

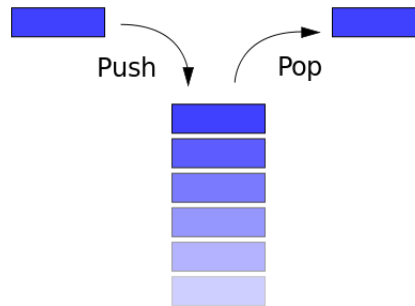
# 15-122: Principles of Imperative Computation

## Recitation 9 Solutions

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### Stacks and Queues

Stacks are a LIFO (last in, first out) data structure. This is just like ordinary references to stack e.g. a stack of books.

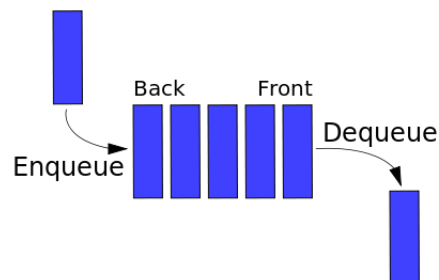


(Picture courtesy of Vegpuff/Wikipedia)

Here is the interface for stacks, as discussed in lecture:

```
1 stack stack_new();
2 bool stack_empty(stack S);
3 void push(stack S, string x);
4 string pop(stack S)
5 /*@requires !stack_empty(S); @*/ ;
```

Queues are a FIFO (first in, first out) data structure. This is just like ordinary references to a queue e.g. the queue (or line) at package pickup in the UC.



(Picture courtesy of Wikipedia)

Here is the interface for queues, as discussed in lecture:

```
1 queue queue_new();
2 bool queue_empty(queue Q);
3 void enq(queue Q, string x);
4 string deq(queue Q)
5 /*@requires !queue_empty(Q); @*/ ;
```

## Checkpoint 0

(NOTE: THIS QUESTION WAS PARTIALLY INCORRECT IN THE HANDOUT; see the revised Checkpoint 1 below for details) Write a function to reverse a queue, using only the functions from the interface.

Below is the general structure of this function. You may not need to fill in all the blanks.

```
1 /* Assume that you have data types stack and queue as described in lecture */
2 void reverse(queue Q)
3 {
4   stack S = stack_new(); //Hint : Allocate a temporary data structure
5   while(!queue_empty(Q))
6   {
7     string temp = deq(Q);
8     push(S, temp);
9   }
10  while(!stack_empty(S))
11  {
12    string temp = pop(S);
13    enq(Q, temp);
14  }
15 }
```

## Checkpoint 1

(NOTE: THIS QUESTION WAS INCORRECT IN THE HANDOUT) Why did we NOT need contracts in the function above?

*Solution:* Since we are clients using a given stack/queue interface, we can assume that, as long as we respect the interface, all stacks/queues are valid and that any operations involving them (which are allowed by the interface) are safe.

## Checkpoint 2

Write a function to calculate the size of a stack without allocating any new data structures. At the end of the function, the stack must be unmodified (Hint : You will need to modify the stack within the function, but must ensure it is the same at the end)

*Solution:*

```
1 //recursive without allocating data structures
2 int stack_size(stack S)
3 //@requires is_stack(S);
4 //@ensures
5 {
6   if (stack_empty(S)) return 0;
7   string x = pop(S);
8   int i = 1 + stack_size(S);
9   push(S, x);
10  return i;
11 }
```

## Clac

*clac* is a relatively simple *postfix*-based programming language. As we read in numbers from the input (which we represent as a queue), we push operands onto a stack and act on them based on the instructions that are in the queue.

Here's an example of *clac* processing some input (you can get this yourself when working on the *clac* assignment by running *clac-ref*).

```
$ clac-ref -trace
Clac top level
clac>> 5 9 2 7 3 + - / dup * %

          stack || queue
          || 5 9 2 7 3 + - / dup * %
          5 || 9 2 7 3 + - / dup * %
          5 9 || 2 7 3 + - / dup * %
          5 9 2 || 7 3 + - / dup * %
          5 9 2 7 || 3 + - / dup * %
          5 9 2 7 3 || + - / dup * %
          5 9 2 10 || - / dup * %
          5 9 -8 || / dup * %
          5 -1 || dup * %
          5 -1 -1 || * %
          5 1 || %
          0 ||

0
```

What's happening here? Well, we push all of the numbers onto the stack after reading them out of the queue. Then, we get to the `+`, so we pop two items (the 7 and the 3) off of the stack, add them, and push their sum, 10, back on. Next, we get to the `-`, pop off the 2 and 10 and subtract them, and get -8, which we push on to the stack. Then, we get to the `/`. We pop 9 and -8 and divide them.  $9/-8$  rounds to -1, so we push that onto the stack. Next, we execute the `dup`, which simply makes the top element of the stack appear twice. We get to the `*`, which multiplies the top two elements, giving us 1. Finally, we get to the `%`.  $5 \% 1 == 0$ , so we push 0. Then, we're out of instructions, so we end and pop the top item off of the stack and print it.

A common source of confusion with *clac* is `if` statements and `else` statements.

When we get to an `if` statement, we pop the top item off of the stack. If it is 0, we skip the next *two* tokens in the queue – we just ignore them. Otherwise (if it's non-zero), we continue processing tokens as normal.

When we get to an `else` statement, we *always* skip the next token in the queue.

So, why are these `if/else` statements? Let's take a look at some *clac* code

NOTE: In *clac* code below, we're using `x` it to mean any arbitrary int – you should fill in an int, like 1, -1, 0, etc, if you're actually running the code.

```
$ clac-ref -trace
Clac top level
clac>> 0 if 2 else 3
```

```
stack || queue
      || 0 if 2 else 3
0 || if 2 else 3
  || 3
  3 ||
```

```
3
clac>> 1 if 2 else 3
```

```
stack || queue
      3 || 1 if 2 else 3
3 1 || if 2 else 3
   3 || 2 else 3
   3 2 || else 3
   3 2 ||
```

```
2
```

Next, let's write a simple clac program: one that calculates absolute value. We can define  $|x|$  as follows:

$$|x| = \begin{cases} x * 1 & \text{if } x \geq 0 \\ x * -1 & \text{if } x < 0 \end{cases}$$

So, if  $x$  is less than 0, we want to multiply it by  $-1$  and otherwise we want to multiply it by 1. If we run the clac command `x 0 <`, then it will result in 1 being on the top of the stack if  $x < 0$  and 0 being on the top of the stack otherwise.

We eventually want to multiply by either 1 or  $-1$ , so we should push the appropriate one of them onto the stack: If  $x < 0$  we multiply by  $-1$ , otherwise we multiply by 1.

So, we add `if -1 else 1` to our command. Now we have

```
x 0 < if -1 else 1
```

This says "if  $x < 0$ , push  $-1$  onto the stack. Otherwise, push 1 onto the stack." This works because when `x 0 <` evaluates to 0 (so  $x \geq 0$ ), we ignore the tokens `-1` and `else`, so we just push 1 onto the stack. If `x 0 <` evaluates to 1 (so  $x < 0$ ), then we push  $-1$  onto the stack and ignore the token 1.

Next, we want to multiply by  $x$ , so we add `*` to the end:

```
x 0 < if -1 else 1 *
```

This doesn't work, though! We popped  $x$  off of the stack when we did the comparison. If we run the above command, we get:

```
Error: Error: not enough elements on stack
```

So, we need to duplicate  $x$  before we compare, so we can still use it later:

```
x dup 0 < if -1 else 1 *
```

That will compute the absolute value of  $x$ .