15-213

Structured Data II
Heterogenous Data
Feb. 15, 2000

Topics

- Structure Allocation
- Alignment
- Unions
- Byte Ordering
- Byte Operations
- IA32/Linux Memory Organization

Basic Data Types

Integral

- Stored & operated on in general registers
- Signed vs. unsigned depends on instructions used

Intel	GAS	Bytes	C
byte	b	1	[unsigned] char
word	w	2	[unsigned] short
double word	l	4	[unsigned] int, char *
quad word		8	

Floating Point

- Stored & operated on in floating point registers

Intel	GAS	Bytes	C
Single	s	4	float
Double	l	8	double
Extended		10	--

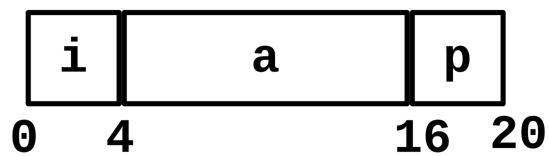
Structures

Concept

- Contiguously-allocated region of memory
- Refer to members within structure by names
- Members may be of different types

```
struct rec {
  int i;
  int a[3];
  int *p;
};
```

Memory Layout



Accessing Structure Member

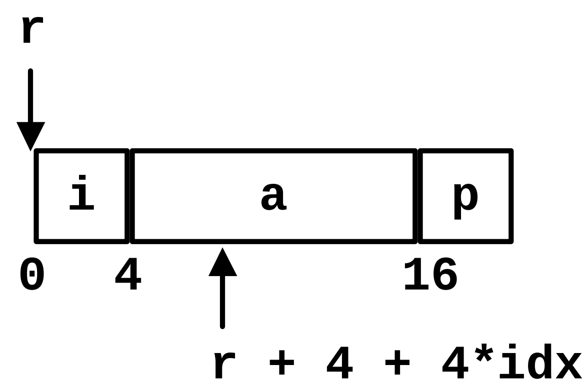
```
void
set_i(struct rec *r,
      int val)
{
  r->i = val;
}
```

Assembly

```
# %eax = val
# %edx = r
movl %eax,(%edx)    # Mem[r] = val
```

Generating Pointer to Structure Member

```
struct rec {
  int i;
  int a[3];
  int *p;
};
```



Generating Pointer to Array Element

- Offset of each structure member determined at compile time

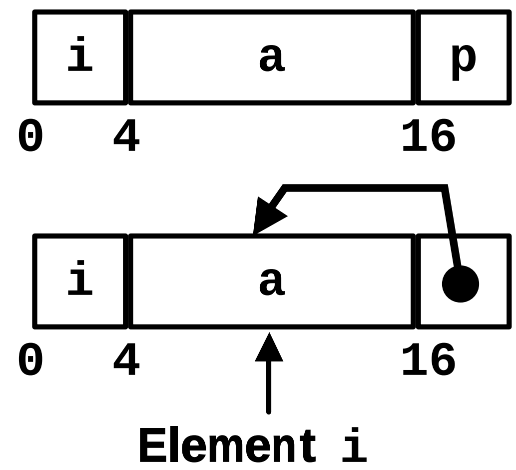
```
int *
find_a
 (struct rec *r, int idx)
{
  return &r->a[idx];
}
```

```
# %ecx = idx
# %edx = r
leal 0(,%ecx,4),%eax   # 4*idx
leal 4(%eax,%edx),%eax # r+4*idx+4
```

Structure Referencing (Cont.)

C Code

```
struct rec {
  int i;
  int a[3];
  int *p;
};
```




```
void
set_p(struct rec *r)
{
  r->p =
   &r->a[r->i];
}
```

```
 # %edx = r
movl (%edx),%ecx        # r->i
leal 0(,%ecx,4),%eax    # 4*(r->i)
leal 4(%edx,%eax),%eax  # r+4+4*(r->i)
movl %eax,16(%edx)      # Update r->p
```

Alignment

Aligned Data

- **Primitive data type requires K bytes**
- **Address must be multiple of K**
- **Required on some machines; advised on IA32**
  - treated differently by Linux and Windows!

Motivation for Aligning Data

- **Memory accessed by (aligned) double or quad-words**
  - Inefficient to load or store datum that spans quad word boundaries
  - Virtual memory very tricky when datum spans 2 pages

Compiler

- **Inserts gaps in structure to ensure correct alignment of fields**

Specific Cases of Alignment

Size of Primitive Data Type:

- **1 byte** (e.g., `char`)
  - no restrictions on address
- **2 bytes** (e.g., `short`)
  - lowest 1 bit of address must be 0_2
- **4 bytes** (e.g., `int`, `float`, `char *`, etc.)
  - lowest 2 bits of address must be 00_2
- **8 bytes** (e.g., `double`)
  - Windows (and most other OS's & instruction sets):
    - » lowest 3 bits of address must be 000_2
  - Linux:
    - » lowest 2 bits of address must be 00_2
    - » i.e. treated the same as a 4 byte primitive data type

