15-251

Great Theoretical Ideas in Computer Science

Graphs

Lecture 18, October 23, 2008

Upcoming Events

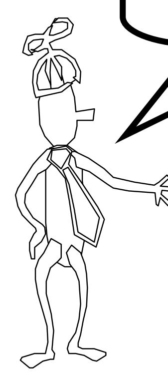
Review Session on Saturday
(5 pm, Wean 5409)

Test on Monday

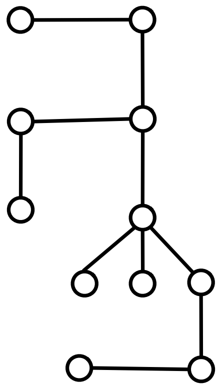
Election Day

What's a tree?

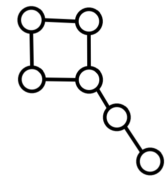
A tree is a connected graph with no cycles



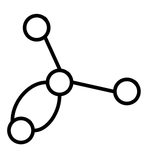
Tree



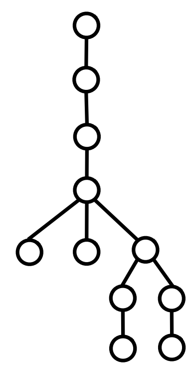
Not a Tree



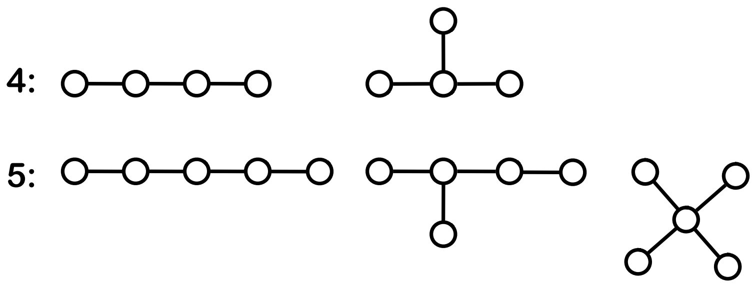
Not aTree



Tree



How Many n-Node Trees?



The Shy People Party

At the shy people party, people enter one-by-one, and as a person comes in, (s)he shakes hand with only one person already at the party.

Prove that at a shy party with n people (n >= 2), at least two people have shaken hands with only one other person.

We'll pass around a piece of paper. Draw a new 8-node tree, and put your name next to it. (There are 23 of them…)

The Shy People Party

Notation

In this lecture:

n will denote the number of nodes in a graph

e will denote the number of edges in a graph

To prove this, it suffices to show

1 => 2 => 3 => 4 => 5 => 1

Theorem: Let G be a graph with n nodes and e edges

The following are equivalent:

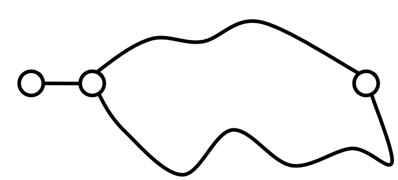
1. G is a tree (connected, acyclic)
2. Every two nodes of G are joined by a unique path
3. G is connected and $n = e + 1$
4. G is acyclic and $n = e + 1$
5. G is acyclic and if any two non-adjacent points are joined by a line, the resulting graph has exactly one cycle

1 => 2

1. G is a tree (connected, acyclic)
2. Every two nodes of G are joined by a unique path

Proof: (by contradiction)

Assume G is a tree that has two nodes connected by two different paths:



Then there exists a cycle!

