

15-213

## Exceptional Control Flow II

March 14, 2002

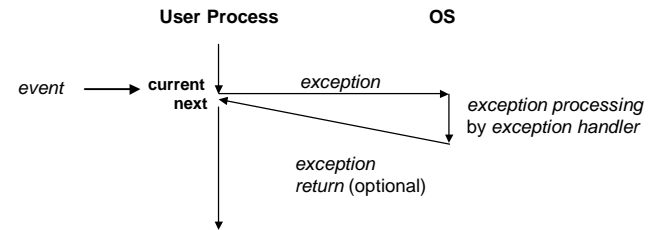
### Topics

- Exceptions
- Process context switches
- Reading: 8.5-8.8
- Problems: 8.19

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## Exceptions

An *exception* is a transfer of control to the OS in response to some *event* (i.e., change in processor state)



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## Role of Exceptions

### Error Handling

- Error conditions detected by hardware and/or OS
  - Divide by zero
  - Invalid pointer reference

### Getting Help from OS

- Initiate I/O operation
- Fetch memory page from disk

### Process Management

- Create illusion that running many programs and services simultaneously

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## The World of Multitasking

### System Runs Many Processes Concurrently

- **Process: executing program**
  - State consists of memory image + register values + program counter
- **Continually switches from one process to another**
  - Suspend process when it needs I/O resource or timer event occurs
  - Resume process when I/O available or given scheduling priority
- **Appears to user(s) as if all processes executing simultaneously**
  - Even though most systems can only execute one process at a time
  - Except possibly with lower performance than if running alone

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## Programmer's Model of Multitasking

### Basic Functions

- `fork()` spawns new process
  - Called once, returns twice
- `exit()` terminates own process
  - Called once, never returns
  - Puts it into "zombie" status
- `wait()` and `waitpid()` wait for and reap terminated children
- `exec1()` and `execve()` replace state of existing process with that of newly started program
  - Called once, never returns

### Programming Challenge

- Understanding the nonstandard semantics of the functions
- Avoiding improper use of system resources
  - Fewer safeguards provided

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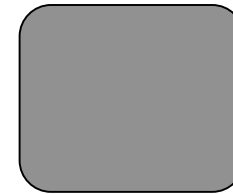
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## Fork Example #4

### Key Points

- Both parent and child can continue forking

```
void fork4()
{
    printf("L0\n");
    if (fork() != 0) {
        printf("L1\n");
        if (fork() != 0) {
            printf("L2\n");
            fork();
        }
    }
    printf("Bye\n");
}
```



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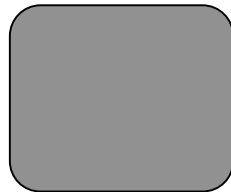
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## Fork Example #5

### Key Points

- Both parent and child can continue forking

```
void fork5()
{
    printf("L0\n");
    if (fork() == 0) {
        printf("L1\n");
        if (fork() == 0) {
            printf("L2\n");
            fork();
        }
    }
    printf("Bye\n");
}
```



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## Zombie Example

```
void fork7()
{
    if (fork() == 0) {
        /* Child */
        printf("Terminating Child, PID = %d\n",
            getpid());
        exit(0);
    } else {
        printf("Running Parent, PID = %d\n",
            getpid());
        while (1)
            ; /* Infinite loop */
    }
}
```

```
linux> ./forks 7 &
[1] 6639
Running Parent, PID = 6639
Terminating Child, PID = 6640
linux> ps
  PID TTY          TIME CMD
 6585 ttyp9        00:00:00 tcsh
 6639 ttyp9        00:00:03 forks
 6640 ttyp9        00:00:00 forks <defunct>
 6641 ttyp9        00:00:00 ps
linux> kill 6639
[1] Terminated
linux> ps
  PID TTY          TIME CMD
 6585 ttyp9        00:00:00 tcsh
 6642 ttyp9        00:00:00 ps
```

- `ps` shows child process as "defunct"
- Killing parent allows child to be reaped

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## Nonterminating Child Example

```
linux> ./forks 8
Terminating Parent, PID = 6675
Running Child, PID = 6676
linux> ps
  PID TTY          TIME CMD
 6585 tttyp9      00:00:00 tcsh
 6676 tttyp9      00:00:06 forks
 6677 tttyp9      00:00:00 ps
linux> kill 6676
linux> ps
  PID TTY          TIME CMD
 6585 tttyp9      00:00:00 tcsh
 6678 tttyp9      00:00:00 ps
```

```
void fork8()
{
    if (fork() == 0) {
        /* Child */
        printf("Running Child, PID = %d\n",
            getpid());
        while (1)
            ; /* Infinite loop */
    } else {
        printf("Terminating Parent, PID = %d\n",
            getpid());
        exit(0);
    }
}
```

- ps shows child process as “defunct”
- Killing parent allows child to be reaped

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## Exec Example

### Task

- Sort a set of files
- E.g., ./sortfiles f1.txt f2.txt f3.txt
- Perform concurrently
  - Using Unix sort command
  - Commands of form sort f1.txt -o f1.txt

```
#include <stdlib.h>
#include <stdio.h>
#include <unistd.h>
#include <sys/types.h>
#include <wait.h>

int main(int argc, char *argv[])
{
    int cnt = invoke(argc, argv);
    complete(cnt);
    return 0;
}
```

### Steps

- Invoke a process for each file
- Complete by waiting for all processes to complete

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## Exec Example (cont.)