Satisfying Alignment with Structures

Offsets Within Structure

- Must satisfy element's alignment requirement

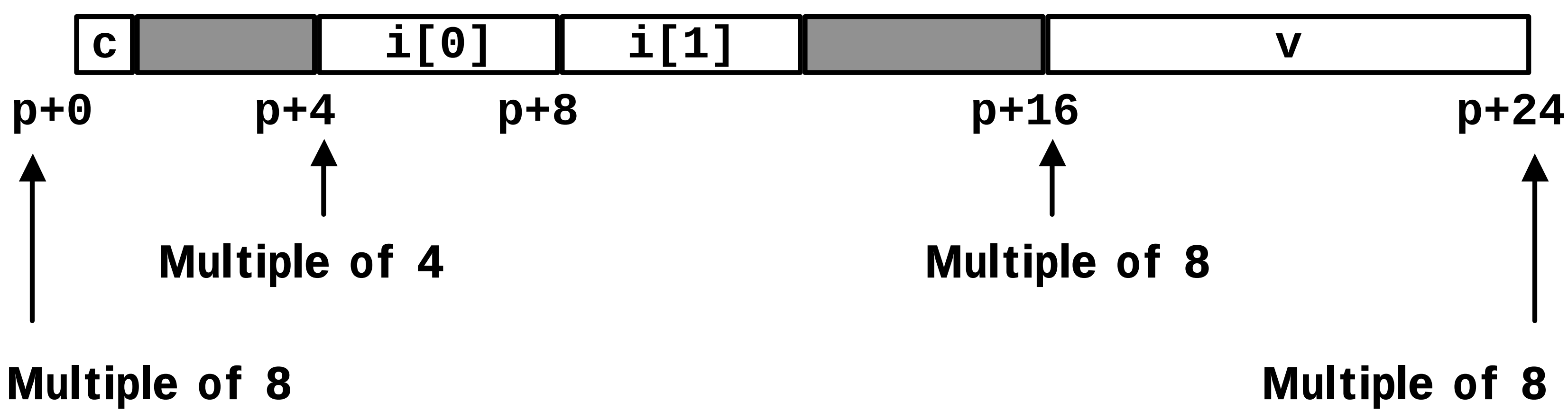
Overall Structure Placement

- Each structure has alignment requirement K
  - Largest alignment of any element
- Initial address & structure length must be multiples of K

Example (under Windows):

- K = 8, due to double element

```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

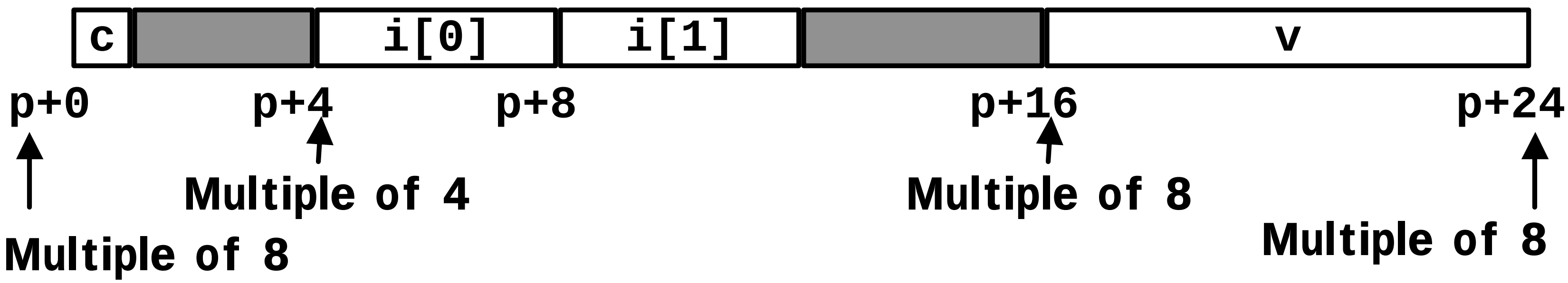


Linux vs. Windows

```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

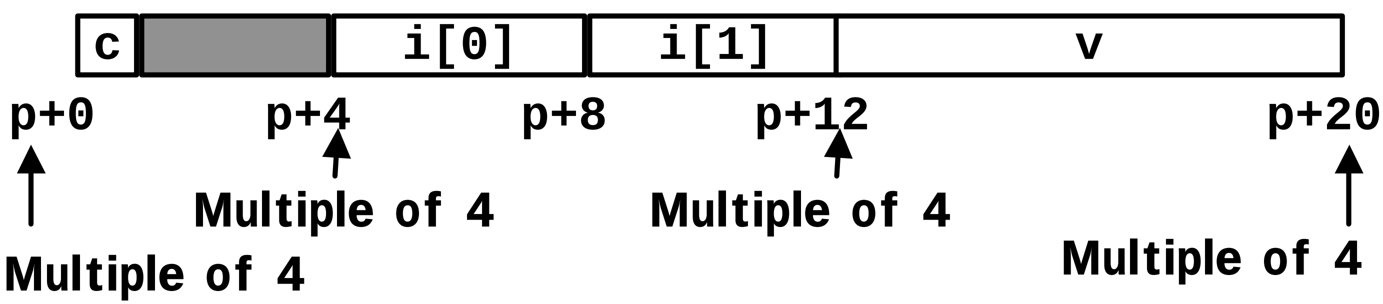
Windows (including Cygwin):

- K = 8, due to double element



Linux:

- K = 4; double treated like a 4-byte data type



Effect of Overall Alignment Requirement

```
struct S2 {
  double x;
  int i[2];
  char c;
} *p;
```

p must be multiple of:
- 8 for Windows
- 4 for Linux

x	i[0]	i[1]	c	

p+0 p+8 p+12 p+16 Windows: p+24
Linux: p+20

```
struct S3 {
  float x[2];
  int i[2];
  char c;
} *p;
```

p must be multiple of 4 (in either OS)

