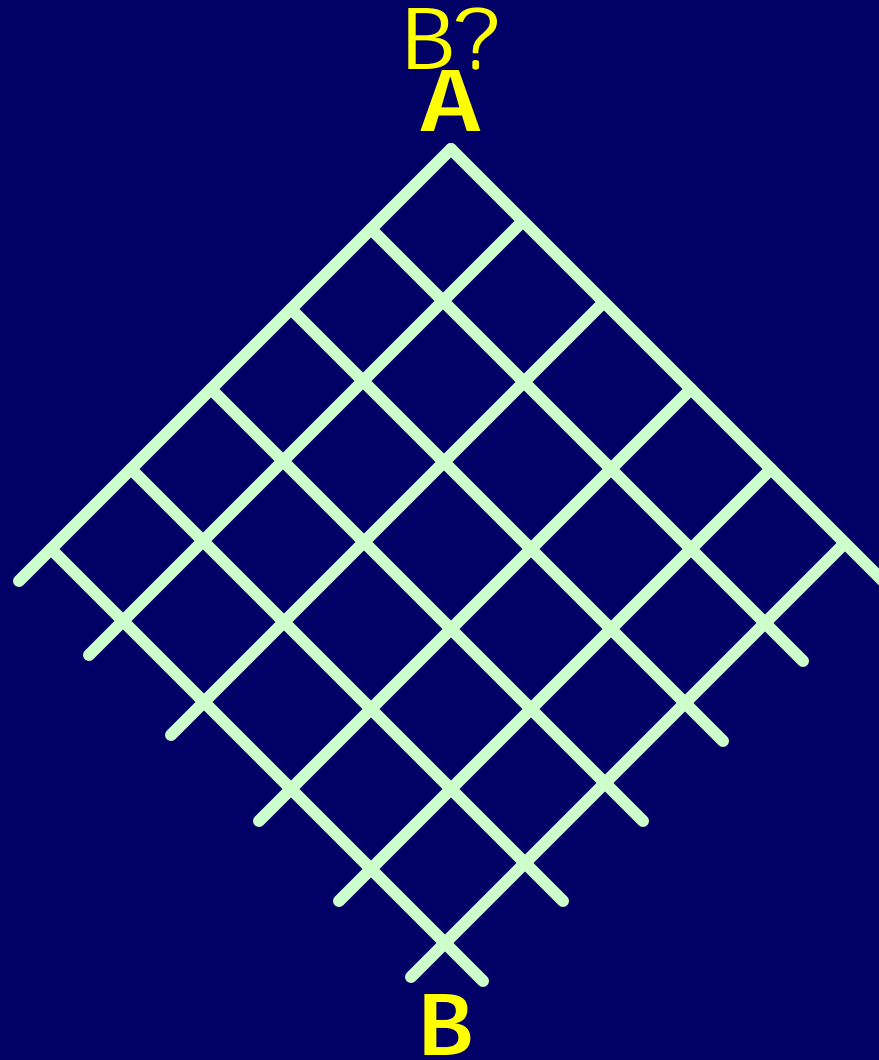
$$1 + 6 + 2 * 15 + 20 + 15 + 6 + 1 = 36$$





