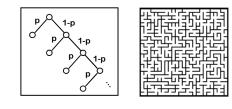
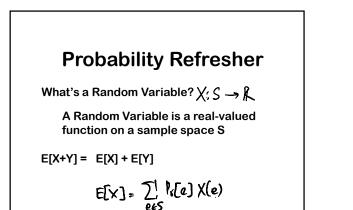
15-251

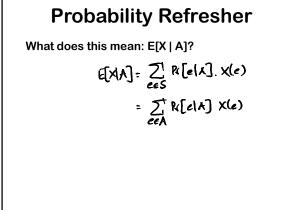
Great Theoretical Ideas in Computer Science

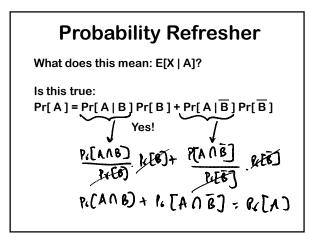
Infinite Sample spaces and Random Walks

Lecture 12 (October 4, 2007)







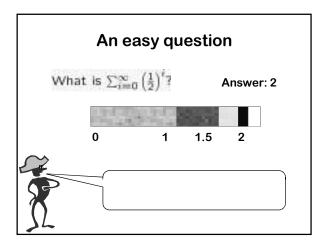


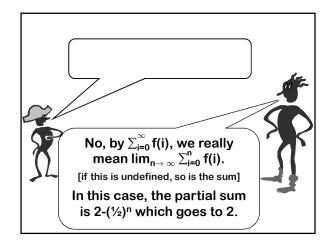


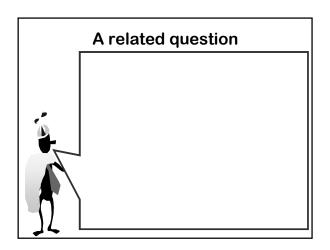
What does this mean: E[X | A]?

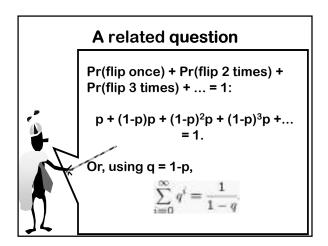
Is this true: Pr[A] = Pr[A | B] Pr[B] + Pr[A | B] Pr[B] Yes!

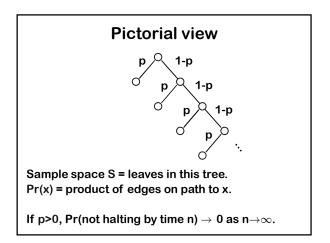
Similarly: E[X] = E[X|A] Pr[A] + E[X $|\overline{A}]$ Pr[\overline{A}]

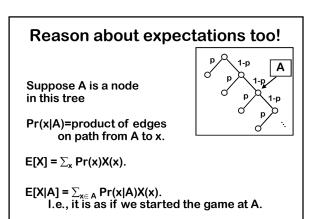


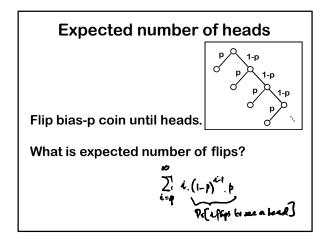


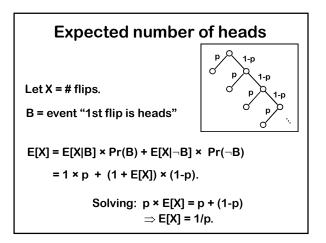












Infinite Probability spaces

Notice we are using infinite probability spaces here, but we really only defined things for <u>finite</u> spaces so far.

Infinite probability spaces can sometimes be weird.

Luckily, in CS we will almost always be looking at spaces that can be viewed as choice trees where $Pr(haven't halted by time t) \rightarrow 0 as t \rightarrow \infty$.

General picture

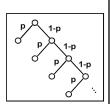
Let sample space S be leaves of a choice tree.

Let S_n = {leaves at depth $\leq n$ }.

For event A, let $A_n = A \cap S_n$.

