System-Level I/O

15-213: Introduction to Computer Systems

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Today

- Unix I/O
- RIO (robust I/O) package
- Metadata, sharing, and redirection
- Standard I/O
- Conclusions and examples

Unix Files

- A Unix *file* is a sequence of *m* bytes:
 - \blacksquare $B_0, B_1, \dots, B_k, \dots, B_{m-1}$
- All I/O devices are represented as files:
 - /dev/sda2 (/usr disk partition)
 - /dev/tty2 (terminal)
- Even the kernel is represented as a file:
 - /dev/kmem (kernel memory image)
 - /proc (kernel data structures)

Unix File Types

Regular file

- File containing user/app data (binary, text, whatever)
- OS does not know anything about the format
 - other than "sequence of bytes", akin to main memory

Directory file

A file that contains the names and locations of other files

Character special and block special files

Terminals (character special) and disks (block special)

FIFO (named pipe)

A file type used for inter-process communication

Socket

A file type used for network communication between processes

Unix I/O

Key Features

- Elegant mapping of files to devices allows kernel to export simple interface called Unix I/O
- Important idea: All input and output is handled in a consistent and uniform way

Basic Unix I/O operations (system calls):

- Opening and closing files
 - open() and close()
- Reading and writing a file
 - read() and write()
- Changing the current file position (seek)
 - indicates next offset into file to read or write

Opening Files

 Opening a file informs the kernel that you are getting ready to access that file

```
int fd; /* file descriptor */
if ((fd = open("/etc/hosts", O_RDONLY)) < 0) {
   perror("open");
   exit(1);
}</pre>
```

- Returns a small identifying integer *file descriptor*
 - fd == -1 indicates that an error occurred
- Each process created by a Unix shell begins life with three open files associated with a terminal:
 - 0: standard input
 - 1: standard output
 - 2: standard error

Closing Files

Closing a file informs the kernel that you are finished accessing that file

```
int fd;    /* file descriptor */
int retval; /* return value */

if ((retval = close(fd)) < 0) {
    perror("close");
    exit(1);
}</pre>
```

- Closing an already closed file is a recipe for disaster in threaded programs (more on this later)
- Moral: Always check return codes, even for seemingly benign functions such as close()

Reading Files

 Reading a file copies bytes from the current file position to memory, and then updates file position

- Returns number of bytes read from file fd into buf
 - Return type ssize_t is signed integer
 - nbytes < 0 indicates that an error occurred
 - Short counts (nbytes < sizeof (buf)) are possible and are not errors!</p>

Writing Files

Writing a file copies bytes from memory to the current file position, and then updates current file position

```
char buf[512];
int fd;     /* file descriptor */
int nbytes;     /* number of bytes read */

/* Open the file fd ... */
/* Then write up to 512 bytes from buf to file fd */
if ((nbytes = write(fd, buf, sizeof(buf)) < 0) {
    perror("write");
    exit(1);
}</pre>
```

- Returns number of bytes written from buf to file fd
 - nbytes < 0 indicates that an error occurred</p>
 - As with reads, short counts are possible and are not errors!

Simple Unix I/O example

Copying standard in to standard out, one byte at a time

```
#include "csapp.h"

int main(void)
{
    char c;

    while(Read(STDIN_FILENO, &c, 1) != 0)
        Write(STDOUT_FILENO, &c, 1);
    exit(0);
}
```

Note the use of error handling wrappers for read and write (Appendix A).

Dealing with Short Counts

Short counts can occur in these situations:

- Encountering (end-of-file) EOF on reads
- Reading text lines from a terminal
- Reading and writing network sockets or Unix pipes

Short counts never occur in these situations:

- Reading from disk files (except for EOF)
- Writing to disk files

One way to deal with short counts in your code:

Use the RIO (Robust I/O) package from your textbook's csapp.c
 file (Appendix B)

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The RIO Package

- RIO is a set of wrappers that provide efficient and robust I/O in apps, such as network programs that are subject to short counts
- RIO provides two different kinds of functions
 - Unbuffered input and output of binary data
 - rio readn and rio writen
 - Buffered input of binary data and text lines
 - rio_readlineb and rio_readnb
 - Buffered RIO routines are thread-safe and can be interleaved arbitrarily on the same descriptor
- Download from http://csapp.cs.cmu.edu/public/code.html
 - → src/csapp.c and include/csapp.h