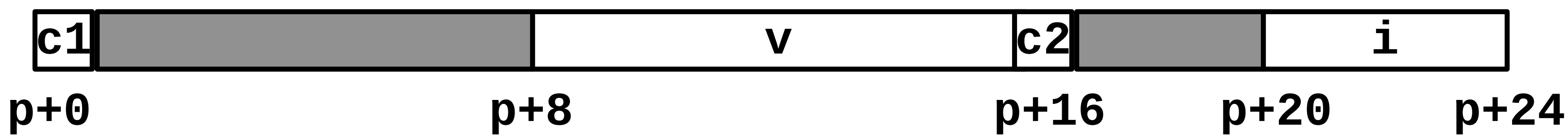
x[0]	x[1]	i[0]	i[1]	c	

p+0 p+4 p+8 p+12 p+16 p+20

Ordering Elements Within Structure

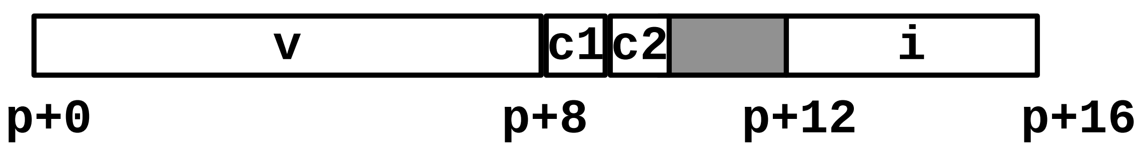
```
struct S4 {
  char c1;
  double v;
  char c2;
  int i;
} *p;
```

10 bytes wasted space in Windows




```
struct S5 {
  double v;
  char c1;
  char c2;
  int i;
} *p;
```

2 bytes wasted space

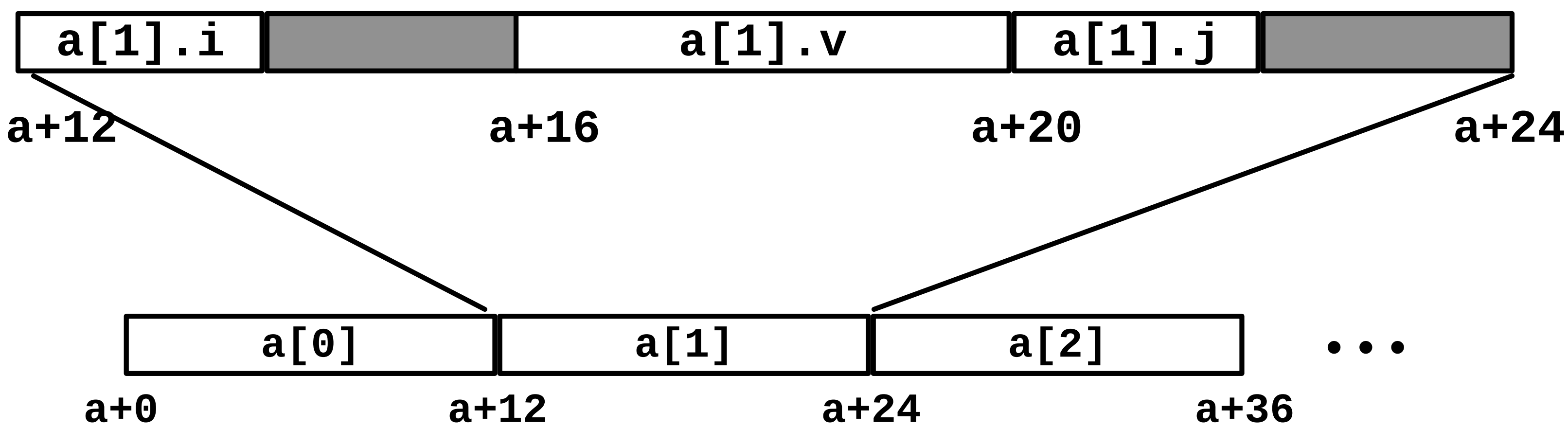


Arrays of Structures

Principle

- Allocated by repeating allocation for array type
- In general, may nest arrays & structures to arbitrary depth

```
struct S6 {
  short i;
  float v;
  short j;
} a[10];
```



