

Floating point encoding. In this problem, you will work with floating point numbers based on the IEEE floating point format. We consider two different 6-bit formats:

Format A:

- There is one sign bit s .
- There are $k = 3$ exponent bits. The bias is $2^{k-1} - 1 = 3$.
- There are $n = 2$ fraction bits.

Format B:

- There is one sign bit s .
- There are $k = 2$ exponent bits. The bias is $2^{k-1} - 1 = 1$.
- There are $n = 3$ fraction bits.

For formats A and B, please write down the binary representation for the following (use round-to-even). Recall that for denormalized numbers, $E = 1 - \text{bias}$. For normalized numbers, $E = e - \text{bias}$.

Value	Format A Bits	Format B Bits
Zero	0 000 00	0 00 000
One		
1/2		
11/8		

Assembly Loop

080483a0 <mystery>:

```

80483a0:      55
80483a1:      89 e5
80483a3:      53
80483a4:      8b 5d 0c
80483a7:      8b 4d 08
80483aa:      85 db
80483ac:      7e 17
80483ae:      31 c0
80483b0:      8b 14 81
80483b3:      f6 c2 01
80483b6:      75 06
80483b8:      83 c2 01
80483bb:      89 14 81
80483be:      83 c0 01
80483c1:      39 d8
80483c3:      75 eb
80483c5:      5b
80483c6:      5d
80483c7:      c3

```

```

push    %ebp
mov     %esp,%ebp
push    %ebx
mov     0xc(%ebp),%ebx
mov     0x8(%ebp),%ecx
test    %ebx,%ebx
jle     80483c5 <mystery+0x25>
xor     %eax,%eax
mov     (%ecx,%eax,4),%edx
test    $0x1,%dl
jne     80483be <mystery+0x1e>
add     $0x1,%edx
mov     %edx, (%ecx,%eax,4)
add     $0x1,%eax
cmp     %ebx,%eax
jne     80483b0 <mystery+0x10>
pop     %ebx
pop     %ebp
ret

```

Assembly Loop

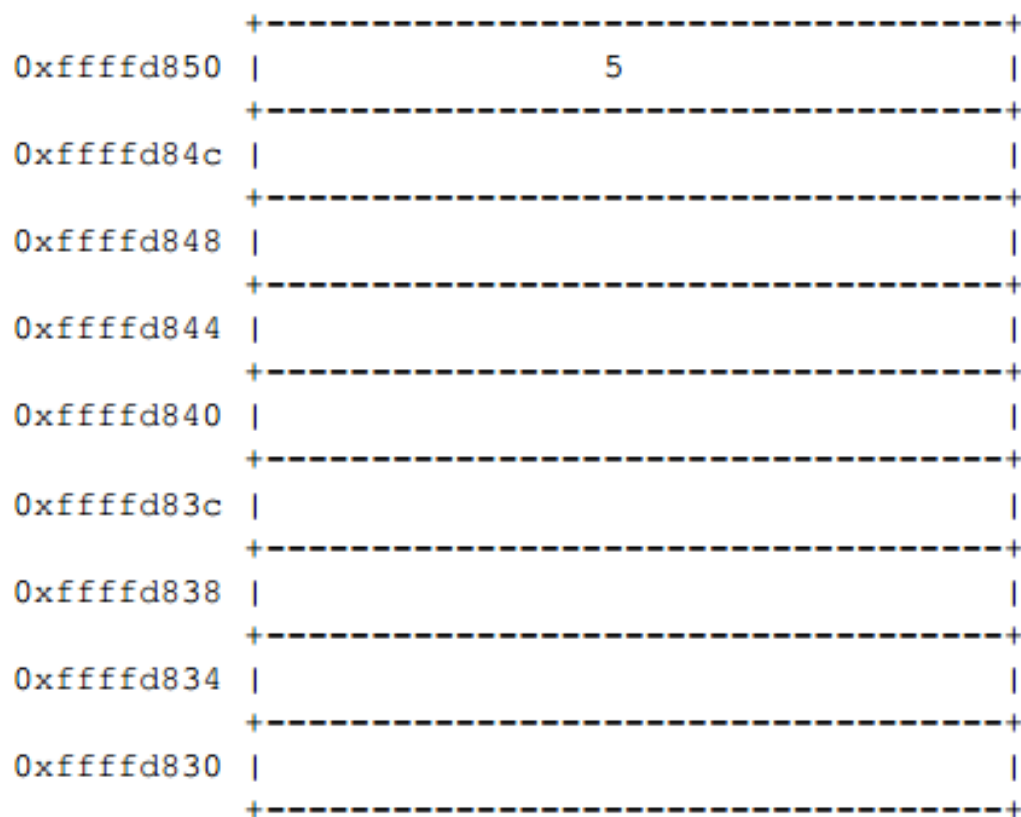
```
void mystery(int *array, int n)
{
    int i;
    for(_____; _____; _____)
    {
        if(_____ == 0)
            _____;
    }
}
```

- Given assembly code of `foo()` and `bar()`, Draw a detailed picture of the stack, starting with the caller invoking `foo(3, 4, 5)`.

Value of `%ebp` when `foo` is called: `0xffffd858`

Return address in function that called `foo`: `0x080483c9`

Stack addresss The diagram starts with the arguments for `foo()`



Stack Diagram

```
int bar (int a, int b) {
    return a + b;
}
```

```
int foo(int n, int m, int c) {
    c += bar(m, n);
    return c;
}
```

08048374 <bar>:

8048374:	55	push	%ebp
8048375:	89 e5	mov	%esp,%ebp
8048377:	8b 45 0c	mov	0xc(%ebp),%eax
804837a:	03 45 08	add	0x8(%ebp),%eax
804837d:	5d	pop	%ebp
804837e:	c3	ret	

0804837f <foo>:

804837f:	55	push	%ebp
8048380:	89 e5	mov	%esp,%ebp
8048382:	83 ec 08	sub	\$0x8,%esp
8048385:	8b 45 08	mov	0x8(%ebp),%eax
8048388:	89 44 24 04	mov	%eax,0x4(%esp)
804838c:	8b 45 0c	mov	0xc(%ebp),%eax
804838f:	89 04 24	mov	%eax,(%esp)
8048392:	e8 dd ff ff ff	call	8048374 <bar>
8048397:	03 45 10	add	0x10(%ebp),%eax
804839a:	c9	leave	
804839b:	c3	ret	

Assembly - Array Access

Find H, J

```
int array1[H][J];
int array2[J][H];

int copy_array(int x, int y) {
    array2[y][x] = array1[x][y];

    return 1;
}
```

Suppose the above C code generates the following x86-64 assembly code:

```
# On entry:
#     %edi = x
#     %esi = y
#
copy_array:
    movslq    %edi,%rdi
    movslq    %esi,%rsi
    movq      %rdi,%rax
    leaq      (%rsi,%rsi,2), %rdx
    salq      $5,%rax
    subq      %rdi,%rax
    leaq      (%rdi,%rdx,2), %rdx
    addq      %rsi,%rax
    movl      array1(,%rax,4), %eax
    movl      %eax,array2(,%rdx,4)
    movl      $1,%eax
    ret
```