Unbuffered RIO Input and Output

- Same interface as Unix read and write
- Especially useful for transferring data on network sockets

```
#include "csapp.h"
ssize_t rio_readn(int fd, void *usrbuf, size_t n);
ssize_t rio_writen(int fd, void *usrbuf, size_t n);
Return: num. bytes transferred if OK, 0 on EOF (rio_readn only), -1 on error
```

- rio_readn returns short count only if it encounters EOF
 - Only use it when you know how many bytes to read
- rio_writen never returns a short count
- Calls to rio_readn and rio_writen can be interleaved arbitrarily on the same descriptor

Implementation of rio readn

```
* rio readn - robustly read n bytes (unbuffered)
 */
ssize t rio readn(int fd, void *usrbuf, size t n)
   size t nleft = n;
   ssize t nread;
   char *bufp = usrbuf;
   while (nleft > 0) {
       if ((nread = read(fd, bufp, nleft)) < 0) {</pre>
           if (errno == EINTR) /* interrupted by sig handler return */
              nread = 0;  /* and call read() again */
           else
              return -1; /* errno set by read() */
       else if (nread == 0)
                              /* EOF */
          break;
       nleft -= nread;
       bufp += nread;
   return (n - nleft); /* return >= 0 */
                                                              csapp
```

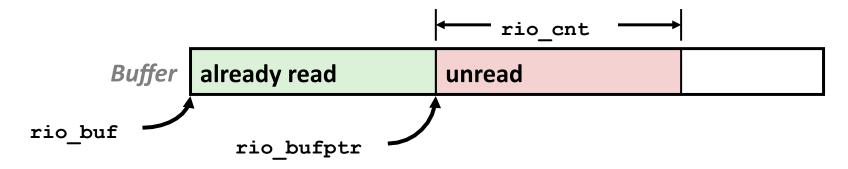
Buffered I/O: Motivation

- Applications often read/write one character at a time
 - getc, putc, ungetc
 - gets, fgets
 - Read line of text on character at a time, stopping at newline
- Implementing as Unix I/O calls expensive
 - read and write require Unix kernel calls
 - > 10,000 clock cycles
- Solution: Buffered read
 - Use Unix read to grab block of bytes
 - User input functions take one byte at a time from buffer
 - Refill buffer when empty

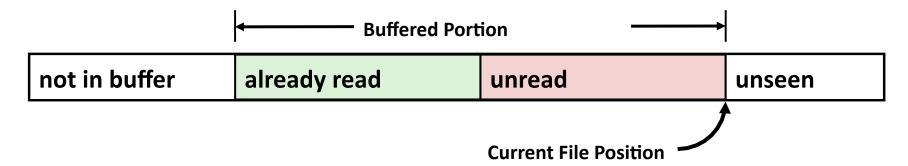
Buffer already read unread

Buffered I/O: Implementation

- For reading from file
- File has associated buffer to hold bytes that have been read from file but not yet read by user code

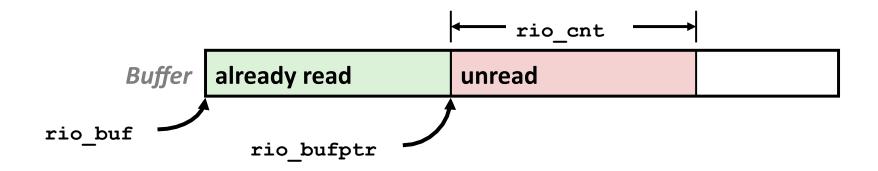


Layered on Unix file:



Buffered I/O: Declaration

All information contained in struct



Buffered RIO Input Functions

 Efficiently read text lines and binary data from a file partially cached in an internal memory buffer

```
#include "csapp.h"

void rio_readinitb(rio_t *rp, int fd);

ssize_t rio_readlineb(rio_t *rp, void *usrbuf, size_t maxlen);

Return: num. bytes read if OK, 0 on EOF, -1 on error
```