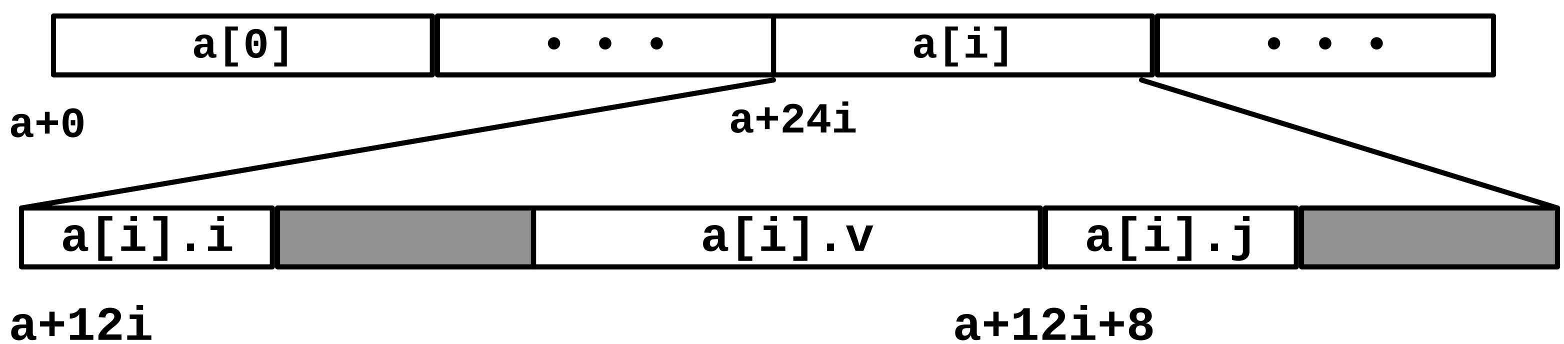
Accessing Element within Array

- **Compute offset to start of structure**
  - Compute 12**i* as 4*(*i*+2*i*)
- **Access element according to its offset within structure**
  - Offset by 8
  - Assembler gives displacement as a + 8
    - » Linker must set actual value

```
struct S6 {
  short i;
  float v;
  short j;
} a[10];
```

```
short get_j(int idx)
{
  return a[idx].j;
}
```

```
 # %eax = idx
 leal (%eax,%eax,2),%eax # 3*idx
 movswl a+8(,%eax,4),%eax
```

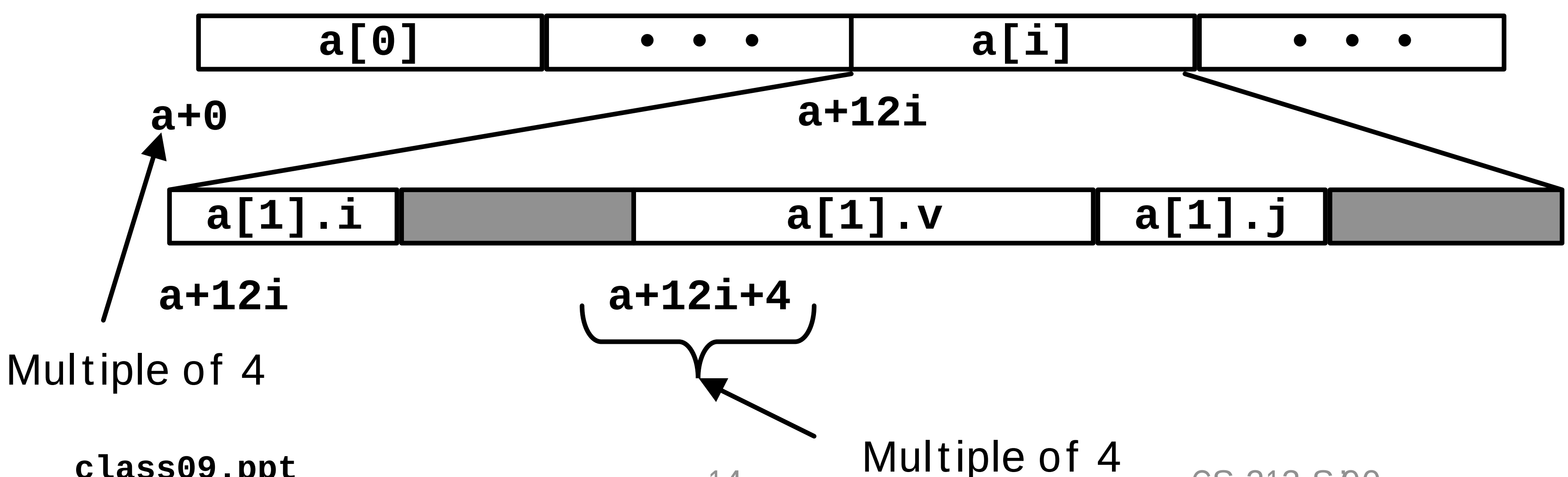


Satisfying Alignment within Structure

Achieving Alignment

- **Starting address of structure array must be multiple of worst-case alignment for any element**
  - **a** must be multiple of 4
- **Offset of element within structure must be multiple of element's alignment requirement**
  - **v**'s offset of 4 is a multiple of 4
- **Overall size of structure must be multiple of worst-case alignment for any element**
  - Structure padded with unused space to be 12 bytes

```
struct S6 {
  short i;
  float v;
  short j;
} a[10];
```

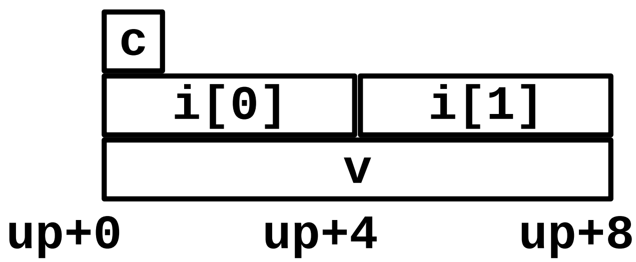


Union Allocation

Principles

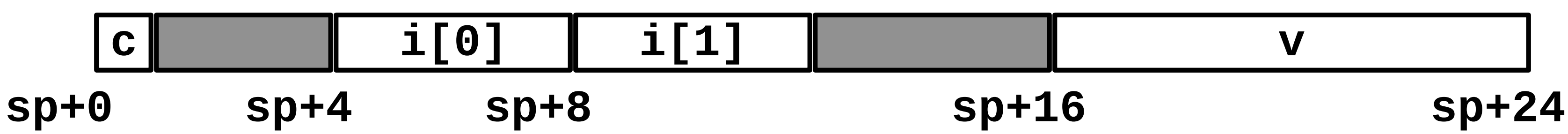
- Overlay union elements
- Allocate according to largest element
- Can only use one field at a time

```
union U1 {
  char c;
  int i[2];
  double v;
} *up;
```




```
struct S1 {
  char c;
  int i[2];
  double v;
} *sp;
```

