

Bits and Bytes

CS 213

Aug. 27, 1998

Topics

- **Why bits?**
- **Representing information as bits**
 - Binary/Hexadecimal
 - Byte representations
 - » numbers
 - » characters and strings
 - » Instructions
- **Bit-level manipulations**
 - Boolean algebra
 - Expressing in C

Why Don't Computers Use Base 10?

Base 10 Number Representation

- That's why fingers are known as "digits"
- Natural representation for financial transactions
 - Floating point number cannot exactly represent \$1.20
- Even carries through in scientific notation
 - 1.5213×10^4

Implementing Electronically

- Hard to store
 - ENIAC (First electronic computer) used 10 vacuum tubes / digit
- Hard to transmit
 - Need high precision to encode 10 signal levels on single wire
- Messy to implement digital logic functions
 - Addition, multiplication, etc.

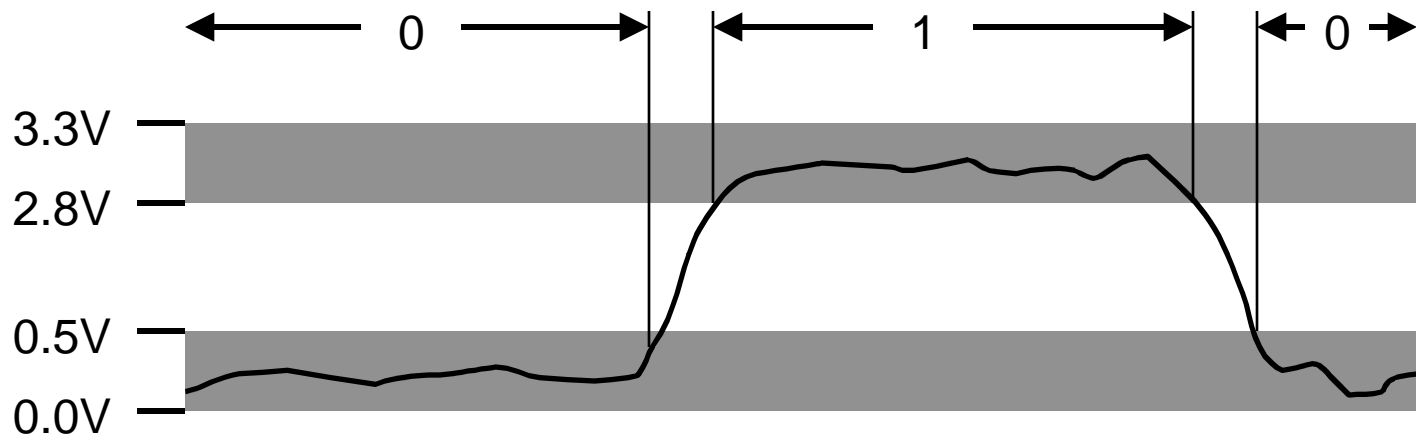
Binary Representations

Base 2 Number Representation

- Represent 15213_{10} as 11101101101101_2
- Represent 1.20_{10} as $1.0011001100110011[0011]\dots_2$
- Represent 1.5213×10^4 as $1.1101101101101_2 \times 2^{13}$

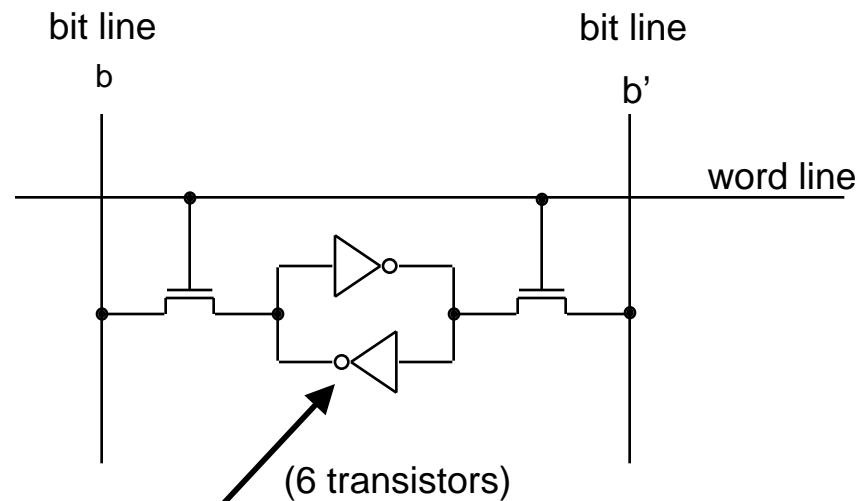
Electronic Implementation

- Easy to store with bistable elements
- Reliably transmitted on noisy and inaccurate wires