- rio_readlineb reads a text line of up to maxlen bytes from file
 fd and stores the line in usrbuf
 - Especially useful for reading text lines from network sockets
- Stopping conditions
 - maxlen bytes read
 - EOF encountered
 - Newline ('\n') encountered

Buffered RIO Input Functions (cont)

```
#include "csapp.h"

void rio_readinitb(rio_t *rp, int fd);

ssize_t rio_readlineb(rio_t *rp, void *usrbuf, size_t maxlen);
ssize_t rio_readnb(rio_t *rp, void *usrbuf, size_t n);

Return: num. bytes read if OK, 0 on EOF, -1 on error
```

- rio readnb reads up to n bytes from file fd
- Stopping conditions
 - maxlen bytes read
 - EOF encountered
- Calls to rio_readlineb and rio_readnb can be interleaved arbitrarily on the same descriptor
 - Warning: Don't interleave with calls to rio_readn

RIO Example

 Copying the lines of a text file from standard input to standard output

```
#include "csapp.h"
int main(int argc, char **argv)
{
   int n;
   rio_t rio;
   char buf[MAXLINE];

   Rio_readinitb(&rio, STDIN_FILENO);
   while((n = Rio_readlineb(&rio, buf, MAXLINE))) != 0)
        Rio_writen(STDOUT_FILENO, buf, n);
   exit(0);
}
```

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File Metadata

- *Metadata* is data about data, in this case file data
- Per-file metadata maintained by kernel
 - accessed by users with the stat and fstat functions

```
/* Metadata returned by the stat and fstat functions */
struct stat {
               st dev; /* device */
   dev t
               st ino; /* inode */
   ino t
               st mode; /* protection and file type */
   mode t
               st_nlink; /* number of hard links */
   nlink t
   uid t
               st uid; /* user ID of owner */
               st gid; /* group ID of owner */
   gid t
               st_rdev; /* device type (if inode device) */
   dev t
   off t st size; /* total size, in bytes */
   unsigned long st blksize; /* blocksize for filesystem I/O */
   unsigned long st blocks; /* number of blocks allocated */
               st atime; /* time of last access */
   time t
   time t st mtime; /* time of last modification */
               st ctime; /* time of last change */
   time t
};
```

Example of Accessing File Metadata

```
/* statcheck.c - Querying and manipulating a file's meta data */
#include "csapp.h"
                                       unix> ./statcheck statcheck.c
                                       type: regular, read: yes
int main (int argc, char **argv)
                                       unix> chmod 000 statcheck.c
                                       unix> ./statcheck statcheck.c
    struct stat stat;
                                       type: regular, read: no
    char *type, *readok;
                                       unix> ./statcheck ...
                                       type: directory, read: yes
    Stat(argv[1], &stat);
                                       unix> ./statcheck /dev/kmem
    if (S ISREG(stat.st mode))
       type = "regular";
                                       type: other, read: yes
    else if (S ISDIR(stat.st mode))
       type = "directory";
    else
       type = "other";
    if ((stat.st mode & S IRUSR)) /* OK to read?*/
       readok = "yes";
    else
       readok = "no";
   printf("type: %s, read: %s\n", type, readok);
    exit(0);
                                                      statcheck.c
```