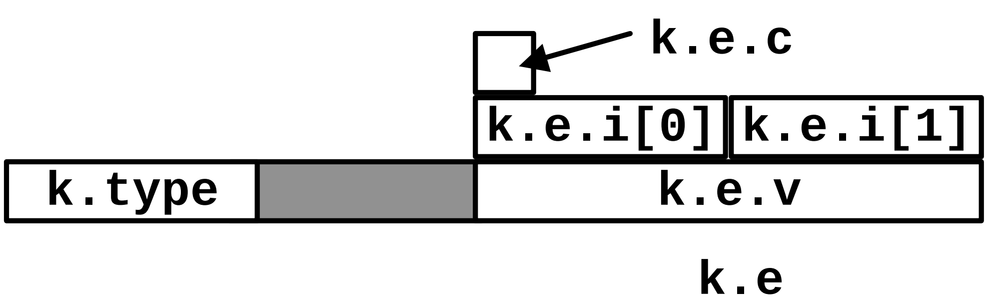
(Windows alignment)



Implementing "Tagged" Union

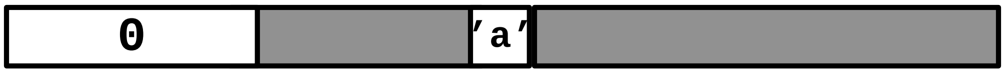
- Structure can hold 3 kinds of data
- Only one form at any given time
- Identify particular kind with flag type

```
typedef enum { CHAR, INT, DBL }
  utype;

typedef struct {
  utype type;
  union {
    char c;
    int i[2];
    double v;
  } e;
} store_ele, *store_ptr;

store_ele k;
```

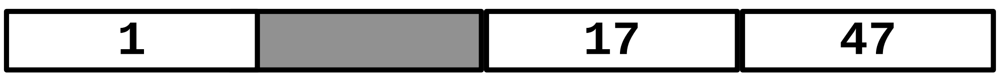


Using "Tagged" Union

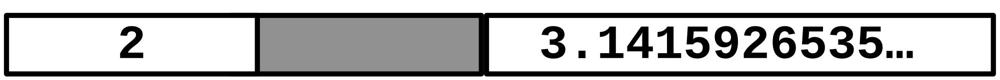
```
store_ele k1;
k1.type = CHAR;
k1.e.c = 'a';
```




```
store_ele k2;
k2.type = INT;
k2.e.i[0] = 17;
k2.e.i[1] = 47;
```

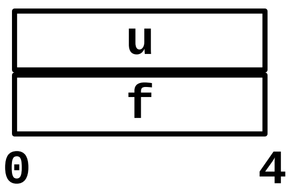



```
store_ele k3;
k3.type = DBL;
k1.e.v =
  3.14159265358979323846;
```



Using Union to Access Bit Patterns

```
typedef union {
  float f;
  unsigned u;
} bit_float_t;
```



- **Get direct access to bit representation of float**
- **`bit2float` generates float with given bit pattern**
  – NOT the same as **`(float) u`**
- **`float2bit` generates bit pattern from float**
  – NOT the same as **`(unsigned) f`**

```
float bit2float(unsigned u)
{
  bit_float_t arg;
  arg.u = u;
  return arg.f;
}
```

```
unsigned float2bit(float f)
{
  bit_float_t arg;
  arg.f = f;
  return arg.u;
}
```

Byte Ordering

Idea

- Long/quad words stored in memory as 4/8 consecutive bytes
- Which is most (least) significant?
- Can cause problems when exchanging binary data between machines

Big Endian

- Most significant byte has lowest address
- IBM 360/370, Motorola 68K, Sparc

Little Endian

- Least significant byte has lowest address
- Intel x86, Digital VAX

Byte Ordering Example

```
union {
  unsigned char c[8];
  unsigned short s[4];
  unsigned int i[2];
  unsigned long l[1];
} dw;
```

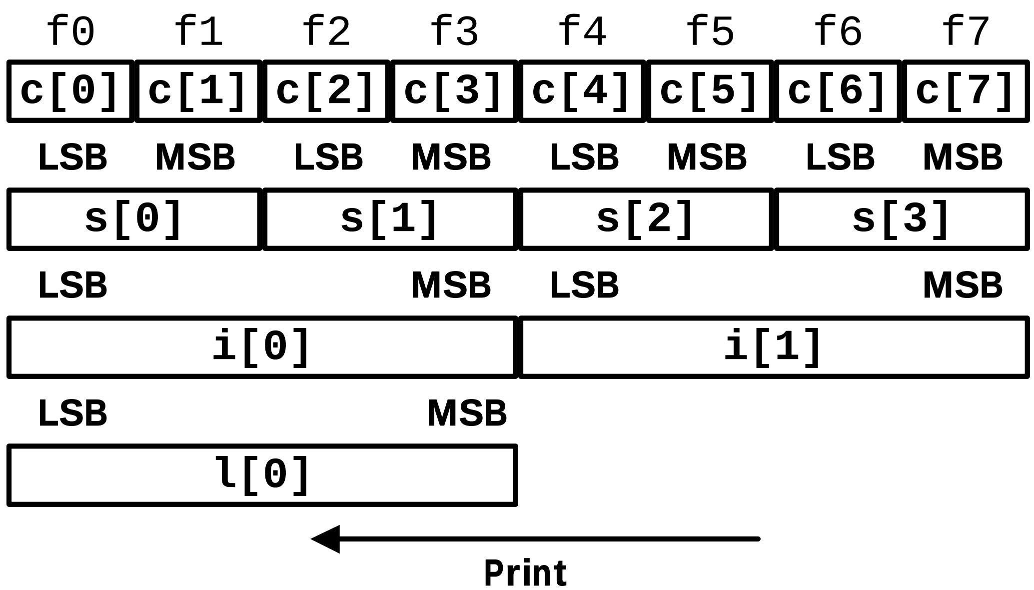
<table>
<tr><td>c[0]</td><td>c[1]</td><td>c[2]</td><td>c[3]</td><td>c[4]</td><td>c[5]</td><td>c[6]</td><td>c[7]</td></tr>
<tr><td colspan="2">s[0]</td><td colspan="2">s[1]</td><td colspan="2">s[2]</td><td colspan="2">s[3]</td></tr>
<tr><td colspan="4">i[0]</td><td colspan="4">i[1]</td></tr>
<tr><td colspan="4">l[0]</td><td colspan="4"></td></tr>
</table>