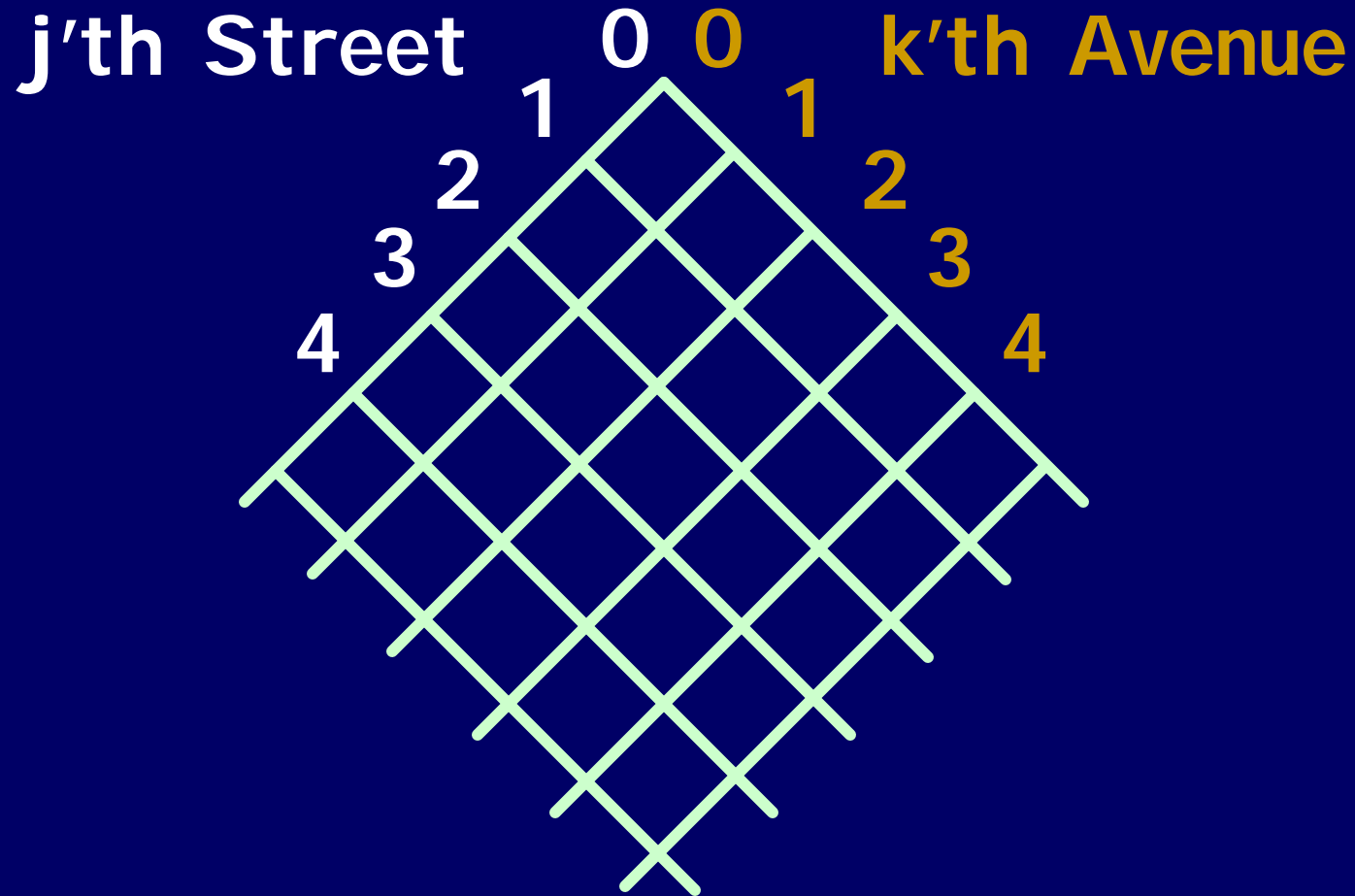
All these properties can be proved inductively and algebraically. We will give combinatorial proofs using the **Manhattan block walking** representation of binomial coefficients.