- **Straightforward implementation of arithmetic functions**

Anatomy of an SRAM Cell



Inverter.

High input	-->	Low Output
Low input	-->	High Output

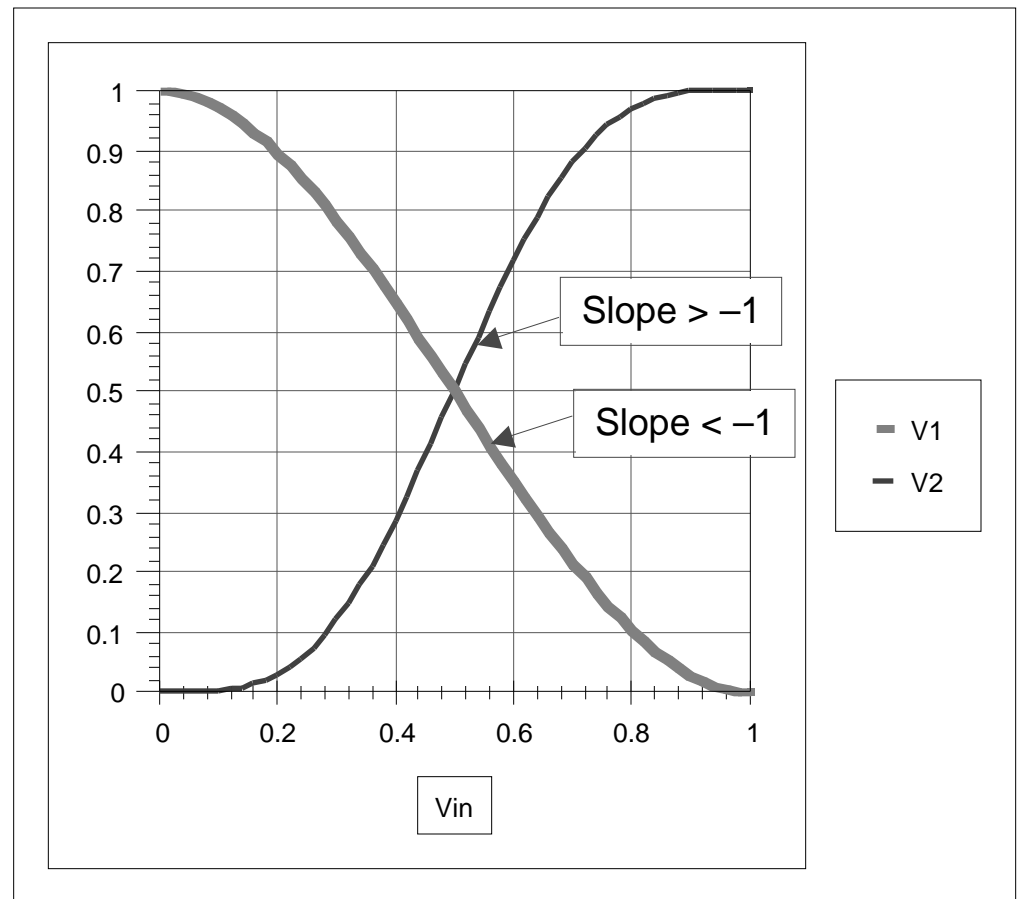
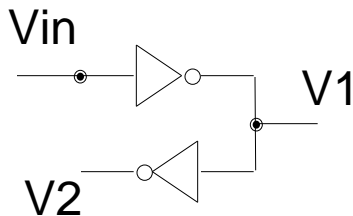
SRAM Cell Principle