Byte Ordering Example (Cont).

```
int j;
for (j = 0; j < 8; j++)
dw.c[j] = 0xf0 + j;

printf("Characters 0-7 ==
[0x%x,0x%x,0x%x,0x%x,0x%x,0x%x,0x%x,0x%x]\n",
    dw.c[0], dw.c[1], dw.c[2], dw.c[3],
    dw.c[4], dw.c[5], dw.c[6], dw.c[7]);

printf("Shorts 0-3 ==
[0x%x,0x%x,0x%x,0x%x]\n",
    dw.s[0], dw.s[1], dw.s[2], dw.s[3]);

printf("Ints 0-1 == [0x%x,0x%x]\n",
    dw.i[0], dw.i[1]);

printf("Long 0 == [0x%lx]\n",
    dw.l[0]);
```

Byte Ordering on x86

Little Endian

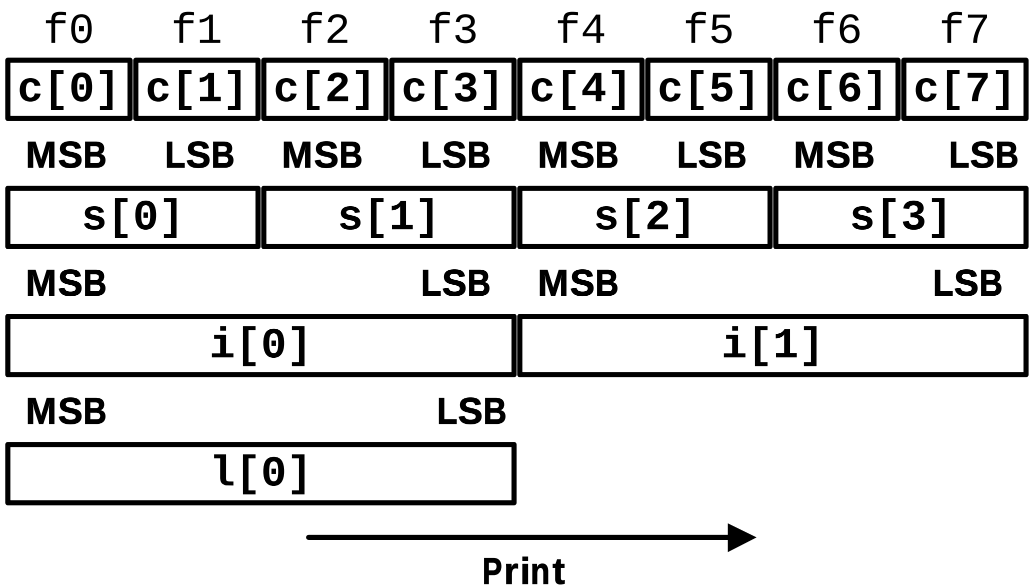


Output on Pentium:

```
Characters 0-7 == [0xf0,0xf1,0xf2,0xf3,0xf4,0xf5,0xf6,0xf7]
Shorts     0-3 == [0xf1f0,0xf3f2,0xf5f4,0xf7f6]
Ints       0-1 == [0xf3f2f1f0,0xf7f6f5f4]
Long       0   == [f3f2f1f0]
```

