Recitation 10: Malloc Lab

Instructors

October 21, 2019
Administrivia

- Malloc checkpoint due **Tuesday, October 29!** yeeT
- Malloc final due the week, **Tuesday, November 5!** yooT
- Malloc Bootcamp:
  - **Sunday, October 27** at Rashid Auditorium, 7-9PM 😎
    - We will cover ✨ fun and flirty ✨ ways to succeed post-malloc checkpoint!
    - Tell your friends to come (if they’re in 213 (if they want to come (don’t force your friends to do things they don’t want to do that’s not what friends are for)))
Checkpoint Submission

Style Grading
- We will grade your checkheap with your checkpoint submission!

Things to Remember:
- Document checkheap
- See writeup for what to include in checkheap
Outline

- Concept
- How to choose blocks
- Metadata
- Debugging / GDB Exercises
What is malloc?

- A function to allocate memory during runtime (dynamic memory allocation).
  - More useful when the size or number of allocations is unknown until runtime (e.g. data structures)

- The heap is a segment of memory addresses reserved almost exclusively for malloc to use.
  - Your code directly manipulates the bytes of memory in this section.
**Concept**

- Overall, malloc does three things:

1. Organizes all blocks and stores information about them in a structured way.
2. Uses the structure made to choose an appropriate location to allocate new memory.
3. Updates the structure when the user frees a block of memory.

This process occurs even for a complicated algorithm like segregated lists.
Concept (Implicit list)

1. Connects and organizes all blocks and stores information about them in a structured way, typically implemented as a singly linked list
Concept (Implicit list)

2. Uses the structure made to choose an appropriate location to allocate new memory.

```
p1 = malloc(3)
p2 = malloc(7)
p3 = malloc(5)
```
Concept (Implicit list)

3. Updates the structure when the user frees a block of memory.

\[
\begin{align*}
p1 &= \text{malloc}(3) \\
p2 &= \text{malloc}(7) \\
p3 &= \text{malloc}(5) \\
\text{free}(p2) &
\end{align*}
\]
Concept (Implicit list)

3. Updates the structure when the user frees a block of memory.

```
p1 = malloc(3)
p2 = malloc(7)
p3 = malloc(5)
free(p2)
```
Goals

- Run as fast as possible
- Waste as little memory as possible

Seemingly conflicting goals, but with the library `malloc` call cleverness you can do very well in both areas!

- The simplest implementation is the implicit list. `mm.c` uses this method.
  - Unfortunately…
This is pretty slow... most explicit list implementations get above 2000 Kops/sec

<table>
<thead>
<tr>
<th>valid</th>
<th>util</th>
<th>ops</th>
<th>msecs</th>
<th>Kops</th>
<th>trace</th>
</tr>
</thead>
<tbody>
<tr>
<td>yes</td>
<td>78.4%</td>
<td>20</td>
<td>0.002</td>
<td>9632</td>
<td>./traces/syn-array-short.rep</td>
</tr>
<tr>
<td>yes</td>
<td>13.4%</td>
<td>20</td>
<td>0.001</td>
<td>25777</td>
<td>./traces/syn-struct-short.rep</td>
</tr>
<tr>
<td>yes</td>
<td>15.2%</td>
<td>20</td>
<td>0.001</td>
<td>24783</td>
<td>./traces/syn-string-short.rep</td>
</tr>
<tr>
<td>yes</td>
<td>73.1%</td>
<td>20</td>
<td>0.001</td>
<td>19277</td>
<td>./traces/syn-mix-short.rep</td>
</tr>
<tr>
<td>yes</td>
<td>16.0%</td>
<td>36</td>
<td>0.001</td>
<td>31192</td>
<td>./traces/NGRAM-FOX1.rep</td>
</tr>
<tr>
<td>yes</td>
<td>73.6%</td>
<td>757</td>
<td>0.145</td>
<td>5237</td>
<td>./traces/syn-mix-realloc.rep</td>
</tr>
</tbody>
</table>
* yes  | 62.0% | 5748| 3.925 | 1464     | ./traces/bdd-aa4.rep            |
* yes  | 58.3% | 87830| 1682.766| 52 | ./traces/bdd-aa32.rep          |
* yes  | 58.0% | 41080| 410.385| 100 | ./traces/bdd-ma4.rep           |
* yes  | 58.1% | 115380| 4636.711| 25 | ./traces/bdd-nq7.rep           |
* yes  | 56.6% | 20547| 26.677| 770 | ./traces/cbit-abs.rep          |
* yes  | 55.8% | 95276| 675.303| 141 | ./traces/cbit-parity.rep       |
* yes  | 58.0% | 89623| 611.511| 147 | ./traces/cbit-satadd.rep       |
* yes  | 49.6% | 50583| 185.382| 273 | ./traces/cbit-xyz.rep          |
* yes  | 40.6% | 32540| 76.919| 423 | ./traces/NGRAM-GULLIVER1.rep   |
* yes  | 42.4% | 127912| 1284.959| 100 | ./traces/NGRAM-GULLIVER2.rep   |
* yes  | 39.4% | 67012| 338.591| 198 | ./traces/NGRAM-MOBY1.rep       |
* yes  | 38.6% | 94828| 701.305| 135 | ./traces/NGRAM-SHAKE1.rep      |
* yes  | 90.9% | 80000| 1455.891| 55 | ./traces/syn-array.rep         |
* yes  | 88.0% | 80000| 915.167| 87  | ./traces/syn-mix.rep           |
* yes  | 74.3% | 80000| 914.366| 87  | ./traces/syn-string.rep        |
* yes  | 75.2% | 80000| 812.748| 98  | ./traces/syn-struct.rep        |