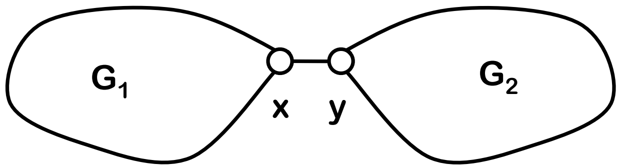
2 => 3 2. Every two nodes of G are joined by a unique path

3. G is connected and $n = e + 1$

Proof: (by induction)

Assume true for every graph with < n nodes

Let G have n nodes and let x and y be adjacent



Let n_1, e_1 be number of nodes and edges in G_1

Then $n = n_1 + n_2 = e_1 + e_2 + 2 = e + 1$

Corollary: Every nontrivial tree has at least two endpoints (points of degree 1)

Proof:

Assume all but one of the points in the tree have degree at least 2

In any graph, sum of the degrees = 2e

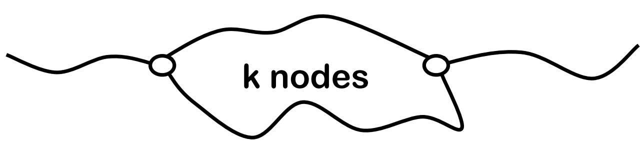
Then the total number of edges in the tree is at least $(2n-1)/2 = n - 1/2 > n - 1$

3 => 4 3. G is connected and $n = e + 1$

4. G is acyclic and $n = e + 1$

Proof: (by contradiction)

Assume G is connected with $n = e + 1$, and G has a cycle containing k nodes

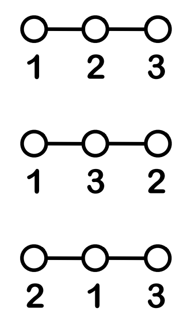


Note that the cycle has k nodes and k edges

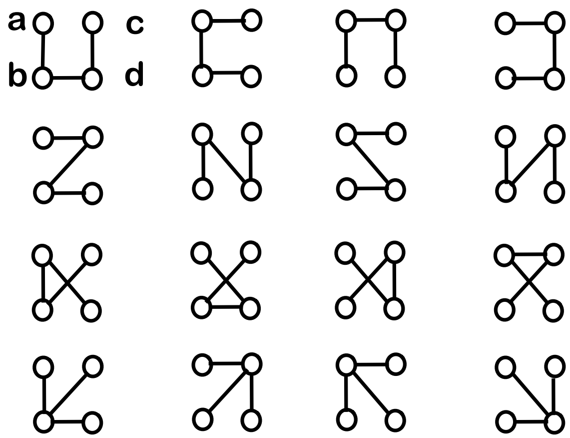
Start adding nodes and edges until you cover the whole graph

Number of edges in the graph will be at least n

How many labeled trees are there with three nodes?



How many labeled trees are there with four nodes?



How many labeled trees are there with n nodes?

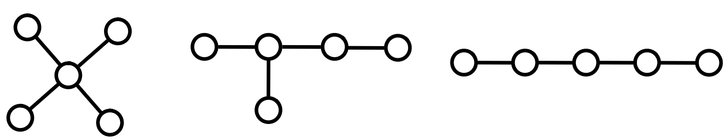
3 labeled trees with 3 nodes

16 labeled trees with 4 nodes

125 labeled trees with 5 nodes

n^{n-2} labeled trees with n nodes

How many labeled trees are there with five nodes?



5 labelings

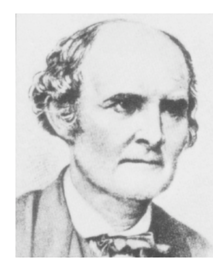
5 x 4 x 3 labelings

5!/ 2 labelings

125 labeled trees

Cayley's Formula

The number of labeled trees on n nodes is n^{n-2}



The proof will use the correspondence principle

Each labeled tree on n nodes

corresponds to

A sequence in $\{1,2,\ldots,n\}^{n-2}$ (that is, n-2 numbers, each in the range [1..n])

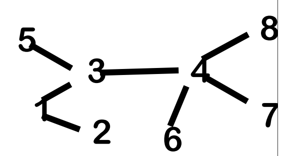
How to reconstruct the unique tree from a sequence S:

Let I = {1, 2, 3, …, n}

Loop until S is empty

- Let i = smallest # in I but not in S
- Let s = first label in sequence S
  - Add edge {i, s} to the tree
  - Delete i from I
  - Delete s from S

Add edge {a,b}, where I = {a,b}



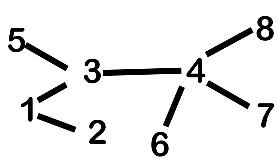
1 3 3 4 4 4

How to make a sequence from a tree?

