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| Optional Material |
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| Optional, but very useful: Manfred Schroeder |
| Fractals, Chaos, Power Laws: Minutes <br> from an Infinite Paradise W.H. Freeman <br> and Company, 1991 |
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- forecasting [CIKM’02]
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$\qquad$ M-tree?
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$\qquad$
Metric trees - analysis
- Problem: How many disk accesses, for an
M-tree?
- Given:
- N (\# of objects)
- C (fanout of disk pages)
- r (radius of range query - BIASED model)
- NOT ENOUGH - what else do we need?
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- A: something about the distribution
[Ciaccia, Patella, Zezula, PODS98]: assumed that the distance distribution is the same, for $\qquad$ every object:

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- A: something about the distribution
- Given our 'fractal' tools, we could try them which one?
- A: Correlation integral [Traina+, ICDE2000]

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- So, what is the \# of disk accesses, for a node of radius $\mathrm{r}_{\mathrm{d}}$, on a query of radius $\mathrm{r}_{\mathrm{q}}$ ?
- $A: \sim\left(r_{d}+r_{q}\right) \ldots$ $\qquad$
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- So, what is the \# of disk accesses, for a node
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$\qquad$ of radius $\mathrm{r}_{\mathrm{d}}$, on a query of radius $\mathrm{r}_{\mathrm{q}}$ ?
- A: $\sim\left(\mathrm{r}_{\mathrm{d}}+\mathrm{r}_{\mathrm{q}}\right)^{\wedge} \boldsymbol{D}$ $\qquad$
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- Normally, D takes O( $\mathrm{N}^{\wedge} 2$ ) time
- Anything faster? suppose we have already built an M-tree
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## $\int^{3 \text { Index }}$ Indexing - Detailed outline

- fractals
- intro
- applications
- disk accesses for R-trees (range queries) $\qquad$
- dimensionality reduction
- selectivity in M-trees
- dim. curse revisited
. "fat fractals"
- quad-tree analysis [Gaede+] $\qquad$
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## (Overview of proofs)

- assume that your points are uniformly distributed in a $d$-dimensional manifold (= hyper-plane)
- derive the formulas
- substitute $d$ for the fractal dimension
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## Reminder: Hausdorff Dimension ( $D_{0}$ )

- $r=$ side length (each dimension)
- $B(r)=\#$ boxes containing points $\propto r^{D 0}$


$$
\log r=-1
$$



$$
\log _{26} B=1
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$$
\log B=2
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- How to determine avg MBR side $l$ ?
$-N=\#$ pts, $C=$ MBR capacity


Hausdorff dimension: $B(r) \propto r^{D 0}$

$$
B(l)=N / C=l^{-\mathrm{D} 0} \Rightarrow l=(N / C)^{-1 / \mathrm{D} 0}
$$

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- For $k$ pts, what radius $\varepsilon$ do we expect?

Correlation dimension: $S(r) \propto r^{D 2}$
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$$
\underset{\operatorname{ses}(2011)}{S(\varepsilon)}=\frac{k}{N-1}=(2 \varepsilon)^{D 2}
$$

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## Asymptotic Formula $P_{\text {all }}^{L \infty}(k) \approx \sum_{j=0}^{h}\left\{\frac{1}{C^{h-j}}+\left[1+\left(\frac{k}{C^{h-j}}\right)^{1 / D}\right]^{D}\right\}$

- NO mention of the embedding dimensionality!!
- Still have dim. curse, but on f.d. $D$

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| Conclusions |  |
| :---: | :---: |
|  | Worst-case theory is over-pessimistic |
|  | High dimensional data can exhibit good performance if correlated, non-uniform |
|  | Many real data sets are self-similar |
|  | Determinant is intrinsic dimensionality <br> - multiple fractal dimensions ( $D_{0}$ and $D_{2}$ ) <br> - indication of how far one can go |
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