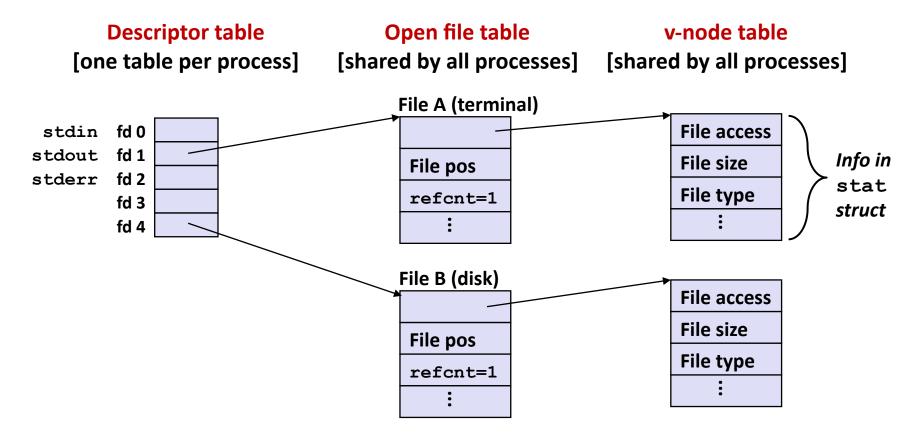
Accessing Directories

- Only recommended operation on a directory: read its entries
 - dirent structure contains information about a directory entry
 - DIR structure contains information about directory while stepping through its entries

```
#include <sys/types.h>
#include <dirent.h>
 DIR *directory;
 struct dirent *de;
 if (!(directory = opendir(dir name)))
     error("Failed to open directory");
 while (0 != (de = readdir(directory))) {
     printf("Found file: %s\n", de->d name);
 closedir(directory);
```

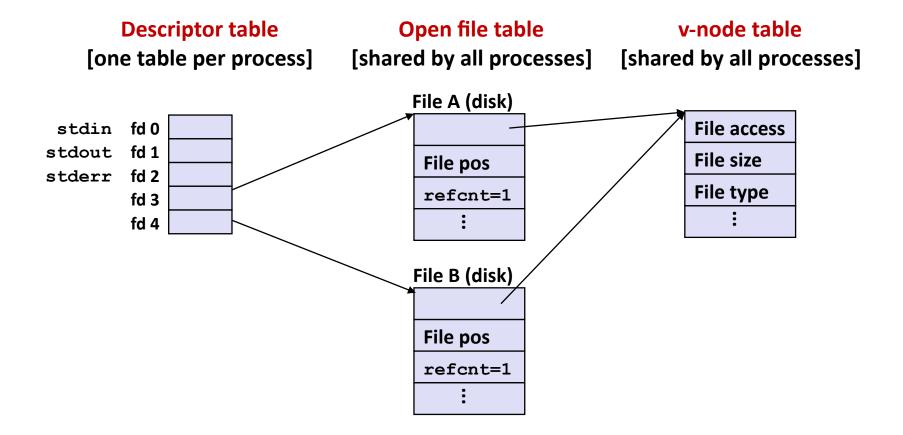
How the Unix Kernel Represents Open Files

Two descriptors referencing two distinct open disk files. Descriptor 1 (stdout) points to terminal, and descriptor 4 points to open disk file



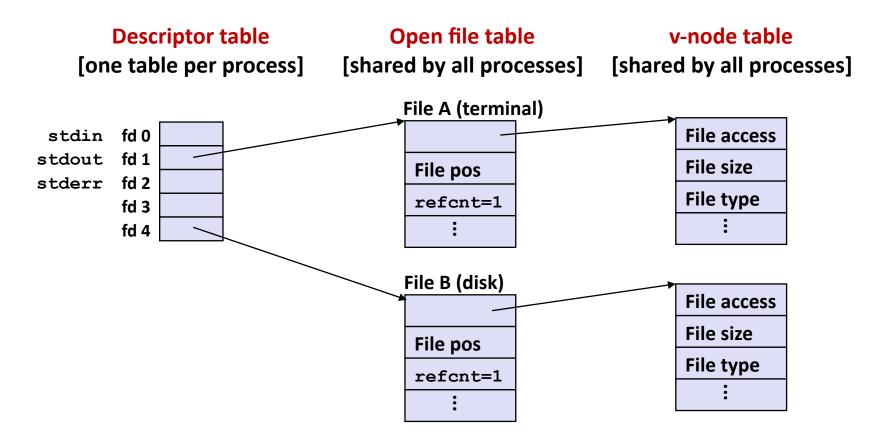
File Sharing

- Two distinct descriptors sharing the same disk file through two distinct open file table entries
 - E.g., Calling **open** twice with the same **filename** argument