Byte Ordering on Sun

Big Endian

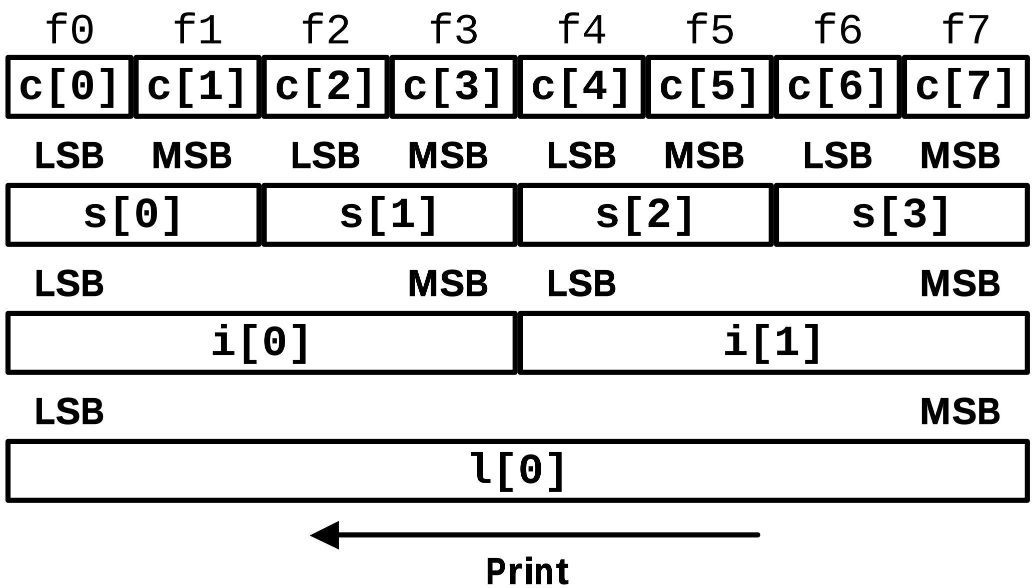


Output on Sun:

```
Characters 0-7 == [0xf0,0xf1,0xf2,0xf3,0xf4,0xf5,0xf6,0xf7]
Shorts     0-3 == [0xf0f1,0xf2f3,0xf4f5,0xf6f7]
Ints       0-1 == [0xf0f1f2f3,0xf4f5f6f7]
Long       0   == [0xf0f1f2f3]
```

Byte Ordering on Alpha

Little Endian



Output on Alpha:

```
Characters 0-7 == [0xf0,0xf1,0xf2,0xf3,0xf4,0xf5,0xf6,0xf7]
Shorts     0-3 == [0xf1f0,0xf3f2,0xf5f4,0xf7f6]
Ints       0-1 == [0xf3f2f1f0,0xf7f6f5f4]
Long       0   == [0xf7f6f5f4f3f2f1f0]
```

Byte-Level Operations

IA32 Support

- Arithmetic and data movement operations have byte-level version

  `movb`, `addb`, `testb`, etc.
- Some registers partially byte-addressable
- Can perform single byte memory references

Compiler

- Parameters and return values of type `char` passed as `int`'s
- Use `movsbl` to sign-extend byte to `int`

%eax	%ah	%al
%edx	%dh	%dl
%ecx	%ch	%cl
%ebx	%bh	%bl

Byte-Level Operation Example

- Compute Xor of characters in string

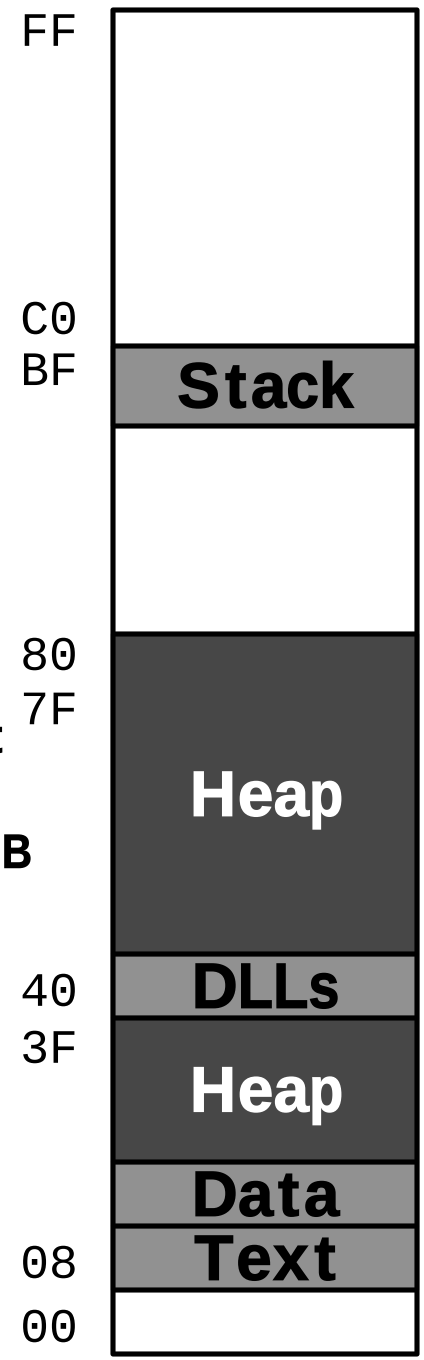
```
char string_xor(char *s)
{
  char result = 0;
  char c;
  do {
    c = *s++;
    result ^= c;
  } while (c);
  return result;
}
```

```
  # %edx = s, %cl = result
  movb $0,%cl      # result = 0
L2:              # loop:
  movb (%edx),%al # *s
  incl %edx        # s++
  xorb %al,%cl     # result ^= c
  testb %al,%al    # al
  jne L2           # If != 0, goto loop
  movsbl %cl,%eax # Sign extend to int
```

Linux Memory Layout

Upper 2 hex digits of address

Red Hat v. 5.2 ~1920MB memory



Stack

- **Runtime stack (8MB limit)**

Heap

- **Dynamically allocated storage**
- **When call `malloc`, `calloc`, `new`**

DLLs

- **Dynamically Linked Libraries**
- **Library routines (e.g., `printf`, `malloc`)**
- **Linked into object code when first executed**