Inverter Amplifies

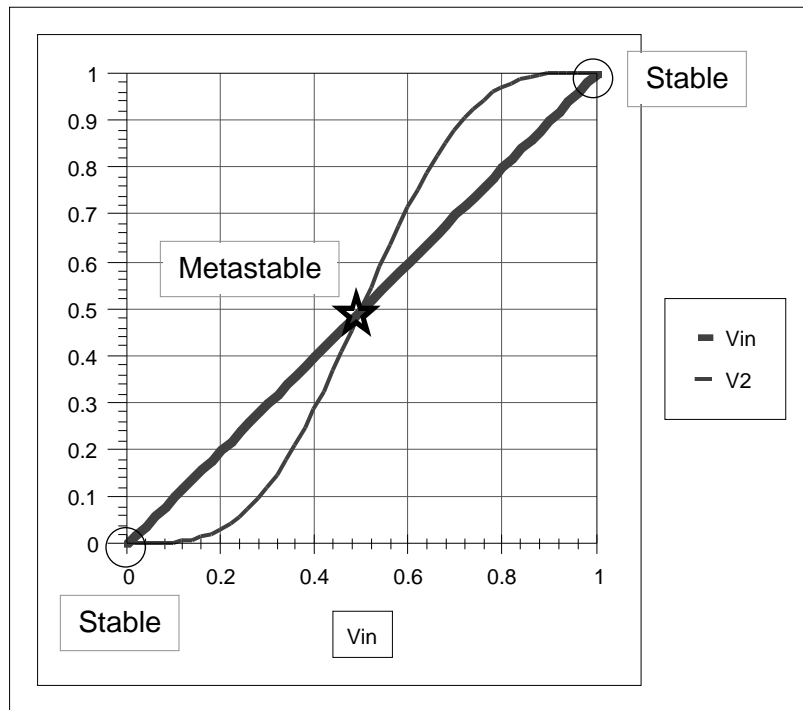
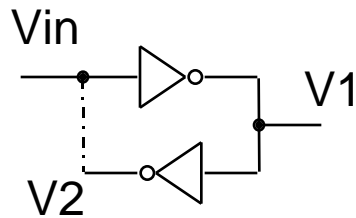
- Negative gain
- Slope < -1 in middle
- Saturates at ends

Inverter Pair Amplifies

- Positive gain
- Slope > 1 in middle
- Saturates at ends



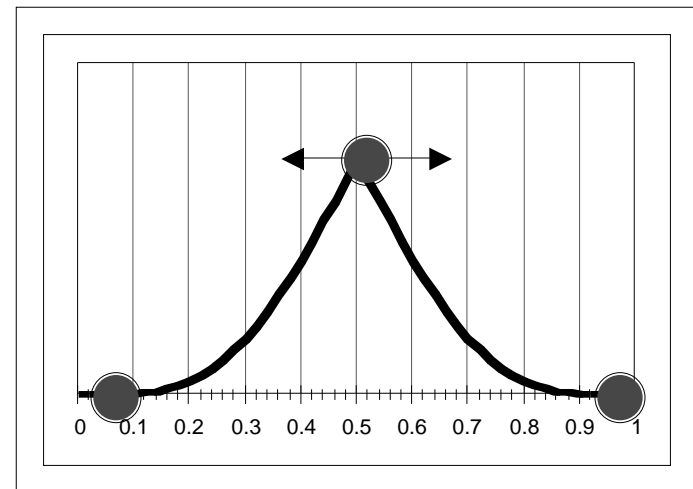
Bistable Element



Stability

- Require $V_{in} = V_2$
- Stable at endpoints
 - recover from perturbation
- Metastable in middle
 - Fall out when perturbed

Ball on Ramp Analogy



Byte-Oriented Memory Organization

Programs Refer to Virtual Addresses

- **Conceptually very large array of bytes**
- **Actually implemented with hierarchy of different memory types**
 - SRAM, DRAM, disk
 - Only allocate for regions actually used by program
- **In Unix and Windows NT, address space private to particular “process”**
 - Program being executed
 - Program can clobber its own data, but not that of others