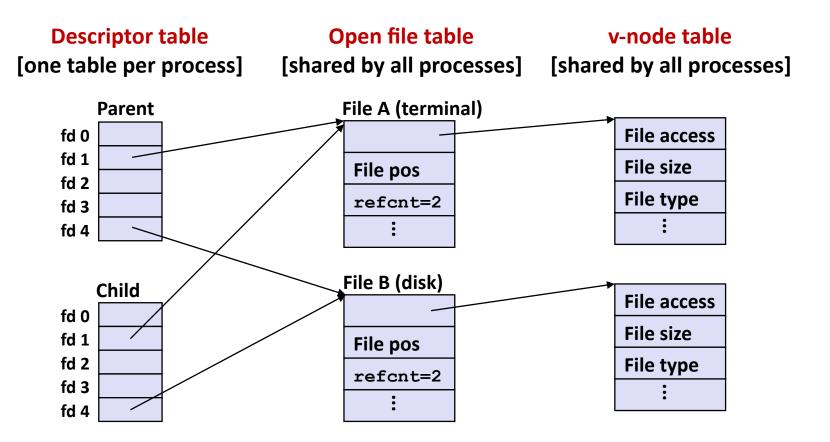
How Processes Share Files: Fork()

- A child process inherits its parent's open files
 - Note: situation unchanged by exec functions (use fcntl to change)
- Before fork() call:



How Processes Share Files: Fork()

- A child process inherits its parent's open files
- After fork():
 - Child's table same as parent's, and +1 to each refent



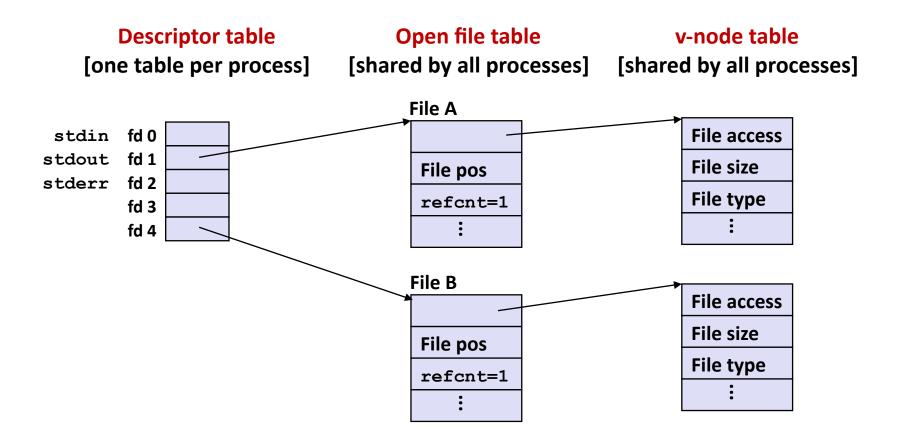
I/O Redirection

- Question: How does a shell implement I/O redirection?
 unix> ls > foo.txt
- Answer: By calling the dup2 (oldfd, newfd) function
 - Copies (per-process) descriptor table entry oldfd to entry newfd

Descriptor table before dup2 (4,1) fd 0 fd 1 fd 2 fd 3 fd 4 b Descriptor table after dup2 (4,1) fd 0 fd 1 fd 1 fd 2 fd 3 fd 4 b

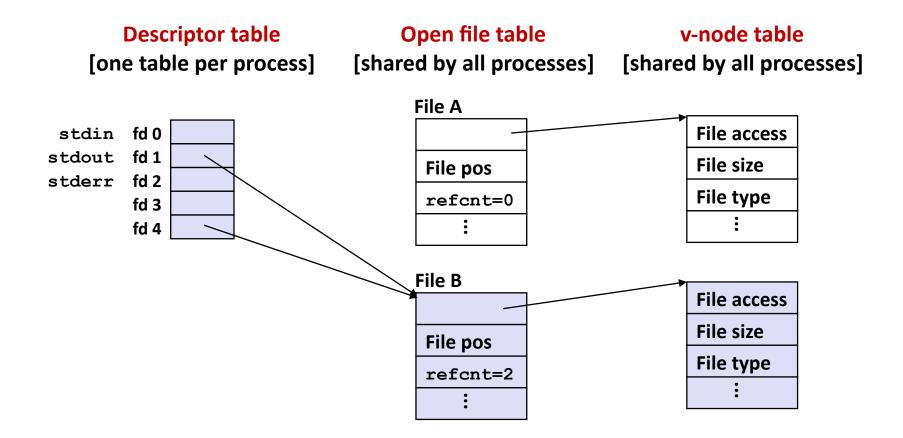
I/O Redirection Example

- Step #1: open file to which stdout should be redirected
 - Happens in child executing shell code, before exec



I/O Redirection Example (cont.)