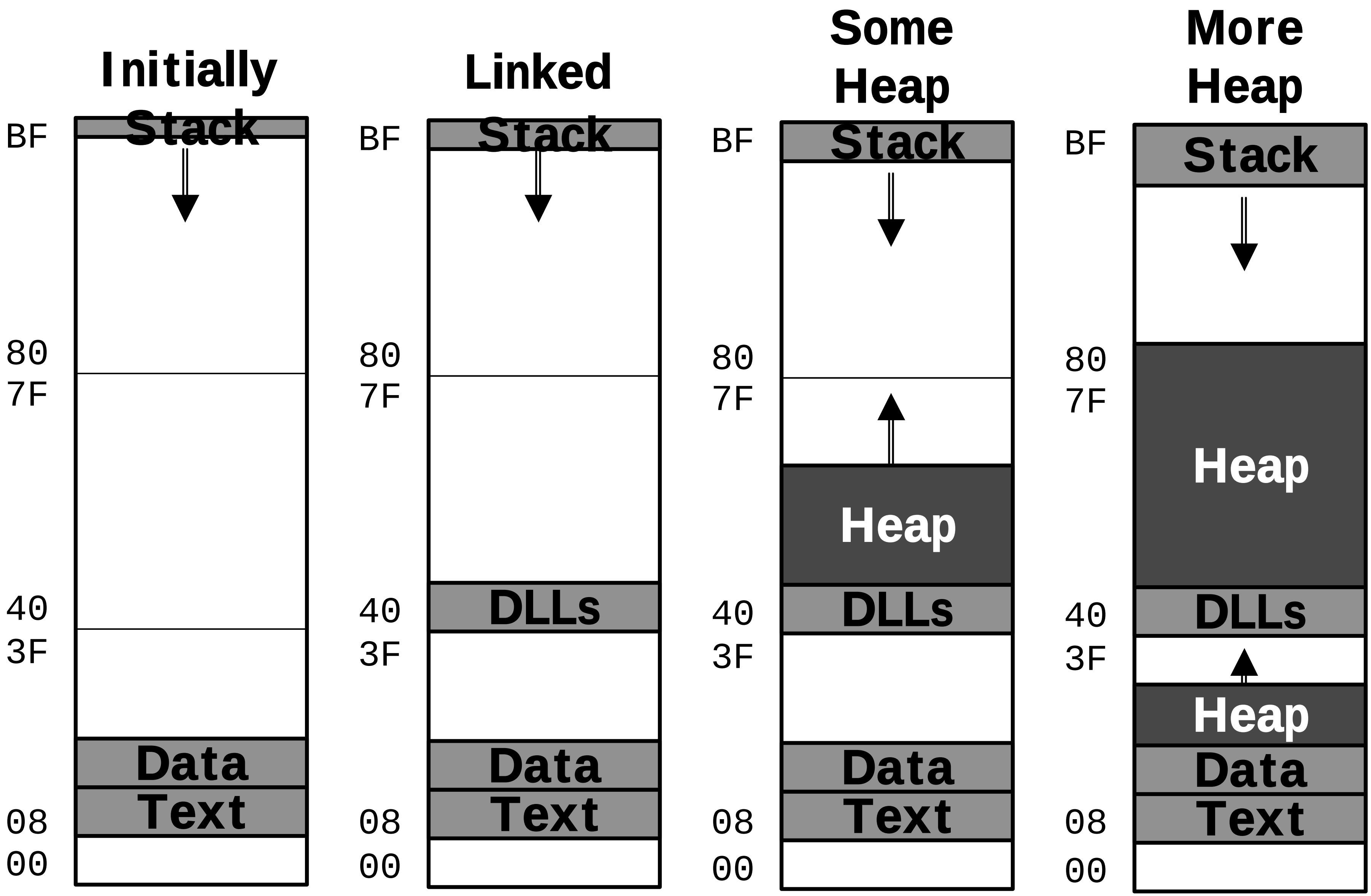
Data

- **Statically allocated data**
- **E.g., arrays & strings declared in code**

Text

- **Executable machine instructions**
- **Read-only**

Linux Memory Allocation



Memory Allocation Example

```
char big_array[1<<24];  /*  16 MB */
char huge_array[1<<28]; /* 256 MB */
int beyond;
char *p1, *p2, *p3, *p4;
int useless() {  return 0; }

int main()
{
 p1 = malloc(1 <<28);  /* 256 MB */
 p2 = malloc(1 << 8);  /* 256 B  */
 p3 = malloc(1 <<28);  /* 256 MB */
 p4 = malloc(1 << 8);  /* 256 B  */
 /* Some print statements ... */
}
```

Dynamic Linking Example

```
(gdb) print malloc
  $1 = {<text variable, no debug info>}
    0x8048454 <malloc>
(gdb) run
  Program exited normally.
(gdb) print malloc
  $2 = {void *(unsigned int)}
    0x40006240 <malloc>
```

Initially

- Code in text segment that invokes dynamic linker
- Address `0x8048454` should be read `0x08048454`

Final

- Code in DLL region

Breakpointing Example

```
(gdb) break main
(gdb) run
  Breakpoint 1, 0x804856f in main ()
(gdb) print $esp
  $3 = (void *) 0xbffffc78
```

Main

- Address `0x804856f` should be read `0x0804856f`

Stack

- Address `0xbffffc78`

Example Addresses

`$esp`	`0xbffffc78`
`p3`	`0x500b5008`
`p1`	`0x400b4008`
Final `malloc`	`0x40006240`
`p4`	`0x1904a640`
`p2`	`0x1904a538`
`beyond`	`0x1904a524`
`big_array`	`0x1804a520`
`huge_array`	`0x0804a510`
`main()`	`0x0804856f`
`useless()`	`0x08048560`
Initial `malloc`	`0x08048454`