Compiler + Run-Time System Control Allocation

- **Where different program objects should be stored**
- **Multiple mechanisms: static, stack, and heap**
- **In any case, all allocation within single virtual address space**

Encoding Byte Values

Byte = 8 bits

- **Binary** 00000000_2 to 11111111_2
- **Decimal:** 0_{10} to 255_{10}
- **Hexadecimal** 00_{16} to FF_{16}
 - Base 16 number representation
 - Use characters '0' to '9' and 'A' to 'F'
 - Write $FA1D37B_{16}$ in C as $0xFA1D37B$
 - » Or $0xfa1d37b$

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

Machine Words

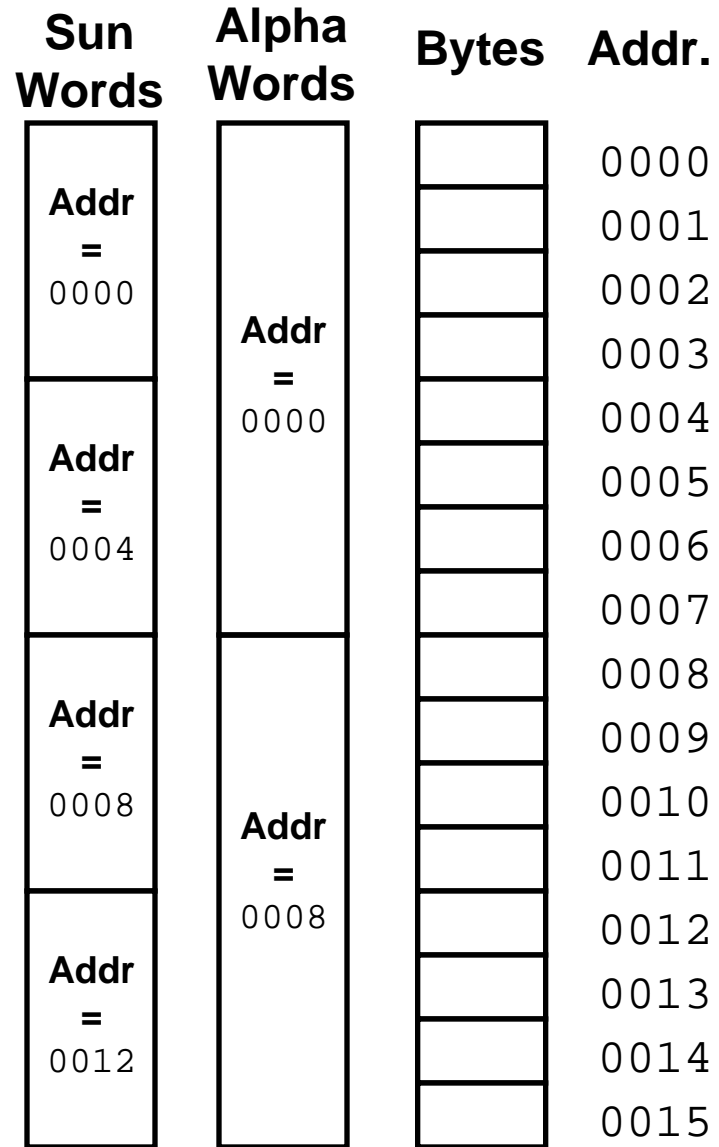
Machine Has “Word Size”

- **Nominal size of integer-valued data**
 - Including addresses
- **Most current machines are 32 bits (4 bytes)**
 - Limits addresses to 4GB
 - Becoming too small for memory-intensive applications
- **Our Alphas are 64 bits (8 bytes)**
 - Potentially address 1.8×10^{19} bytes
 - Although current machines cannot do this
 - » Limit is 4Terabytes
- **Machines support multiple data formats**
 - Fractions or multiples of word size
 - Always integral number of bytes

Word-Oriented Memory Organization

Addresses Specify Byte Locations

- Address of first byte in word
- Addresses of successive words differ by 4 (Sun) or 8 (Alpha)



Data Representations

Sizes of C Objects (in Bytes)

