

15-213

“The course that gives CMU its Zip!”

Code Optimization I

Feb. 12, 2008

Topics

- Machine-Independent Optimizations
 - Basic optimizations
 - Optimization blockers

Harsh Reality

There's more to performance than asymptotic complexity

Constant factors matter too!

- Easily see 10:1 performance range depending on how code is written
- Must optimize at multiple levels:
 - algorithm, data representations, procedures, and loops

Must understand system to optimize performance

- How programs are compiled and executed
- How to measure program performance and identify bottlenecks
- How to improve performance without destroying code modularity and generality

Optimizing Compilers

Provide efficient mapping of program to machine

- register allocation
- code selection and ordering (scheduling)
- dead code elimination
- eliminating minor inefficiencies

Don't (usually) improve asymptotic efficiency

- up to programmer to select best overall algorithm
- big-O savings are (often) more important than constant factors
 - but constant factors also matter

Have difficulty overcoming “optimization blockers”

- potential memory aliasing
- potential procedure side-effects

Limitations of Optimizing Compilers

Operate under fundamental constraint

- Must not cause any change in program behavior under any possible condition
- Often prevents it from making optimizations when would only affect behavior under pathological conditions.

Behavior that may be obvious to the programmer can be obfuscated by languages and coding styles

- e.g., Data ranges may be more limited than variable types suggest

Most analysis is performed only within procedures

- Whole-program analysis is too expensive in most cases

Most analysis is based only on *static* information

- Compiler has difficulty anticipating run-time inputs

When in doubt, the compiler must be conservative

Machine-Independent Optimizations

Optimizations that you or the compiler should do regardless of processor / compiler

Code Motion

- Reduce frequency with which computation performed
 - If it will always produce same result
 - Especially moving code out of loop

```
void set_row(double *a, double *b,
             long i, long n)
{
    long j;
    for (j = 0; j < n; j++)
        a[n*i+j] = b[j];
}
```



```
long j;
int ni = n*i;
for (j = 0; j < n; j++)
    a[ni+j] = b[j];
```

Compiler-Generated Code Motion

```
void set_row(double *a, double *b,
            long i, long n)
{
    long j;
    for (j = 0; j < n; j++)
        a[n*i+j] = b[j];
}
```

```
long j;
long ni = n*i;
double *rowp = a+ni;
for (j = 0; j < n; j++)
    *rowp++ = b[j];
```

Where are the FP operations?

```
set_row:
    xorl    %r8d, %r8d          # j = 0
    cmpq    %rcx, %r8          # j:n
    jge     .L7                # if >= goto done
    movq    %rcx, %rax          # n
    imulq    %rdx, %rax         # n*i outside of inner loop
    leaq     (%rdi,%rax,8), %rdx # rowp = A + n*i*8
.L5:
    movq     (%rsi,%r8,8), %rax # t = b[j]
    incq     %r8                # j++
    movq     %rax, (%rdx)       # *rowp = t
    addq     $8, %rdx           # rowp++
    cmpq     %rcx, %r8          # j:n
    jl      .L5                # if < goto loop
.L7:
    rep ; ret                   # done:
                                # return
```

Reduction in Strength

- Replace costly operation with simpler one

- Shift, add instead of multiply or divide

$16 * x \quad \rightarrow \quad x \ll 4$

- Utility machine dependent
- Depends on cost of multiply or divide instruction
- On Pentium IV, integer multiply requires 10 CPU cycles

- Recognize sequence of products

```
for (i = 0; i < n; i++)  
    for (j = 0; j < n; j++)  
        a[n*i + j] = b[j];
```



```
int ni = 0;  
for (i = 0; i < n; i++) {  
    for (j = 0; j < n; j++)  
        a[ni + j] = b[j];  
    ni += n;  
}
```

Share Common Subexpressions

- Reuse portions of expressions
- Compilers often not very sophisticated in exploiting arithmetic properties

```
/* Sum neighbors of i,j */
up =    val[(i-1)*n + j  ];
down =  val[(i+1)*n + j  ];
left =  val[i*n          + j-1];
right = val[i*n          + j+1];
sum = up + down + left + right;
```

3 multiplications: $i*n$, $(i-1)*n$, $(i+1)*n$

```
int inj = i*n + j;
up =    val[inj - n];
down =  val[inj + n];
left =  val[inj - 1];
right = val[inj + 1];
sum = up + down + left + right;
```