If $\text{lim}_{n \rightarrow \infty} \text{Pr}(S_n) \text{=} 1$, can define:

 $Pr(A)=lim_{n\to\infty}Pr(A_n).$

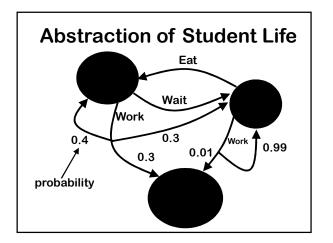


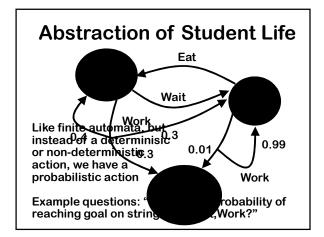
Setting that doesn't fit our model

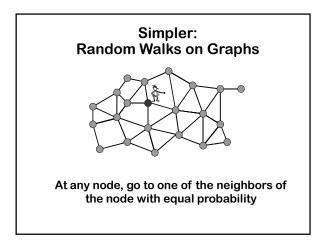
Event: "Flip coin until #heads > 2*#tails."

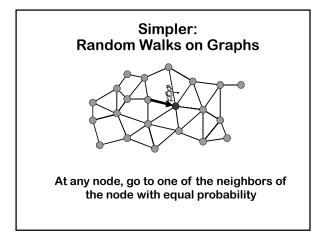
There's a reasonable chance this will <u>never</u> stop...