C Data Type	Alpha	Sun, PC, Mac, etc.
<code>int</code>	4	4
<code>long int</code>	8	4
<code>char</code>	1	1
<code>short</code>	2	2
<code>float</code>	4	4
<code>double</code>	8	8
<code>char *</code>	8	4

» Or any other pointer

Byte Ordering

- **Alphas, PC's are "Little Endian" machines**
 - Least significant byte has lowest address
- **Sun's, Mac's are "Big Endian" machines**
 - Least significant byte has highest address

Examining Data Representations

Code to Print Byte Representation of Data

- Cast pointer to unsigned char * creates byte array

```
typedef unsigned char *pointer;

void show_bytes(pointer start, int len)
{
    int i;
    for (i = 0; i < len; i++)
        printf("0x%p\t0x%x\n",
              start+i, start[i]);
    printf("\n");
}
```

Printf directives:

%p: Print pointer

%x: Print Hexadecimal

show_bytes Execution Example

```
int a = 15213;  
printf("int a = 15213;\n");  
show_bytes((pointer) &a, sizeof(int));
```

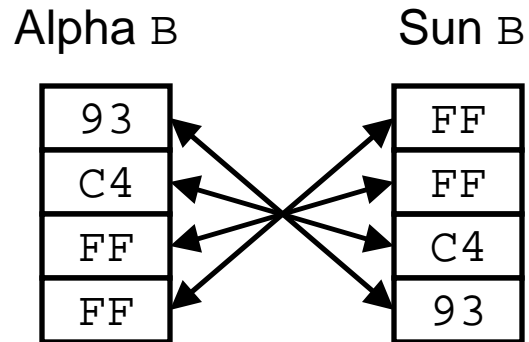
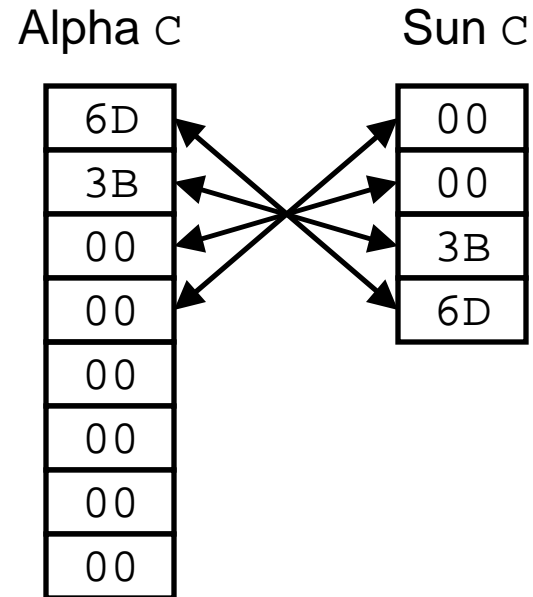
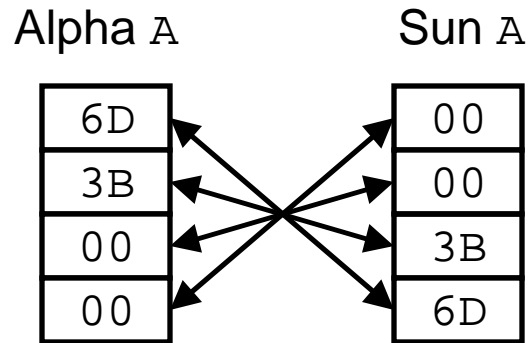
Result:

```
int a = 15213;  
0x11ffffcb8  0x6d  
0x11ffffcb9  0x3b  
0x11ffffcba  0x0  
0x11ffffcbb  0x0
```

Representing Integers

```
int A = 15213;  
int B = -15213;  
long int C = 15213;
```

Decimal:	15213
Binary:	0011 1011 0110 1101
Hex:	3 B 6 D



Two's complement representation
(Covered next lecture)

Representing Pointers

```
int B = -15213;
int *P = &B;
```

Alpha P

A0
FC
FF
FF
01
00
00
00

Alpha Address									
Hex:	1	F	F	F	F	F	C	A	0
Binary:	0001	1111	1111	1111	1111	1111	1100	1010	0000