- Step #2: call dup2 (4,1)
 - cause fd=1 (stdout) to refer to disk file pointed at by fd=4



Fun with File Descriptors (1)

```
#include "csapp.h"
int main(int argc, char *argv[])
    int fd1, fd2, fd3;
    char c1, c2, c3;
    char *fname = arqv[1];
    fd1 = Open(fname, O RDONLY, 0);
    fd2 = Open(fname, O RDONLY, 0);
    fd3 = Open(fname, O RDONLY, 0);
    Dup2 (fd2, fd3);
    Read(fd1, &c1, 1);
    Read(fd2, &c2, 1);
    Read(fd3, &c3, 1);
    printf("c1 = %c, c2 = %c, c3 = %c\n", c1, c2, c3);
    return 0;
                                              ffiles1.c
```

■ What would this program print for file containing "abcde"?

Fun with File Descriptors (2)

```
#include "csapp.h"
int main(int argc, char *argv[])
{
    int fd1;
    int s = qetpid() & 0x1;
    char c1, c2;
    char *fname = argv[1];
    fd1 = Open(fname, O RDONLY, 0);
   Read(fd1, &c1, 1);
    if (fork()) { /* Parent */
        sleep(s);
        Read(fd1, &c2, 1);
        printf("Parent: c1 = %c, c2 = %c\n", c1, c2);
    } else { /* Child */
        sleep(1-s);
        Read(fd1, &c2, 1);
        printf("Child: c1 = %c, c2 = %c\n", c1, c2);
    return 0;
                                            ffiles2.c
```

What would this program print for file containing "abcde"?

Fun with File Descriptors (3)

```
#include "csapp.h"
int main(int argc, char *argv[])
{
    int fd1, fd2, fd3;
    char *fname = argv[1];
    fd1 = Open(fname, O_CREAT|O_TRUNC|O_RDWR, S_IRUSR|S_IWUSR);
    Write(fd1, "pqrs", 4);
    fd3 = Open(fname, O_APPEND|O_WRONLY, 0);
    Write(fd3, "jklmn", 5);
    fd2 = dup(fd1); /* Allocates descriptor */
    Write(fd2, "wxyz", 4);
    Write(fd3, "ef", 2);
    return 0;
}
```

What would be the contents of the resulting file?

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Standard I/O Functions

- The C standard library (libc.so) contains a collection of higher-level standard I/O functions
 - Documented in Appendix B of K&R.
- Examples of standard I/O functions:
 - Opening and closing files (fopen and fclose)
 - Reading and writing bytes (fread and fwrite)
 - Reading and writing text lines (fgets and fputs)
 - Formatted reading and writing (fscanf and fprintf)

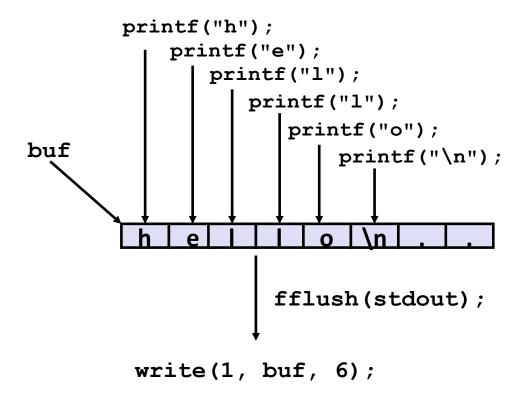
Standard I/O Streams

- Standard I/O models open files as streams
 - Abstraction for a file descriptor and a buffer in memory.
 - Similar to buffered RIO
- C programs begin life with three open streams (defined in stdio.h)
 - stdin (standard input)
 - stdout (standard output)
 - stderr (standard error)

```
#include <stdio.h>
extern FILE *stdin; /* standard input (descriptor 0) */
extern FILE *stdout; /* standard output (descriptor 1) */
extern FILE *stderr; /* standard error (descriptor 2) */
int main() {
   fprintf(stdout, "Hello, world\n");
}
```

