

CMU - SCS  
15-415/15-615 Database Applications  
Spring 2013, C. Faloutsos  
Homework 6: Query Optimization + Schema  
Refinement  
**Deadline: 1:30pm on Tuesday, 3/26/2013**

**Reminders - IMPORTANT:**

- Like all homeworks, it has to be done **individually**.
- Please **typeset** your answers.
- Please submit your answers in **hard copy, in class**, 1:30pm, on Tuesday, 03/26/2013 .
- As before, for ease of grading, please solve each of the three questions on a **separate** page, and type your **name and andrew ID on each** of the three pages.

**Reminders - FYI:**

- Weight: 5% of homework grade.
- The points of this homework add up to 100.
- Rough time estimates: 2-4 hours.

## Q1. Query Optimization - Selectivities, 30 pts - SUBMIT ON SEPARATE PAGE

Consider the tables `WORKS_AT(SSN, gymID)` and `GYM(gymID, name)`. Notice that `gymID` is not a candidate key for the table `GYM`.

`WORKS_AT(SSN, gymID)` consists of  $N_1 = 100,000$  tuples and has

- $V(\text{SSN}, \text{WORKS\_AT}) = 50,000$  distinct values of SSN
- $V(\text{gymID}, \text{WORKS\_AT}) = 20,000$  distinct values of gymID.

`GYM(gymID, name)` consists of  $N_1 = 40,000$  tuples and has

- $V(\text{gymID}, \text{GYM}) = 20,000$  distinct values of gymID
- $V(\text{name}, \text{GYM}) = 30,000$  distinct values of name.

For all the computations below:

- Please give numerical answers, accurate up to the **fourth** significant digit.
- No need for explanations, unless explicitly requested.

Q1.1 [5 pts] Estimate the number of qualifying tuples of the query:

```
SELECT *
FROM WORKS_AT
WHERE SSN = 123456789;
```

Q1.2 [5 pts] Can SSN be a candidate key for the table `WORKS_AT`? Give a short explanation for your answer.

Q1.3 [5 pts] Estimate the number of qualifying tuples of the query:

```
SELECT *
FROM GYM
WHERE name = "Gym_planet";
```

Q1.4 [5 points] Estimate the number of qualifying tuples of the query:

```
SELECT *
FROM WORKS_AT
WHERE SSN = 123456789 AND gymID=101;
```

Q1.5 [5 points] Notice that `gymID` is *not* a candidate key for the table `GYM`. Estimate the number of qualifying tuples of the query:

```
SELECT SSN, GYM.gymID, name
FROM WORKS_AT JOIN GYM
WHERE GYM.gymID = WORKS_AT.gymID;
```

Q1.6 [5 points] Estimate the number of qualifying tuples of the query:

```
SELECT WA1.SSN, WA2.SSN
FROM WORKS_AT AS WA1 JOIN WORKS_AT AS WA2
WHERE WA1.gymID = WA2.gymID;
```

## Q2. Functional Dependencies, 30pts - SUBMIT ON SEPARATE PAGE

**Q2.1** Consider the relation schema  $R = \{P, Q, S, T, U, V\}$  and the set of functional dependencies  $FD =$

$$PQ \rightarrow S \quad (1)$$

$$PS \rightarrow Q \quad (2)$$

$$PT \rightarrow U \quad (3)$$

$$Q \rightarrow T \quad (4)$$

$$QS \rightarrow P \quad (5)$$

$$U \rightarrow V \quad (6)$$

Answer the following questions. Notice that:

- For Yes/No or True/False questions, you may just give binary answers. Explanations are optional and will be used for partial credit. Wrong answers, with no, or wrong explanations, will get **negative** points.
- For the rest of the questions, please give *short* justifications.

2.1.1 [2 pts] Yes/No. Is  $FD$  a minimum cover?

2.1.1 [4 pts] Yes/No. Is the decomposition  $\{PQ, QS, PQTU, UV\}$  lossless?

2.1.2 [4 pts] Somebody claims that the decomposition  $\{PQ, QS, PQTU, UV\}$  is not dependency-preserving. If you agree with the statement, give all the missing dependencies. If you disagree, just state so.

2.1.3 [5 pts] Yes/No. Is the decomposition  $\{PQS, PSTU, PTV\}$  lossless?