How many shortest routes from A to



$$\begin{pmatrix} 10 \\ 5 \end{pmatrix}$$

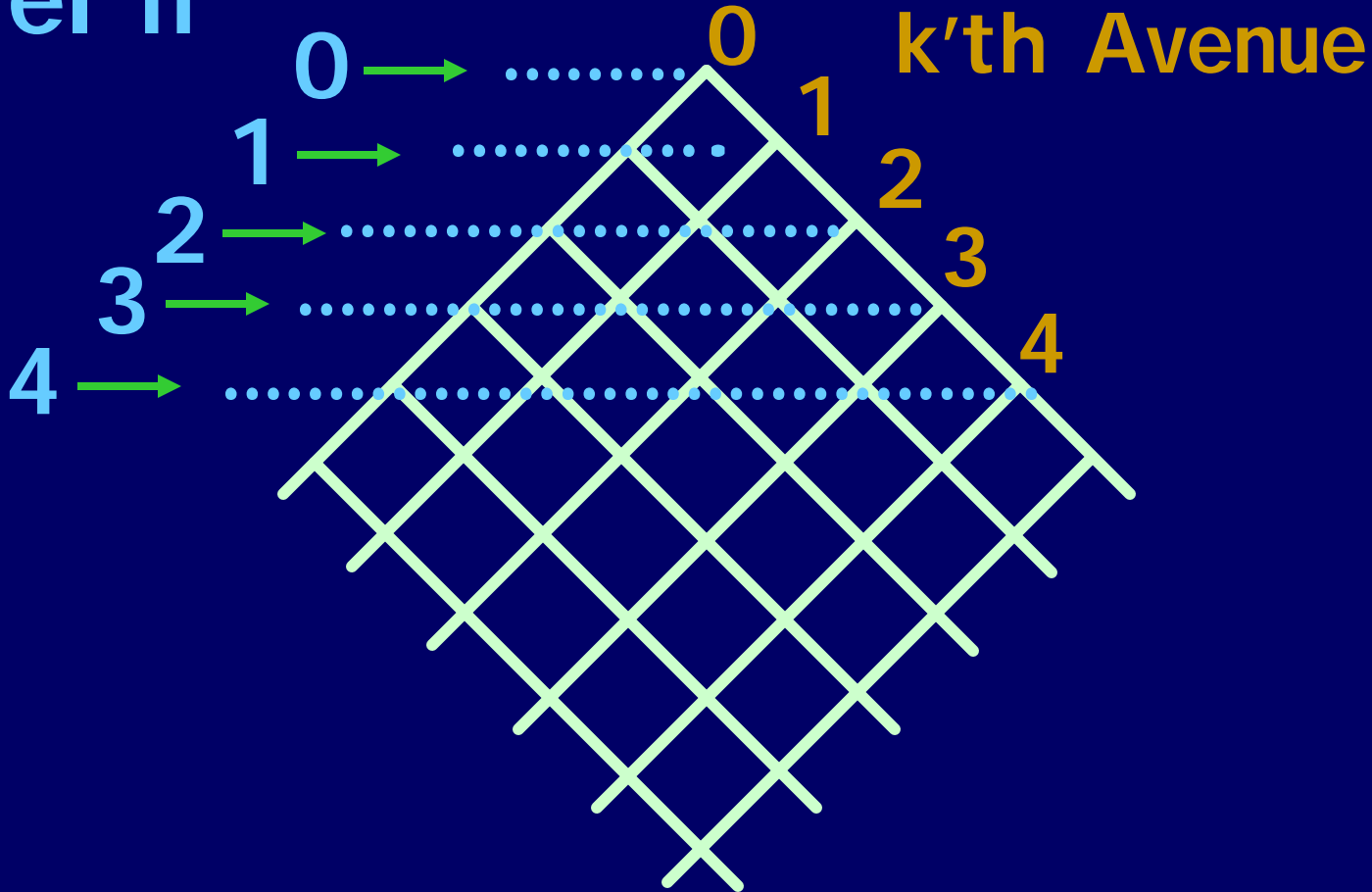
Manhattan



There are $\binom{j+k}{k}$ shortest routes from $(0,0)$ to (j,k) .