Loop through i from 1 to n-2

- Let L be the degree-1 node with the lowest label
- Define the i^{th} element of the sequence as the label of the node adjacent to L
- Delete the node L from the tree

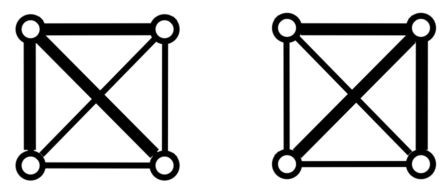
Example:



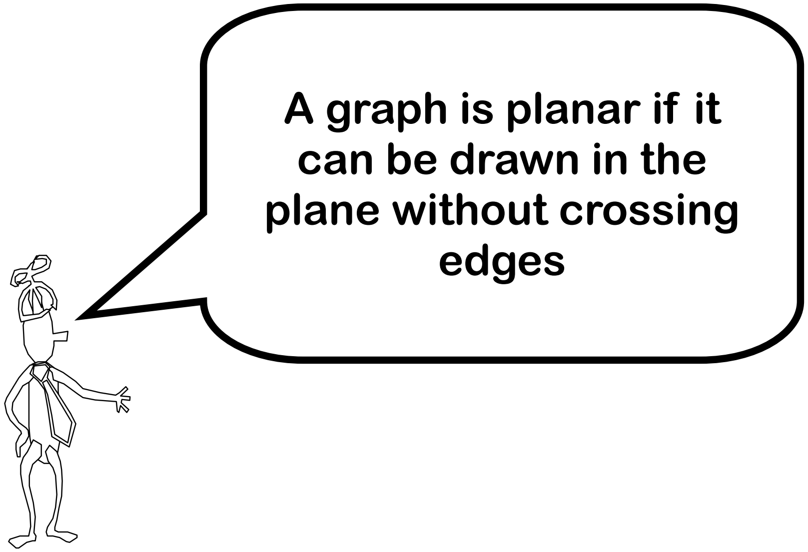
1 3 3 4 4 4

Spanning Trees

A spanning tree of a graph G is a tree that touches every node of G and uses only edges from G

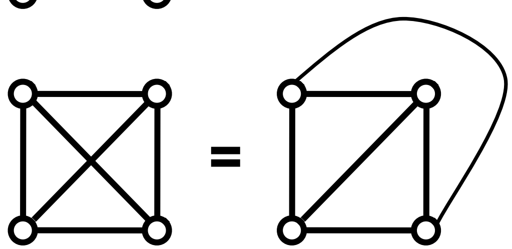


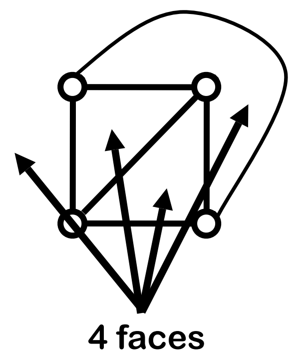
Every connected graph has a spanning tree



