
$\qquad$
$\qquad$


## Outline

Goal: 'Find similar / interesting things'

- Intro to DB $\qquad$
- Indexing - similarity search
- Data Mining


## Indexing - Detailed outline

- primary key indexing
- secondary key / multi-key indexing $\qquad$
spatial access methods
- problem dfn $\qquad$
- z -ordering
- R-trees
- ...
- text
${ }_{15} \cdot . .226$


## Spatial Access Methods problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer spatial queries (like??)


Spatial Access Methods problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
- point queries
- range queries
- k -nn queries
- spatial joins ('all pairs' queries)



## Spatial Access Methods problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
- point queries
- range queries
- k-nn queries
- spatial joins ('all pairs' queries)

$\qquad$
$\qquad$


## Spatial Access Methods problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
- point queries
- range queries
- k -nn queries
- spatial joins ('all pairs' queries)



## Spatial Access Methods problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
- point queries
- range queries
- k-nn queries
- spatial joins ('all pairs' queries)



## Spatial Access Methods problem

- Given a collection of geometric objects (points, lines, polygons, ...)
- organize them on disk, to answer
- point queries
- range queries
- k -nn queries
- spatial joins ('all pairs' within $\varepsilon$ )

$\qquad$
$\qquad$
$\qquad$
$\qquad$




## Indexing - Detailed outline

- primary key indexing
- secondary key / multi-key indexing
- spatial access methods
- problem dfn $\qquad$
- z-ordering
- R-trees
- text
$\stackrel{\text { • }}{15-826}$
Copyright: C. Faloutsos (2007)


## SAMs: solutions

- z-ordering
- R-trees
- (grid files)

Q: how would you organize, e.g., $n$-dim $\qquad$ points, on disk? ( $C$ points per disk page)

$\qquad$
$\qquad$
15-826

## z-ordering

Q: how would you organize, e.g., n-dim points, on disk? ( $C$ points per disk page) $\qquad$
Hint: reduce the problem to 1-d points(!!)
Q1: why? $\qquad$
A:
Q2: how?


$\qquad$
$\qquad$
$\qquad$

## z-ordering

Q2: how?
A: assume finite granularity (e.g., $2^{32} \times 2^{32}$; $\qquad$ $4 \times 4$ here)
Q2.1: how to map n-d cells to 1-d cells? $\qquad$

Q2.1: how to map $n$-d cells to 1-d cells?


Copyright: C. Faloutsos (2007)
$\qquad$
$\qquad$

$\qquad$







## z-ordering

Drill: z-value of magenta cell, with the three methods? $\qquad$

$\qquad$
$\qquad$
$\qquad$
15-826


Drill: z-value of magenta cell, with the three methods?



$\qquad$
$\qquad$
$\qquad$
$\qquad$

## $3^{\text {z.ordering - usage } \& \text { algo's }}$

Q1: How to store on disk?
A:
Q2: How to answer range queries etc


15-826

```
z-ordering - usage & algo's
Q1: How to store on disk?
A: treat z-value as primary key; feed to B-tree PGH
```

SF

$\qquad$
$\qquad$
$\qquad$
$\qquad$

$\qquad$



$\qquad$
$\qquad$
$\qquad$

```
CMUSCS
```


## z-ordering - usage \& algo's

Q2', range queries - how to break a query $\qquad$ into ranges?


9,11-15

## z-ordering - usage \& algo's

Q2'’: range queries - how to break a query $\qquad$ into ranges?
A: recursively, quadtree-style; decompose $\qquad$ only non-full quadrants

$\qquad$
$\qquad$
$\qquad$

$\qquad$
$\qquad$
$\qquad$

## cnuscs <br> z-ordering - Detailed outline

- spatial access methods
- z -ordering
- main idea - 3 methods
- use w/ B-trees; algorithms (range, knn queries ...)
- non-point (eg., region) data $\qquad$
- analysis; variations
- R-trees $\qquad$
- ...




## ${ }^{\text {cnuscs }}$ <br> z-ordering - Detailed outline

- spatial access methods
- z -ordering
- main idea - 3 methods
- use w/ B-trees; algorithms (range, knn queries ...)
$\Rightarrow$ - non-point (eg., region) data $\qquad$
- analysis; variations
- R-trees $\qquad$
- ...

15-826

## z-ordering - regions

$\mathrm{Q}: \mathrm{z}$-value for a region?