1 multiplication: $i*n$

```
leaq    1(%rsi), %rax    # i+1
leaq    -1(%rsi), %r8    # i-1
imulq   %rcx, %rsi       # i*n
imulq   %rcx, %rax       # (i+1)*n
imulq   %rcx, %r8       # (i-1)*n
addq    %rdx, %rsi       # i*n+j
addq    %rdx, %rax       # (i+1)*n+j
addq    %rdx, %r8       # (i-1)*n+j
```

```
imulq   %rcx, %rsi      # i*n
addq    %rdx, %rsi      # i*n+j
movq    %rsi, %rax      # i*n+j
subq    %rcx, %rax      # i*n+j-n
leaq    (%rsi,%rcx), %rcx # i*n+j+n
```


Optimization Blocker #1: Procedure Calls

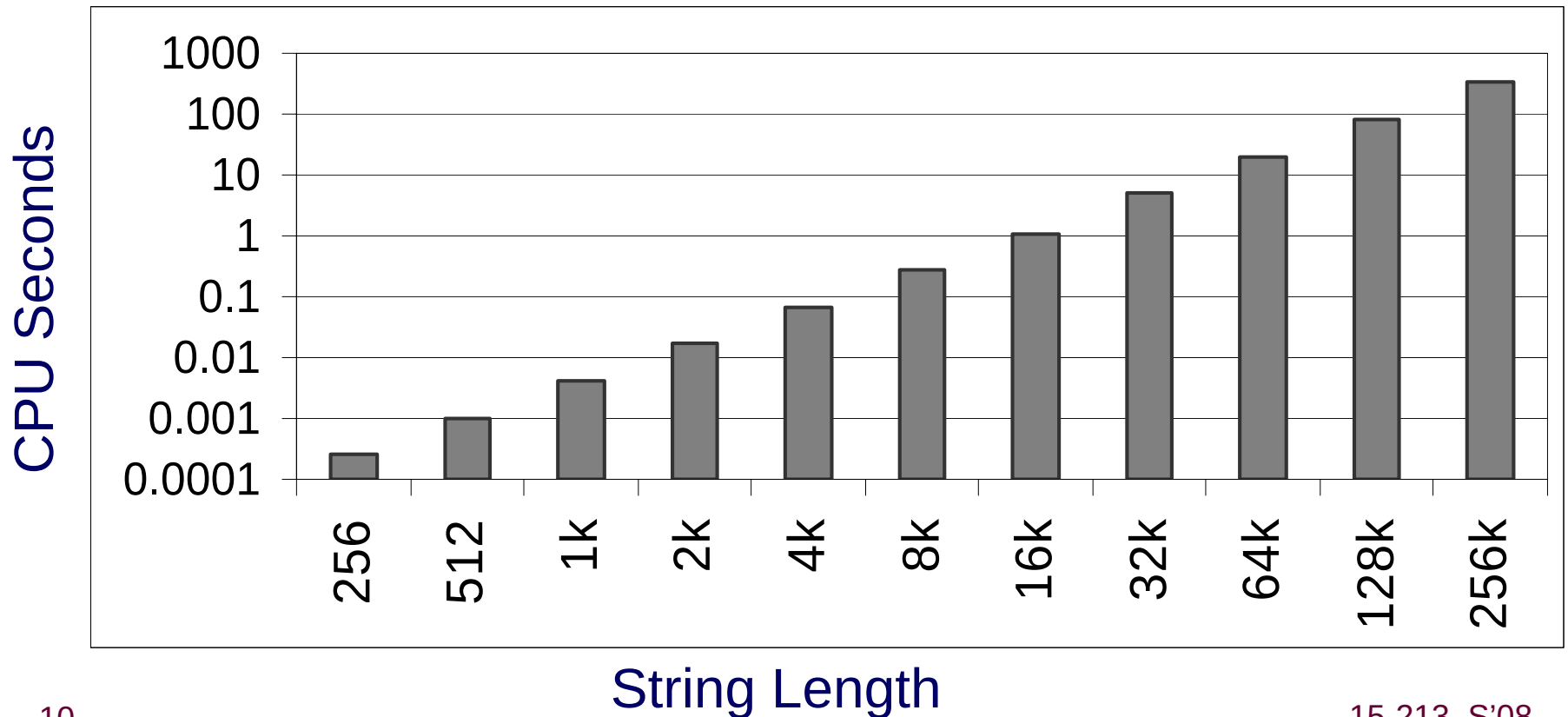
Procedure to Convert String to Lower Case

```
void lower(char *s)
{
    int i;
    for (i = 0; i < strlen(s); i++)
        if (s[i] >= 'A' && s[i] <= 'Z')
            s[i] -= ('A' - 'a');
}
```

