





### MOTIVATION

- Software most complex part of today's safety critical embedded systems
- Most embedded systems are <u>legacy designs</u>
- Written in a low level language such as **ANSI-C** or even assembly language
- **Bigger problem than new, high level designs**
- Existing tools do not address the verification problem

SCREENSHOT

- Goal: Verify legacy code with respect to
  - A formal specification
  - A high level design
  - Safety properties

 $\frac{1}{2}$ 

- 'Most embedded systems are legacy designs,
- *i.e., written in a low level language such as*
- ANSI-C or even assembly language. These
- systems are a bigger problem than the new,
- high level designs, but the existing tools do not
  - address the verification problem."

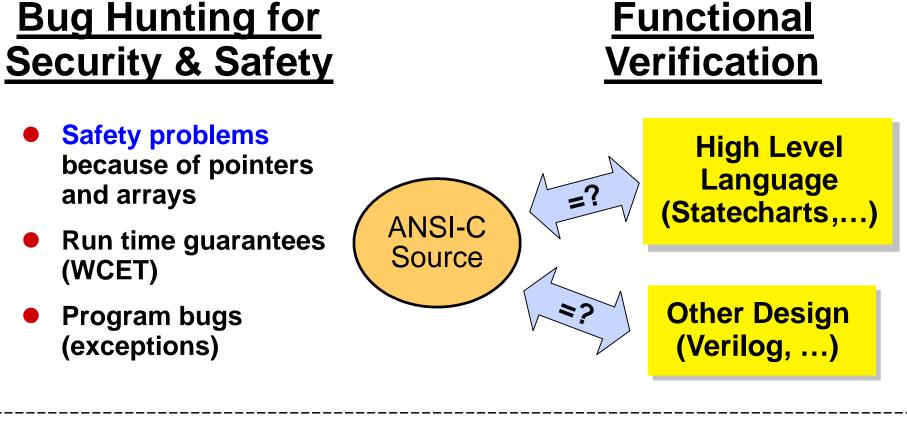
# TOOL OVERVIEW

## ANSI-C BMC

- **Problem: Fixpoint computation is too expensive for** software
- Idea:
  - Unwind program into equation
  - Check equation using SAT
- Advantages:
  - Completely automated
  - Allows full set of ANSI-C, including full treatment of pointers and dynamic memory
- **Properties:** 
  - Simple assertions
  - Security (Pointers/Arrays)
  - Run time guarantees (WECT)



CBMC-GUI			
File Tools Options Info			
<pre>Elle Tools Options Into test_pointer.c fptr4.c voru ('p)(); int input; p=&amp;b if(input) p=&amp;a (*p)(); assert(g==2);</pre>			
Output Errors Watch Debug			
Name	Value		
g	0 (000000000000000000000000000000000000	000000000000000000000000000000000000000	
main::input main::p	-21474 5a CBMC-GUI		
initial of the second s	Eile Tools Options		
fotr	r4.log test_pointer.c fptr4.c	test_pointer.c fptr4.c loop2.c bound.c	
	for (k =		
	Output Errors Wat	ch Debug	
	Name	Value	
	main::x	{ 1, 0, 0, 0, 0, 0, 0, 0, 255, 255 }	
	main::sum	256 (000000000000000000000000000000000000	
	main::k	10 (00000000000000000000000000000000000	
		bound.log 21 of 22 steps	



CBMC

### Motivation:

• assertions often not expressive enough

• E.g.: complex computations

• One wants specification that can be inspected  $\Rightarrow$  We use PVS language

- Problems:
  - Basically everything about PVS language is undecidable, including type consistency
  - PVS language highly compact due to overloading
  - Requires complex resolver and type checker
- Good news: both resolver and type checker implemented (first time outside of PVS!)

# DONE

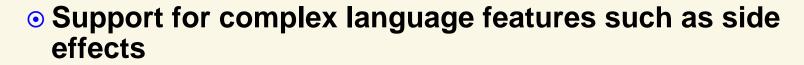
- Implemented tool that automatically detects
  - Buffer overruns
  - Pointer bugs
  - Worst Case Execution time
  - <u>No false positives, no false negatives!</u>
- Tool takes ANSI-C as input • Support for <u>all</u> integer operators

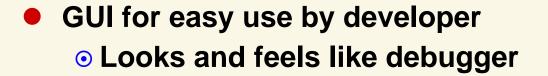
# **CURRENT PROJECT**

- Verify Safety Properties of a part of a train controller provided by GE
  - Termination / WCET
  - Correctness of pointer constructs
  - The code uses two channels for redundancy: Check that they compute the same result

## FUTURE WORK

- Interval abstraction for floating point aritmetic
- Concurrent ANSI-C programs (SpecC)
- Object oriented languages (C++, Java)
- Statechart-like specification language









The code contains x86 assembly language