$\qquad$
$\qquad$
$\qquad$
$\qquad$
$\qquad$


## z-ordering - regions

Q: How to search (range etc queries) - eg 'red' range query
A: break query in z -values; check B -tree


## cmuscs

## z-ordering - regions

Almost identical to range queries for point
$\qquad$ data, except for the "don't cares" - i.e.,

$\qquad$
$\qquad$
$\qquad$
$\qquad$

## z-ordering - regions

Almost identical to range queries for point
$\qquad$ data, except for the "don't cares" - i.e., $\mathrm{z} 1=1100$ ?? $11^{* *}=\mathrm{z} 2$ $\qquad$
Specifically: does z1 contain/avoid/intersect z2? $\qquad$
Q : what is the criterion to decide?

## z-ordering - regions

$$
\mathrm{z} 1=1100 \text { ?? } 11^{* *}=\mathrm{z} 2
$$

Specifically: does z1 contain/avoid/intersect z2?
Q: what is the criterion to decide?
A: Prefix property: let r 1 , r 2 be the corresponding regions, and let rl be the smallest (=> z1 has fewest '*'s). Then:

15-826
Copyright: C. Faloutsos (2007)

## z-ordering - regions

- r2 will either contain completely, or avoid completely r1.
- it will contain r1, if $z 2$ is the prefix of $z 1$



## z-ordering - regions

Drill (True/False). Given:

- z1=011001** $\qquad$
- z2=01******
- z3 $=0100^{* * * *}$ $\qquad$
T/F r2 contains r1
T/F r3 contains r1
T/F r3 contains r2 $\qquad$


## z-ordering - regions

Drill (True/False). Given:

- z1=011001**
- $\mathrm{z} 2=01 * * * * * *$
- $\mathrm{z} 3=0100^{* * * *}$

T/F r2 contains r1-TRUE (prefix property) T/F r3 contains r1-FALSE (disjoint)
T/F r3 contains r2-FALSE (r2 contains r3)

## $7^{\text {z-ordering - regions }}$

Drill (True/False). Given:

- z1=011001**
- z2=01******
- z3= 0100****







$\qquad$
$\qquad$
$\qquad$


Z-ordering - variations
Q: function for the Hilbert curve $(h=f(x, y))$ ?
A: bit-shuffling, followed by post-processing,
to account for rotations. Linear on \# bits.

| See textbook, for pointers to |
| :--- |
| code/algorithms (eg., [Jagadish, 90]) |

${ }^{15.826}$ Copyright c. Faloustos (2007)


- spatial access methods
- z-ordering
- main idea - 3 methods
- use w/ B-trees; algorithms (range, knn queries ...)
- non-point (eg., region) data $\qquad$
- analysis; variations
- R-trees $\qquad$
$\qquad$
15-826


## z-ordering - analysis

Q: How many pieces ('quad-tree blocks') per region? $\qquad$
A: proportional to perimeter (surface etc)


$\qquad$
$\qquad$
$\qquad$


Z-ordering - analysis
Q: So, is Hilbert really better?
A: 27\% fewer runs, for 2-d (similar for 3-d)
Q: are there formulas for \#runs, \#of quadtree
blocks etc?
A: Yes ([Jagadish; Moon+ etc] see textbook)
${ }^{15.826}$

## z-ordering - fun observations

Hilbert and z-ordering curves: "space filling curves": eventually, they visit every point in n -d space - therefore:

order-1
15-826

order-2
$\ldots$ order $(\mathrm{n}+1)$
Copyright: C. Faloutsos (2007)


## $\int_{\text {z-ordering - fun observations }}$

... they show that the plane has as many points as a line (-> headaches for 1900's mathematics/topology). (fractals, again!)


15-826
Copyright: C. Faloutsos (2007)

## z-ordering - fun observations

Observation \#2: Hilbert (like) curve for video encoding [Y. Matias+, CRYPTO ‘87]: $\qquad$
Given a frame, visit its pixels in randomized hilbert order; compress; and transmit



```
8 cmuscs
```


## Indexing - Detailed outline

```
- primary key indexing
- secondary key / multi-key indexing
- spatial access methods
- problem dfn
- z -ordering
- R-trees
- text
15-826

\section*{Conclusions}
- z-ordering is a great idea ( n -d points -> 1-d points; feed to B-trees) \(\qquad\)
- used by TIGER system and (most probably) by other GIS products \(\qquad\)
- works great with low-dim points```