16 16 | 59.1% | 1148359| 14732.604| 78  |                                  |

Average utilization = 59.1%. Average throughput = 78 Kops/sec

Checkpoint Perf index = 20.0 (util) + 0.0 (thru) = 20.0/100
Allocation methods in a nutshell

- Implicit list: a list is implicitly formed by jumping between blocks, using knowledge about their sizes.
  - Implicit list: a list is implicitly formed by jumping between blocks, using knowledge about their sizes.

- Explicit list: Free blocks explicitly point to other blocks, like in a linked list.
  - Understanding explicit lists requires understanding implicit lists

- Segregated list: Multiple linked lists, each containing blocks in a certain range of sizes.
  - Understanding segregated lists requires understanding explicit lists
Choices

- **What kind of implementation to use?**
  - Implicit list, explicit list, segregated lists, binary tree methods, etc.
  - You can use specialized strategies depending on the size of allocations.
  - Adaptive algorithms are fine, though not necessary to get 100%.
    - Don’t hard-code for individual trace files - you’ll get no credit/code deductions!

- **What fit algorithm to use?**
  - Best fit: choose the smallest block that is big enough to fit the requested allocation size.
  - First fit / next fit: search linearly starting from some location, and pick the first block that fits.
  - Which is faster? Which uses less memory?
  - “Good enough” fit: a blend between the two.

- **This lab has many more ways to get an A+ than, say, Cache Lab Part 2**
Finding a Best Block

- Suppose you have implemented the explicit list approach
  - You were using best fit with explicit lists

  - Will your memory utilization score improve?

  *Note: you don’t have to implement seglists and run mdriver to answer this. That’s, uh, hard to do within one recitation session.*

  - What other advantages does segregated lists provide?

- Losing memory because of the way you choose your free blocks is called **external fragmentation**.
Metadata

- All blocks need to store some data about themselves in order for `malloc` to keep track of them (e.g. headers)
  - This takes memory too…
  - Losing memory for this reason is called **internal fragmentation**.

- **What data might a block need?**
  - Does it depend on the `malloc` implementation you use?
  - Is it different between free and allocated blocks?

- **Can we use the extra space in free blocks?**
  - Or do we have to leave the space alone?

- **How can we overlap two different types of data at the same location?**
In a perfect world...