2.1.4 [1 pts] True/False. The decomposition  $\{PQS, PSTU, PTV\}$  is not dependency-preserving, because it does not preserve  $U \rightarrow V$ .

2.1.5 [2 pts] True/False. The decomposition  $\{PQS, PSTU, PTV\}$  is not dependency-preserving, because it does not preserve  $U \rightarrow V$  nor  $Q \rightarrow T$ .

2.1.6 [2 pts] True/False. The decomposition  $\{PQS, PSTU, PTV\}$  is dependency-preserving.

**Q2.2** Consider now the same relation schema  $R = \{P, Q, S, T, U, V\}$  with **different**, simpler, set of functional dependencies  $FD' =$

$$Q \rightarrow ST \quad (7)$$

$$P \rightarrow T \quad (8)$$

$$PS \rightarrow T \quad (9)$$

$$QU \rightarrow V \quad (10)$$

Answer the following questions. Again, **negative** points for wrong, binary answers.

2.2.1 [1 pts] True/False. The attribute closure  $\{P\}^+$  is  $\{P, S, T\}$ .

2.2.2 [1 pts] True/False. The attribute closure  $\{P\}^+$  is  $\{P, T\}$ .

- 2.2.3 [1 pts] True/False. The attribute closure  $\{P,Q\}^+$  is  $\{P,T,Q,S\}$ .
- 2.2.4 [1 pts] True/False. The attribute closure  $\{P,Q\}^+$  is  $\{P,S,T\}$ .
- 2.2.4 [1 pts] True/False. The attribute closure  $\{P,Q\}^+$  is  $\{P,T,Q,S,U,V\}$ .
- 2.2.5 [1 pts] True/False. The dependency  $Q \rightarrow S$  can be deduced from  $FD'$ .
- 2.2.6 [2 pts] True/False. The dependency  $QU \rightarrow TUV$  can be deduced from  $FD'$ .
- 2.2.7 [2 pts] True/False. All the candidate keys of  $R$  are  $\{P,Q\}$ .

### Q3. BCNF and 3NF, 40pts - SUBMIT ON SEPARATE PAGE

Consider the relation schema  $R = \{P, Q, S, T, U, V\}$  and the functional dependencies  $FD =$

$$PQ \rightarrow S \quad (11)$$

$$PS \rightarrow Q \quad (12)$$

$$PT \rightarrow U \quad (13)$$

$$Q \rightarrow T \quad (14)$$

$$QS \rightarrow P \quad (15)$$

$$U \rightarrow V \quad (16)$$

Consider also the relation schemas

- $R1 = \{P, Q, S\}$
- $R2 = \{P, Q, S, U, V\}$  and
- $R3 = \{P, Q, S, T\}$

As before, **negative points** for wrong, binary answers; explanations are optional, unless explicitly requested.

Q3.1 [2 pts] Write the projection of the FDs on  $R1$ .

Q3.2 [2 pts] True/False. The set of dependencies  $FD$  given above (11-16) is a minimal cover.

Q3.3 [4 pts] True/False.  $R1$  is in 3NF.

Q3.4 [4 pts] True/False.  $R1$  is in BCNF.

Q3.5 [2 pts]. Write the projection of the FDs on  $R2$ .

Q3.6 [4 pts] True/False. All the candidate keys of  $R2$  are  $\{PQU, QSU\}$ .

Q3.7 [4 pts]. True/False.  $R2$  is in BCNF.

Q3.8 [4 pts] True/False. Consider the decomposition of  $R2$   $\{PQU, PQS, UV\}$ . The new relations are in BCNF.

Q3.9 [2 pts]. Write the projection of the FDs on  $R3$ .

Q3.10 [2 pts] True/False. The candidate keys of  $R3$  are  $\{PQ, QS, PS\}$ .

Q3.11 [4 pts]  $R3$  is *not* in BCNF. Give all the dependencies of  $FD$  that violate the BCNF.

Q3.12 [2 pts] True/False.  $R3$  is in 1NF.

Q3.13 [4 pts] True/False. Consider the decomposition of  $R3$  to  $\{PQS, QT\}$ . The new relations are in BCNF.

\_\_\_\_\_ End of homework questions \_\_\_\_\_