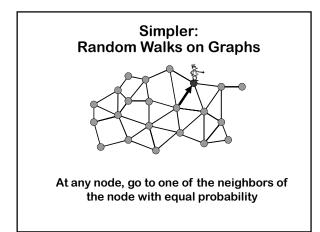
How to walk home drunk

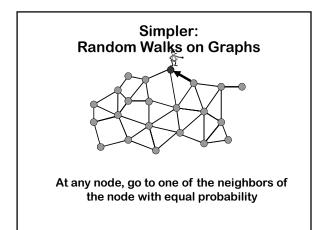


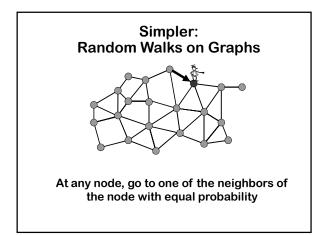


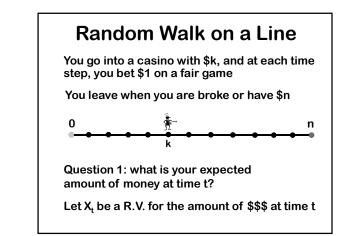


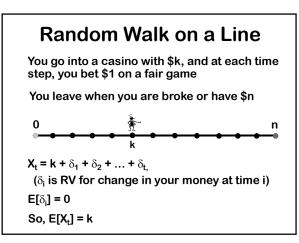


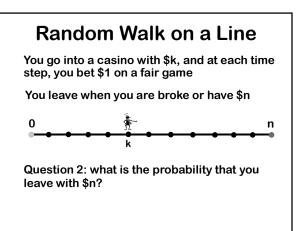


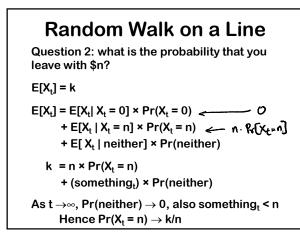


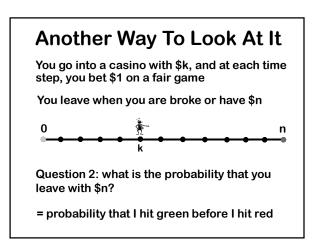


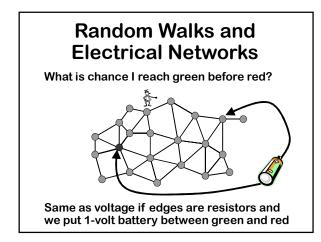


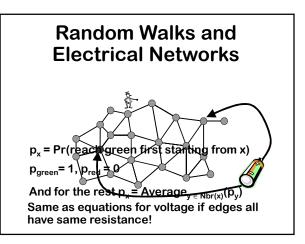


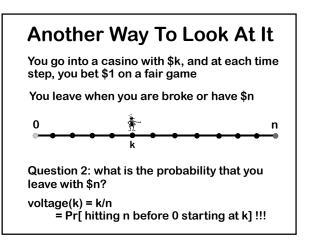


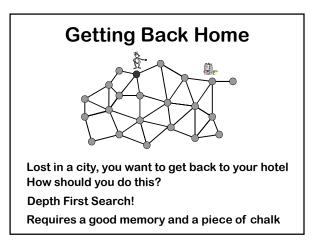


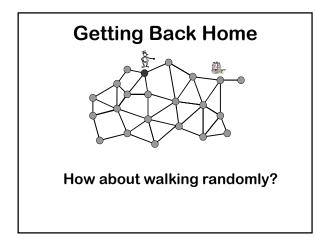


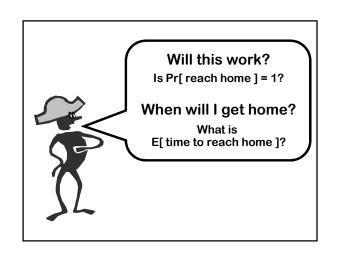


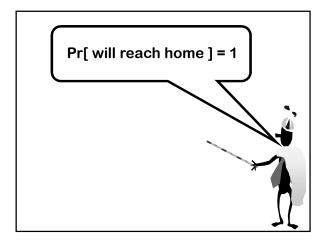












We Will Eventually Get Home

Look at the first n steps

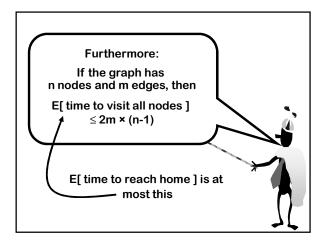
There is a non-zero chance \boldsymbol{p}_1 that we get home

Also, $p_1 \ge (1/n)^n$

Suppose we fail

Then, wherever we are, there is a chance $p_2 \geq (1/n)^n$ that we hit home in the next n steps from there

Probability of failing to reach home by time kn = $(1 - p_1)(1 - p_2) \dots (1 - p_k) \rightarrow 0$ as $k \rightarrow \infty$

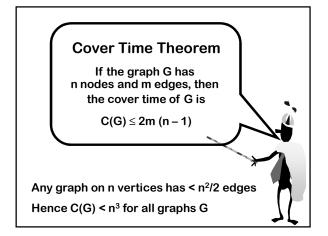


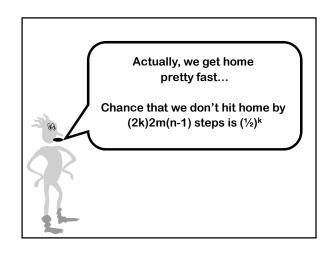
Cover Times

Cover time (from u) $C_u = E$ [time to visit all vertices | start at u]

Cover time of the graph C(G) = max_u { C_u }

(worst case expected time to see all vertices)







True or False:

If the average income of people is \$100 then more than 50% of the people can be earning more than \$200 each

False! else the average would be higher!!!

Markov's Inequality

If X is a non-negative r.v. with mean E[X], then

 $\Pr[X > 2 E[X]] \leq \frac{1}{2}$

 $\Pr[X > k E[X]] \leq 1/k$



Markov's Inequality

Non-neg random variable X has expectation A = E[X]

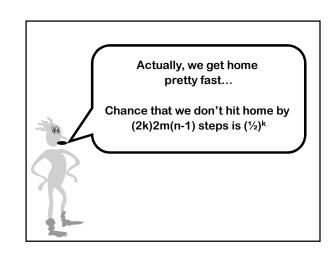
 \geq E[X | X > 2A] Pr[X > 2A] (since X is non-neg)

Also, E[X | X > 2A] > 2A

 \Rightarrow A \ge 2A × Pr[X > 2A]

$$\Rightarrow \frac{1}{2} \ge \Pr[X > 2A]$$

 $\Pr[X > k \times expectation] \le 1/k$



An Averaging Argument

Suppose I start at u

E[time to hit all vertices | start at $u] \leq C(G)$

Hence, by Markov's Inequality:

Pr[time to hit all vertices > 2C(G) | start at u] $\leq \frac{1}{2}$

So Let's Walk Some Mo!