Sun P

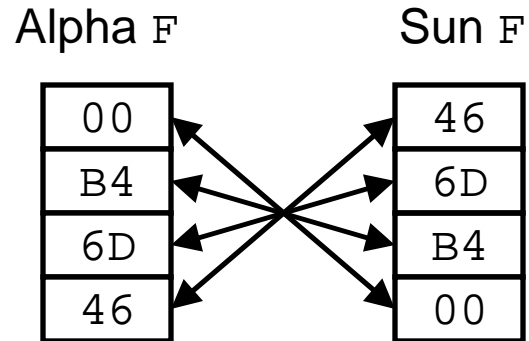
EF
FF
FB
2C

Sun Address								
Hex:	E	F	F	F	F	B	2	C
Binary:	1110	1111	1111	1111	1111	1011	0010	1100

Different compilers & machines assign different locations to objects

Representing Floats

Float F = 15213.0;



IEEE Single Precision Floating Point Representation

Hex:	4	6	6	D	B	4	0	0
Binary:	0100	0110	0110	1101	1011	0100	0000	0000
15213:			1110	1101	1011	01		



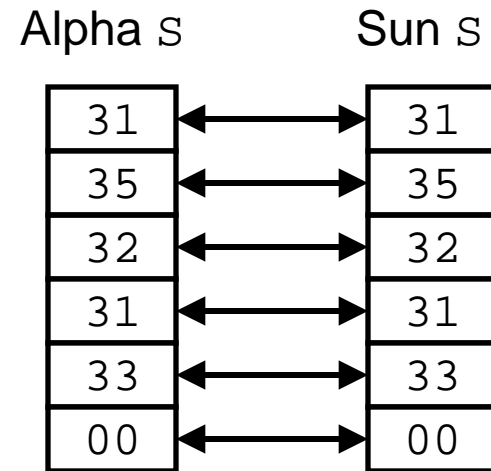
Not same as integer representation, but consistent across machines

Representing Strings

Strings in C

- Represented by array of characters
- Each character encoded in ASCII format
 - Standard 7-bit encoding of character set
 - Other encodings exist, but uncommon
 - Character “0” has code $0x30$
 - » Digit i has code $0x30+i$
- String should be null-terminated
 - Final character = 0

```
char s[6] = "15213";
```



Compatibility

- Byte ordering not an issue
 - Data are single byte quantities
- Text files generally platform independent
 - Except for different conventions of line termination character!

Machine-Level Code Representation

Encode Program as Sequence of Instructions

- **Each simple operation**
 - Arithmetic operation
 - Read or write memory
 - Conditional branch
- **Instructions encoded as bytes**
 - Alpha's, Sun's, Mac's use 4 byte instructions
 - » Reduced Instruction Set Computer (RISC)
 - PC's use variable length instructions
 - » Complex Instruction Set Computer (CISC)
- **Different instruction types and encodings for different machines**
 - Most code not binary compatible

Programs are Byte Sequences Too!

Representing Instructions

```
int sum(int x, int y)
{
    return x+y;
}
```

Alpha sum

00
00
30
42
01
80
FA
6B

Sun sum

81
C3
E0
08
90
02
00
09

- **For this example, both use two 4-byte instructions**
 - Use differing numbers of instructions in other cases

Different machines use totally different instructions and encodings

Boolean Algebra

Developed by George Boole in 19th Century

- Algebraic representation of logic
 - Encode “True” as 1 and “False” as 0

And

- $A \& B = 1$ when both $A=1$ and $B=1$

$\&$	0	1
0	0	0
1	0	1

Or

- $A | B = 1$ when either $A=1$ or $B=1$

	0	1
0	0	1
1	1	1

Not

- $\sim A = 1$ when $A=0$

\sim	
0	1
1	0

Exclusive-Or (Xor)