- Extracted from 213 lab submissions, Fall, 1998

Lower Case Conversion Performance

- Time quadruples when double string length
- Quadratic performance



Convert Loop To Goto Form

```
void lower(char *s)
{
    int i = 0;
    if (i >= strlen(s))
        goto done;
loop:
    if (s[i] >= 'A' && s[i] <= 'Z')
        s[i] -= ('A' - 'a');
    i++;
    if (i < strlen(s))
        goto loop;
done:
}
```

- `strlen` executed every iteration

Calling Strlen

```
/* My version of strlen */
size_t strlen(const char *s)
{
    size_t length = 0;
    while (*s != '\0') {
        s++;
        length++;
    }
    return length;
}
```

Strlen performance

- Only way to determine length of string is to scan its entire length, looking for null character.

Overall performance, string of length N

- N calls to strlen
- Require times N, N-1, N-2, ..., 1
- Overall $O(N^2)$ performance

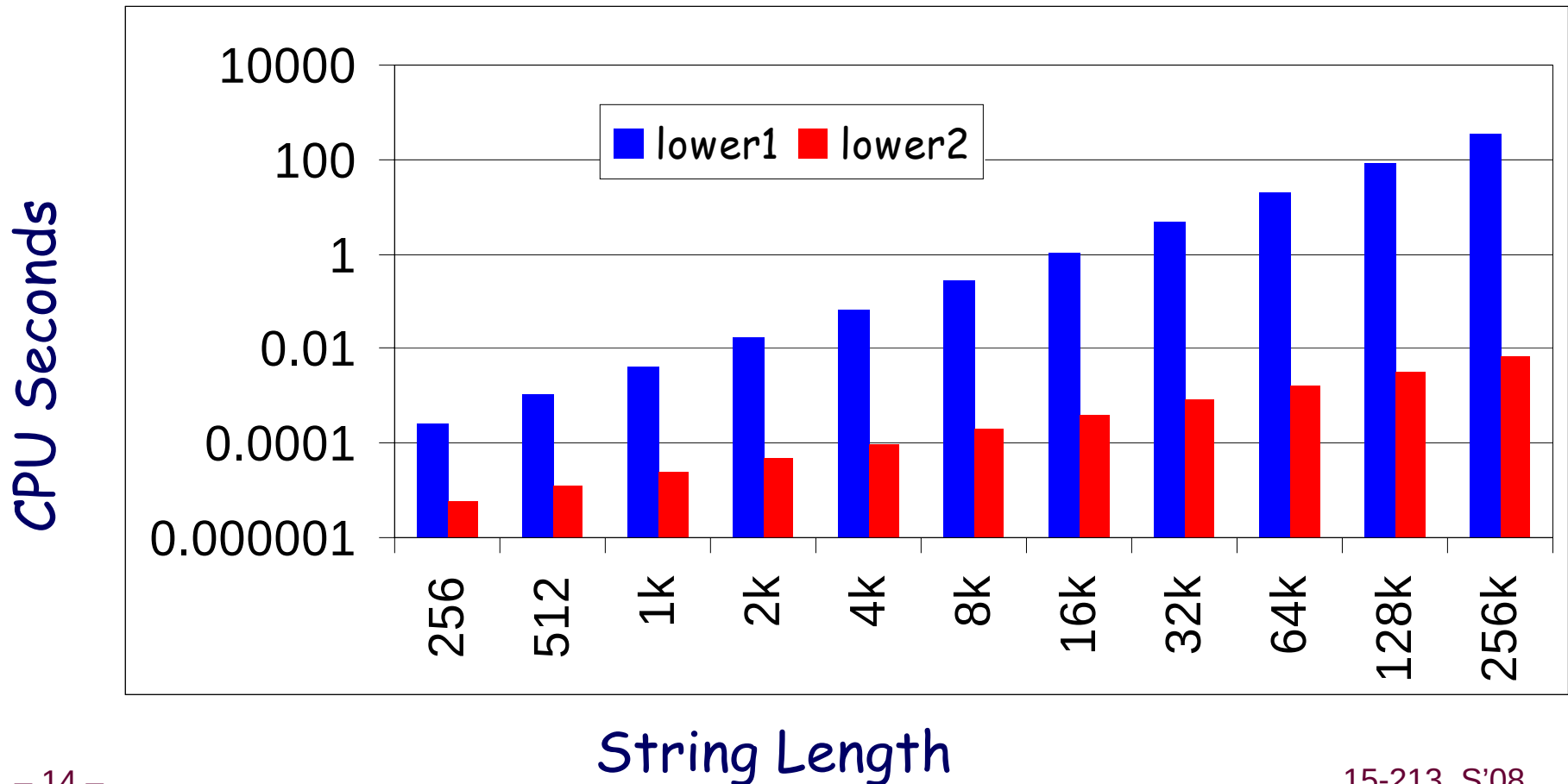
Improving Performance

```
void lower(char *s)
{
    int i;
    int len = strlen(s);
    for (i = 0; i < len; i++)
        if (s[i] >= 'A' && s[i] <= 'Z')
            s[i] -= ('A' - 'a');
}
```