Buffering in Standard I/O

Standard I/O functions use buffered I/O



■ Buffer flushed to output fd on "\n" or fflush() call

Standard I/O Buffering in Action

■ You can see this buffering in action for yourself, using the always fascinating Unix strace program:

```
#include <stdio.h>
int main()
{
    printf("h");
    printf("e");
    printf("l");
    printf("l");
    printf("o");
    printf("\n");
    fflush(stdout);
    exit(0);
}
```

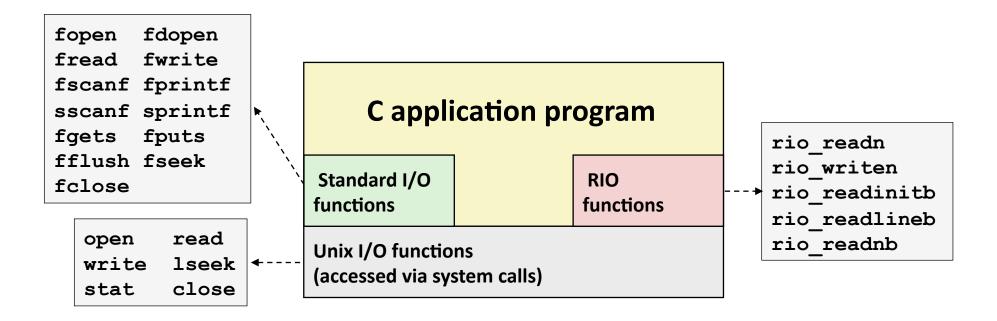
```
linux> strace ./hello
execve("./hello", ["hello"], [/* ... */]).
...
write(1, "hello\n", 6) = 6
...
exit_group(0) = ?
```

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Unix I/O vs. Standard I/O vs. RIO

Standard I/O and RIO are implemented using low-level Unix I/O



Which ones should you use in your programs?

Pros and Cons of Unix I/O

Pros

- Unix I/O is the most general and lowest overhead form of I/O.
 - All other I/O packages are implemented using Unix I/O functions.
- Unix I/O provides functions for accessing file metadata.
- Unix I/O functions are async-signal-safe and can be used safely in signal handlers.

Cons

- Dealing with short counts is tricky and error prone.
- Efficient reading of text lines requires some form of buffering, also tricky and error prone.
- Both of these issues are addressed by the standard I/O and RIO packages.

Pros and Cons of Standard I/O

Pros:

- Buffering increases efficiency by decreasing the number of read and write system calls
- Short counts are handled automatically

Cons:

- Provides no function for accessing file metadata
- Standard I/O functions are not async-signal-safe, and not appropriate for signal handlers.
- Standard I/O is not appropriate for input and output on network sockets
 - There are poorly documented restrictions on streams that interact badly with restrictions on sockets (CS:APP2e, Sec 10.9)

Choosing I/O Functions

- General rule: use the highest-level I/O functions you can
 - Many C programmers are able to do all of their work using the standard I/O functions

■ When to use standard I/O

When working with disk or terminal files

When to use raw Unix I/O

- Inside signal handlers, because Unix I/O is async-signal-safe.
- In rare cases when you need absolute highest performance.

When to use RIO

- When you are reading and writing network sockets.
- Avoid using standard I/O on sockets.

Aside: Working with Binary Files

- Binary File Examples
 - Object code, Images (JPEG, GIF),
- Functions you shouldn't use on binary files
 - Line-oriented I/O such as fgets, scanf, printf, rio_readlineb
 - Different systems interpret $0 \times 0 A$ (\n') (newline) differently:
 - Linux and Mac OS X: LF(0x0a) [\n']
 - HTTP servers & Windoes: CR+LF(0x0d 0x0a) ['\r\n']
 - Use rio_readn or rio_readnb instead
 - String functions
 - strlen, strcpy
 - Interprets byte value 0 (end of string) as special

For Further Information

The Unix bible:

- W. Richard Stevens & Stephen A. Rago, Advanced Programming in the Unix Environment, 2nd Edition, Addison Wesley, 2005
 - Updated from Stevens's 1993 classic text.

Stevens is arguably the best technical writer ever.

- Produced authoritative works in:
 - Unix programming
 - TCP/IP (the protocol that makes the Internet work)
 - Unix network programming
 - Unix IPC programming

■ Tragically, Stevens died Sept. 1, 1999

But others have taken up his legacy