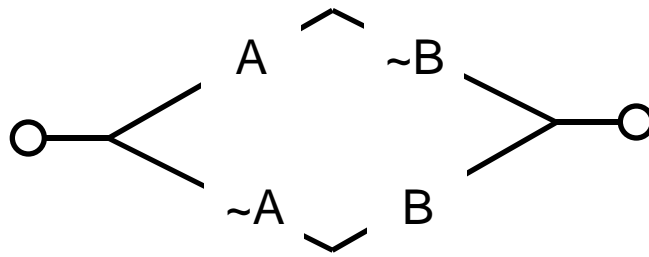
- $A \wedge B = 1$ when either $A=1$ or $B=1$, but not both

\wedge	0	1
0	0	1
1	1	0

Application of Boolean Algebra

Applied to Digital Systems by Claude Shannon

- 1937 MIT Master's Thesis
- Reason about networks of relay switches
 - Encode closed switch as 1, open switch as 0



Connection when
 $A \& \sim B \mid \sim A \& B$
 $= A \wedge B$

Properties of & and | Operations

Integer Arithmetic

- $\mathbb{Z}, +, *, -, 0, 1$ forms a “ring”
- Addition is “sum” operation
- Multiplication is “product” operation
- $-$ is additive inverse
- 0 is identity for sum
- 1 is identity for product

Boolean Algebra

- $\{0,1\}, |, \&, \sim, 0, 1$ forms a “Boolean algebra”
- Or is “sum” operation
- And is “product” operation
- \sim is “complement” operation (not additive inverse)
- 0 is identity for sum
- 1 is identity for product

Properties of Rings & Boolean Algebras

Boolean Algebra

- *Commutativity*

$$A \mid B = B \mid A$$

$$A \& B = B \& A$$

- *Associativity*

$$(A \mid B) \mid C = A \mid (B \mid C)$$

$$(A \& B) \& C = A \& (B \& C)$$

- *Product distributes over sum*

$$A \& (B \mid C) = (A \& B) \mid (A \& C)$$

- *Sum and product identities*

$$A \mid 0 = A$$

$$A \& 1 = A$$

- *Zero is product annihilator*

$$A \& 0 = 0$$

- *Cancellation of negation*

$$\sim(\sim A) = A$$

Integer Ring

$$A + B = B + A$$

$$A * B = B * A$$

$$(A + B) + C = A + (B + C)$$

$$(A * B) * C = A * (B * C)$$

$$A * (B + C) = A * B + B * C$$

$$A + 0 = A$$

$$A * 1 = A$$

$$A * 0 = 0$$

$$-(-A) = A$$

Ring Boolean Algebra

Boolean Algebra

- Boolean: *Sum distributes over product*

$$A \mid (B \& C) = (A \mid B) \& (A \mid C)$$

- Boolean: *Idempotency*

$$A \mid A = A$$

– “A is true” or “A is true” = “A is true”

$$A \& A = A$$

- Boolean: *Absorption*

$$A \mid (A \& B) = A$$

– “A is true” or “A is true and B is true” = “A is true”

$$A \& (A \mid B) = A$$

- Boolean: *Laws of Complements*

$$A \mid \sim A = 1$$

– “A is true” or “A is false”

- Ring: *Every element has additive inverse*

$$A \mid \sim A = 0$$

Integer Ring

$$A + (B * C) = (A + B) * (B + C)$$

$$A + A = A$$

$$A * A = A$$

$$A + (A * B) = A$$

$$A * (A + B) = A$$

$$A + \sim A = 1$$

$$A + \sim A = 0$$

Properties of & and ^

Boolean Ring

- $\{0,1\}$, \wedge , $\&$, \neg , 0 , 1
- **Identical to integers mod 2**
- \neg is identity operation: $\neg(\neg A) = A$
 $\neg A \wedge A = 0$

Property

- Commutative sum
- Commutative product
- Associative sum
- Associative product
- Prod. over sum
- 0 is sum identity
- 1 is prod. identity
- 0 is product annihilator
- Additive inverse

Boolean Ring

$$A \wedge B = B \wedge A$$
$$A \& B = B \& A$$
$$(A \wedge B) \wedge C = A \wedge (B \wedge C)$$
$$(A \& B) \& C = A \& (B \& C)$$
$$A \& (B \wedge C) = (A \& B) \wedge (A \& C)$$
$$A \wedge 0 = 0$$
$$A \& 1 = A$$
$$A \& 0 = 0$$
$$A \wedge A = 0$$