- Move call to `strlen` outside of loop
- Since result does not change from one iteration to another
- Form of code motion

Lower Case Conversion Performance

- Time doubles when double string length
- Linear performance of lower2



Optimization Blocker: Procedure Calls

Why couldn't compiler move strlen out of inner loop?

- Procedure may have side effects
 - Alters global state each time called
- Function may not return same value for given arguments
 - Depends on other parts of global state
 - Procedure lower could interact with strlen

Warning:

- Compiler treats procedure call as a black box
- Weak optimizations near them

Remedies:

- Use of inline functions
- Do your own code motion

```
int lencnt = 0;
size_t strlen(const char *s)
{
    size_t length = 0;
    while (*s != '\0') {
        s++; length++;
    }
    lencnt += length;
    return length;
}
```

Memory Matters

```
/* Sum rows is of n X n matrix a
   and store in vector b */
void sum_rows1(double *a, double *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
        b[i] = 0;
        for (j = 0; j < n; j++)
            b[i] += a[i*n + j];
    }
}
```

```
# sum_rows1 inner loop
.L53:
    addsd    (%rcx), %xmm0           # FP add
    addq     $8, %rcx
    decq     %rax
    movsd    %xmm0, (%rsi,%r8,8)     # FP store
    jne      .L53
```

- Code updates `b[i]` on every iteration
- Why couldn't compiler optimize this away?

Memory Aliasing

```
/* Sum rows is of n X n matrix a
   and store in vector b */
void sum_rows1(double *a, double *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
        b[i] = 0;
        for (j = 0; j < n; j++)
            b[i] += a[i*n + j];
    }
}
```

```
double A[9] =
{ 0, 1, 2,
  4, 8, 16},
  32, 64, 128};

double B[3] = A+3;

sum_rows1(A, B, 3);
```

Value of B:

init: [4, 8, 16]

i = 0: [3, 8, 16]

i = 1: [3, 22, 16]

i = 2: [3, 22, 224]

- Code updates `b[i]` on every iteration
- Must consider possibility that these updates will affect program behavior

Removing Aliasing

```
/* Sum rows is of n X n matrix a
   and store in vector b */
void sum_rows2(double *a, double *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
        double val = 0;
        for (j = 0; j < n; j++)
            val += a[i*n + j];
        b[i] = val;
    }
}
```

```
# sum_rows2 inner loop
.L66:
    addsd    (%rcx), %xmm0    # FP Add
    addq     $8, %rcx
    decq     %rax
    jne      .L66
```

- No need to store intermediate results

Unaliased Version

```
/* Sum rows is of n X n matrix a
   and store in vector b */
void sum_rows2(double *a, double *b, long n) {
    long i, j;
    for (i = 0; i < n; i++) {
        double val = 0;
        for (j = 0; j < n; j++)
            val += a[i*n + j];
        b[i] = val;
    }
}
```

```
double A[9] =
{ 0, 1, 2,
  4, 8, 16},
  32, 64, 128};

double B[3] = A+3;

sum_rows1(A, B, 3);
```

Value of B:

init: [4, 8, 16]

i = 0: [3, 8, 16]

i = 1: [3, 27, 16]

i = 2: [3, 27, 224]

■ Aliasing still creates interference

Optimization Blocker: Memory Aliasing

Aliasing

- Two different memory references specify single location
- Easy to have happen in C
 - Since allowed to do address arithmetic
 - Direct access to storage structures
- Get in habit of introducing local variables
 - Accumulating within loops
 - **Your way of telling compiler not to check for aliasing**

Machine-Independent Opt. Summary

Code Motion

- Compilers are good at this for simple loop/array structures
- Don't do well in the presence of procedure calls and memory aliasing

Reduction in Strength

- Shift, add instead of multiply or divide
 - *Compilers are (generally) good at this*
 - *Exact trade-offs machine-dependent*
- Keep data in registers (local variables) rather than memory
 - *Compilers are not good at this, since concerned with aliasing*
 - *Compilers do know how to allocate registers (no need for register declaration)*

Share Common Subexpressions

- Compilers have limited algebraic reasoning capabilities