Setting up the blocks, metadata, lists… etc (500 LoC)
+ Finding and allocating the right blocks (500 LoC)
+ Updating your heap structure when you free (500 LoC) =

```
[dalud@angelshark:~/.s17/malloclabcheckpoint-handout] $ ./mdriver
Found benchmark throughput 13056 for cpu type Intel(R)Xeon(R)CPUE5520@2.27
Throughput targets: min=6528, max=11750, benchmark=13056

Results for mm malloc:
valid  util  ops  msecs  Kops  trace
yes  78.1%  20  0.004  5595 ./traces/syn-array-short.rep
yes  3.2%  20  0.004  5273 ./traces/syn-struct-short.rep
* yes  96.0%  80000  17.176  4658 ./traces/syn-array.rep
* yes  93.2%  80000  6.154  12999 ./traces/syn-mix.rep
* yes  86.4%  80000  3.717  21521 ./traces/syn-string.rep
* yes  85.6%  80000  3.649  21924 ./traces/syn-struct.rep
16 16  74.2%  1148359  55.949  20525

Average utilization = 74.2%. Average throughput = 20525 Kops/sec
Perf index = 60.0 (util) + 40.0 (thru) = 100.0/100
```
In reality...

- Setting up the blocks, metadata, lists… etc (500 LoC)
- Finding and allocating the right blocks (500 LoC)
- Updating your heap structure when you free (500 LoC)
- One bug, somewhere lost in those 1500 LoC =

```
[dalud@angelshark:~/.../15213/s17/malloclabcheckpoint-handout] $ ./mdriver
Found benchmark throughput 13056 for cpu type Intel(R)Xeon(R)CPUE5520@2.27
Throughput targets: min=6528, max=11750, benchmark=13056
......Segmentation fault
[dalud@angelshark:~/.../15213/s17/malloclabcheckpoint-handout] $  
```
LARGE AMOUNT OF CODE COMPILLES THE FIRST TIME.

NOT SURE IF GENIUS OR HUGE EXPLOSION COMING.
Common errors you might see

- **Garbled bytes**
  - Problem: overwriting data in an allocated block
  - Solution: remembering data lab and the good ol’ days finding where you’re overwriting by stepping through with gdb

- **Overlapping payloads**
  - Problem: having unique blocks whose payloads overlap in memory
  - Solution: literally print debugging everywhere finding where you’re overlapping by stepping through with gdb

- **Segmentation fault**
  - Problem: accessing invalid memory
  - Solution: crying a little finding where you’re accessing invalid memory by stepping through with gdb
Try running $ make

- If you look closely, our code compiles your malloc implementation with the \(-O3\) flag.
- This is an optimization flag. \(-O3\) makes your code run as efficiently as the compiler can manage, but also makes it horrible for debugging (almost everything is “optimized out”).

For malloclab, we’ve provide you a driver, `mdriver-dbg`, that not only enables debugging macros, but compiles your code with \(-O0\). This allows more useful information to be displayed in GDB.
The *Real* Activity: GDB Practice

- Using GDB well in malloclab can save you \(\text{HOURS}^{1,2}\) of debugging time
  - Average 20 hours using GDB for “B” on malloclab
  - Average 23 hours not using GDB for “B” on malloclab

* Average time is based on Summer 2016 survey results

- Form pairs

  wget https://www.cs.cmu.edu/~213/activities/f19-rec-malloc.tar
  tar xvf f19-rec-malloc.tar
  cd f19-rec-malloc
  make

- Two buggy mdrivers
Debugging Guidelines

If you have this problem...

- Ran into segfault
- Trace results don’t match yours
- Don’t know what trace output should be

You might want to...

- Locate a segfault:
  - run
  - <>
  - backtrace
  - list
- Reproduce results of a trace:
  - Run with gdb
    - gdb args
Debugging mdriver

$ gdb --args ./mdriver -c traces/syn-mix-short.rep
(gdb) run
(gdb) backtrace
(gdb) list

Optional: Type Ctrl-X Ctrl-A to see the source code. Don’t linger there for long, since this visual mode is buggy. Type that key combination again to go back to console mode.

(Or use cgdb - see Piazza post)

1) What function is listed on the top of backtrace?
2) What line of code crashed?
3) How did that line cause the crash?
Debugging mdriver

- `(gdb) x /gx block`
  - Shows the memory contents within the block
  - In particular, look for the header.

- `(gdb) print *block`
  - Shows struct contents

  *Alternative:* `(gdb) print *(block_t *) <address>`

Remember the output from `(gdb) bt`?

- `(gdb) frame 1`
  - Jumps to the function one level down the call stack (aka the function that called `write_footer`)
  - Ctrl-X, Ctrl-A again if you want to see visuals

What was the caller function? What is its purpose?