Relations Between Operations

DeMorgan's Laws

- Express & in terms of |, and vice-versa

$$A \& B = \sim(\sim A | \sim B)$$

» A and B are true if and only if neither A nor B is false

$$A | B = \sim(\sim A \& \sim B)$$

» A or B are true if and only if neither A is false nor B is false

Exclusive-Or using Inclusive Or

$$A \wedge B = (\sim A \& B) | (A \& \sim B)$$

» Exactly one of A and B is true

$$A \wedge B = (A | B) \& \sim(A \& B)$$

» Either A is true, or B is true, but not both

General Boolean Algebras

Operate on Bit Vectors

- Operations applied bitwise

01101001	01101001	01101001	01101001
<u>& 01010101</u>	<u> 01010101</u>	<u>^ 01010101</u>	<u>~ 01010101</u>
01000001	01111101	00111100	10101010

Representation of Sets

- Width w bit vector represents subsets of $\{0, \dots, w-1\}$

- $a_j = 1$ if $j \in A$

- 01101001 { 0, 3, 5, 6 }

- 01010101 { 0, 2, 4, 6 }

- **& Intersection** 01000001 { 0, 6 }
- **| Union** 01111101 { 0, 2, 3, 4, 5, 6 }
- **^ Symmetric difference** 00111100 { 2, 3, 4, 5 }
- **~ Complement** 10101010 { 1, 3, 5, 7 }

Bit-Level Operations in C

Operations $\&$, $|$, \sim , \wedge Available in C

- Apply to any “integer” data type
 - long int, int, short, char
- View arguments as bit vectors
- Arguments applied bit-wise

Examples (Char data type)

- $\sim 0x41 \rightarrow 0xBE$
 $\sim 01000001_2 \rightarrow 10111110_2$
- $\sim 0x00 \rightarrow 0xFF$
 $\sim 00000000_2 \rightarrow 11111111_2$

- $0x69 \ \& \ 0x55 \rightarrow 0x41$
 $01101001_2 \ \& \ 01010101_2 \rightarrow 01000001_2$
- $0x69 \ | \ 0x55 \rightarrow 0x7D$
 $01101001_2 \ | \ 01010101_2 \rightarrow 01111101_2$

Contrast: Logic Operations in C

Contrast to Logical Operators

- `&&`, `||`, `!`
 - View 0 as “False”
 - Anything nonzero as “True”
 - Always return 0 or 1

Examples (char data type)

- `!0x41 --> 0x00`
- `!0x00 --> 0x01`
- `!!0x41 --> 0x01`

- `0x69 && 0x55 --> 0x01`
- `0x69 || 0x55 --> 0x01`

Shift Operations

Left Shift: $x \ll y$

- **Shift bit-vector x left y positions**
 - Throw away extra bits on left
 - Fill with 0's on right

Argument x	01100010
$\ll 3$	00010000
Log. $\gg 2$	00011000
Arith. $\gg 2$	00011000

Right Shift: $x \gg y$

- **Shift bit-vector x right y positions**
 - Throw away extra bits on right
- **Logical shift**
 - Fill with 0's on left
- **Arithmetic shift**
 - Replicate most significant bit on left
 - Useful with two's complement integer representation

Argument x	10100010
$\ll 3$	00010000
Log. $\gg 2$	00101000
Arith. $\gg 2$	11101000

Cool Stuff with Xor

- Bitwise Xor is form of addition
- With extra property that every value is its own additive inverse
 - $A \oplus A = 0$

```
void funny(int *x, int *y)
{
    *x = *x ^ *y;    /* #1 */
    *y = *x ^ *y;    /* #2 */
    *x = *x ^ *y;    /* #3 */
}
```

Step	*x	*y
Begin	A	B
1	$A \oplus B$	B
2	$A \oplus B$	$(A \oplus B) \oplus B = A \oplus (B \oplus B) = A \oplus 0 = A$
3	$(A \oplus B) \oplus A = (B \oplus A) \oplus A = B \oplus (A \oplus A) = B \oplus 0 = B$	A
End	B	A