http://www.planarity.net

Examples of Planar Graphs





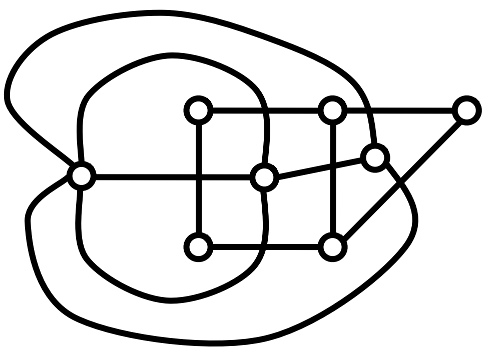
Faces

A planar graph splits the plane into disjoint faces

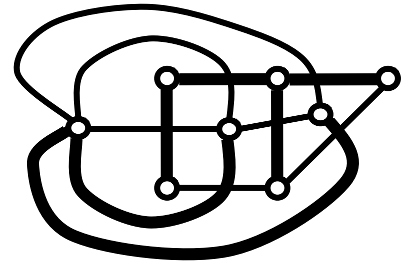
Euler's Formula

If G is a connected planar graph with n vertices, e edges and f faces, then $n - e + f = 2$

Rather than using induction, we'll use the important notion of the dual graph



Dual = put a node in every face, and an edge between every adjacent face



Let G* be the dual graph of G

Let T be a spanning tree of G

Let T* be the graph where there is an edge in dual graph for each edge in G – T

Then T* is a spanning tree for G*

$n = e_T + 1$

$f = e_{T^*} + 1$

$n + f = e_T + e_{T^*} + 2$

$= e + 2$

Corollary: Let G be a simple planar graph with n > 2 vertices. Then:

1. G has a vertex of degree at most 5
2. G has at most 3n – 6 edges

Proof of 1:

In any graph, (sum of degrees) = 2e

Assume all vertices have degree ≥ 6

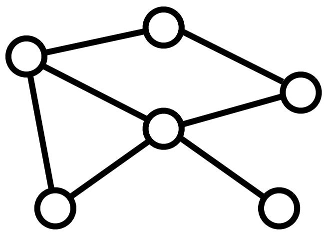
Then $e \geq 3n$

Furthermore, since G is simple, $3f \leq 2e$

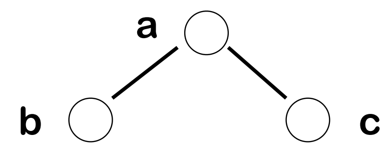
So $3n + 3f \leq 3e \Rightarrow 3(n-e+f) \leq 0$, contradiction.

Graph Coloring

A coloring of a graph is an assignment of a color to each vertex such that no neighboring vertices have the same color



Instructions	Live variables
	a
b = a+2	
	a,b
c = b*b	
	a,c
b = c+1	
	a,b
return a*b	



Graph Coloring

Arises surprisingly often in CS

Register allocation: assign temporary variables to registers for scheduling instructions. Variables that interfere, or are simultaneously active, cannot be assigned to the same register

Theorem: Every planar graph can be 6-colored

Proof Sketch (by induction):

Assume every planar graph with less than n vertices can be 6-colored

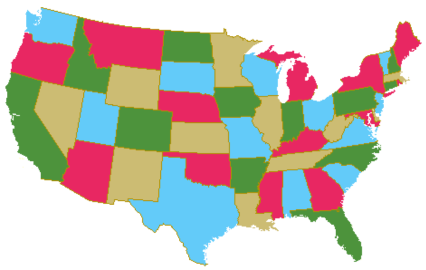
Assume G has n vertices

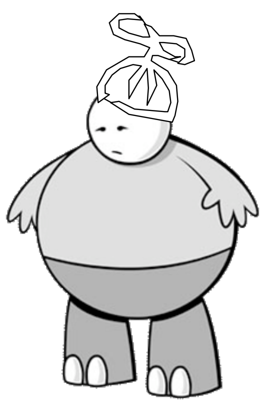
Since G is planar, it has some node v with degree at most 5

Remove v and color by Induction Hypothesis

Not too difficult to give an inductive proof of 5-colorability, using same fact that some vertex has degree ≤ 5

4-color theorem remains challenging!





Here's What You Need to Know…

Trees

- Counting Trees
- Different Characterizations
- Cayley's formula

Planar Graphs

- Definition
- Euler's Theorem
- Coloring Planar Graphs