Manhattan

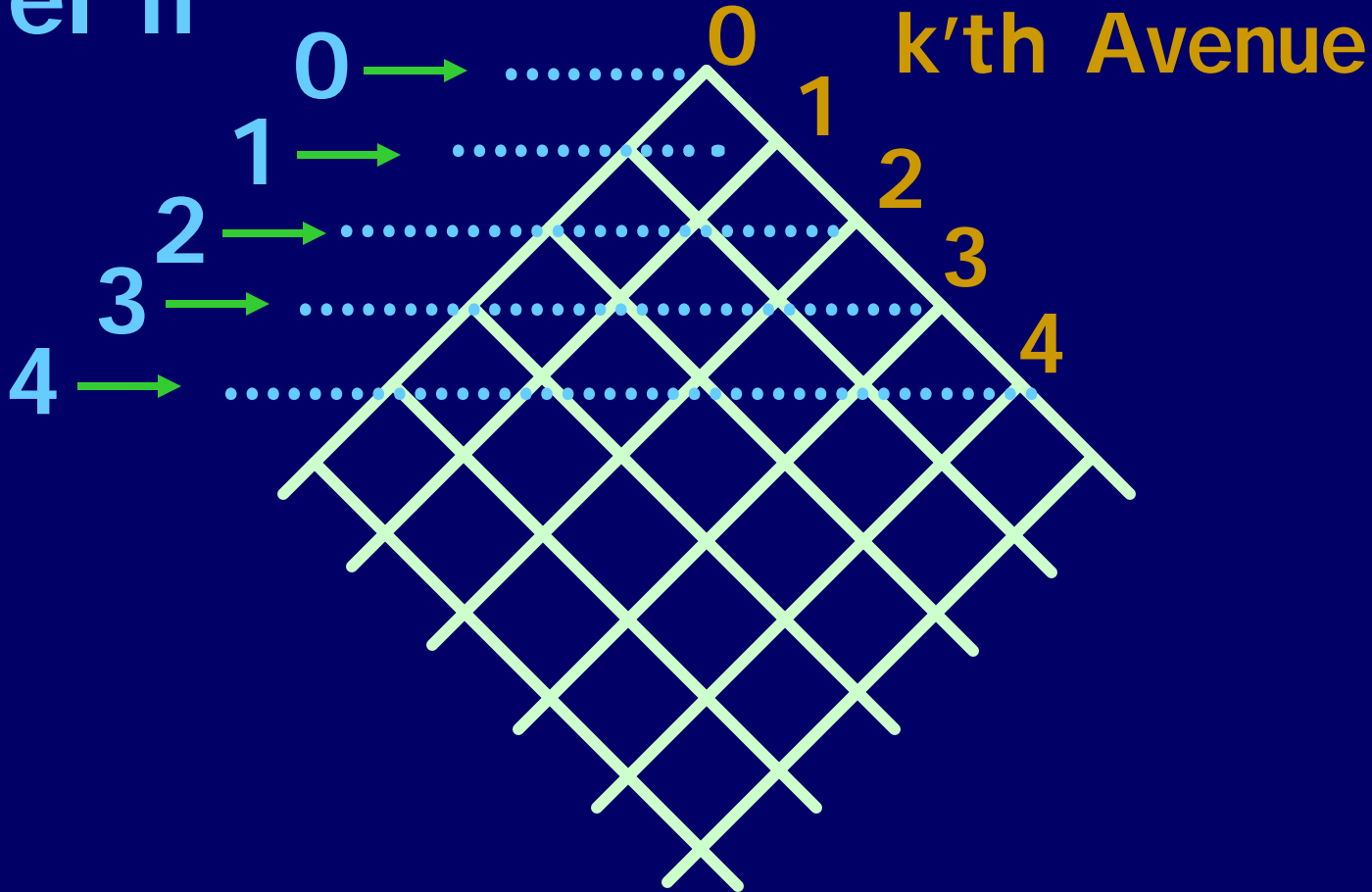
Level n



There are $\binom{n}{k}$ shortest routes from $(0,0)$ to $(n-k,k)$.