- Use fork and execl to spawn set of sorting processes

```
int invoke(int argc, char *argv[])
{
    int i;
    for (i = 1; i < argc; i++) {
        /* Fork off a new process */
        if (fork() == 0) {
            /* Child: Invoke sort program */
            printf("Process %d sorting file %s\n", getpid(), argv[i]);
            if (execl("/bin/sort", "sort",
                argv[i], "-o", argv[i], 0) < 0) {
                perror("sort");
                exit(1);
            }
            /* Never reach this point */
        }
    }
    return argc-1;
}
```

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## Exec Example (cont.)

- Use wait to wait for and reap terminating children

```
void complete(int cnt)
{
    int i, child_status;
    for (i = 0; i < cnt; i++) {
        pid_t wpid = wait(&child_status);
        if (WIFEXITED(child_status))
            printf("Process %d completed with status %d\n",
                wpid, WEXITSTATUS(child_status));
        else
            printf("Process %d terminated abnormally\n", wpid);
    }
}
```

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## Signals

### Signals

- Software events generated by OS and processes
  - an OS abstraction for exceptions and interrupts
- Sent from the kernel or a process to other processes.
- Different signals are identified by small integer ID's
- Only information in a signal is its ID and the fact that it arrived.

Num.	Name	Default	Description
2	SIGINT	Terminate	Interrupt from keyboard (ctrl-c)
9	SIGKILL	Terminate	Kill program (cannot override or ignore)
11	SIGSEGV	Terminate & Dump	Segmentation violation
14	SIGALRM	Terminate	Timer signal
17	SIGCHLD	Ignore	Child stopped or terminated

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## Sending Signals

### Unix kill Program

- Sends arbitrary signal to process
- e.g., `/bin/kill -s 9 pid`
  - sends SIGKILL to specified process

### Function kill

- Send signal to another process
- ```
kill(pid, signal)
```

```
linux> ./forks 8
Terminating Parent, PID = 6675
Running Child, PID = 6676
linux> ps
  PID TTY          TIME CMD
 6585 tttyp9    00:00:00 tcsh
 6676 tttyp9    00:00:06 forks
 6677 tttyp9    00:00:00 ps
linux> /bin/kill -s 9 6676
linux> ps
  PID TTY          TIME CMD
 6585 tttyp9    00:00:00 tcsh
 6678 tttyp9    00:00:00 ps
```

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## Kill Example

```
void fork12()
{
    pid_t pid[N];
    int i;
    int child_status;
    for (i = 0; i < N; i++)
        if ((pid[i] = fork()) == 0) {
            /* Child: Infinite Loop */
            while(1);
        }
    for (i = 0; i < N; i++) {
        printf("Killing process %d\n", pid[i]);
        kill(pid[i], SIGINT);
    }
    for (i = 0; i < N; i++) {
        pid_t wpid = wait(&child_status);
        if (WIFEXITED(child_status))
            printf("Child %d terminated with exit status %d\n",
                wpid, WEXITSTATUS(child_status));
        else
            printf("Child %d terminated abnormally\n", wpid);
    }
}
```

- Use kill to forcibly terminate children

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## Handling Signals

### Every Signal Type has Default Behavior

- Typically terminate or ignore

### Can Override by Declaring Special Signal Handler Function

- `signal(sig, handler)`
  - Indicates that signals of type `sig` should invoke function handler
  - Handler returns to point where exception occurred

```
void int_handler(int sig)
{
    printf("Process %d received signal %d\n", getpid(), sig);
    exit(0);
}

void fork13()
{
    pid_t pid[N];
    int i, child_status;
    signal(SIGINT, int_handler);
    . . .
}
```

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## Signal Handler Funkiness

### Signals are not Queued

```
int ccount = 0;
void child_handler(int sig)
{
    int child_status;
    pid_t pid = wait(&child_status);
    ccount--;
    printf("Received signal %d from process %dn",
        sig, pid);
}

void fork14()
{
    pid_t pid[N];
    int i, child_status;
    ccount = N;
    signal(SIGCHLD, child_handler);
    for (i = 0; i < N; i++)
        if ((pid[i] = fork()) == 0) {
            /* Child: Exit */
            exit(0);
        }
    while (ccount > 0)
        pause(); /* Suspend until signal occurs */
}
```

- For each signal type, just have single bit indicating whether or not signal has occurred
- Even if multiple processes have sent this signal

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## Living with Nonqueuing Signals

### Must Check for All Possible Signal Sources

- Typically loop with wait

```
void child_handler2(int sig)
{
    int child_status;
    pid_t pid;
    while ((pid = wait(&child_status)) > 0) {
        ccount--;
        printf("Received signal %d from process %dn", sig, pid);
    }
}

void fork15()
{
    ...
    signal(SIGCHLD, child_handler2);
    ...
}
```