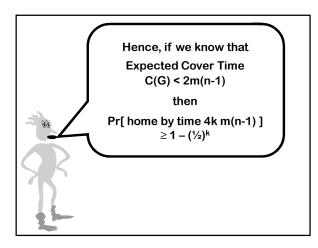
Pr [time to hit all vertices > 2C(G) | start at u] $\leq \frac{1}{2}$

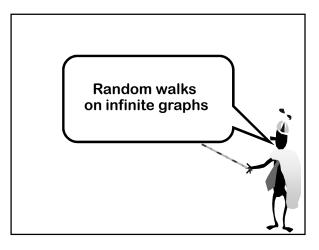
Suppose at time 2C(G), I'm at some node with more nodes still to visit

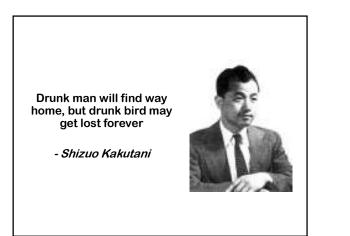
Pr [haven't hit all vertices in 2C(G) <u>more</u> time | start at v] $\leq \frac{1}{2}$

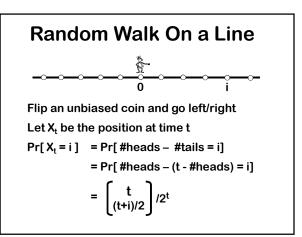
Chance that you failed both times $\leq \frac{1}{4} = (\frac{1}{2})^2$

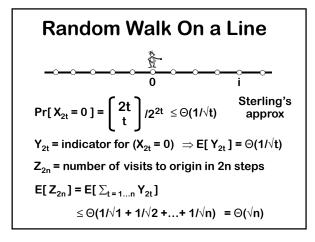
Hence, Pr[havent hit everyone in time k × 2C(G)] \leq (1/2)^k



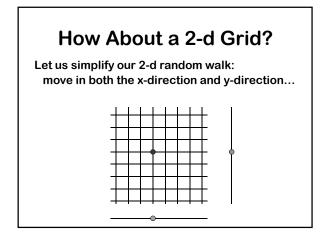


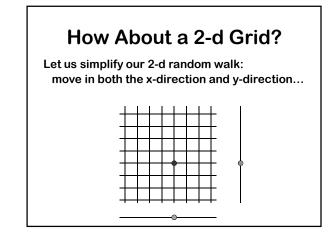


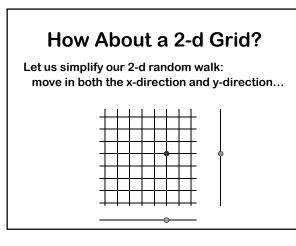






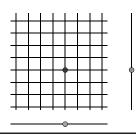


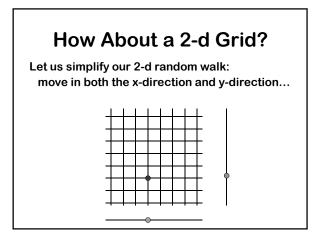




How About a 2-d Grid?

Let us simplify our 2-d random walk: move in both the x-direction and y-direction...





In The 2-d Walk

Returning to the origin in the grid ⇔ both "line" random walks return to their origins

Pr[visit origin at time t] = $\Theta(1/\sqrt{t}) \times \Theta(1/\sqrt{t})$ = $\Theta(1/t)$

E[# of visits to origin by time n] $= <math>\Theta(1/1 + 1/2 + 1/3 + ... + 1/n) = \Theta(\log n)$

But In 3D

Pr[visit origin at time t] = $\Theta(1/\sqrt{t})^3 = \Theta(1/t^{3/2})$ lim_{n→∞} E[# of visits by time n] < K (constant) Hence Pr[never return to origin] > 1/K