Manhattan

Level n

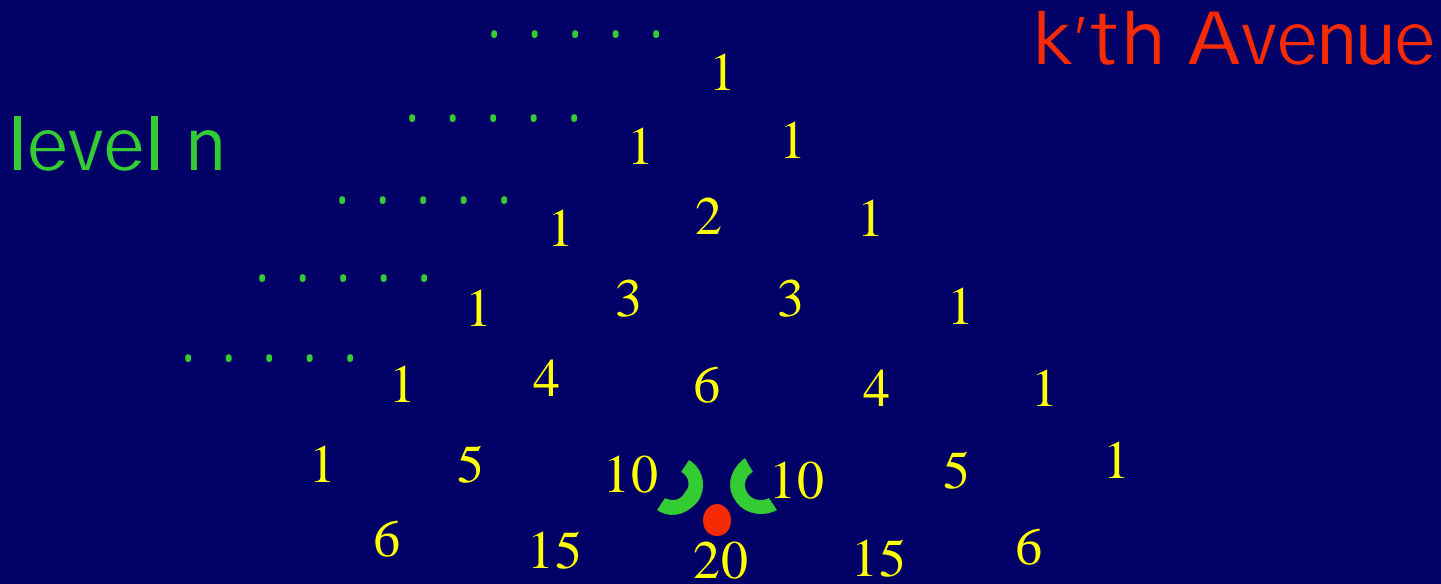


There are $\binom{n}{k}$ shortest routes from $(0,0)$ to Level n and k^{th} Avenue.