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## A program that reacts to externally generated events (ctrl-c)

```
#include <stdlib.h>
#include <stdio.h>
#include <signal.h>

static void handler(int sig) {
    printf("You think hitting ctrl-c will stop the bomb?\n");
    sleep(2);
    printf("Well...\n");
    fflush(stdout);
    sleep(1);
    printf("OK\n");
    exit(0);
}

main() {
    signal(SIGINT, handler); /* installs ctrl-c handler */
    while(1) {
    }
}
```

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## A program that reacts to internally generated events

```
#include <stdio.h>
#include <signal.h>

int beeps = 0;

/* SIGALRM handler */
void handler(int sig) {
    printf("BEEP\n");
    fflush(stdout);

    if (++beeps < 5)
        alarm(1);
    else {
        printf("BOOM!\n");
        exit(0);
    }
}
```

```
main() {
    signal(SIGALRM, handler);
    alarm(1); /* send SIGALRM in
              1 second */

    while (1) {
        /* handler returns here */
    }
}
```

```
bass> a.out
BEEP
BEEP
BEEP
BEEP
BEEP
BOOM!
bass>
```

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## Nonlocal jumps: setjmp()/longjmp()

**Powerful (but dangerous) user-level mechanism for transferring control to an arbitrary location.**

- controlled to way to break the procedure call/return discipline
- useful for error recovery

```
int setjmp(jmp_buf j)
```

- must be called before longjmp
- identifies a return site for a subsequent longjmp.
- Called once, returns one or more times

**Implementation:**

- remember where you are by storing the current register context, stack pointer, and PC value in jmp\_buf.
- return 0

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## setjmp/longjmp (cont)

```
void longjmp(jmp_buf j, int i)
```

- meaning:
  - return from the setjmp remembered by jump buffer j again...
  - ...this time returning i
- called after setjmp
- Called once, but never returns

**longjmp Implementation:**

- restore register context from jump buffer j
- set %eax (the return value) to i
- jump to the location indicated by the PC stored in jump buf j.

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## setjmp/longjmp example

```
#include <setjmp.h>
jmp_buf buf;

main() {
    if (setjmp(buf) != 0) {
        printf("back in main due to an error\n");
    } else {
        printf("first time through\n");
        p1(); /* p1 calls p2, which calls p3 */
    }
    ...
    p3() {
        <error checking code>
        if (error)
            longjmp(buf, 1)
    }
}
```

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## Putting it all together: A program that restarts itself when ctrl-c'd

```
#include <stdio.h>
#include <signal.h>
#include <setjmp.h>

sigjmp_buf buf;

void handler(int sig) {
    siglongjmp(buf, 1);
}

main() {
    signal(SIGINT, handler);

    if (!sigsetjmp(buf, 1))
        printf("starting\n");
    else
        printf("restarting\n");
}
```

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```
while(1) {
    sleep(1);
    printf("processing...\n");
}
```

```
base> a.out
starting
processing...
processing...
restarting ← Ctrl-c
processing...
processing...
restarting ← Ctrl-c
processing...
processing...
restarting ← Ctrl-c
processing...
```

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## Limitations of Long Jumps

### Works Within Stack Discipline

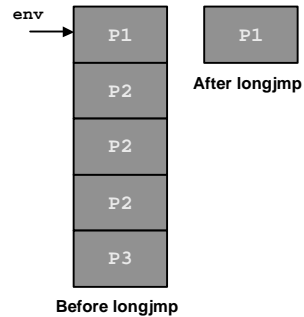
- Can only long jump to environment of function that has been called but not yet completed

```
jmp_buf env;

P1()
{
    if (setjmp(env)) {
        /* Long Jump to here */
    } else {
        P2();
    }
}

P2()
{
    . . . P2(); . . . P3();
}

P3()
{
    longjmp(env, 1);
}
```



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## Limitations of Long Jumps (cont.)

### Works Within Stack Discipline

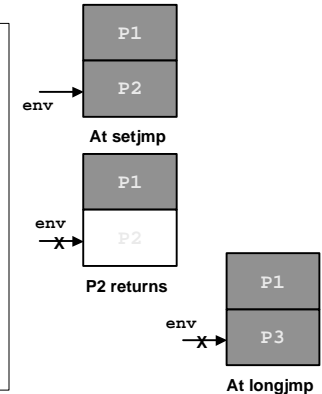
- Can only long jump to environment of function that has been called but not yet completed

```
jmp_buf env;

P1()
{
    P2(); P3();
}

P2()
{
    if (setjmp(env)) {
        /* Long Jump to here */
    }
}

P3()
{
    longjmp(env, 1);
}
```



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## Summary

### Signals Provide Process-Level Exception Handling

- Can generate with `kill`
- Can define effect by declaring signal handler

### Some Caveats

- Very high overhead
  - >10,000 clock cycles
  - Only use for exceptional conditions
- Don't have queues
  - Just one bit of status for each signal type

### Long Jumps Provide Exceptional Control Flow Within Process

- Within constraints of stack discipline

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