- Was it writing to `block` or `block_next` when it crashed?
Thought process while debugging

- `write_footer` crashed because it got the wrong address for the footer...
- The address was wrong because the header of the block was some garbage value
  - Since `write_footer` uses `get_size(block)` after all
- But why in the world does the header contain garbage??
  - The crash happened in `place`, which basically splits a free block into two and uses the first one to store things.
  - Hm, `block_next` would be the new block created after the split? The one on the right?
  - The header would be in the middle of the original free block actually. Wait, but I wrote a new header before I wrote the footer!
    - Right? …Oh, I didn’t. Darn.
Heap consistency checker

- mm-2.c activates debug mode, and so mm_checkheap runs at the beginning and end of many of its functions.

- The next bug will be a total nightmare to find without this heap consistency checker*.

*Even though the checker in mm-2.c is short and buggy
Now you try debugging this - second example!

$ gdb --args ./mdriver-2 -c traces/syn-array-short.rep

(gdb) run

Yikes… what error are we getting?
Now you try debugging this - second example!

$ gdb --args ./mdriver-2 -c traces/syn-array-short.rep
(gdb) run

Yikes... what error are we getting?

~g a r b l e d b y t e s~

* an accurate representation of what’s actually going on in your blocks
Now you try debugging this - second example!

(gdb) watch *0x8000026d0 /* Track from first garbled payload */
(gdb) run
(gdb) continue
(gdb) continue /* Keep going until coalesce_block */
(gdb) backtrace
(gdb) list

Ah, it seems like nothing’s amiss…
Running with mdriver-2-dbg...

- Let’s run it with mdriver-2-dbg, which has a lower optimization - will give us more insight, like the stack trace below

(in the same gdb session)

(gdb) file mdriver-2-dbg
(gdb) run
(gdb) continue
...
(gdb) list
Running with mdriver-2-dbg...

- Now try printing out the values of prev_alloc / next_alloc...

  (gdb) print prev_alloc
  (gdb) $1 = <optimized out>

%rip, they’re optimized out! We have to change the optimization level to get what we truly want.
Running with mdriver-2-dbg...

- Go into your Makefile (vim Makefile) => change “COPT_DBG = -O0” so that all local variables are preserved

```
$ make clean
$ make
$ gdb --args ./mdriver-2-dbg -d 2 -c traces/syn-mix-short.rep
(gdb) b mm-2.c:450 /* Cut to the chase... */
(gdb) run
(gdb) continue
...
(gdb) print next_alloc
(gdb) $1 = true /* SUCCESS! */
```
Strategy - Suggested Plan for Completing Malloc

0. *Start writing your checkheap!*

1. Get an explicit list implementation to work with proper coalescing and splitting

3. Get to a segregated list implementation to improve utilization

4. Work on optimizations (each has its own challenges!)
   - Remove footers
   - Decrease minimum block size
   - Reduce header sizes
Strategy - Suggested Plan for Completing Malloc

0. **Start writing your checkheap!**  
Keep writing your checkheap!

1. Get an explicit list implementation to work with proper coalescing and splitting  
Keep writing your checkheap!

3. Get to a segregated list implementation to improve utilization  
Keep writing your checkheap!

4. Work on optimizations (each has its own challenges!)
   - Remove footers  
   - Decrease minimum block size  
   - Reduce header sizes  
   Keep writing your checkheap!
MallocLab Checkpoint

- Due *next Tuesday*!

- Checkpoint should take a bit less than half of the time you spend overall on the lab.

- Read the write-up. Slowly. Carefully.

- Use GDB - watch, backtrace

- Ask us for debugging help
  - Only after you implement mm_checkheap though! You gotta learn how to understand your own code - help us help you!

please write checkheap or we will scream
Appendix: Advanced GDB Usage

- **backtrace**: Shows the call stack
- **up/down**: Lets you go up/down one level in the call stack
- **frame**: Lets you go to one of the levels in the call stack
- **list**: Shows source code
- **print <expression>**: Runs any valid C command, even something with side effects like `mm_malloc(10)` or `mm_checkheap(1337)`
- **watch <expression>**: Breaks when the value of the expression changes
- **break <function / line> if <expression>**: Only stops execution when the expression holds true
- **Ctrl-X Ctrl-A or cgdb for visualization**