level n

k'th Avenue

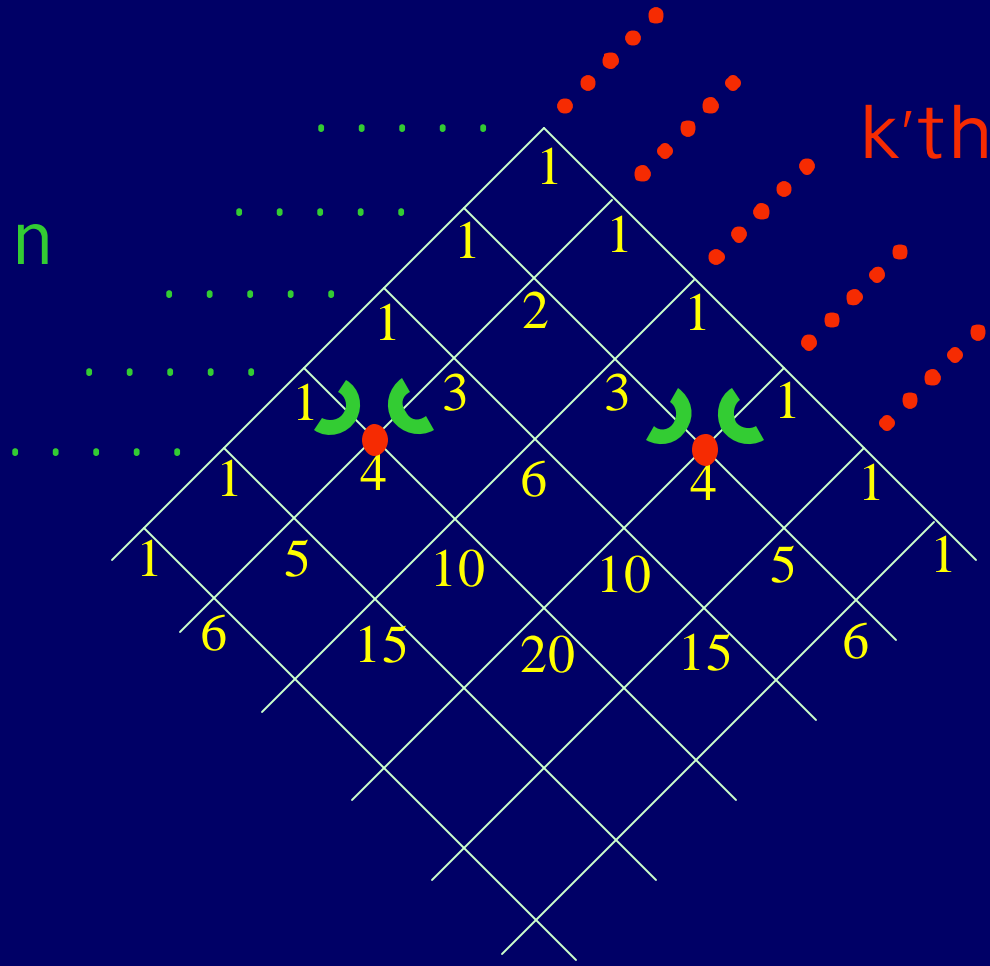




$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}$$

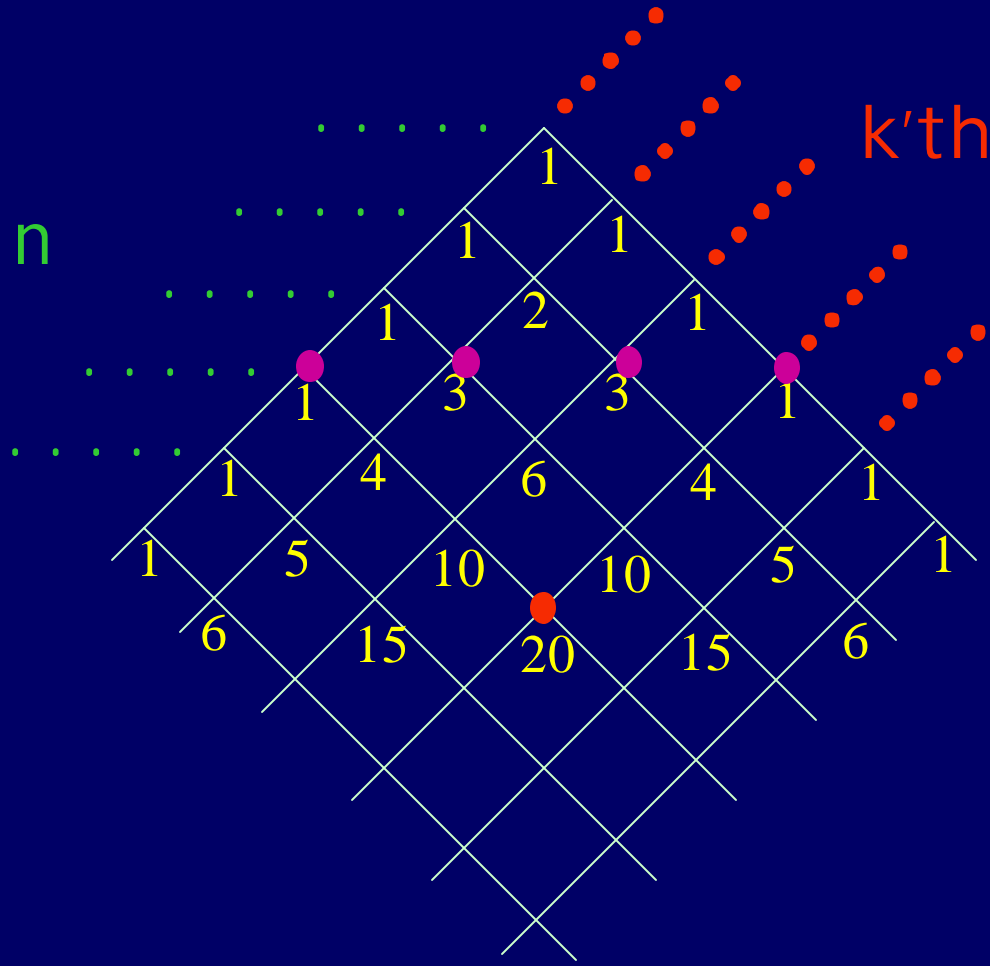
level n

k'th Avenue



$$\sum_{k \text{ even}} \binom{n}{k} = 2^{n-1}$$

level n



k'th Avenue

$$\sum_{k=0}^n \binom{n}{k}^2 = \binom{2n}{n}$$

By convention:

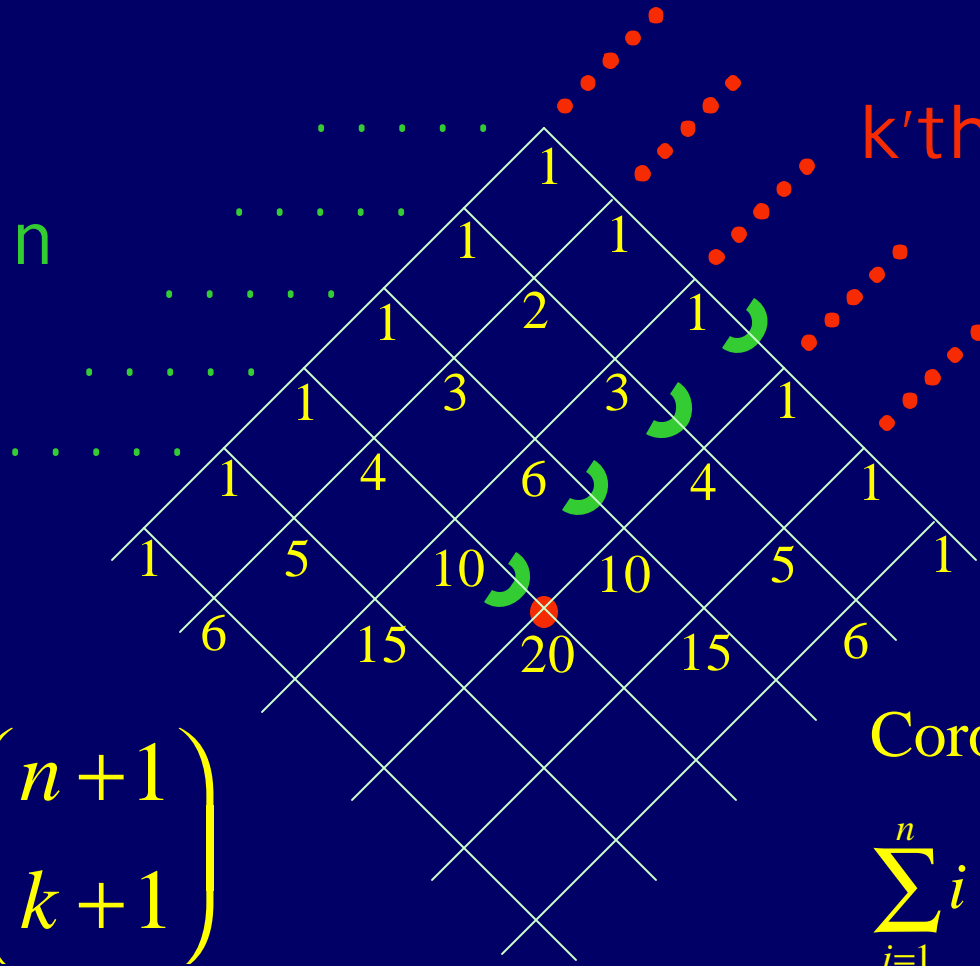
$$0! = 1 \quad (\text{empty product} = 1)$$

$$\binom{n}{k} = 1 \quad \text{if } k = 0$$

$$\binom{n}{k} = 0 \quad \text{if } k < 0 \text{ or } k > n$$

level n

k'th Avenue



$$\sum_{i=1}^n \binom{i}{k} = \binom{n+1}{k+1}$$

Corollary ($k = 1$)

$$\sum_{i=1}^n i = \binom{n+1}{2} = \frac{n(n+1)}{2}$$

Application (Al-Karaji):

$$\begin{aligned}\sum_{i=0}^n i^2 &= 1^2 + 2^2 + 3^2 + \cdots + n^2 \\ &= (1 \cdot 0 + 1) + (2 \cdot 1 + 2) + (3 \cdot 2 + 3) + \cdots + (n(n-1) + n) \\ &= 1 \cdot 0 + 2 \cdot 1 + 3 \cdot 2 + \cdots + n(n-1) + \sum_{i=1}^n i \\ &= 2 \left[\binom{2}{2} + \binom{3}{2} + \binom{4}{2} + \binom{5}{2} + \cdots + \binom{n}{2} \right] + \binom{n+1}{2} \\ &= 2 \binom{n+1}{3} + \binom{n+1}{2} = \frac{(2n+1)(n+1)n}{6}\end{aligned}$$

Vector Programs

Let's define a (parallel) programming language called VECTOR that operates on possibly infinite vectors of numbers. Each variable V^i can be thought of as:

< * , * , * , * , * , * , >

0 1 2 3 4 5

Vector Programs

Let k stand for a scalar constant

$\langle k \rangle$ will stand for the vector $\langle k, 0, 0, 0, \dots \rangle$

$$\langle 0 \rangle = \langle 0, 0, 0, 0, \dots \rangle$$

$$\langle 1 \rangle = \langle 1, 0, 0, 0, \dots \rangle$$

$V! + T!$ means to add the vectors position-wise.

$$\langle 4, 2, 3, \dots \rangle + \langle 5, 1, 1, \dots \rangle = \langle 9, 3, 4, \dots \rangle$$

Vector Programs

RIGHT($V!$) means to shift every number in $V!$ one position to the **right** and to place a 0 in position 0.

$$\text{RIGHT}(\langle 1, 2, 3, \dots \rangle) = \langle 0, 1, 2, 3, \dots \rangle$$

Vector Programs

Example:

State

$V^i := \langle 6 \rangle;$

$V^i = \langle 6, 0, 0, 0, \dots \rangle$

$V^i := \text{RIGHT}(V^i) + \langle 42 \rangle;$

$V^i = \langle 42, 6, 0, 0, \dots \rangle$

$V^i := \text{RIGHT}(V^i) + \langle 2 \rangle;$

$V^i = \langle 2, 42, 6, 0, \dots \rangle$

$V^i := \text{RIGHT}(V^i) + \langle 13 \rangle;$

$V^i = \langle 13, 2, 42, 6, \dots \rangle$

$V^i = \langle 13, 2, 42, 6, 0, 0, 0, \dots \rangle$

Vector Programs

Example:

Stare

$V^! := \langle 1 \rangle;$

$V^! = \langle 1, 0, 0, 0, \dots \rangle$

Loop n times:

$V^! = \langle 1, 1, 0, 0, \dots \rangle$

$V^! := V^! + \text{RIGHT}(V^!);$

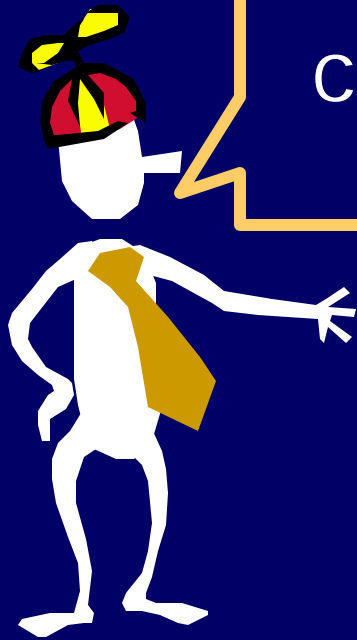
$V^! = \langle 1, 2, 1, 0, \dots \rangle$

$V^! = \langle 1, 3, 3, 1, \dots \rangle$

$V^! = n^{\text{th}}$ row of Pascal's triangle.



$x^1 + x^2 + x^3$



Vector programs
can be implemented
by polynomials!

Programs -----> Polynomials

The vector $V! = \langle a_0, a_1, a_2, \dots \rangle$ will be represented by the polynomial:

$$P_V = \sum_{i=0}^{i=\infty} a_i X^i$$

Formal Power Series

The vector $V^i = \langle a_0, a_1, a_2, \dots \rangle$ will be represented by the **formal power series**:

$$P_V = \sum_{i=0}^{i=\infty} a_i X^i$$

$$V! = \langle a_0, a_1, a_2, \dots \rangle$$

$$P_V = \sum_{i=0}^{i=\infty} a_i X^i$$

$\langle 0 \rangle$ is represented by 0

$\langle k \rangle$ is represented by k

$V! + T!$ is represented by $(P_V + P_T)$

RIGHT($V!$) is represented by $(P_V X)$

Vector Programs

Example:

$V^1 := \langle 1 \rangle;$

$P_V := 1;$

Loop n times:

$V^i := V^i + \text{RIGHT}(V^i);$

$P_V := P_V + P_V X;$

$V^n = n^{\text{th}}$ row of Pascal's triangle.

Vector Programs

Example:

$V^1 := \langle 1 \rangle;$

$P_V := 1;$

Loop n times:

$V^i := V^i + \text{RIGHT}(V^i);$

$P_V := P_V (1 + X);$

$V^i = n^{\text{th}}$ row of Pascal's triangle.

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